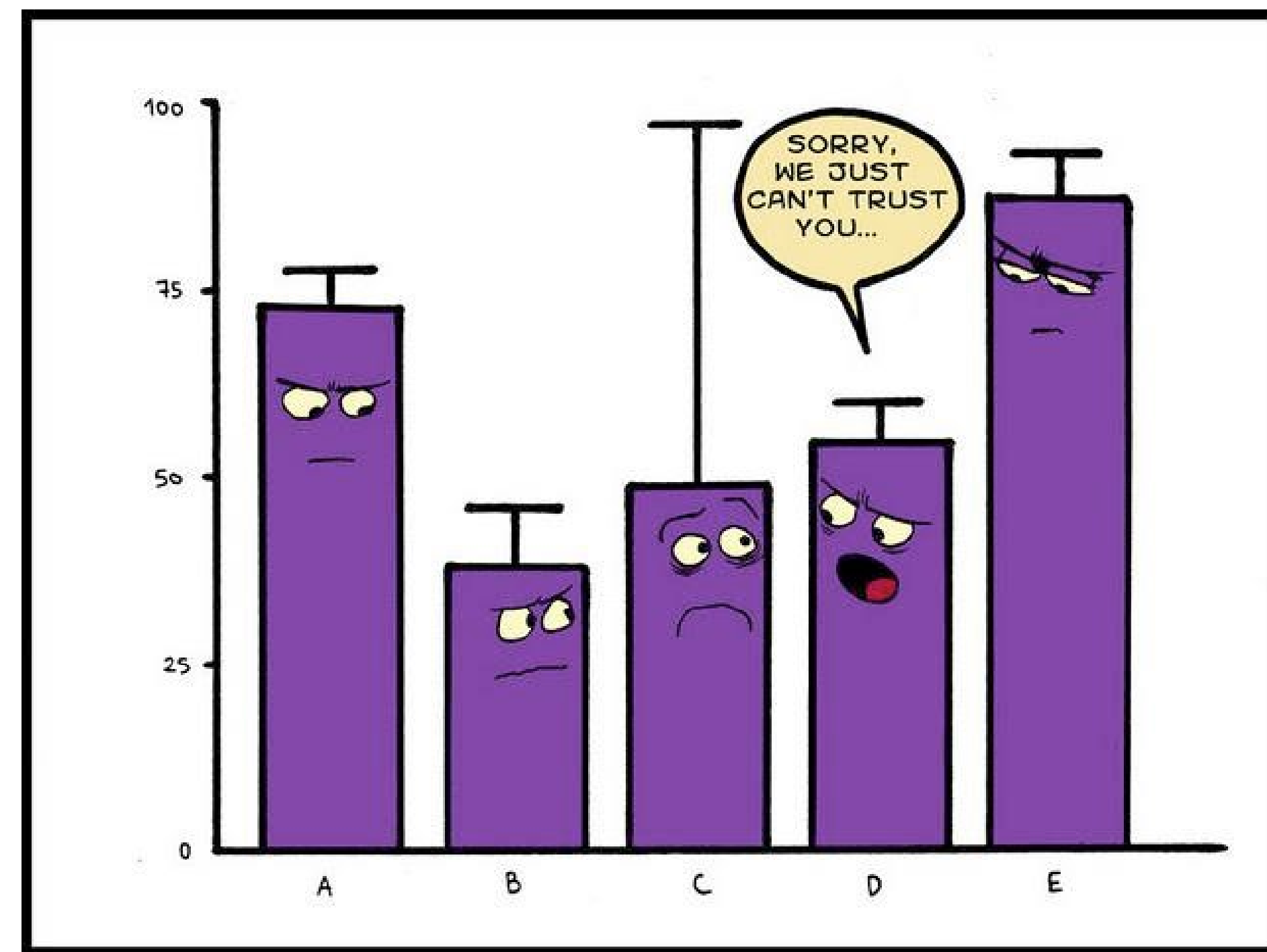


# New Directions in Automated Mechanism Design

Vincent Conitzer; joint work with:



Michael  
Albert  
(Duke → UVA)



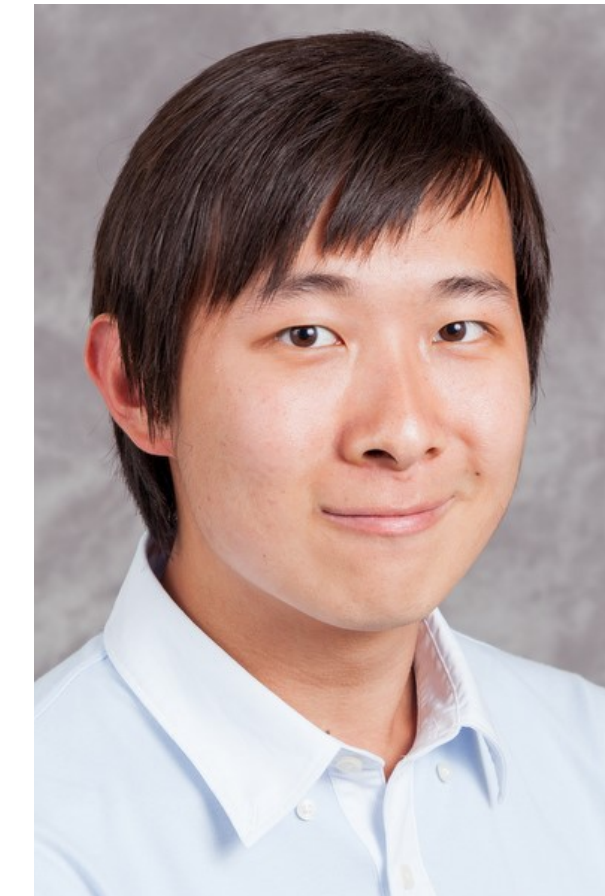
Giuseppe  
(Pino) Lopomo  
(Duke)



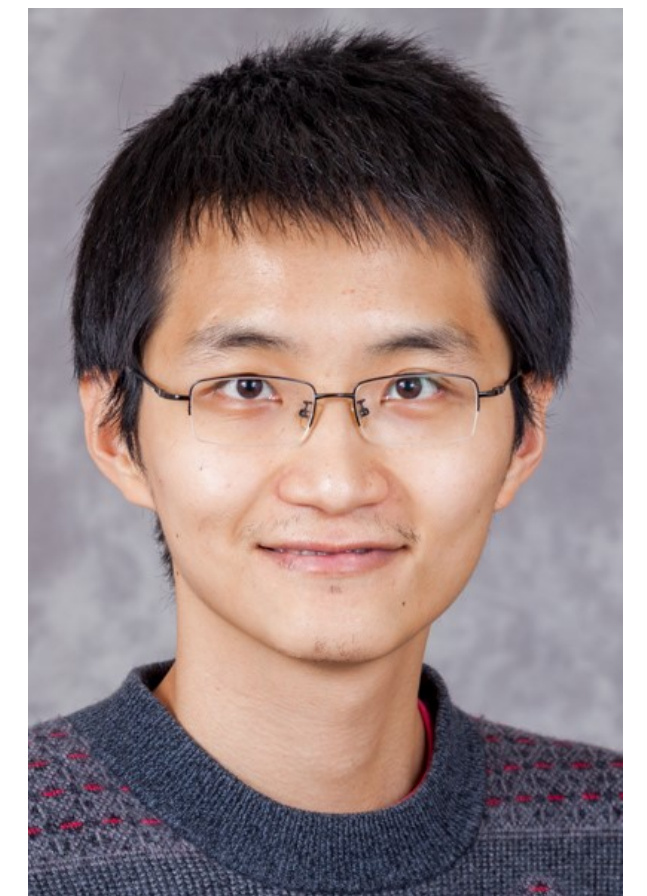
Peter  
Stone  
(UT Austin)



Andrew  
Kephart  
(Duke →  
KeepTruckin)



Hanrui  
Zhang  
(Duke)



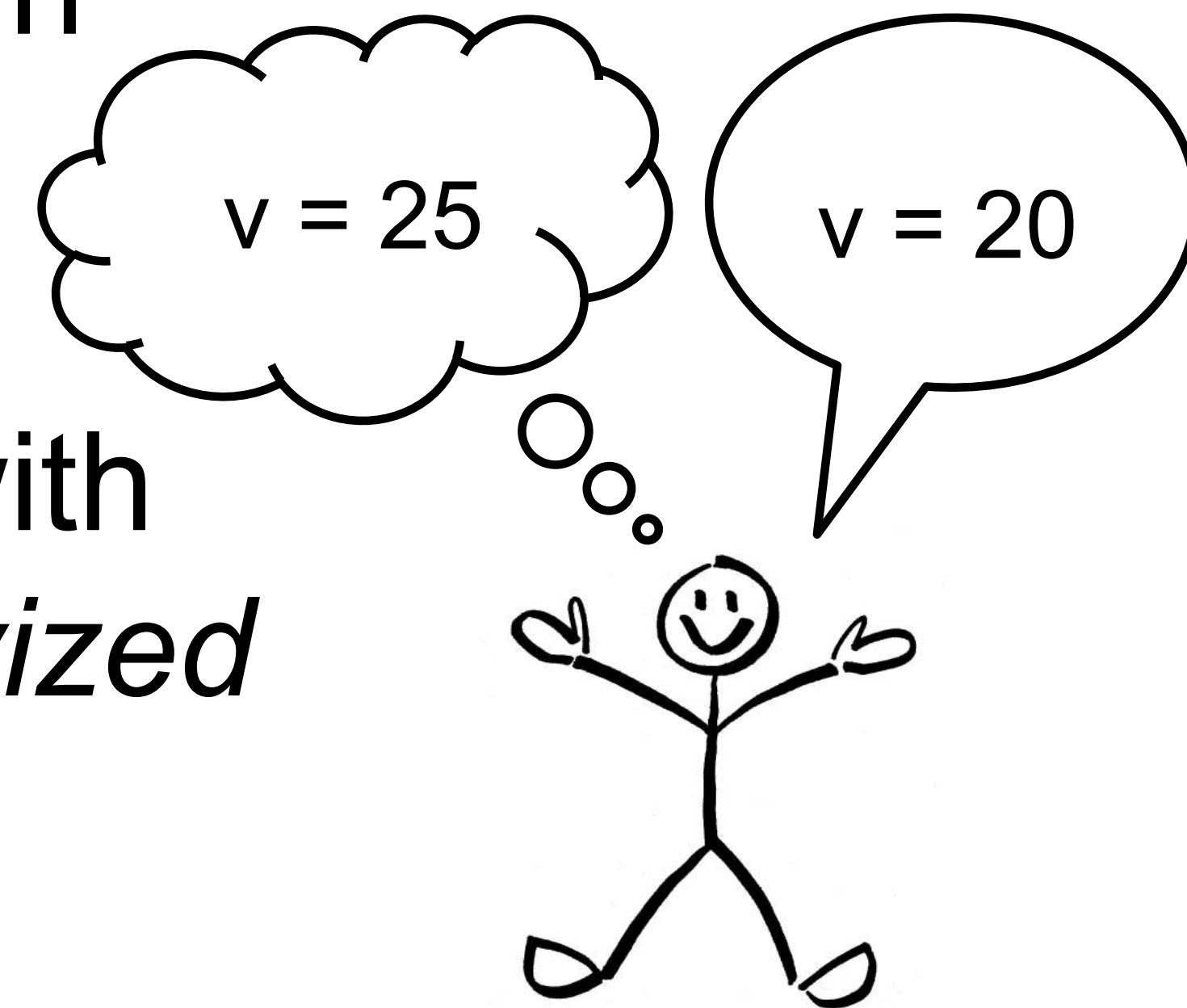
Yu Cheng  
(Duke → IAS  
→ UIC)

# Mechanism design

Make decisions based on the preferences (or other information) of one or more agents (as in social choice)

Focus on *strategic* (game-theoretic) agents with *privately held* information; have to be *incentivized* to reveal it truthfully

Popular approach in design of auctions, matching mechanisms, ...



# Sealed-bid auctions (on a single item)



Bidder  $i$  determines how much the item is worth to her ( $v_i$ )

Writes a bid ( $v'_i$ ) on a piece of paper

How would you bid? How much would I make?

