

# Automated Mechanism Design for Correlated Valuations

Vincent Conitzer; joint work with:



**Michael  
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(Duke → UVA)



**Giuseppe  
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(Duke)



**Peter  
Stone**  
(UT Austin)

**Paper:**

**Mechanism Design for  
Correlated Valuations: Efficient  
Methods for Revenue  
Maximization.** Michael Albert,  
Vincent Conitzer, Giuseppe  
Lopomo, and Peter Stone.  
*Operations Research,*  
forthcoming.

# Automated mechanism design input

**Instance** is given by

Set of possible *outcomes*

Set of *agents*

For each agent

set of possible *types*

*probability distribution* over these types

*Objective function*

Gives a value for each outcome for each combination of agents' types

E.g., social welfare, revenue

*Restrictions* on the mechanism

Are *payments* allowed?

Is *randomization* over outcomes allowed?

What versions of *incentive compatibility (IC)* & *individual rationality (IR)* are used?

# How hard is designing an optimal *deterministic* mechanism (without reporting costs)?

[C. & Sandholm UAI'02, ICEC'03, EC'04]

<b>NP-complete</b> (even with 1 reporting agent):	Solvable in <b>polynomial time</b> (for any <i>constant</i> number of agents):
<ol style="list-style-type: none"><li>1. Maximizing social welfare (no payments)</li><li>2. Designer's own utility over outcomes (no payments)</li><li>3. General (linear) objective that doesn't regard payments</li><li>4. Expected revenue</li></ol>	<ol style="list-style-type: none"><li>1. Maximizing social welfare (not regarding the payments) (<b>VCG</b>)</li></ol>

1 and 3 hold even with no IR constraints

# Positive results (randomized mechanisms)

[C. & Sandholm UAI'02, ICEC'03, EC'04]

- Use linear programming

- Variables:

$p(o \mid \theta_1, \dots, \theta_n)$  = probability that outcome  $o$  is chosen given types  $\theta_1, \dots, \theta_n$

(maybe)  $\pi_i(\theta_1, \dots, \theta_n)$  =  $i$ 's payment given types  $\theta_1, \dots, \theta_n$

- Strategy-proofness constraints: for all  $i, \theta_1, \dots, \theta_n, \theta_i'$ :

$$\sum_o p(o \mid \theta_1, \dots, \theta_n) u_i(\theta_i, o) + \pi_i(\theta_1, \dots, \theta_n) \geq$$

$$\sum_o p(o \mid \theta_1, \dots, \theta_i', \dots, \theta_n) u_i(\theta_i, o) + \pi_i(\theta_1, \dots, \theta_i', \dots, \theta_n)$$

- Individual-rationality constraints: for all  $i, \theta_1, \dots, \theta_n$ :

$$\sum_o p(o \mid \theta_1, \dots, \theta_n) u_i(\theta_i, o) + \pi_i(\theta_1, \dots, \theta_n) \geq 0$$

- Objective (e.g., sum of utilities)

$$\sum_{\theta_1, \dots, \theta_n} p(\theta_1, \dots, \theta_n) \sum_i (\sum_o p(o \mid \theta_1, \dots, \theta_n) u_i(\theta_i, o) + \pi_i(\theta_1, \dots, \theta_n))$$

- Also works for BNE incentive compatibility, ex-interim individual rationality notions, other objectives, etc.

- For deterministic mechanisms, can still use mixed integer programming: require probabilities in  $\{0, 1\}$

–Remember typically designing the optimal deterministic mechanism is NP-hard

# A simple example

One item for sale (free disposal)

2 agents, IID valuations: uniform over  $\{1, 2\}$

Maximize expected revenue under ex-interim

IR, Bayes-Nash equilibrium

How much can we get?

(What is optimal expected welfare?)

		<i>Agent 2's valuation</i>	
		1	2
<i>Agent 1's valuation</i>	1	0.25	0.25
	2	0.25	0.25
		probabilities	

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probabilities

Status: OPTIMAL

Objective: obj = 1.5 (MAXimum)

[nonzero variables:]

Our old AMD

solver [C. & Sandholm, 2002, 2003]

gives:

p_t_1_1_o3	1	(probability of disposal for (1, 1))
p_t_2_1_o1	1	(probability 1 gets the item for (2, 1))
p_t_1_2_o2	1	(probability 2 gets the item for (1, 2))
p_t_2_2_o2	1	(probability 2 gets the item for (2, 2))
pi_2_2_1	2	(1's payment for (2, 2))
pi_2_2_2	4	(2's payment for (2, 2))

# A slightly different distribution

One item for sale (free disposal)

2 agents, valuations drawn as on right

Maximize expected revenue under ex-interim

IR, Bayes-Nash equilibrium

How much can we get?

(What is optimal expected welfare?)

		<i>Agent 2's valuation</i>	
		1	2
<i>Agent 1's valuation</i>	1	0.251	0.250
	2	0.250	0.249

probabilities

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probabilities

**Status: OPTIMAL**

**Objective: obj = 1.749 (MAXimum)**

*[some of the nonzero payment variables:]*

**pi\_1\_1\_2            62501**

**pi\_2\_1\_2            -62750**

**pi\_2\_1\_1            2**

**pi\_1\_2\_2            3.992**

*You'd better be really sure about your distribution!*



# A nearby distribution without correlation

One item for sale (free disposal)

2 agents, valuations IID: 1 w/ .501, 2 w/ .499

Maximize expected revenue under ex-interim

IR, Bayes-Nash equilibrium

How much can we get?

(What is optimal expected welfare?)

		<i>Agent 2's valuation</i>	
		1	2
<i>Agent 1's valuation</i>	1	0.251001	0.249999
	2	0.249999	0.249001
		probabilities	

# A nearby distribution without correlation

One item for sale (free disposal)

2 agents, valuations IID: 1 w/ .501, 2 w/ .499

Maximize expected revenue under ex-interim

IR, Bayes-Nash equilibrium

How much can we get?

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**Status:** OPTIMAL

**Objective:** obj = 1.499 (MAXimum)

		<i>Agent 2's valuation</i>	
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		probabilities	

# Cremer-McLean [1985]

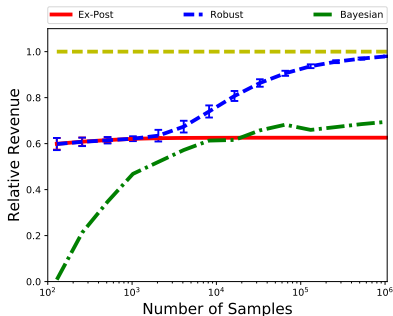
For every agent, consider the following matrix  $\Gamma$  of conditional probabilities, where  $\Theta$  is the set of types for the agent and  $\Omega$  is the set of **signals** (joint types for other agents, or something else observable to the auctioneer)

$$\Gamma = \begin{bmatrix} \pi(1|1) & \cdots & \pi(|\Omega||1) \\ \vdots & \ddots & \vdots \\ \pi(1||\Theta|) & \cdots & \pi(|\Omega|||\Theta|) \end{bmatrix}$$

If  $\Gamma$  has rank  $|\Theta|$  for every agent then the auctioneer can **allocate efficiently** and **extract the full surplus as revenue** (!!)

# Preview of Results

Automated mechanism design procedure that is both *computationally efficient* and *sample efficient* for **maximizing revenue** when there is *sufficient* correlation.



# Outline

- **Known Distributions** (AAAI 2015; 2016)
- **Robust Mechanism Design for Revenue Maximizing Mechanisms** (AAMAS 2017, AAAI 2017, to appear in OR 2021)

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# Problem Description

- A monopolistic seller with one item
- A single bidder with type  $\theta \in \Theta$  and valuation  $v(\theta)$
- An external signal  $\omega \in \Omega$  and distribution  $\pi(\theta, \omega)$