**First price:** Highest bid wins, pays bid

**Second price:** Highest bid wins, pays next-highest bid

**First price with reserve:** Highest bid wins iff it exceeds  $r$ , pays bid

**Second price with reserve:** Highest bid wins iff it exceeds  $r$ , pays next highest bid or  $r$  (whichever is higher)

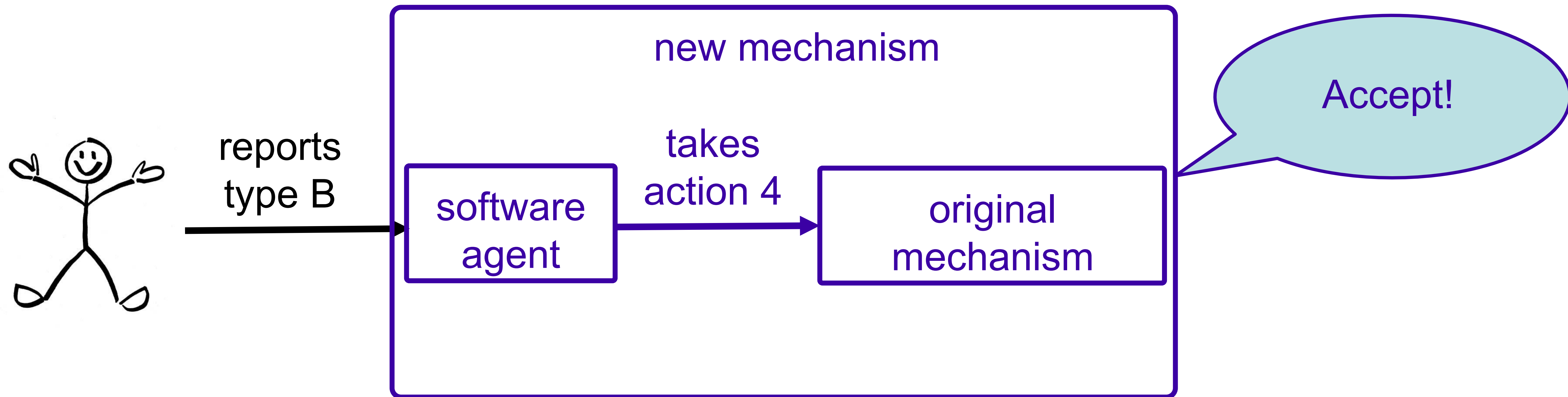
# Revelation Principle

Anything you can achieve, you can also achieve with a truthful (AKA incentive compatible) mechanism.



# Revelation Principle

Anything you can achieve, you can also achieve with a truthful (AKA incentive compatible) mechanism.



# Automated mechanism design input

**Instance** is given by

Set of possible *outcomes*

Set of *agents*

For each agent

set of possible *types*

*probability distribution* over these types

*Objective function*

Gives a value for each outcome for each combination of agents' types

E.g., social welfare, revenue

*Restrictions* on the mechanism

Are *payments* allowed?

Is *randomization* over outcomes allowed?

What versions of *incentive compatibility (IC)* & *individual rationality (IR)* are used?

# How hard is designing an optimal *deterministic* mechanism (without reporting costs)?

[C. & Sandholm UAI'02, ICEC'03, EC'04]

<b>NP-complete</b> (even with 1 reporting agent):	Solvable in <b>polynomial time</b> (for any <i>constant</i> number of agents):
<ol style="list-style-type: none"><li>1. Maximizing social welfare (no payments)</li><li>2. Designer's own utility over outcomes (no payments)</li><li>3. General (linear) objective that doesn't regard payments</li><li>4. Expected revenue</li></ol>	<ol style="list-style-type: none"><li>1. Maximizing social welfare (not regarding the payments) (<b>VCG</b>)</li></ol>

1 and 3 hold even with no IR constraints

# Positive results (randomized mechanisms)

[C. & Sandholm UAI'02, ICEC'03, EC'04]

- Use linear programming

- Variables:

$p(o \mid \theta_1, \dots, \theta_n)$  = probability that outcome  $o$  is chosen given types  $\theta_1, \dots, \theta_n$

(maybe)  $\pi_i(\theta_1, \dots, \theta_n)$  =  $i$ 's payment given types  $\theta_1, \dots, \theta_n$

- Strategy-proofness constraints: for all  $i, \theta_1, \dots, \theta_n, \theta_i'$ :

$$\sum_o p(o \mid \theta_1, \dots, \theta_n) u_i(\theta_i, o) + \pi_i(\theta_1, \dots, \theta_n) \geq$$

$$\sum_o p(o \mid \theta_1, \dots, \theta_i', \dots, \theta_n) u_i(\theta_i, o) + \pi_i(\theta_1, \dots, \theta_i', \dots, \theta_n)$$

- Individual-rationality constraints: for all  $i, \theta_1, \dots, \theta_n$ :

$$\sum_o p(o \mid \theta_1, \dots, \theta_n) u_i(\theta_i, o) + \pi_i(\theta_1, \dots, \theta_n) \geq 0$$

- Objective (e.g., sum of utilities)

$$\sum_{\theta_1, \dots, \theta_n} p(\theta_1, \dots, \theta_n) \sum_i (\sum_o p(o \mid \theta_1, \dots, \theta_n) u_i(\theta_i, o) + \pi_i(\theta_1, \dots, \theta_n))$$

- Also works for BNE incentive compatibility, ex-interim individual rationality notions, other objectives, etc.

- For deterministic mechanisms, can still use mixed integer programming: require probabilities in  $\{0, 1\}$

–Remember typically designing the optimal deterministic mechanism is NP-hard



# A simple example

One item for sale (free disposal)

2 agents, IID valuations: uniform over {1, 2}

Maximize expected revenue under ex-interim

IR, Bayes-Nash equilibrium

How much can we get?

(What is optimal expected welfare?)

		Agent 2's valuation	
		1	2
Agent 1's valuation	1	0.25	0.25
	2	0.25	0.25

probabilities

Status: OPTIMAL

Objective: obj = 1.5 (MAXimum)

[nonzero variables:]

Our old AMD

solver [C. & Sandholm, 2002, 2003]

gives:

$p_{t_1_1_o3}$	1	(probability of disposal for (1, 1))
$p_{t_2_1_o1}$	1	(probability 1 gets the item for (2, 1))
$p_{t_1_2_o2}$	1	(probability 2 gets the item for (1, 2))
$p_{t_2_2_o2}$	1	(probability 2 gets the item for (2, 2))
$pi_{2_2_1}$	2	(1's payment for (2, 2))
$pi_{2_2_2}$	4	(2's payment for (2, 2))

# A slightly different distribution

One item for sale (free disposal)

2 agents, valuations drawn as on right

Maximize expected revenue under ex-interim

IR, Bayes-Nash equilibrium

How much can we get?

(What is optimal expected welfare?)

		<i>Agent 2's valuation</i>	
		1	2
<i>Agent 1's valuation</i>	1	0.251	0.250
	2	0.250	0.249
		probabilities	

Status: OPTIMAL

Objective: obj = 1.749 (MAXimum)

[some of the nonzero payment variables:]

pi\_1\_1\_2 62501

pi\_2\_1\_2 -62750

pi\_2\_1\_1 2

pi\_1\_2\_2 3.992

*You'd better be really sure about your distribution!*

# A nearby distribution without correlation

One item for sale (free disposal)

2 agents, valuations IID: 1 w/ .501, 2 w/ .499

Maximize expected revenue under ex-interim

IR, Bayes-Nash equilibrium

How much can we get?

(What is optimal expected welfare?)

**Status:** OPTIMAL

**Objective:** obj = 1.499 (MAXimum)

		<i>Agent 2's valuation</i>	
		1	2
<i>Agent 1's valuation</i>	1	0.251001	0.249999
	2	0.249999	0.249001
		probabilities	

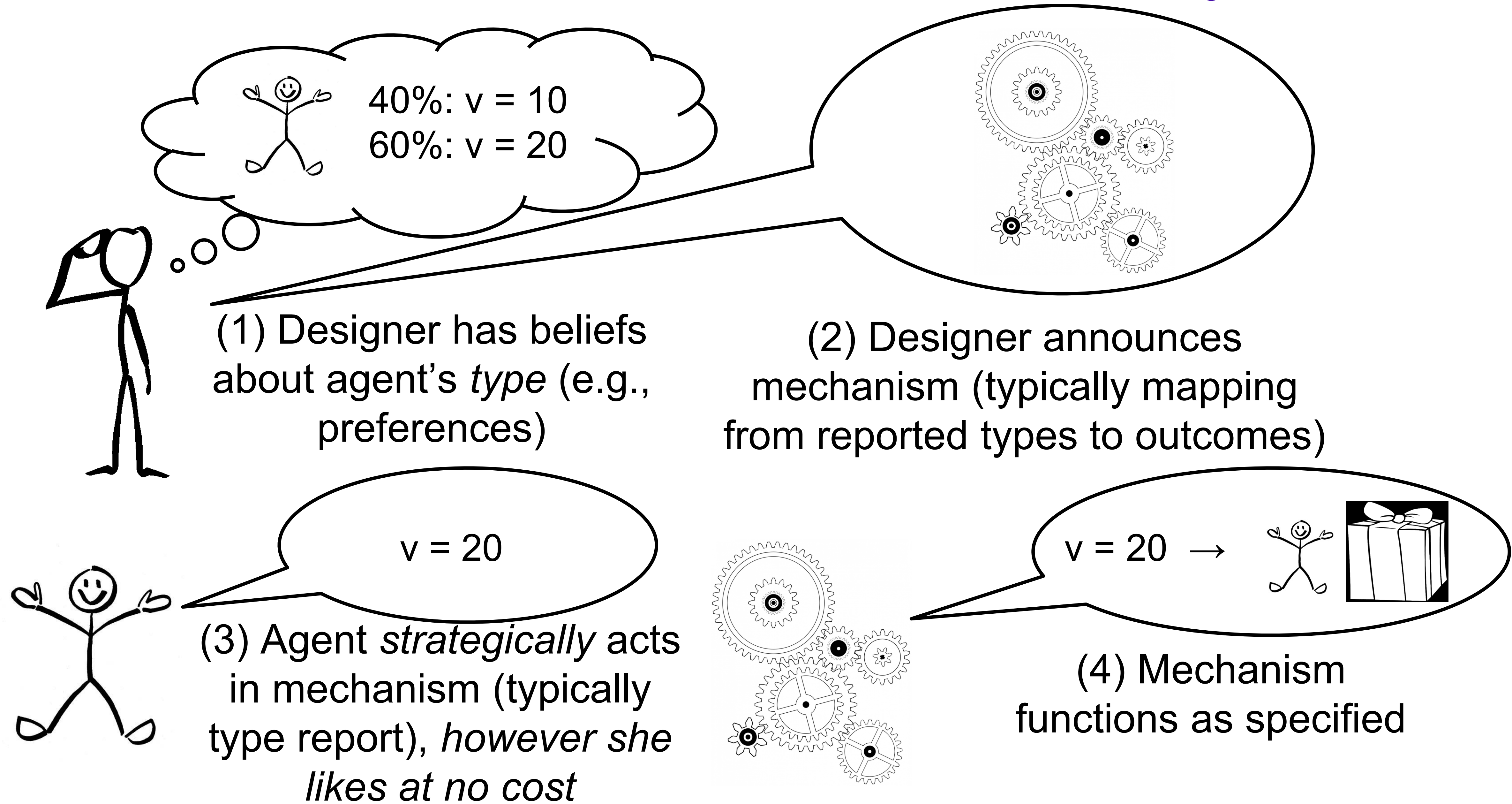
# Cremer-McLean [1985]

For every agent, consider the following matrix  $\Gamma$  of conditional probabilities, where  $\Theta$  is the set of types for the agent and  $\Omega$  is the set of **signals** (joint types for other agents, or something else observable to the auctioneer)

$$\Gamma = \begin{bmatrix} \pi(1|1) & \cdots & \pi(|\Omega||1) \\ \vdots & \ddots & \vdots \\ \pi(1||\Theta|) & \cdots & \pi(|\Omega|||\Theta|) \end{bmatrix}$$

If  $\Gamma$  has rank  $|\Theta|$  for every agent then the auctioneer can **allocate efficiently** and **extract the full surplus as revenue (!!)**

# Standard setup in mechanism design



The mechanism may have more information about the specific agent!

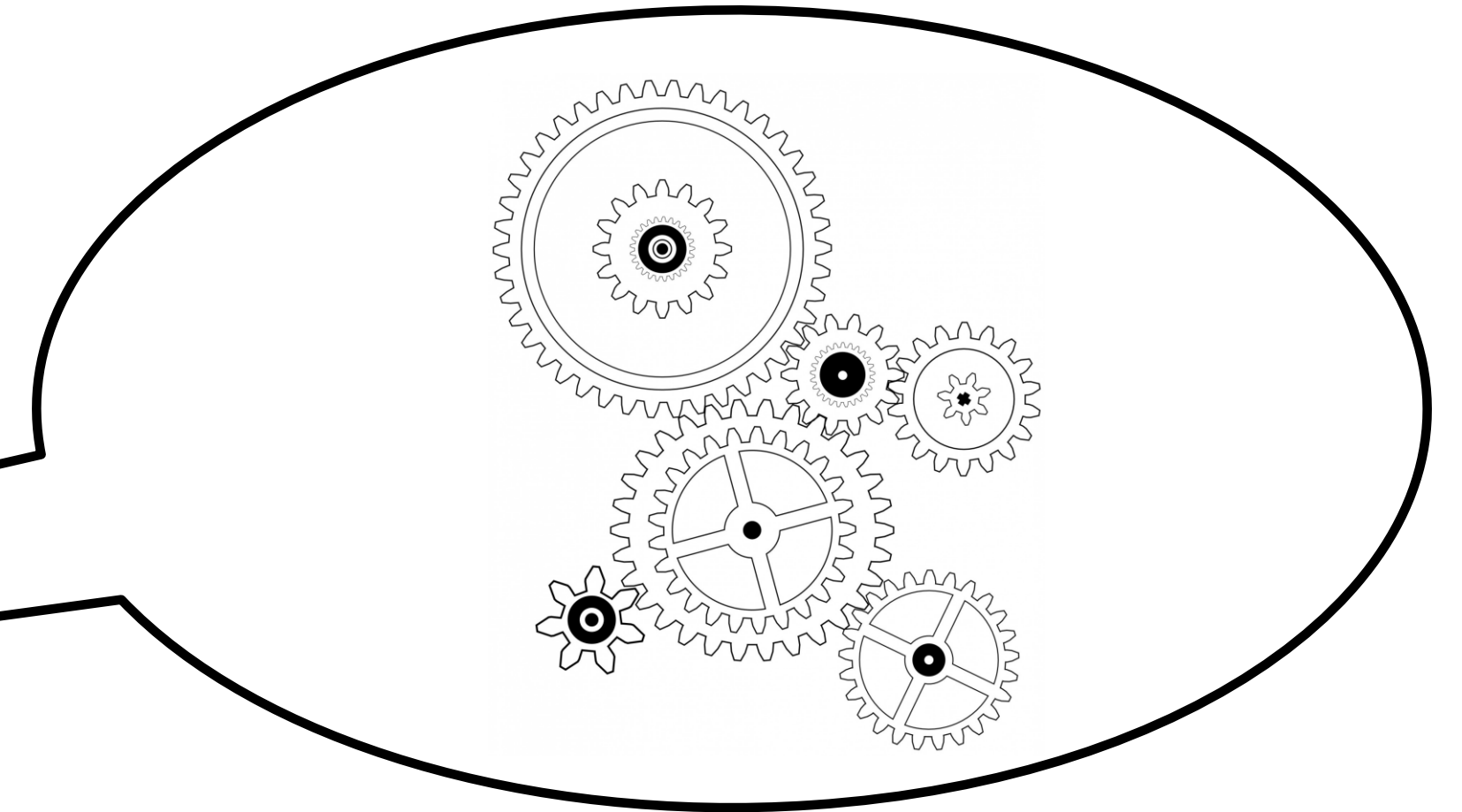
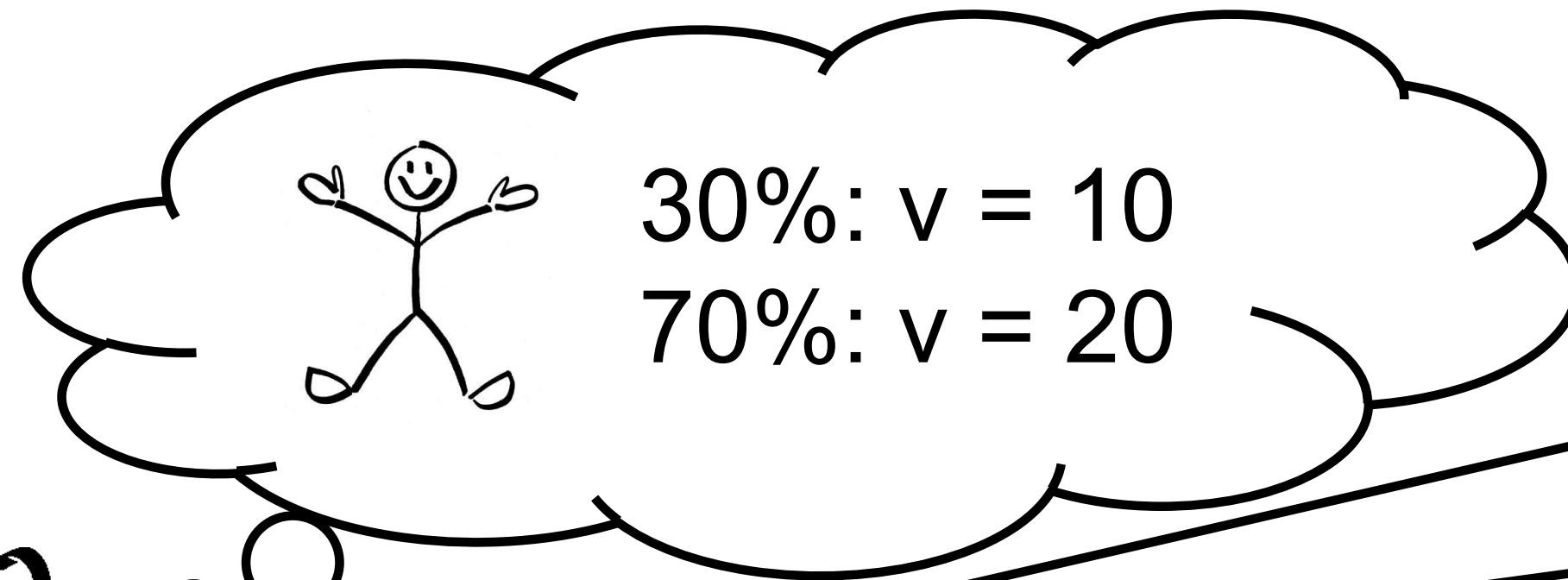
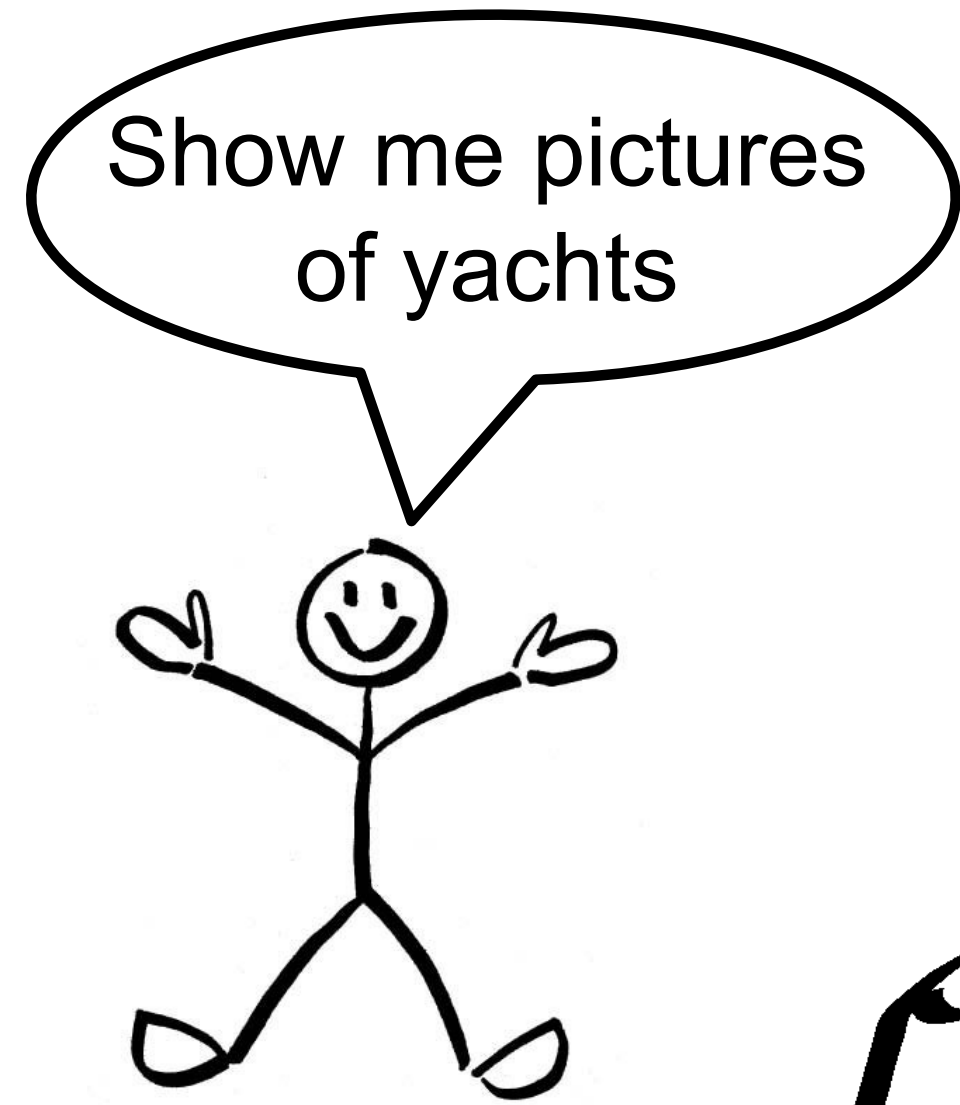
## **application**

online marketplaces  
selling insurance  
university admissions  
webpage ranking