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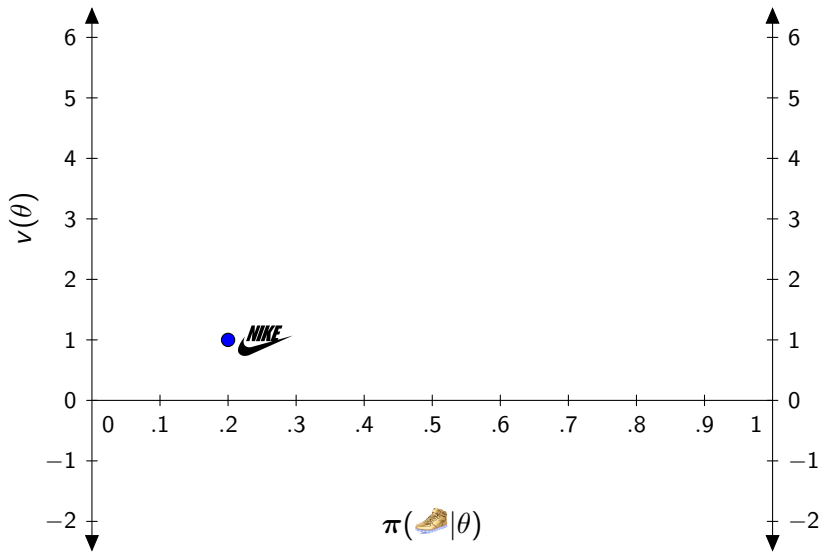
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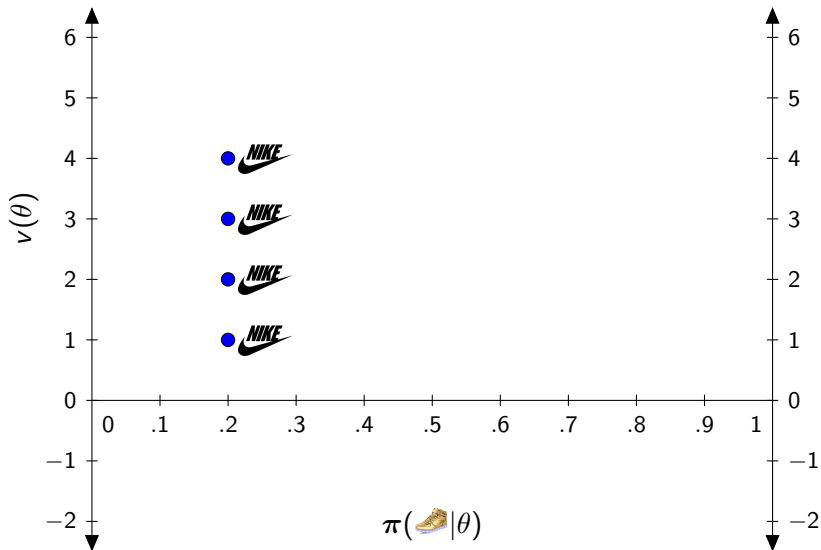
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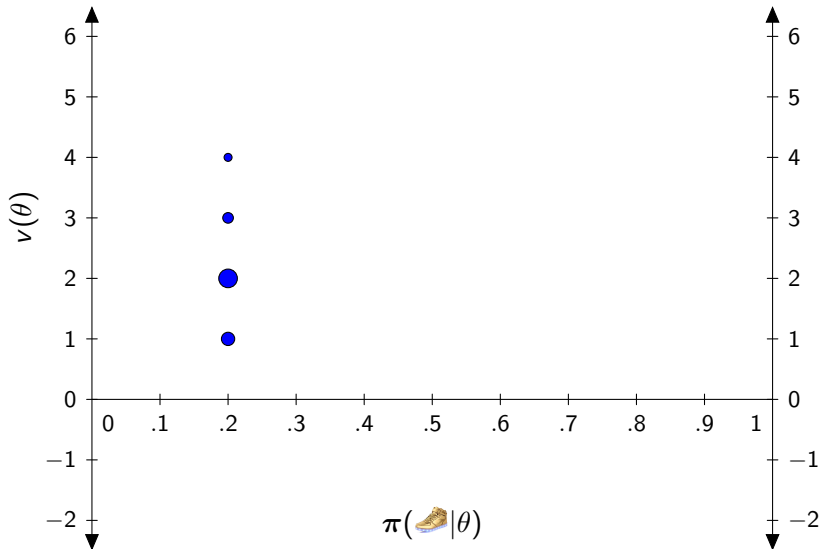
# Mechanism Design with Correlated Distributions



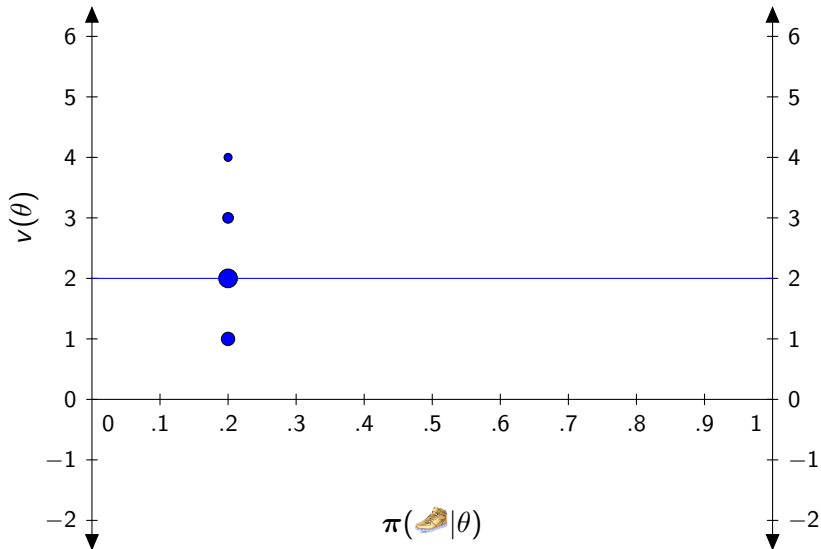
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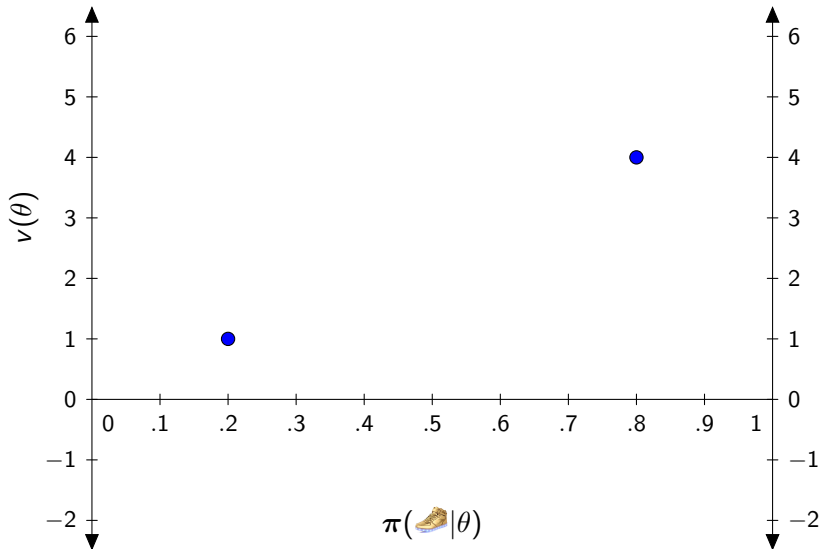
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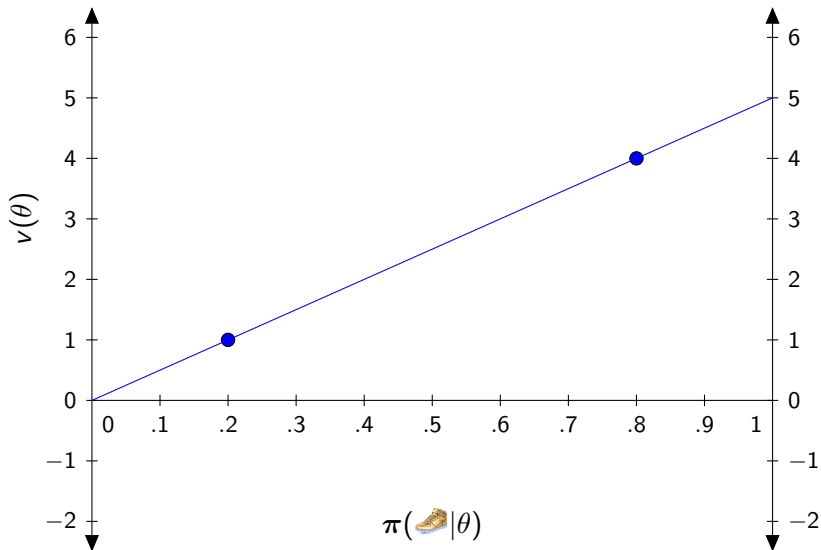
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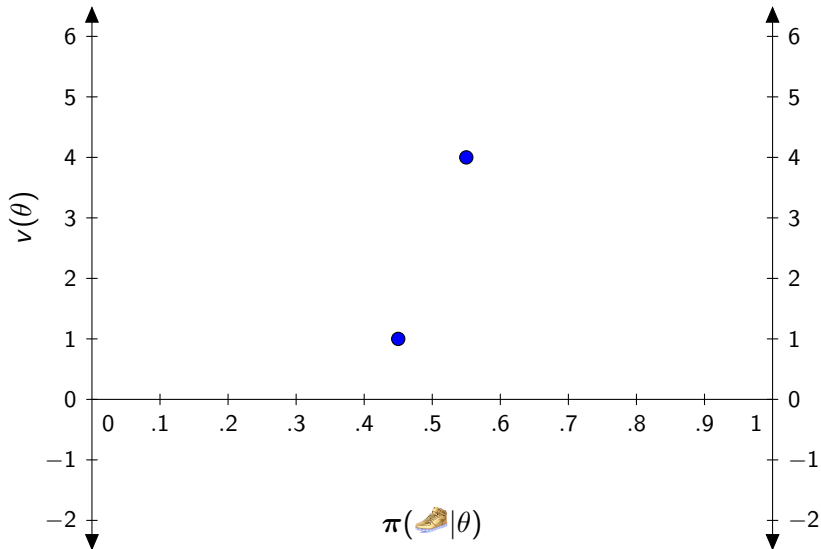


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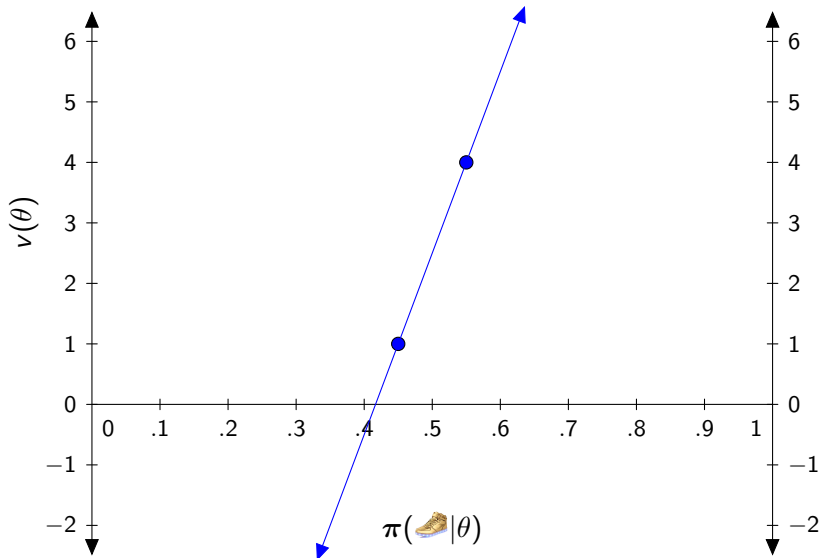