## **information**

actions taken online  
driving record  
courses taken  
links to page

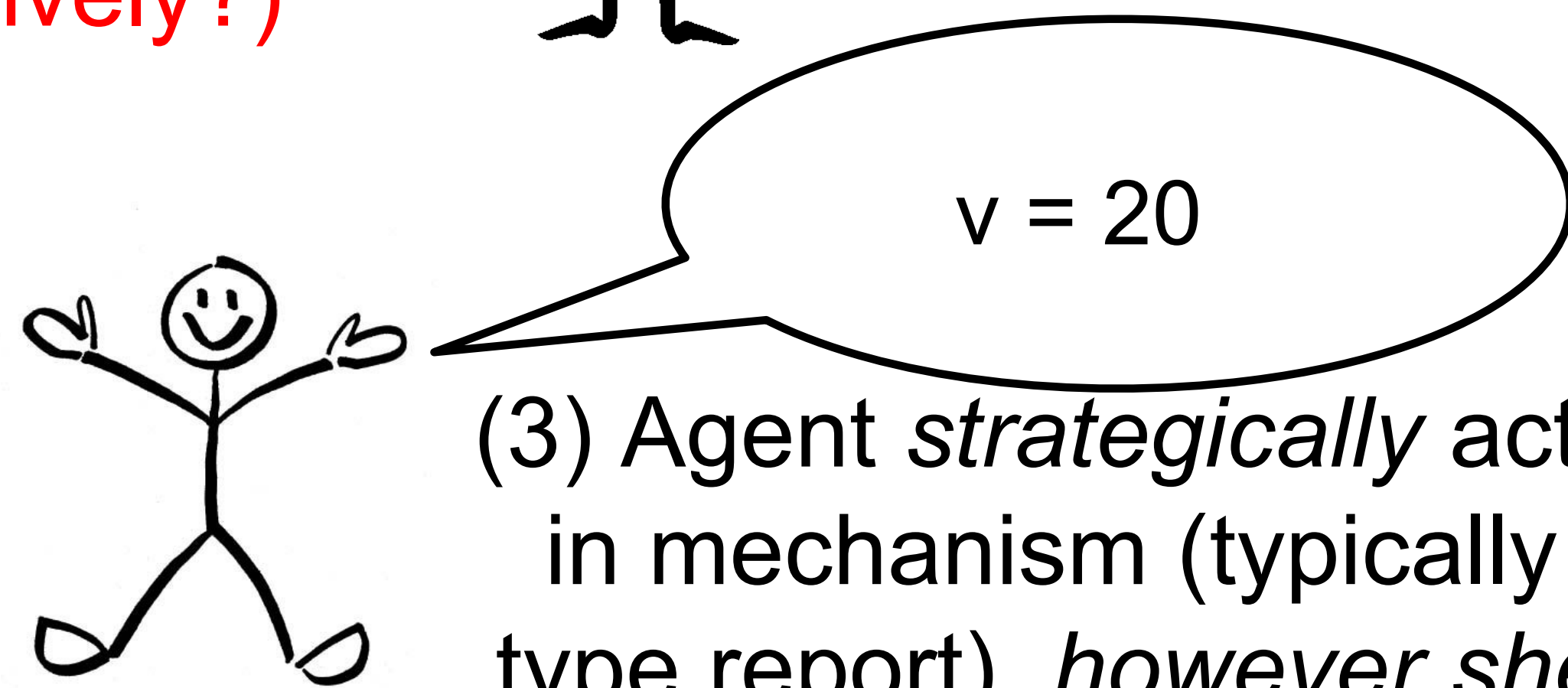
# Attempt 1 at fixing this



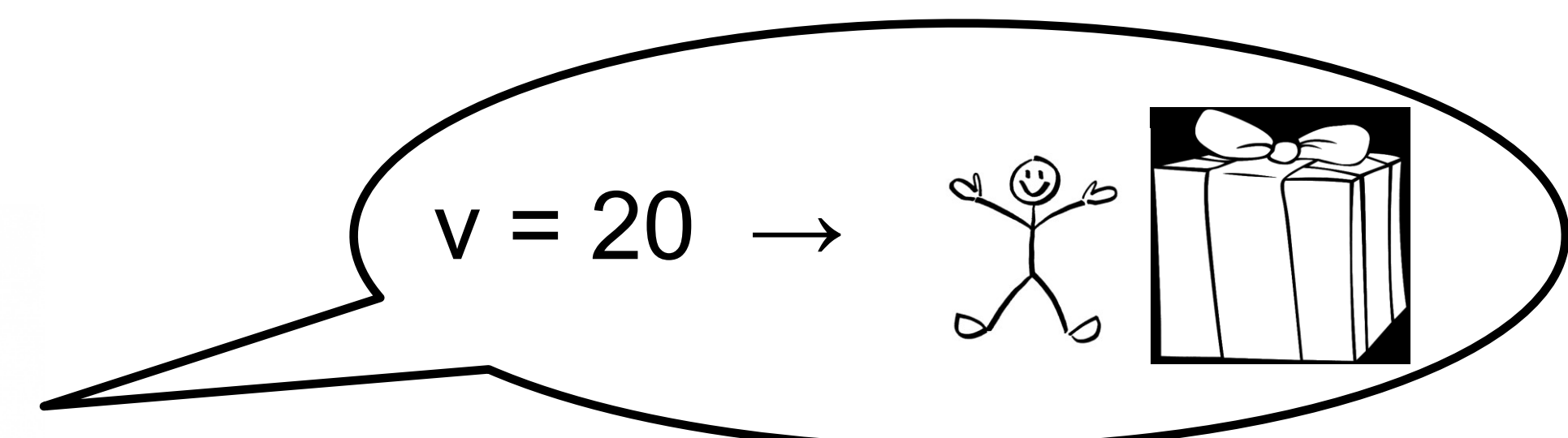
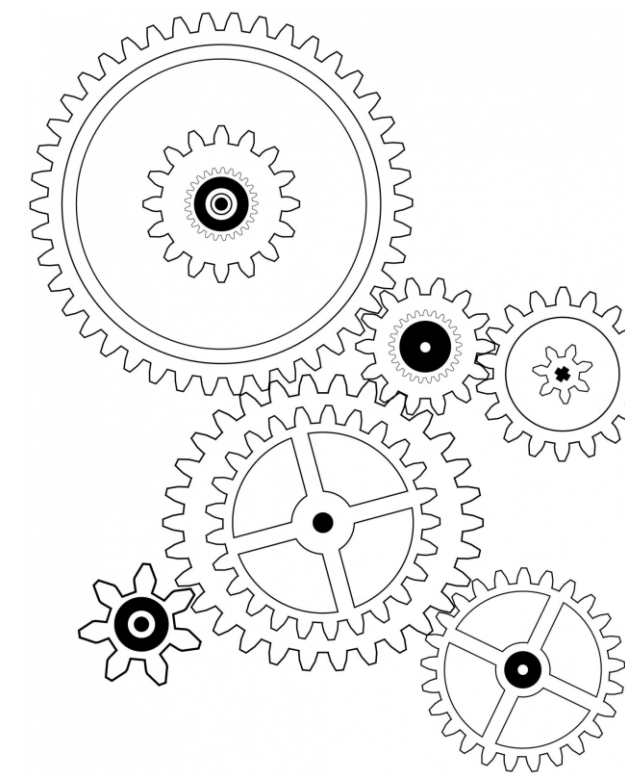
(0) Agent acts in the world  
(naively?)

(1) Designer obtains beliefs about agent's *type* (e.g., preferences)

(2) Designer announces mechanism (typically mapping from reported types to outcomes)

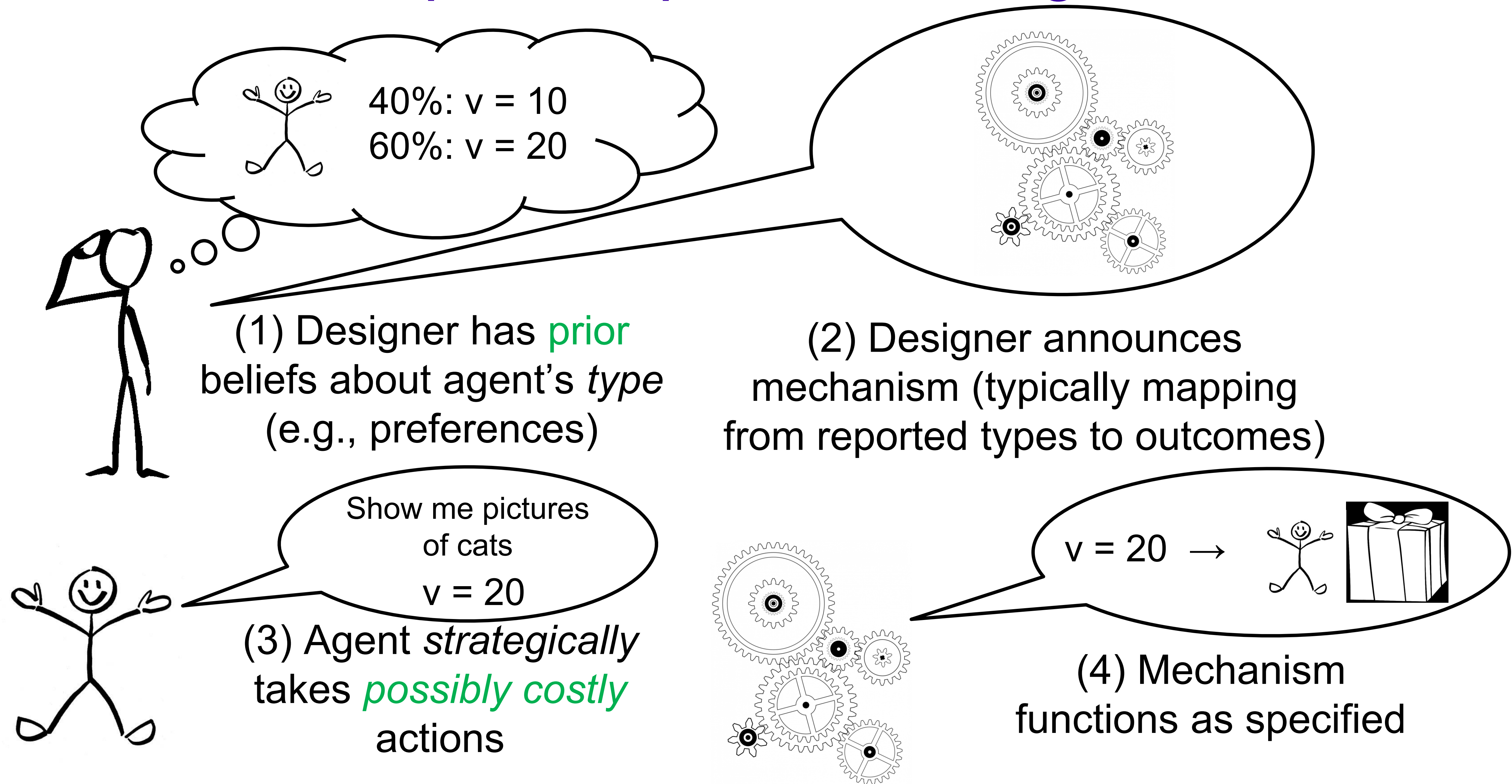


(3) Agent *strategically* acts in mechanism (typically type report), *however she likes at no cost*



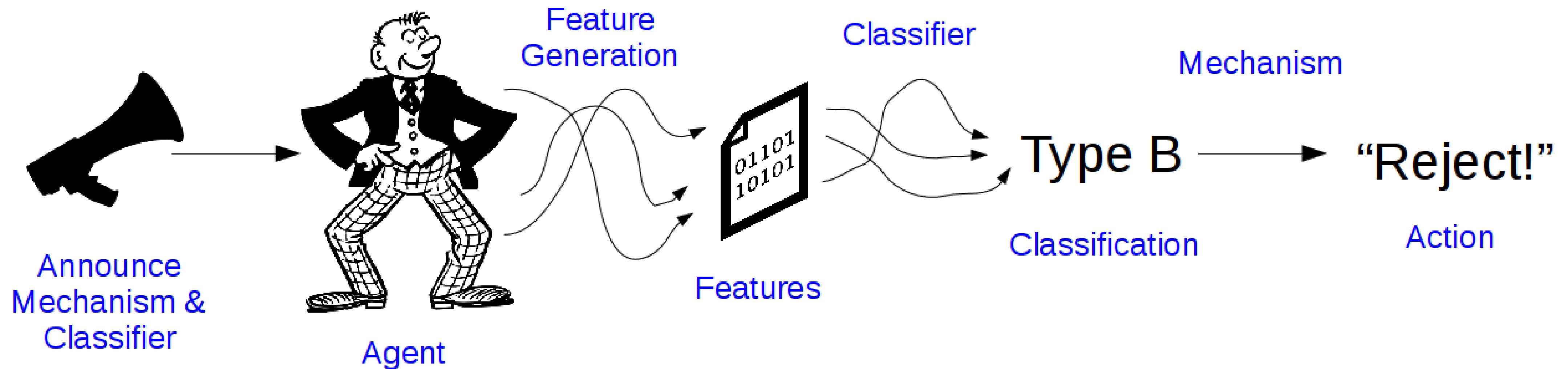
(4) Mechanism functions as specified

# Attempt 2: Sophisticated agent





# Machine learning view



See also later work by Hardt, Megiddo, Papadimitriou, Wootters [2015/2016]

# From Ancient Times...

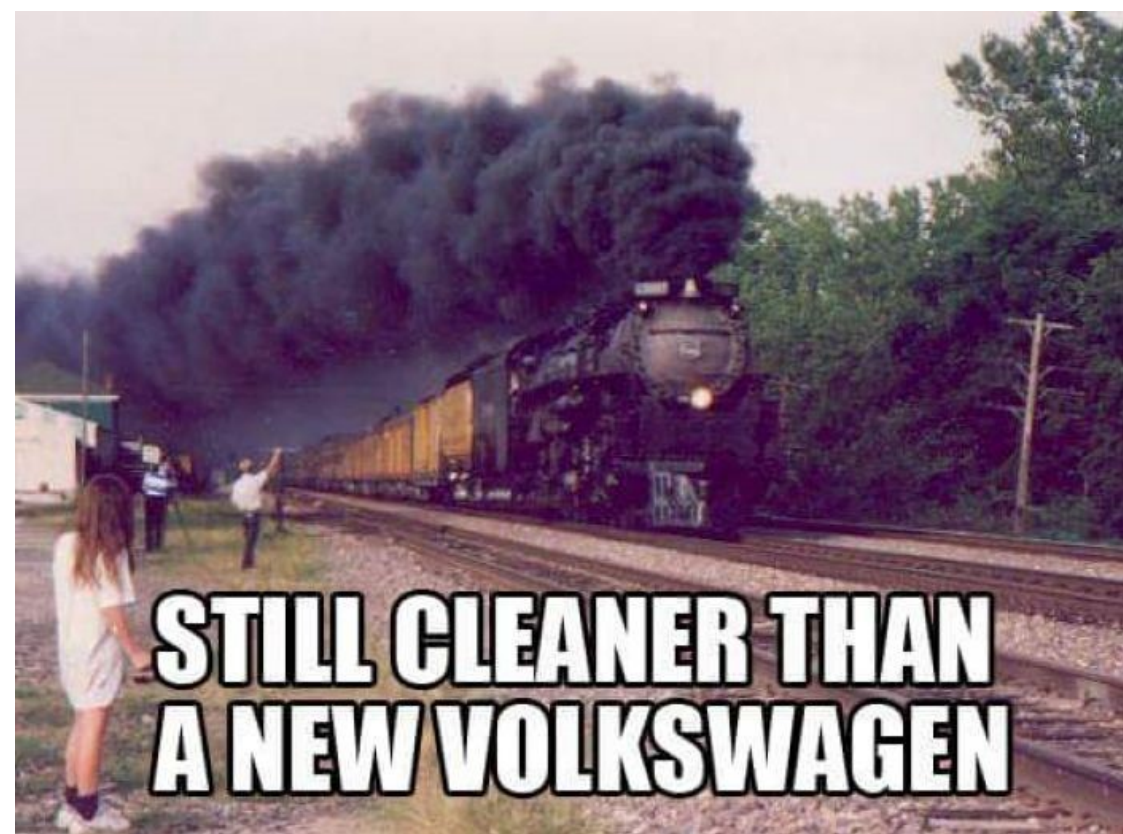
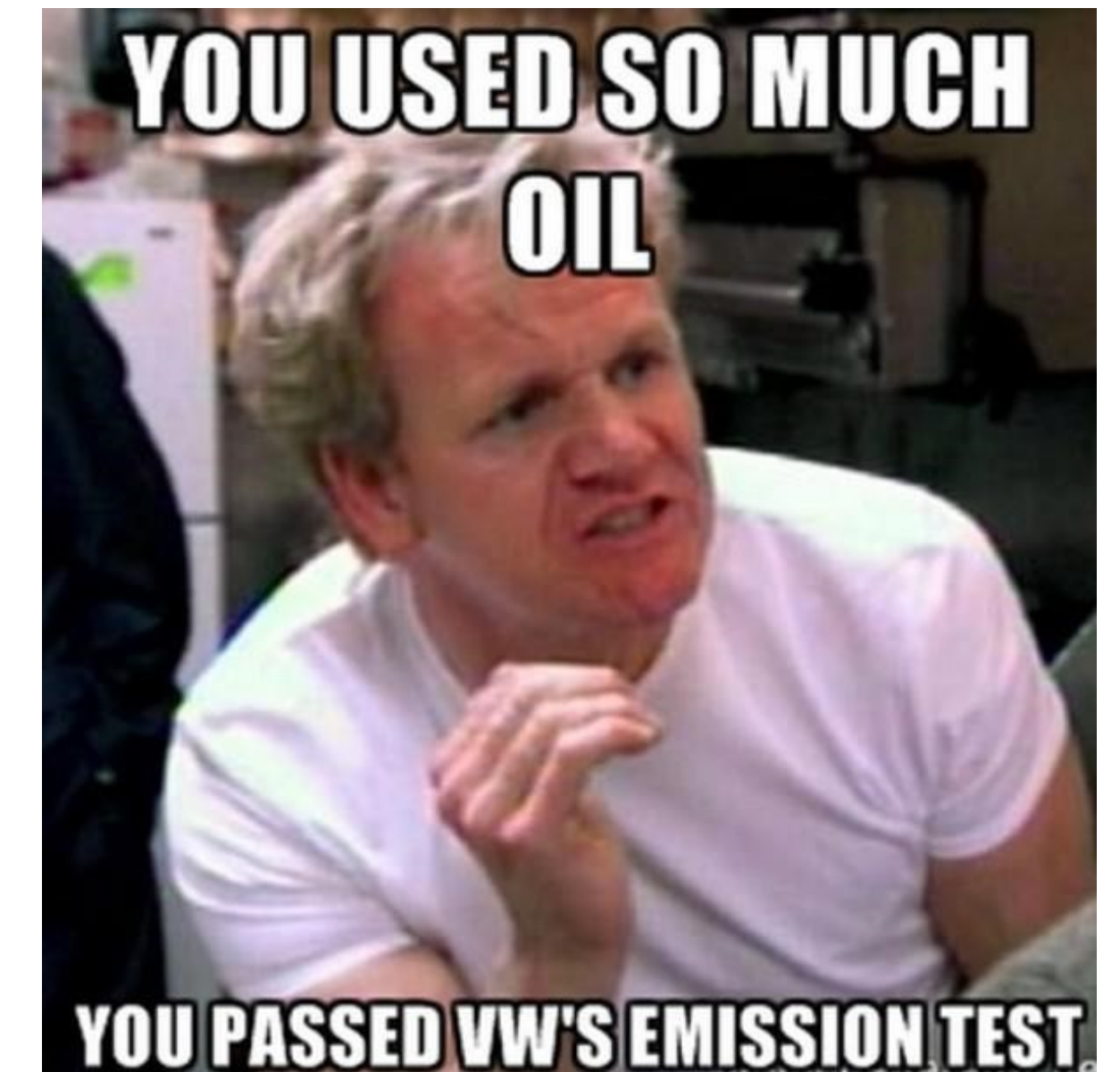
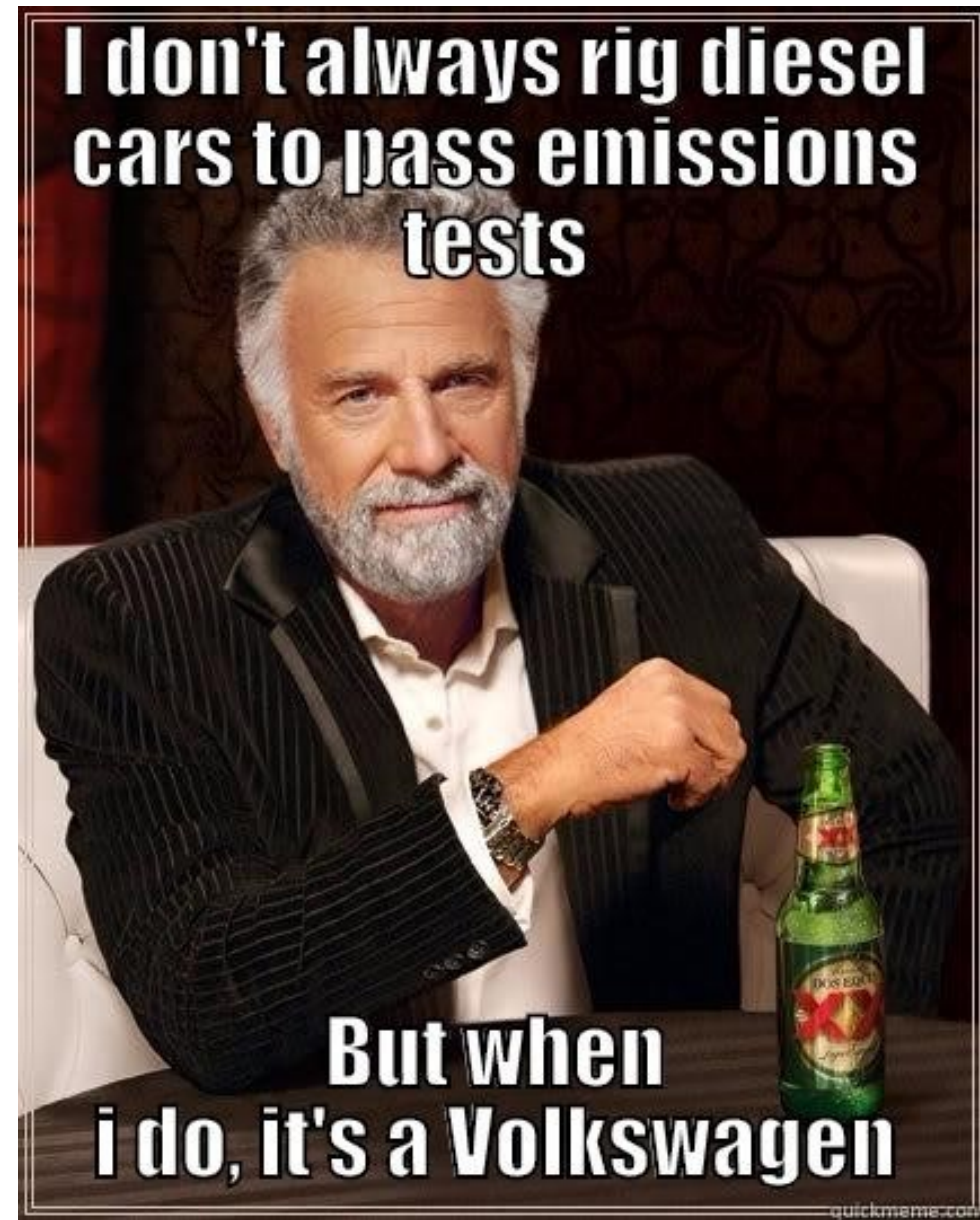


Jacob and Esau



Trojan Horse

# ... to Modern Times



# Illustration: Barbara Buying Fish From Fred

<u>Types (<math>t \in T</math>)</u>	<u>Actions (<math>a \in A</math>)</u>	<u>Choice Function (<math>F : T \rightarrow A</math>)</u>
<i>fresh</i>	<i>accept</i>	<i>fresh</i> $\rightarrow$ <i>accept</i>
<i>ok</i>	<i>reject</i>	<i>ok</i> $\rightarrow$ <i>accept</i>
<i>rotten</i>		<i>rotten</i> $\rightarrow$ <i>reject</i>

Classifications ( $\hat{t} \in \hat{T}$ ): *frêsh*, *ôk*, *rottên*



# ... continued

Effort Function ( $E : T \times \hat{T} \rightarrow \mathbb{R}$ ):

	<i>frêsh</i>	<i>ôk</i>	<i>rotten</i>
<i>fresh</i>	0	0	0
<i>ok</i>	10	0	0
<i>rotten</i>	30	10	0



Valuation Function ( $V : T \times A \rightarrow \mathbb{R}$ ):

$$V(\cdot, \text{accept}) = 20, \quad V(\cdot, \text{reject}) = 0$$

Mechanism  $M : \hat{T} \rightarrow A$

First Try:  $M = \text{frêsh} \rightarrow \text{accept}, \text{ôk} \rightarrow \text{accept}, \text{rotten} \rightarrow \text{reject}$

# ... continued

Effort Function ( $E : T \times \hat{T} \rightarrow \mathbb{R}$ ):

	<i>frêsh</i>	<i>ôk</i>	<i>rotten</i>
<i>fresh</i>	0	0	0
<i>ok</i>	10	0	0
<i>rotten</i>	30	10	0



Valuation Function ( $V : T \times A \rightarrow \mathbb{R}$ ):

$$V(\cdot, \textit{accept}) = 20, \quad V(\cdot, \textit{reject}) = 0$$

Mechanism  $M : \hat{T} \rightarrow A$

**First Try:**  $M = \textit{frêsh} \rightarrow \textit{accept}, \textit{ôk} \rightarrow \textit{accept}, \textit{rotten} \rightarrow \textit{reject}$

**Better:**  $M^* = \textit{frêsh} \rightarrow \textit{accept}, \textit{ôk} \rightarrow \textit{reject}, \textit{rotten} \rightarrow \textit{reject}.$

# Comparison With Other Models

	$\hat{f}resh$	$\hat{o}k$	$\hat{r}otten$
$fresh$	0	0	0
$ok$	0	0	0
$rotten$	0	0	0

Standard Mechanism Design

	$\hat{f}resh$	$\hat{o}k$	$\hat{r}otten$
$fresh$	0	$\infty$	0
$ok$	$\infty$	0	0
$rotten$	0	0	0

Mechanism Design with Partial Verification

	$\hat{f}resh$	$\hat{o}k$	$\hat{r}otten$
$fresh$	$\infty$	0	0
$ok$	1.2	5	$-\infty$
$rotten$	-3	0	0

Mechanism Design with Signaling Costs

Green and Laffont. [Partially verifiable information and mechanism design](#). RES 1986

Auletta, Penna, Persiano, Ventre. [Alternatives to truthfulness are hard to recognize](#). AAMAS 2011

# Question

Given:

Types ( $t \in T$ )	Actions ( $a \in A$ )	Choice Function ( $F : T \rightarrow A$ )
<i>fresh</i>	<i>accept</i>	<i>fresh</i> $\rightarrow$ <i>accept</i>
<i>ok</i>	<i>reject</i>	<i>ok</i> $\rightarrow$ <i>accept</i>
<i>rotten</i>		<i>rotten</i> $\rightarrow$ <i>reject</i>

Classifications ( $\hat{t} \in \hat{T}$ ): *frêsh*, *ôk*, *rotten*



Effort Function ( $E : T \times \hat{T} \rightarrow \mathbb{R}$ ):

	<i>frêsh</i>	<i>ôk</i>	<i>rotten</i>
<i>fresh</i>	0	0	0
<i>ok</i>	10	0	0
<i>rotten</i>	30	10	0



Valuation Function ( $V : T \times A \rightarrow \mathbb{R}$ ):

$$V(\cdot, \text{accept}) = 20, \quad V(\cdot, \text{reject}) = 0$$

Then:

Does there exist a

Mechanism  $M : \hat{T} \rightarrow A$

which implements the choice function?