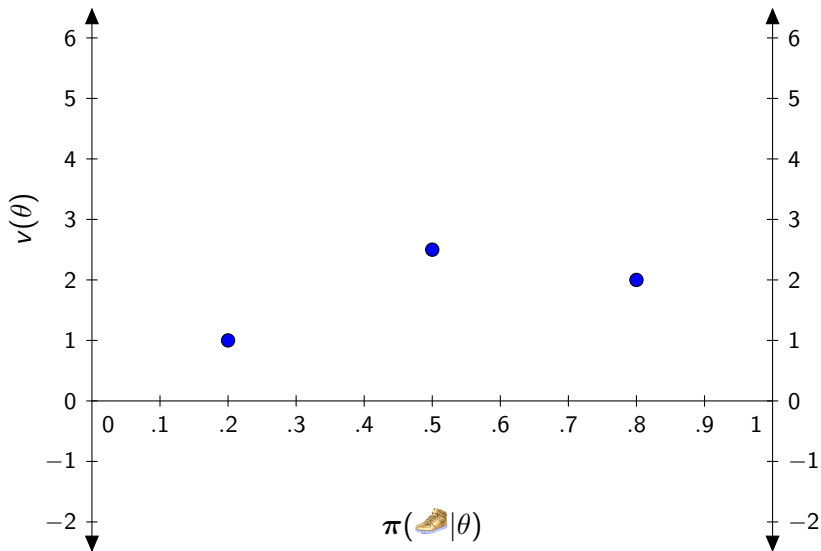
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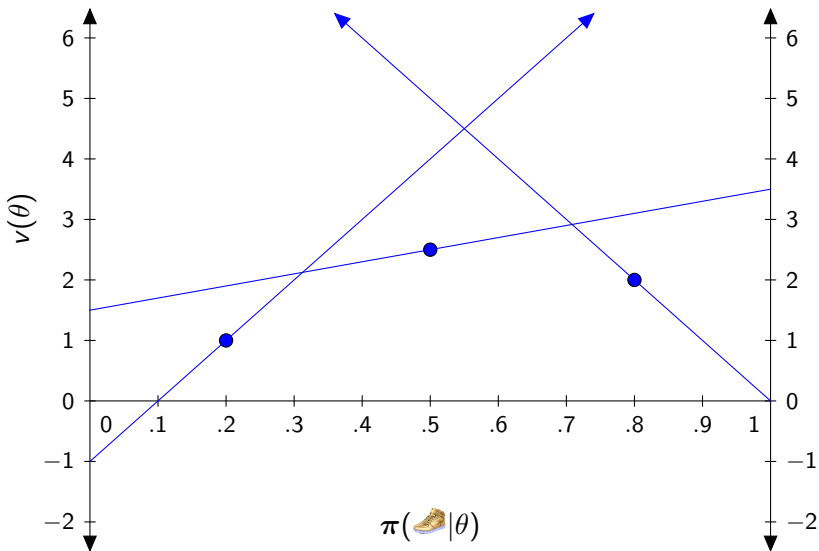
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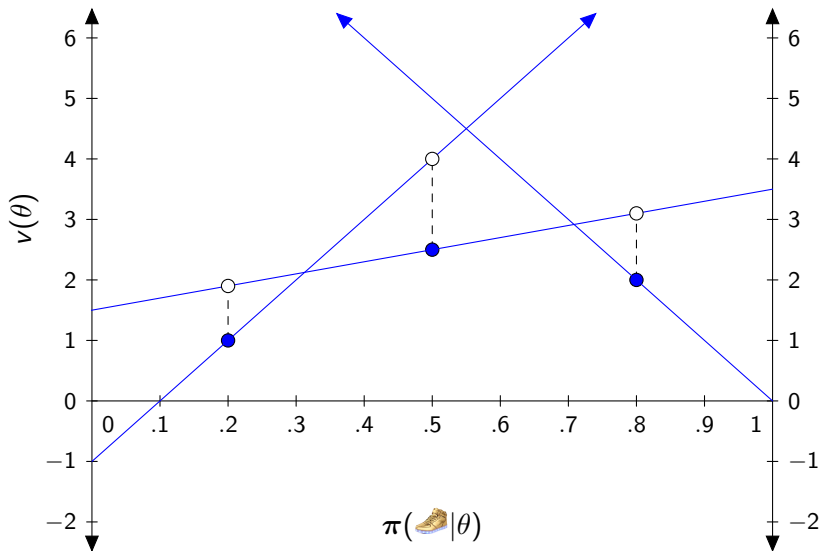
# Necessary and Sufficient Condition



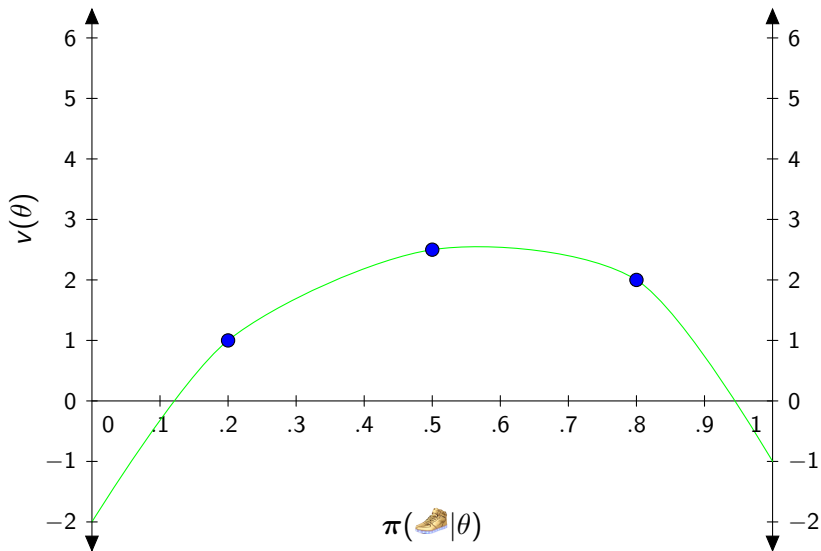
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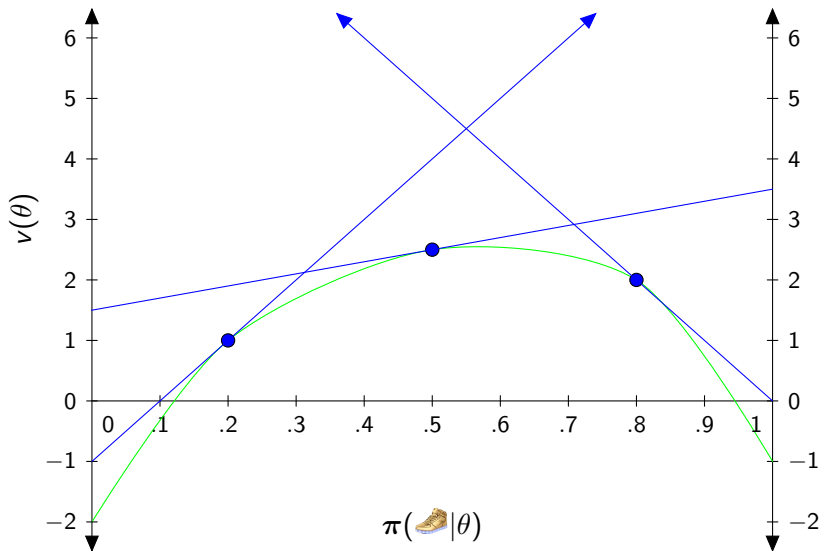
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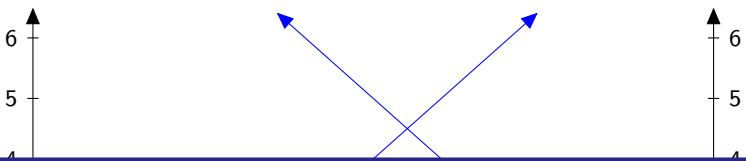
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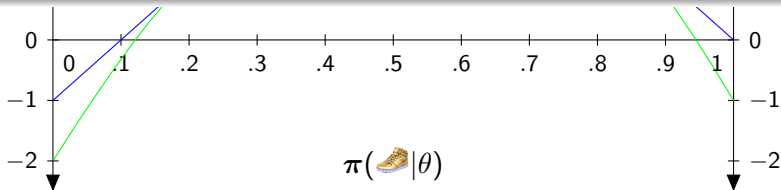


# Necessary and Sufficient Condition



Theorem: Full Surplus Extraction with a Bayesian Mechanism (AAAI 2016)

For a given  $(\pi, \Theta, \Omega)$ , full surplus extraction is possible for a Bayesian mechanism if and only if there exists a concave function  $G : \mathbb{R}^{|\Omega|} \rightarrow \mathbb{R}$  such that  $G(\pi(\bullet|\theta)) = v(\theta)$ . [Full Discussion](#)





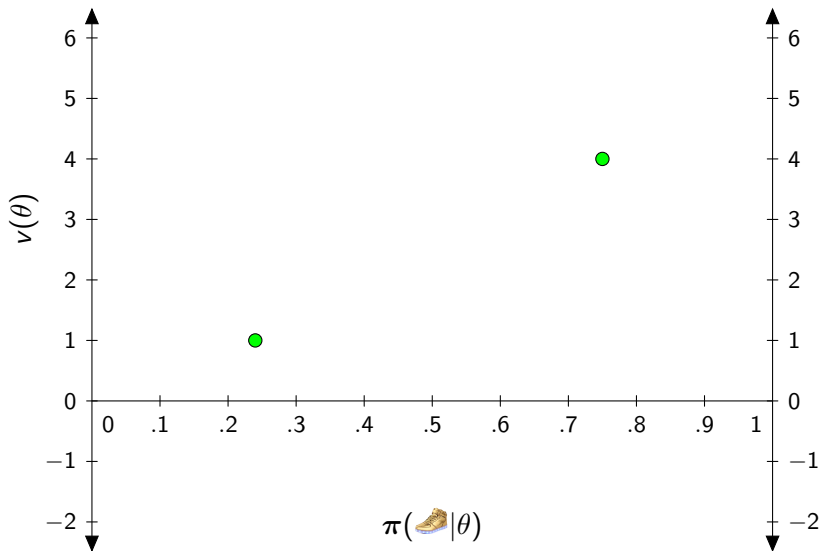
*What if we have access to samples and there is “sufficient” correlation?*

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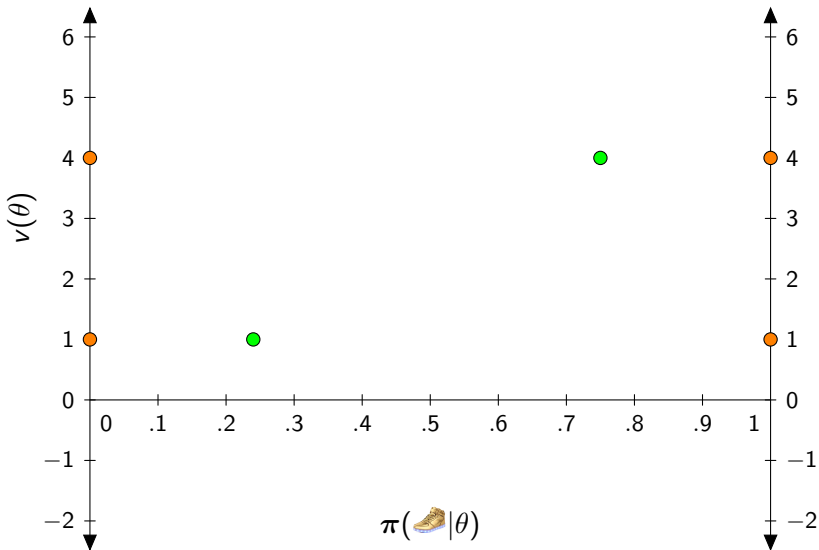
### Theorem: Learning is Impossible (AAMAS 17)

*For any finite number of samples, there exists a distribution for which the optimal learned mechanism is no better than the ex-post mechanism.* [▶ Discussion](#)

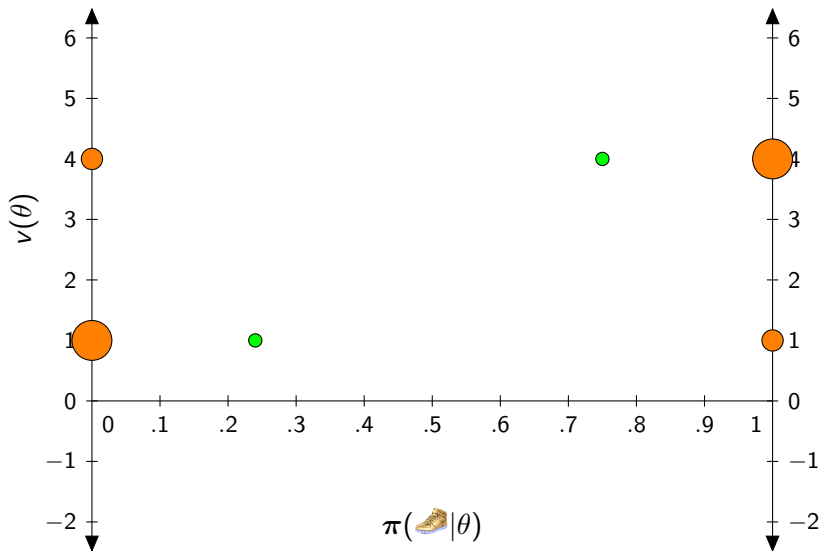
# Consistent Distributions



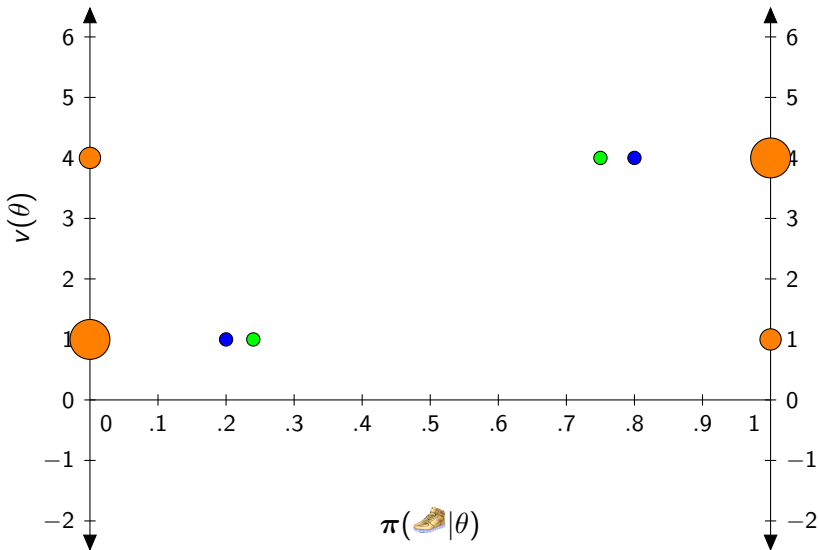
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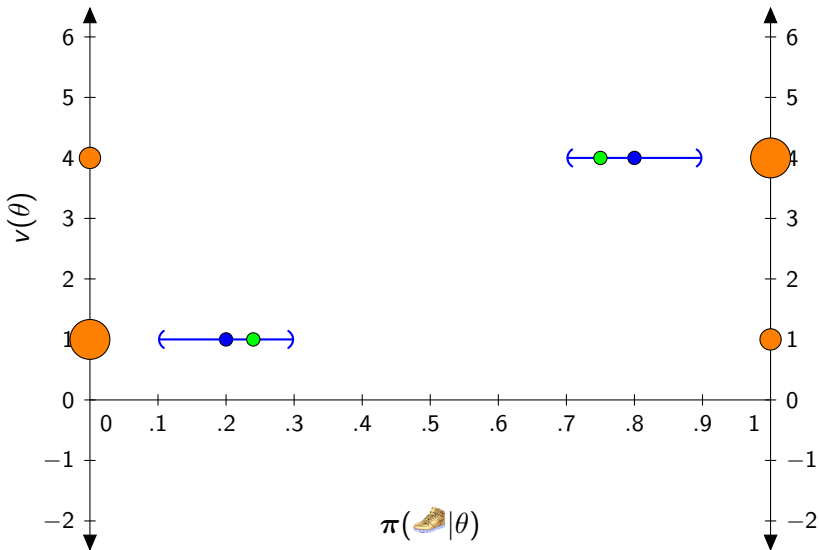
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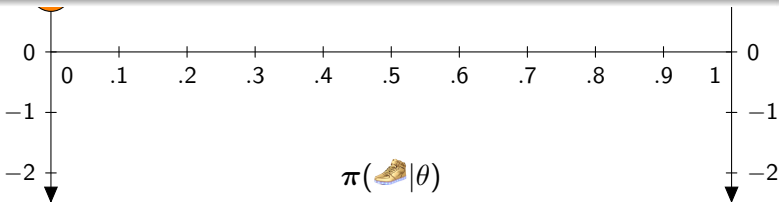


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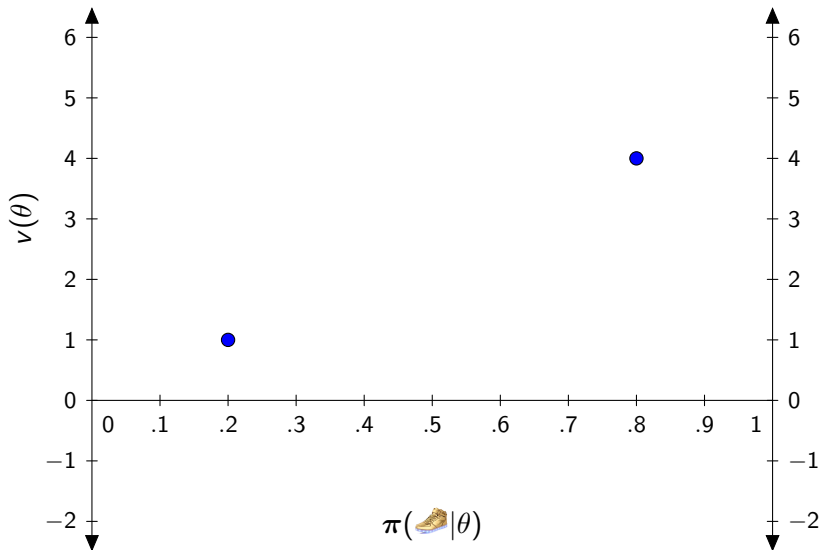
## Definition: Set of Consistent Distributions

A set of distributions,  $\mathcal{P}(\hat{\pi})$ , is a *consistent set of distributions* for the estimated distribution  $\hat{\pi}$  if the true distribution,  $\pi$ , is guaranteed to be in  $\mathcal{P}(\hat{\pi})$  and  $\hat{\pi} \in \mathcal{P}(\hat{\pi})$ .