**NP-complete!**

Auletta, Penna, Persiano, Ventre. [Alternatives to truthfulness are hard to recognize](#). AAMAS 2011



# Results



with Andrew  
Kephart  
(AAMAS 2015)

		<b>Transfers (T)</b>		<b>No Transfers (NT)</b>	
		<i>Two Outcomes (TO)</i>	<i>Injective SCF (FI)</i>	<i>Two Outcomes (TO)</i>	<i>Injective SCF (FI)</i>
<b>Free Utilities (FU)</b>	<i>Unrestricted Costs (U)</i>	NP-c	<b>NP-c</b>	NP-c	<b>NP-c</b>
	$\{0, \infty\}$ <i>Costs (ZI)</i>	NP-c	<b>NP-c</b>	NP-c	<b>P</b>
<b>Targeted Utilities (TU)</b>	<i>Unrestricted Costs (U)</i>	<b>NP-c</b>	<b>P</b>	NP-c	<b>P</b>
	$\{0, \infty\}$ <i>Costs (ZI)</i>	<b>NP-c</b>	<b>P</b>	NP-c	<b>P</b>

Non-bolded results are from:

Auletta, Penna, Persiano, Ventre. [Alternatives to truthfulness are hard to recognize](#). AAMAS 2011

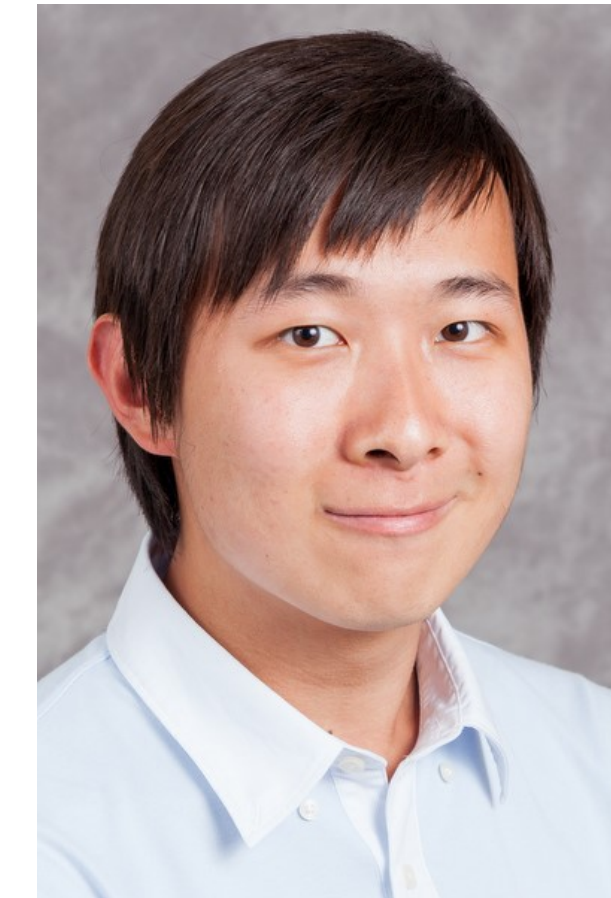
Hardness results fundamentally rely on **revelation principle failing** – conditions under which revelation principle still holds in [Green & Laffont '86](#) and [Yu '11](#) (partial verification), and [Kephart & C. EC'16](#) (costly signaling).

# When Samples Are Strategically Selected

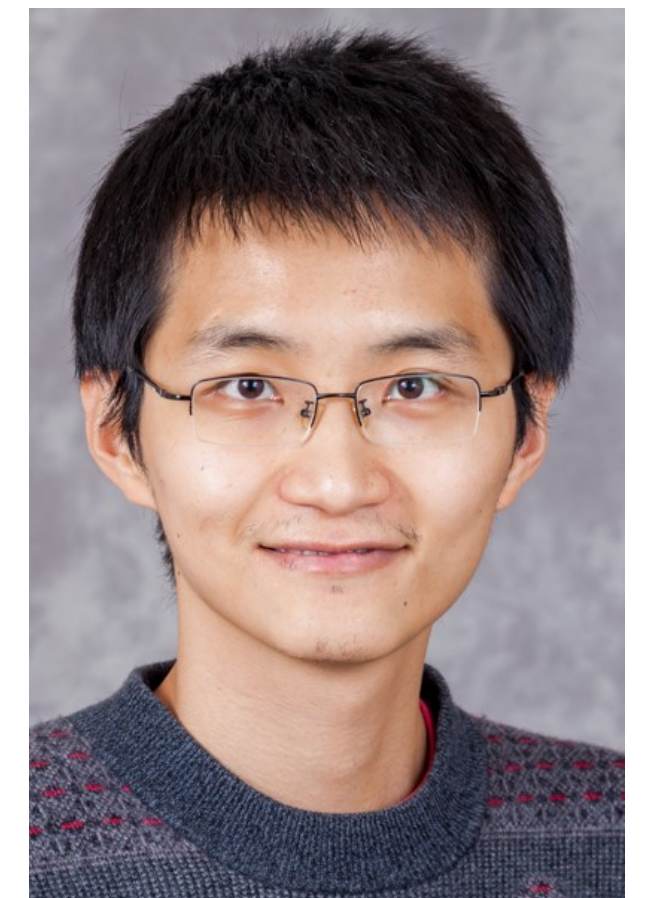
ICML 2019, with



Bob, Professor of Rocket Science

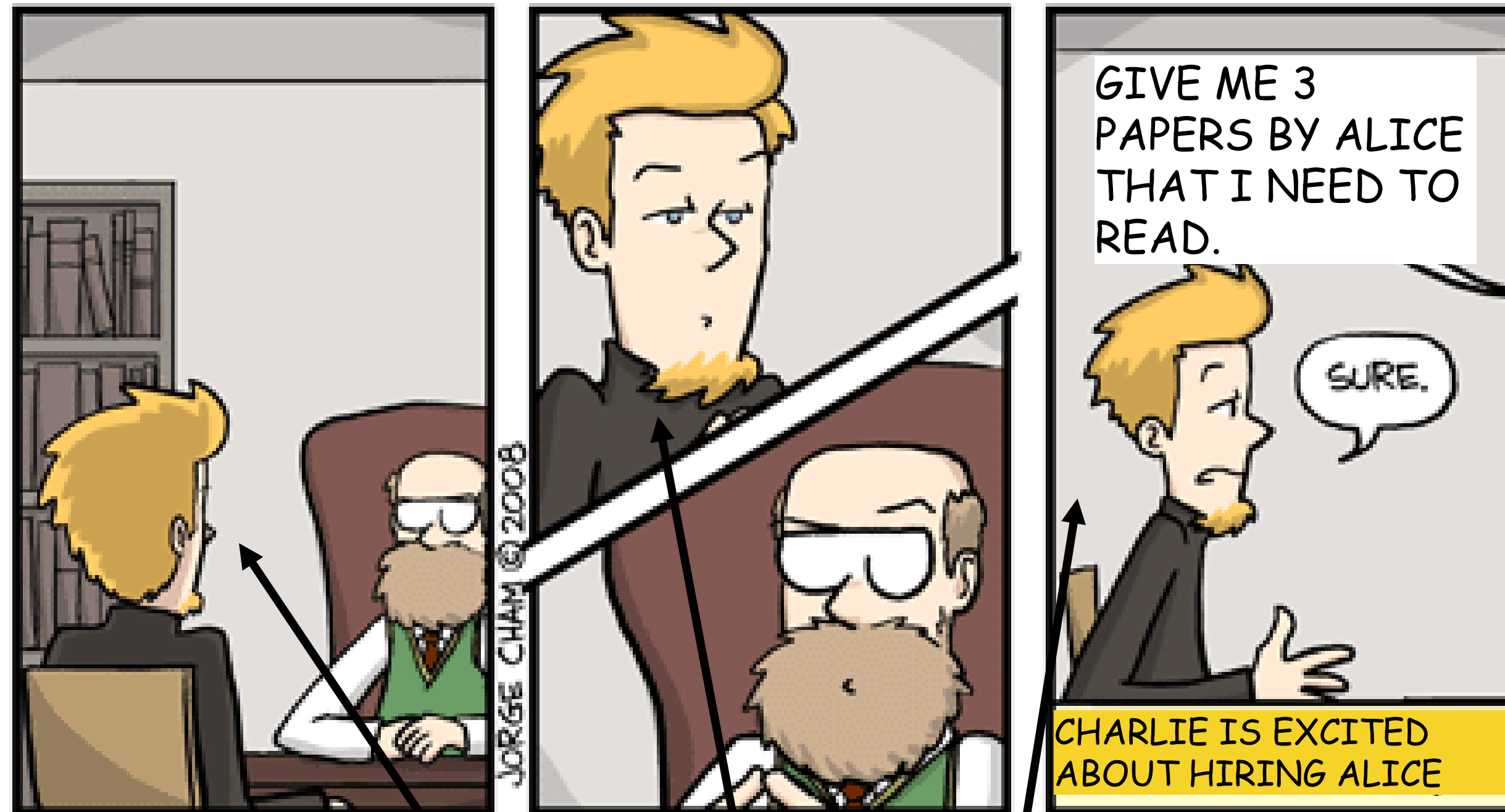


Hanrui  
Zhang  
(Duke)



Yu Cheng  
(Duke → IAS  
→ UIC)

# Academic hiring...



Charlie, Bob's student

# Academic hiring...

I NEED TO CHOOSE THE BEST 3 PAPERS TO CONVINCe BOB, SO THAT HE WILL HIRE ALICE.

CHARLIE WILL DEFINITELY PICK THE BEST 3 PAPERS BY ALICE, AND I NEED TO CALIBRATE FOR THAT.



# The general problem

A **distribution (Alice)** over paper qualities  $\theta \in \{g, b\}$  arrives, which can be either a **good** one ( $\theta = g$ ) or a **bad** one ( $\theta = b$ )



Alice, the postdoc applicant

# The general problem

The **principal (Bob)** announces a **policy**, according to which he decides, based on the **report** of the **agent (Charlie)**, whether to **accept  $\theta$  (hire Alice)**



# The general problem

The **agent (Charlie)** has access to  **$n(=15)$  iid samples (papers)** from  $\theta$  (Alice), from which he **chooses  $m(=3)$**  as his **report**



# The general problem

The agent (**Charlie**) sends his **report** to the principal, aiming to convince the principal (**Bob**) to accept  $\theta$  (**Alice**)





# The general problem

The **principal (Bob)** observes the **report** of the **agent (Charlie)**, and makes the decision according to the policy announced

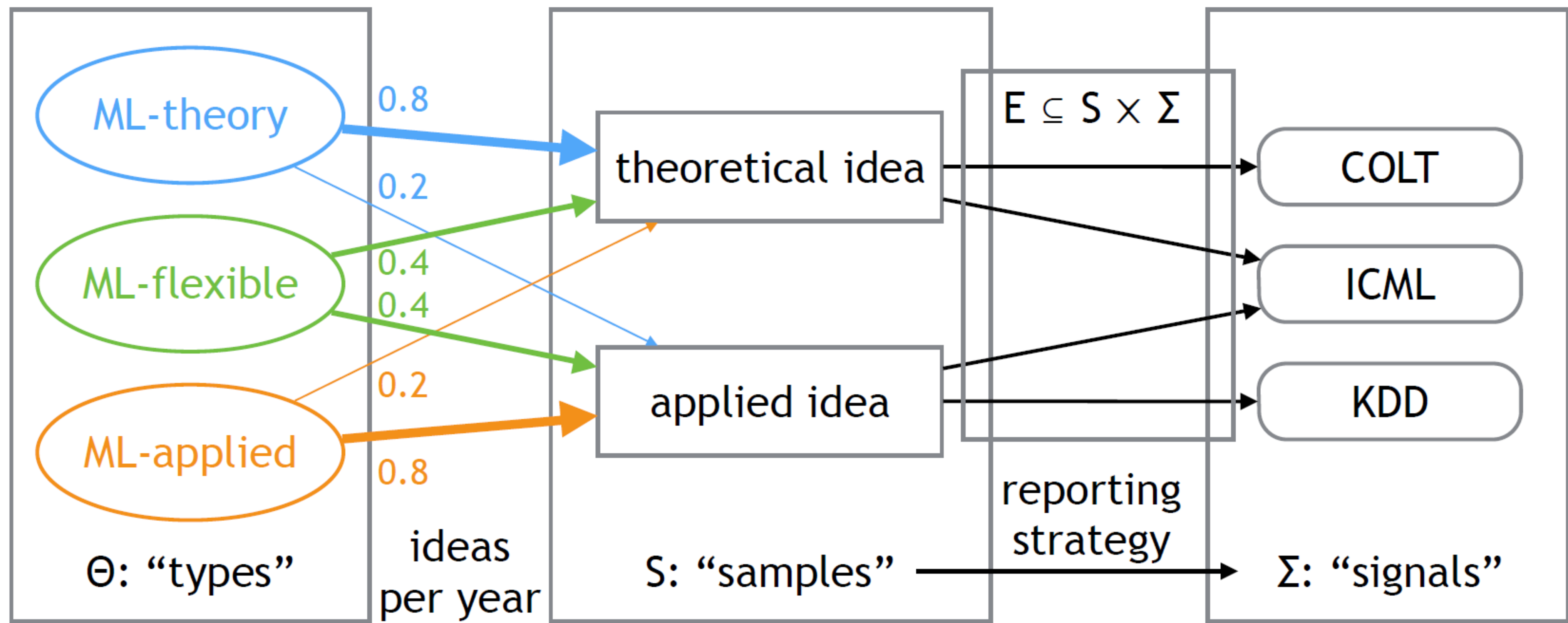


# Questions

How does strategic selection affect the principal's policy?