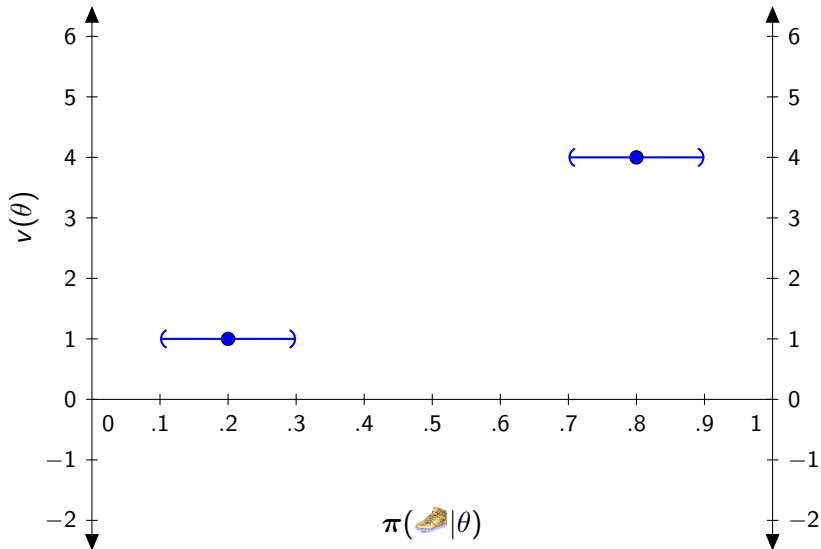




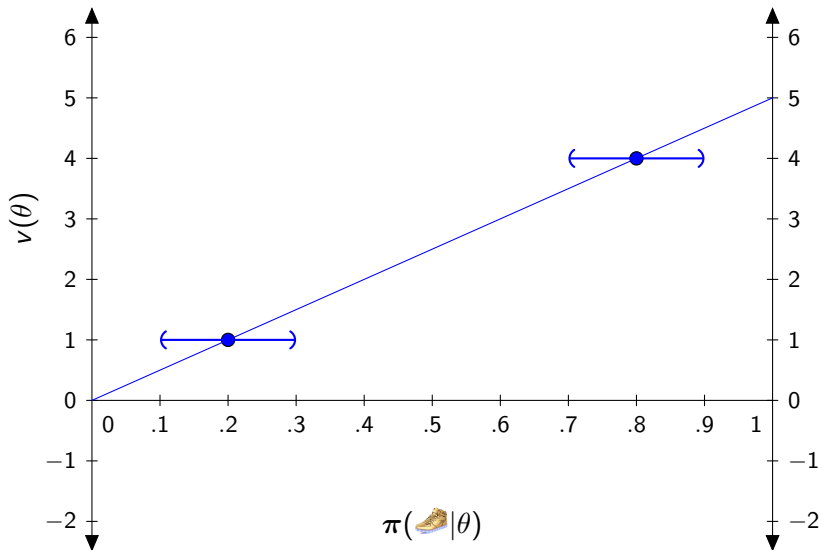
# Robust Mechanism Design



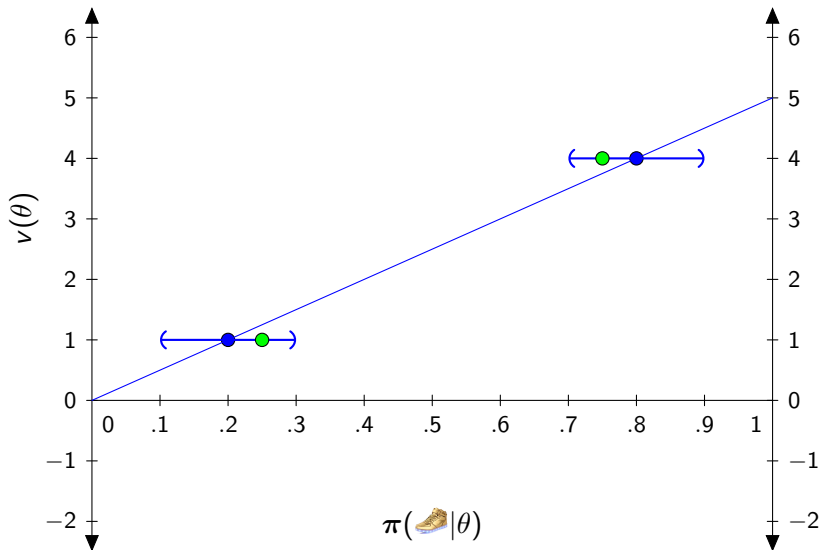
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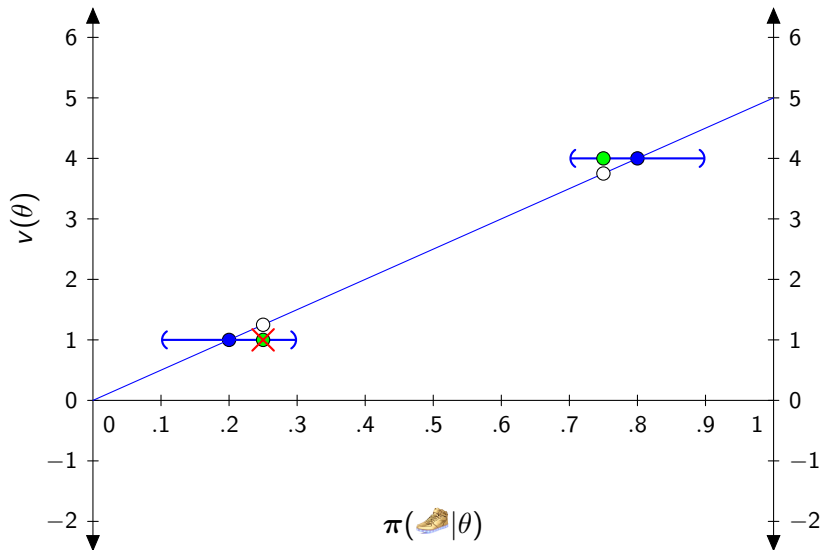
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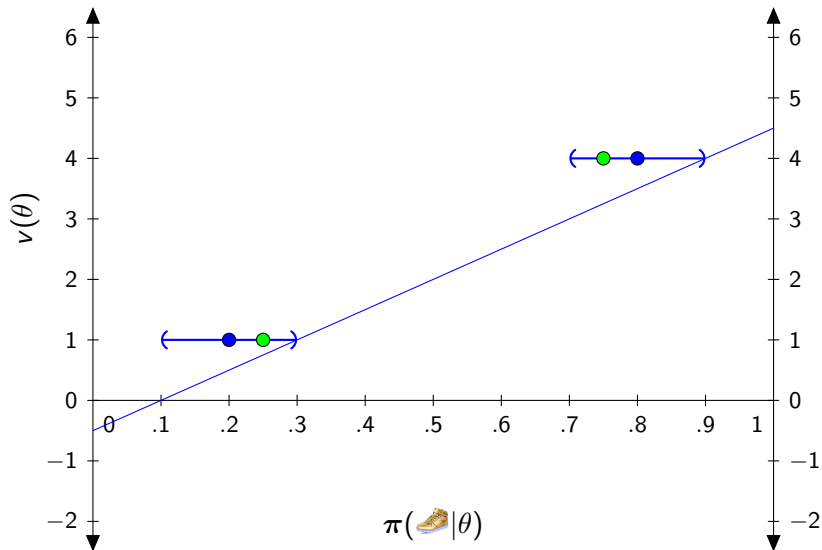
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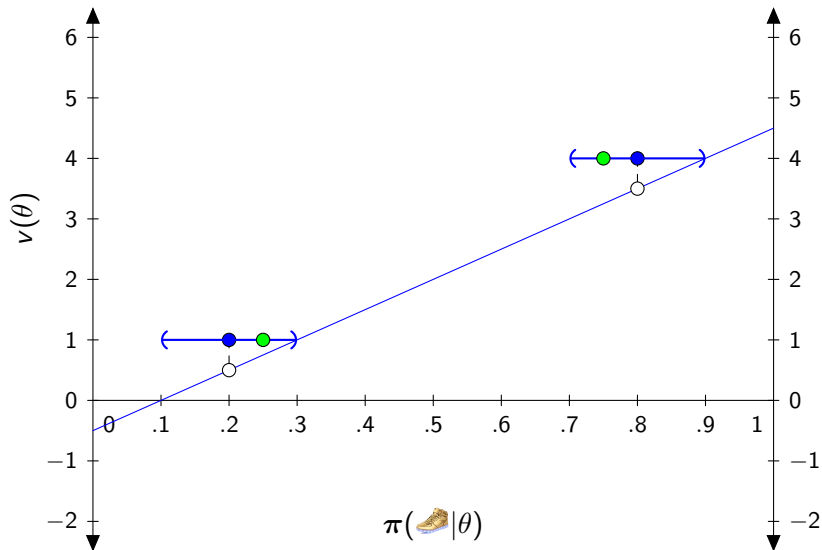
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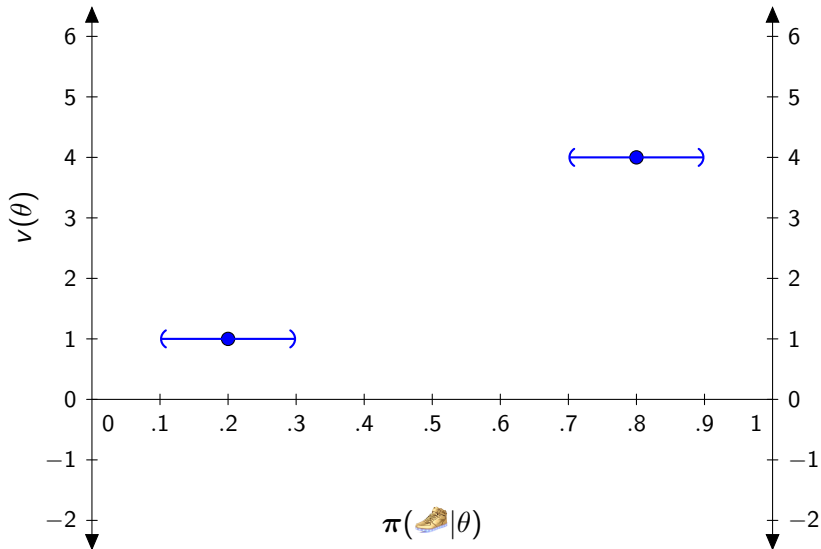
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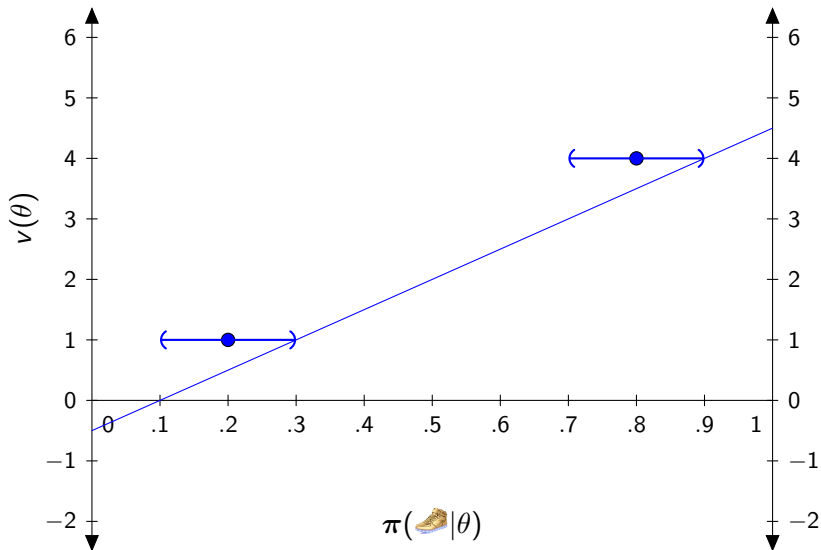


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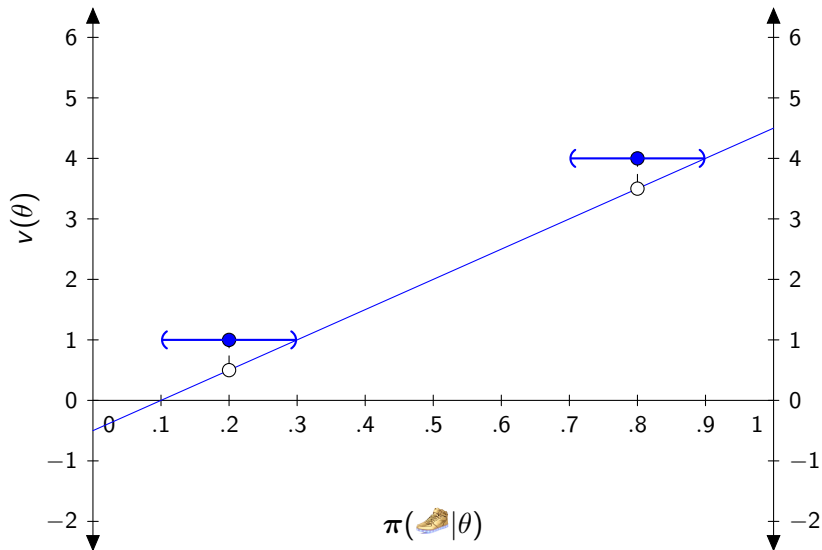




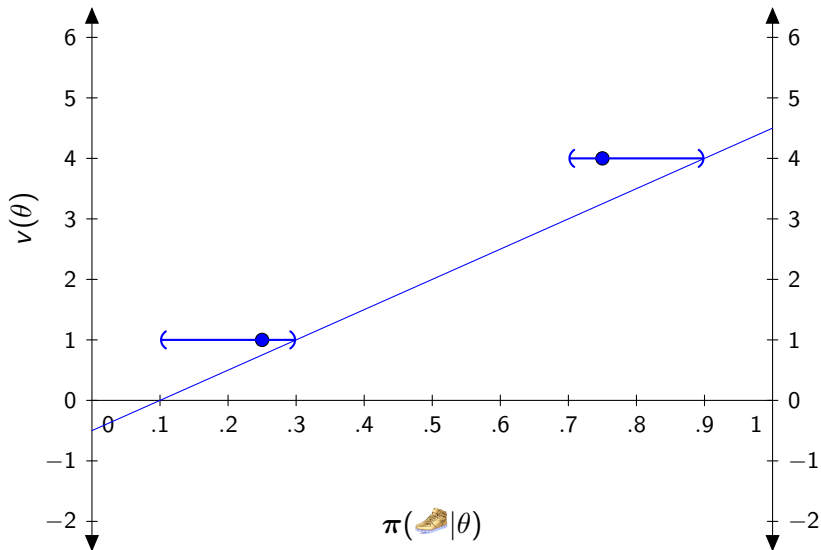
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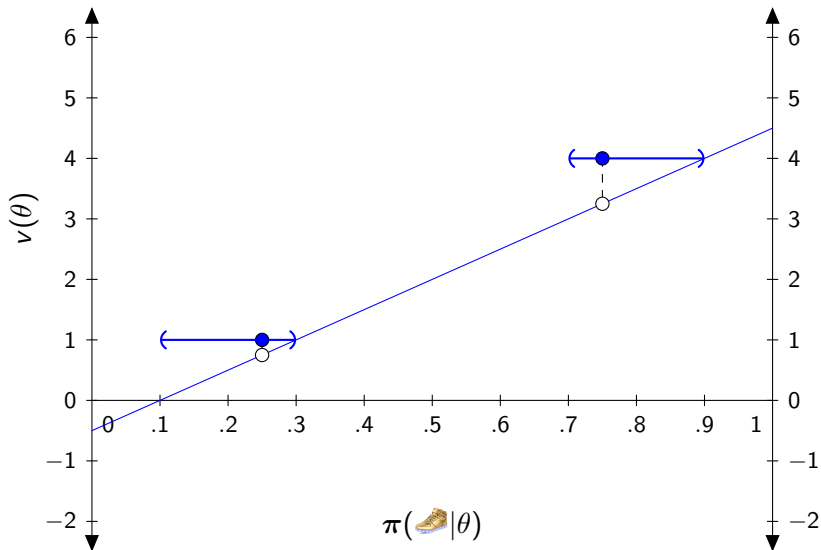
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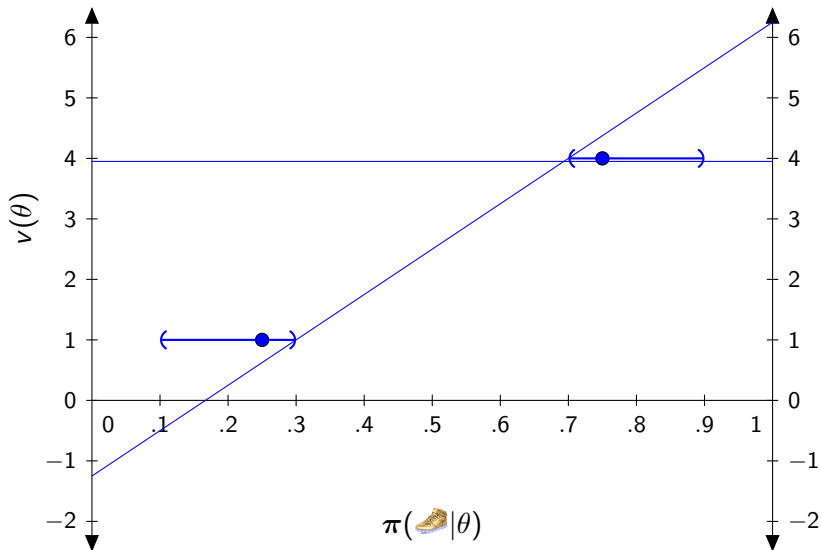
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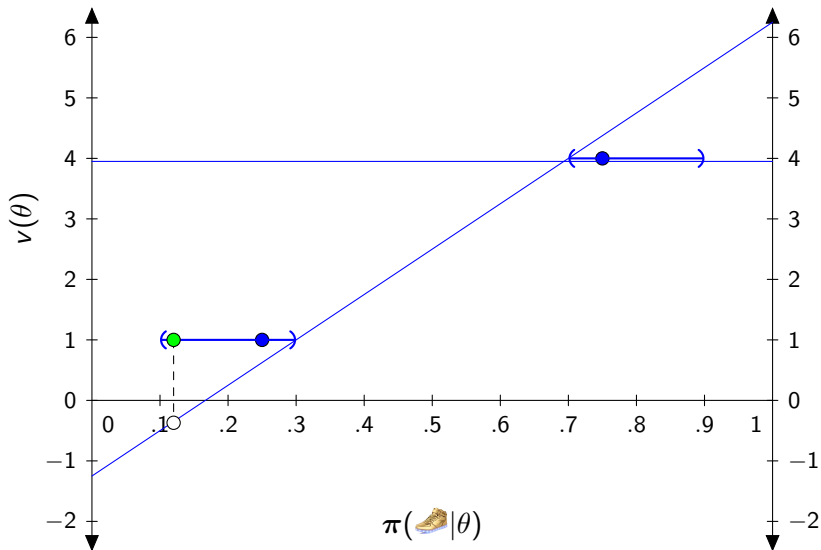
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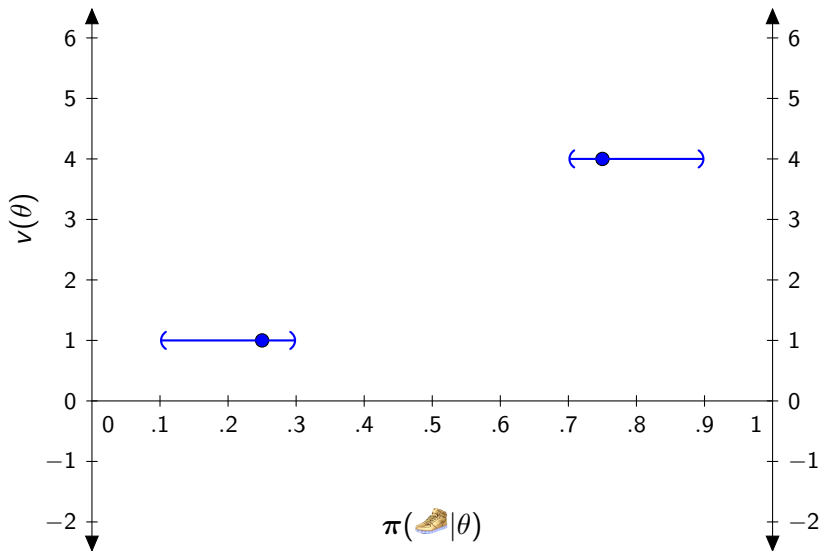
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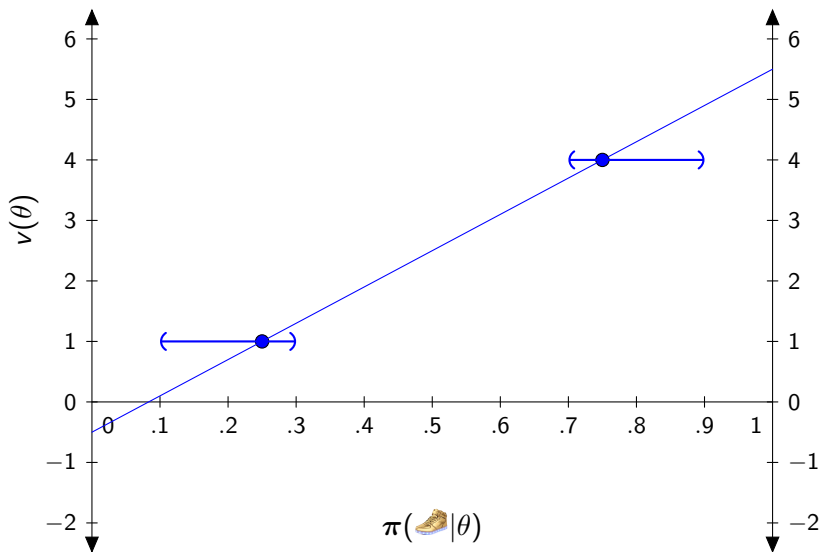
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# Illustration of the Algorithm