Is it easier or harder to classify based on strategic samples, compared to when the principal has access to iid samples?

Should the principal ever have a diversity requirement (e.g., at least 1 mathematical paper and at least 1 experimental paper), or only go by total quality according to a single metric?



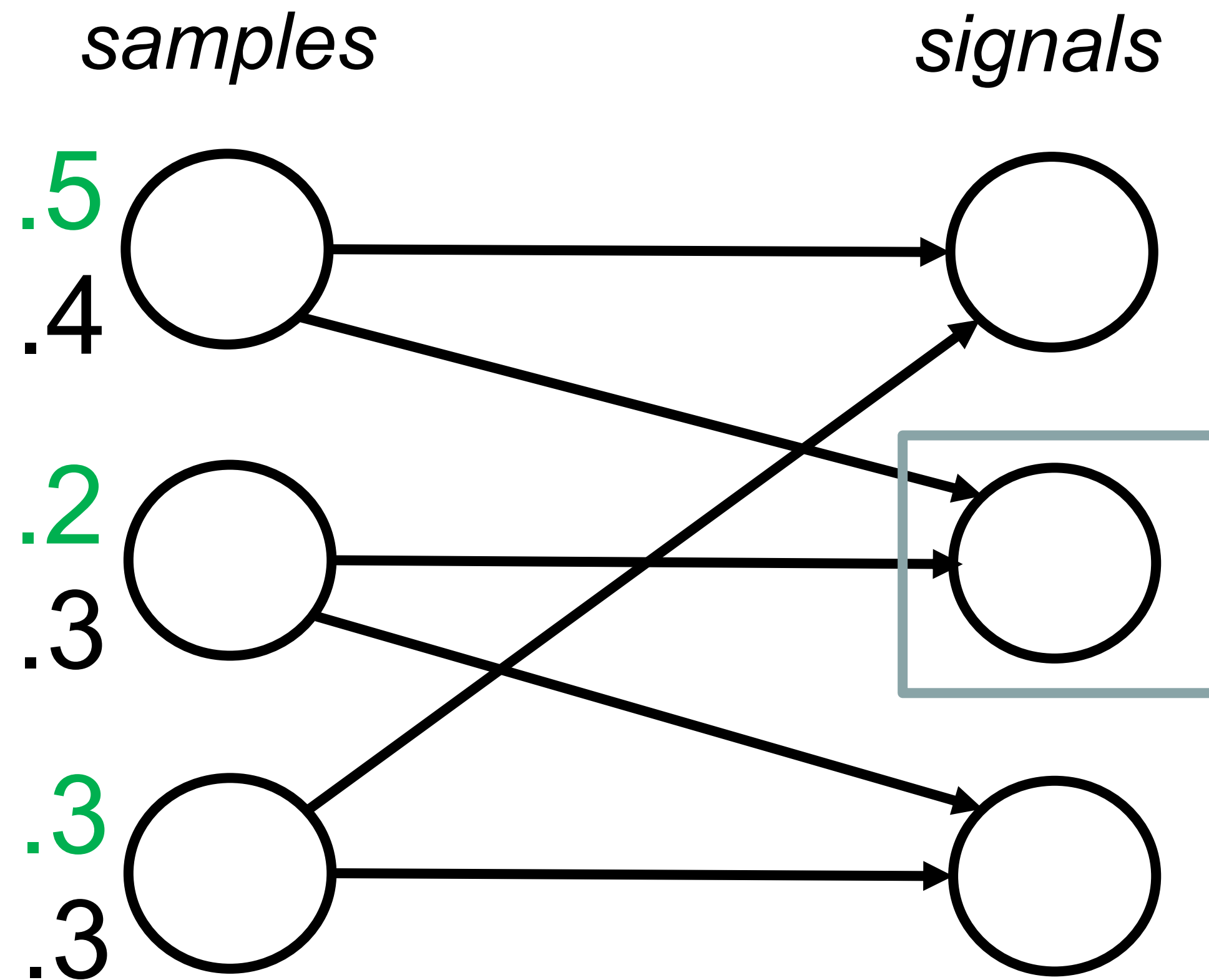
Agent's problem:

- “How do I distinguish myself from other types?”
- “How many samples do I need for that?”

Principal's problem:

- “How do I tell ML-flexible agents from others?”
- “At what point in their career can I reliably do that?”

# One **good** and one **bad** distribution



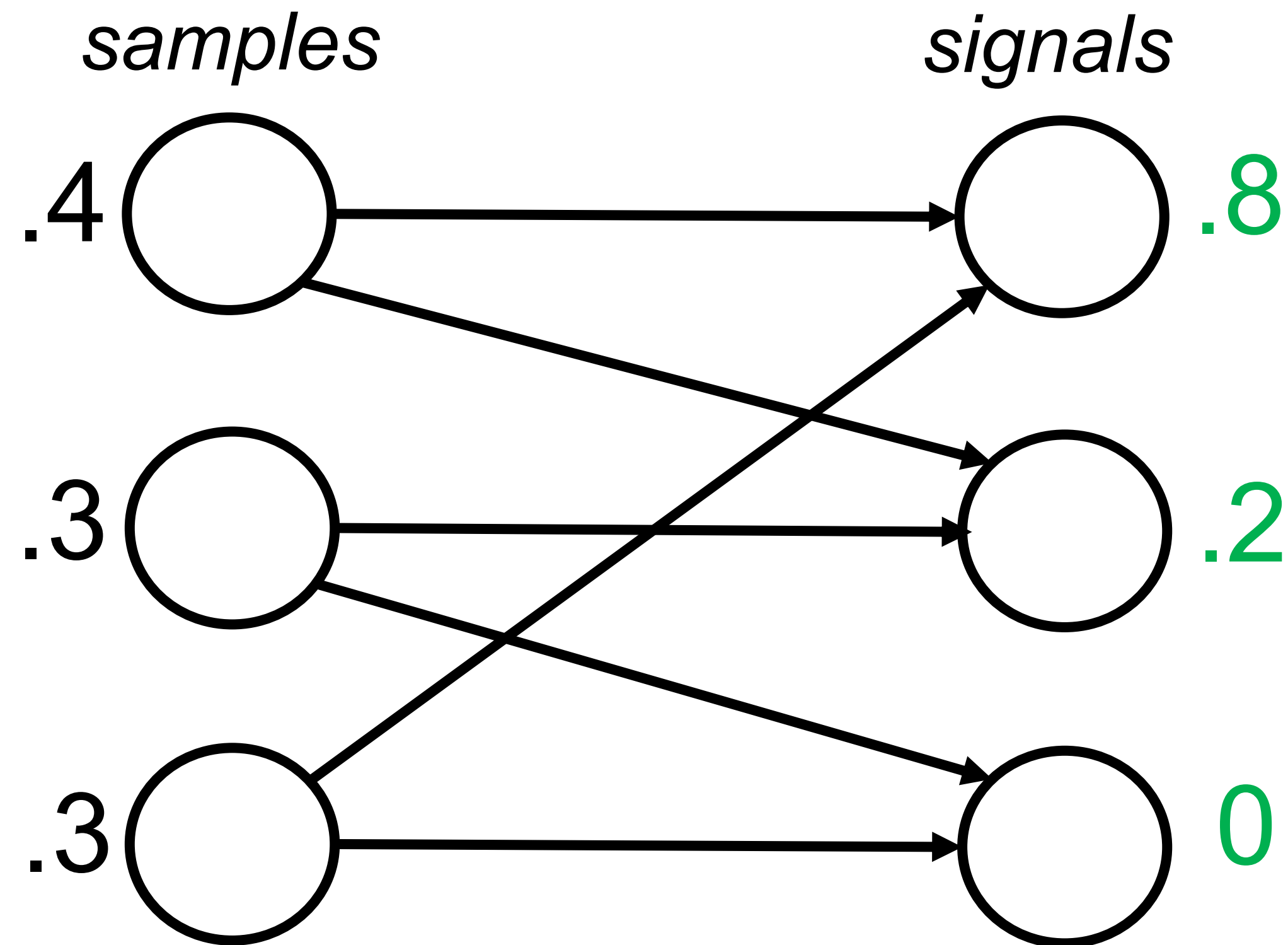
This subset covers  
*.5+.2=.7 good mass* and  
*.4+.3=.7 bad mass*, so it  
doesn't work. (What  
does?)

Pick a subset of the right-hand side (to accept) that maximizes  
(*green mass covered* - *black mass covered*)

If positive, can (eventually) distinguish; otherwise not.

NP-hard in general.