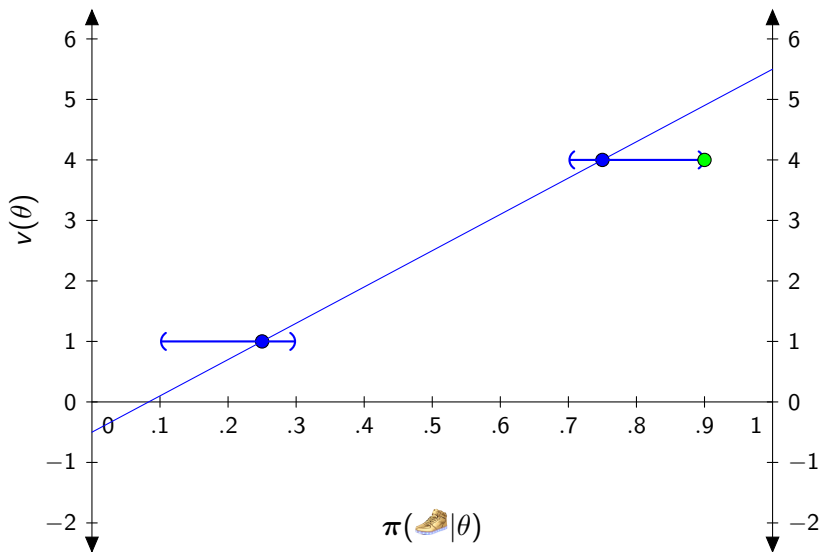


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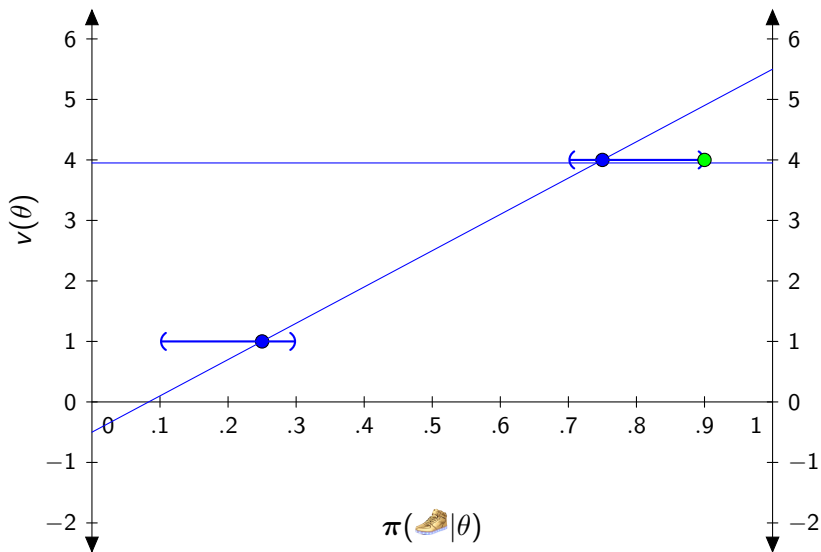