# But if we know the strategy for the good distribution (revelation principle holds):



Can place good mass on the signals side because we know the strategy

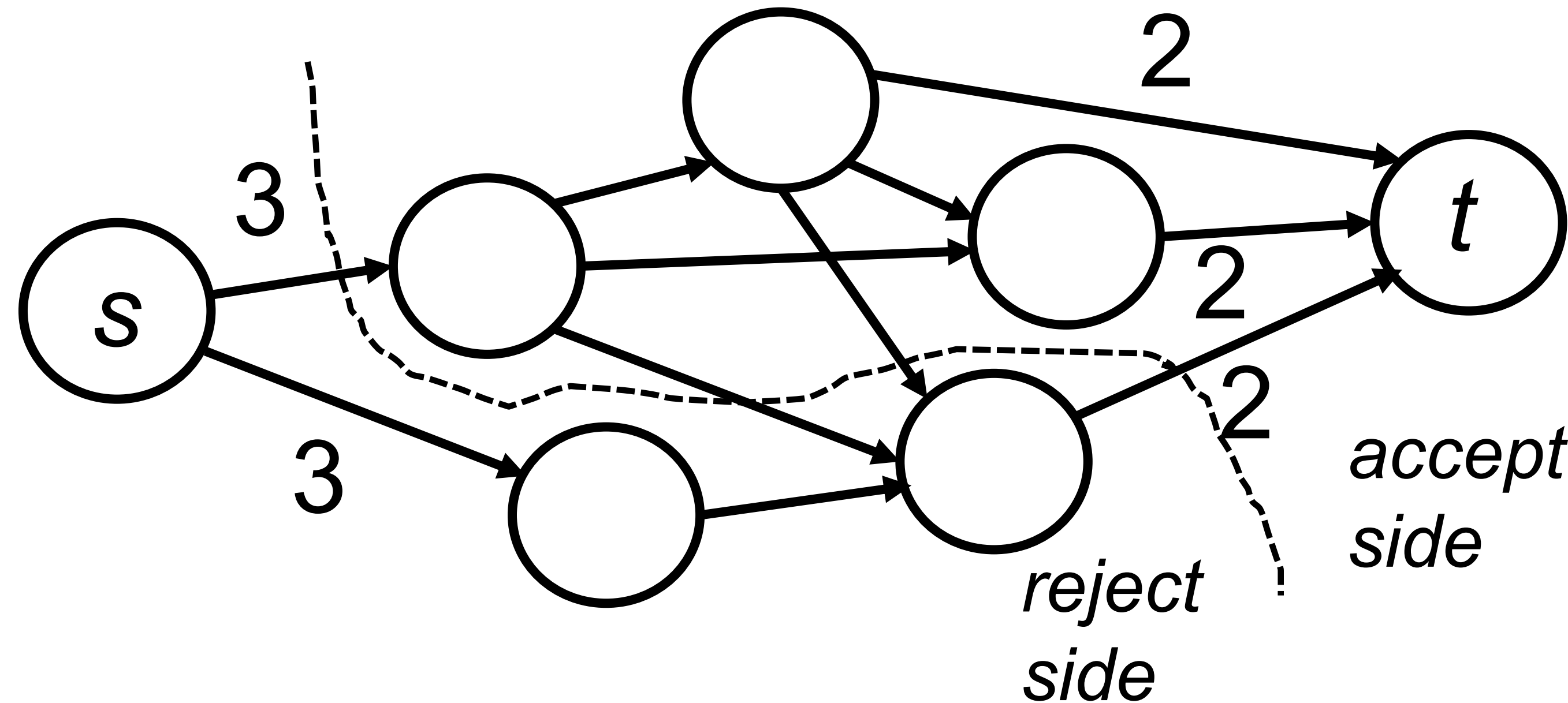
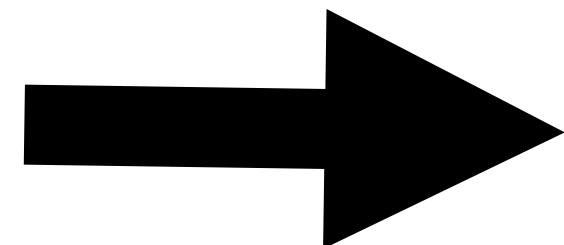
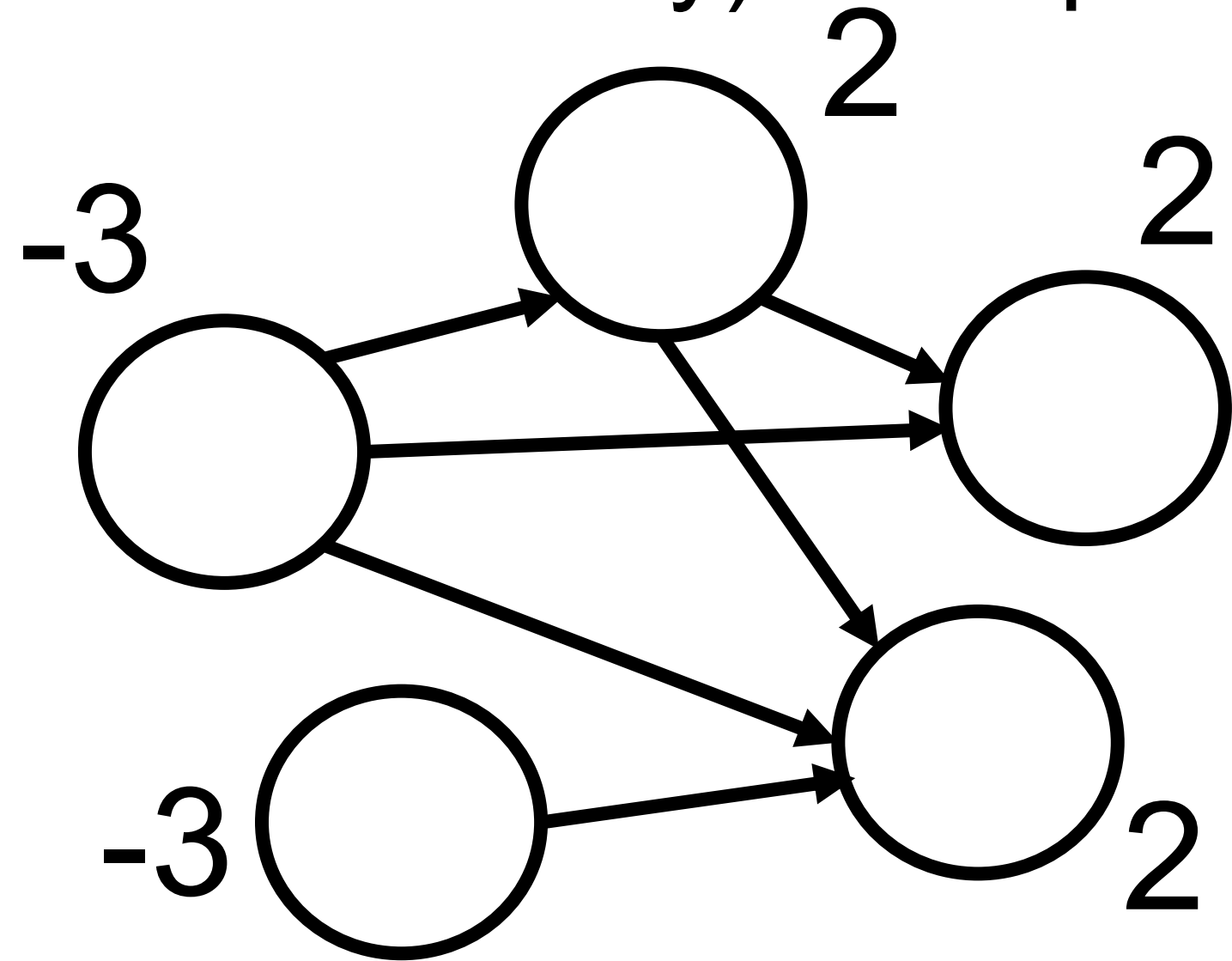
Solve as maximum flow/matching from left to right with capacities on vertices  
Duality gives set of signals to accept (~Hall's marriage theorem)

# Optimization: reduction to min cut

(when revelation principle holds)

types are vertices; edges imply ability to (cost-effectively) misreport

edges between types have capacity  $\infty$



In sampling case, can check existence of edges with previous technique

Values are  $P(\text{type}) \cdot \text{value}(\text{type})$

Can be generalized to more outcomes than accept/reject, if types have the same utility over them.

# Conclusion

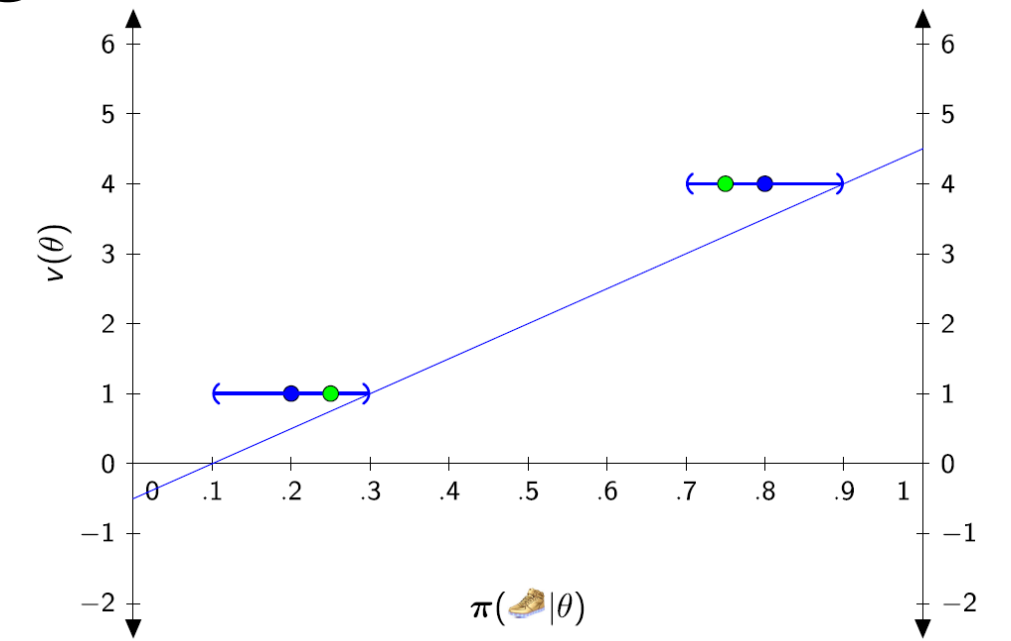
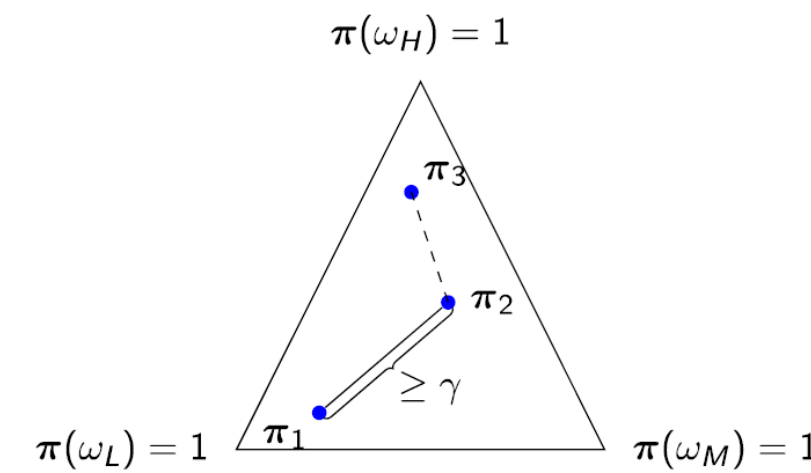
## First part:

When considering correlation, **small** changes can have a **huge** effect  
 Automatically designing **robust** mechanisms addresses this  
 Combines well with **learning** (under some conditions)

0.251	0.250
0.250	0.249

v.

0.251001	0.249999
0.249999	0.249001

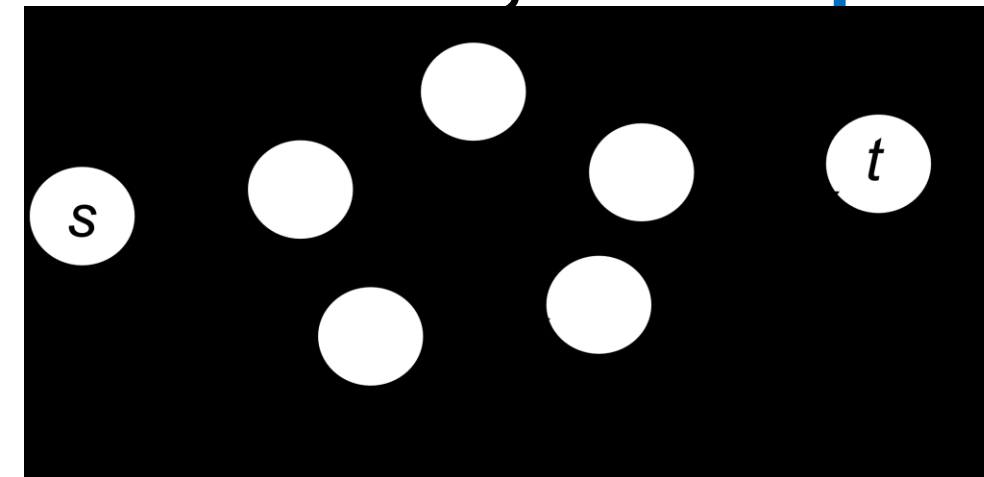
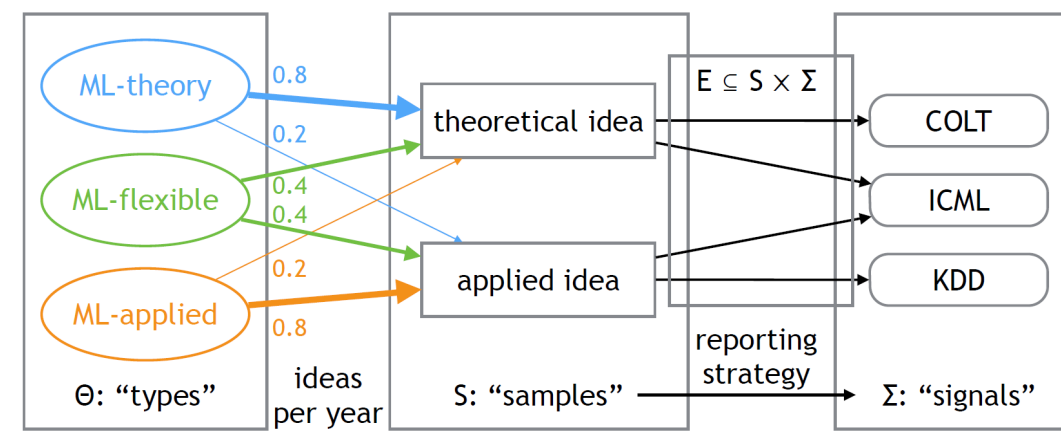


## Second part:

With costly or limited misreporting, **revelation principle** can fail  
 Causes **computational hardness** in general  
 Sometimes agents report based on their **samples**  
 Some **efficient algorithms** for the infinite limit case; **sample bounds**

**Effort Function** ( $E : T \times \hat{T} \rightarrow \mathbb{R}$ ):

	<i>fresh</i>	<i>ok</i>	<i>rotten</i>
<i>fresh</i>	0	0	0
<i>ok</i>	10	0	0
<i>rotten</i>	30	10	0



# Thank you for your attention!