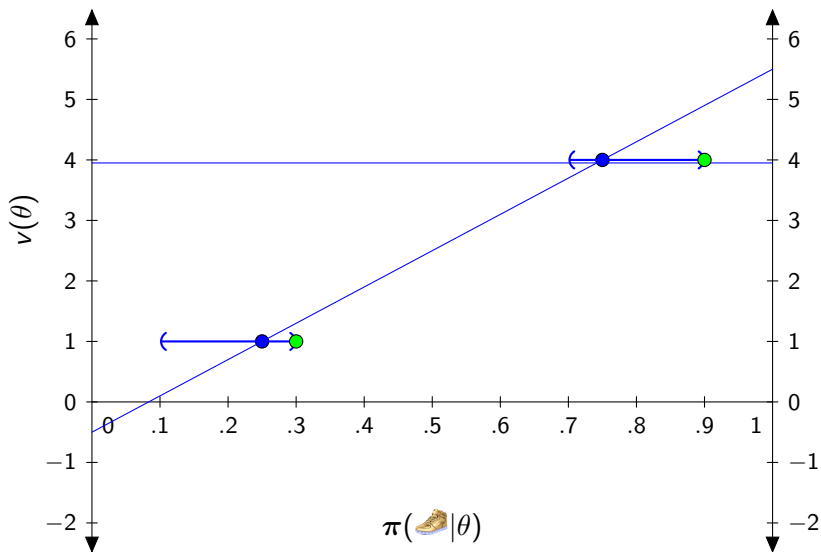
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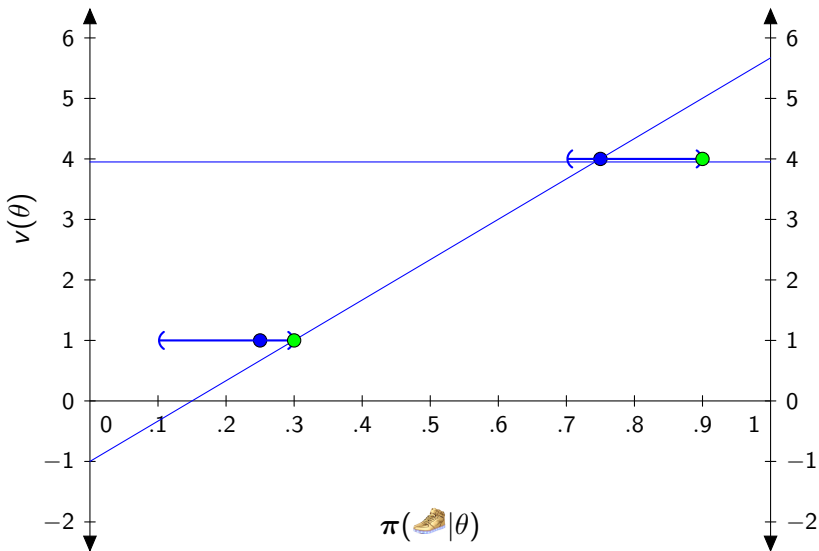
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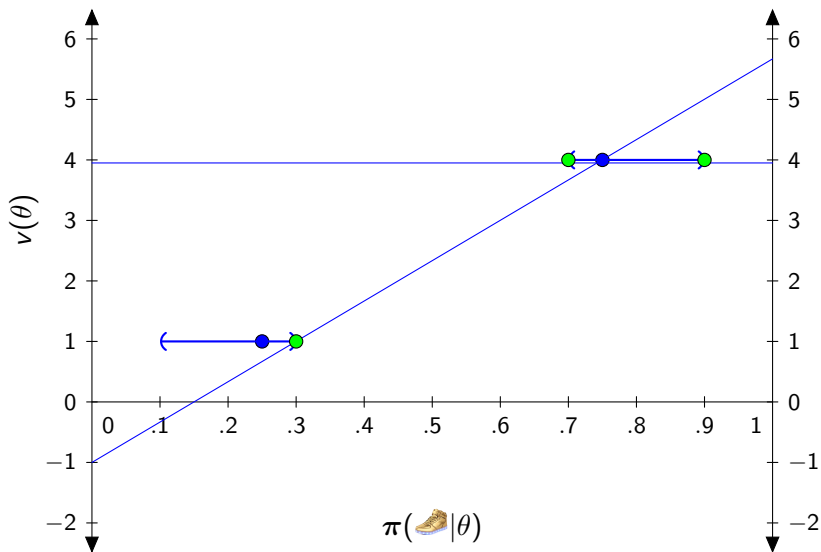
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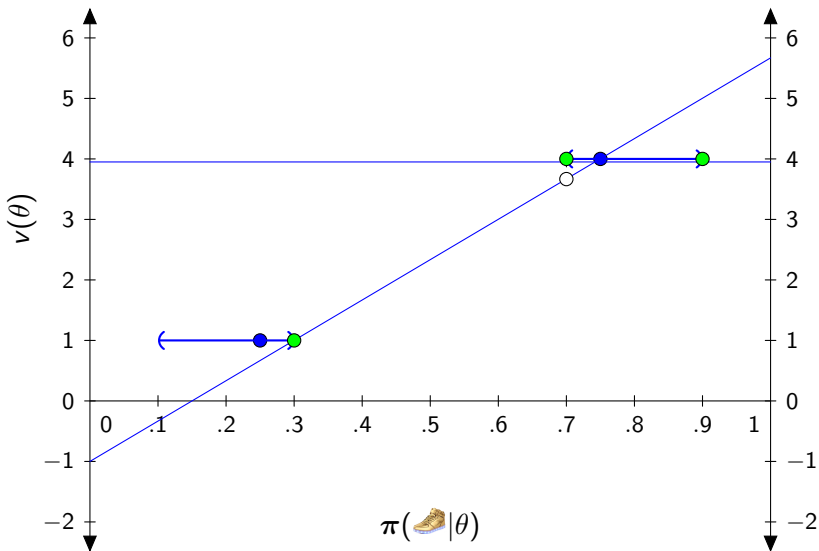
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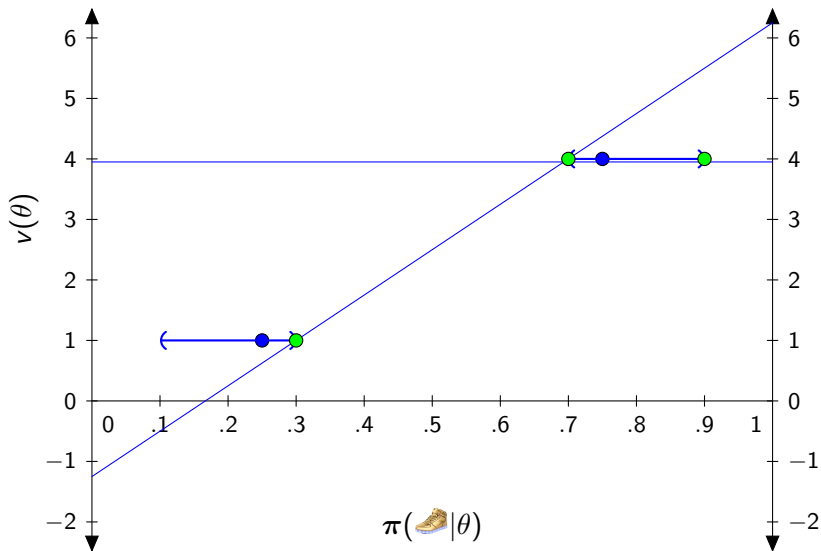
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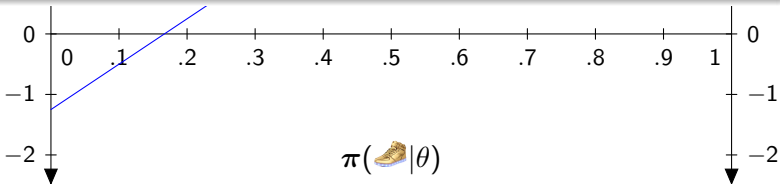


# Illustration of the Algorithm



**Theorem: Polynomial Complexity of the Optimal Robust Mechanism**

*The optimal robust mechanism can be calculated in **time polynomial** in the number of types of the bidder and external signal.*





# $\epsilon$ -Robust Mechanism Design

*Robust is not sufficient*

- All results and intuition for robust mechanism design carries over to restricted  $\epsilon$ -robust mechanism design

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## *Robust is not sufficient*

### Definition: Set of $\epsilon$ -Consistent Distributions

A set of distributions,  $\mathcal{P}_\epsilon(\hat{\pi})$ , is an  $\epsilon$ -consistent set of distributions for the estimated distribution  $\hat{\pi}$  if the true distribution,  $\pi$ , is in  $\mathcal{P}_\epsilon(\hat{\pi})$  with probability  $1 - \epsilon$  and  $\hat{\pi} \in \mathcal{P}_\epsilon(\hat{\pi})$ .

- All results and intuition for robust mechanism design carries over to restricted  $\epsilon$ -robust mechanism design

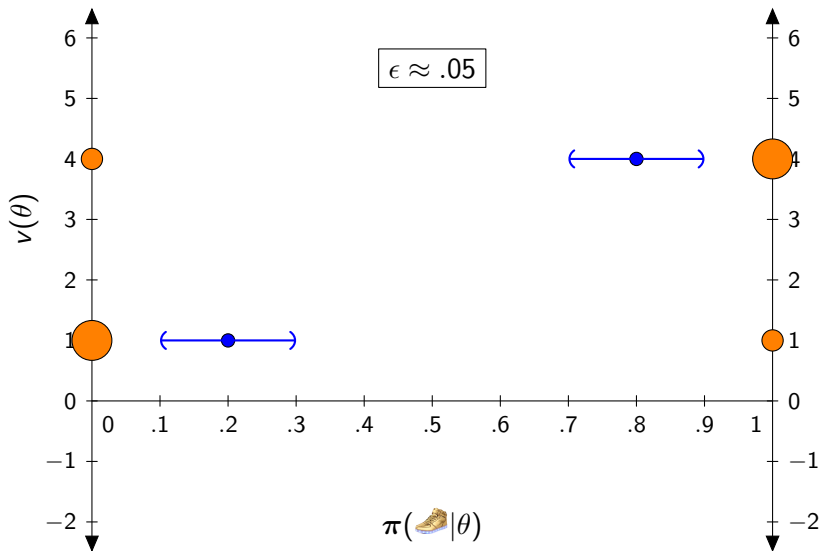
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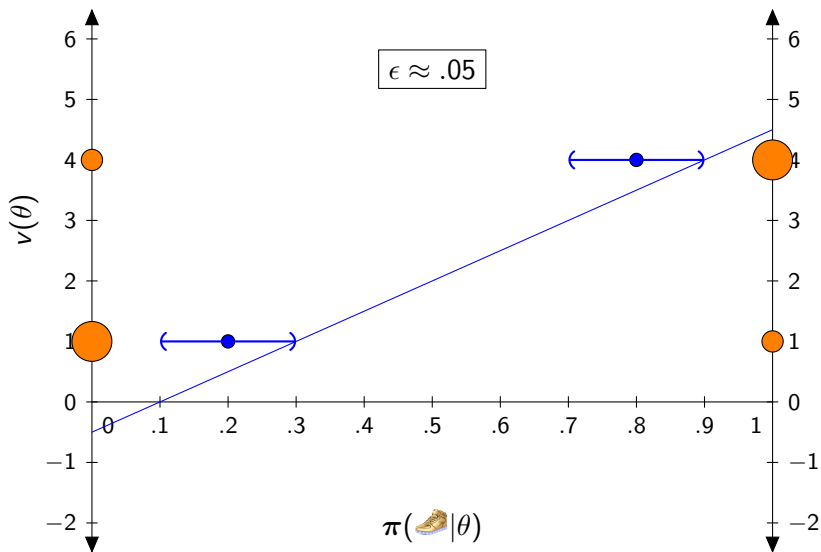
## *Robust is not sufficient*

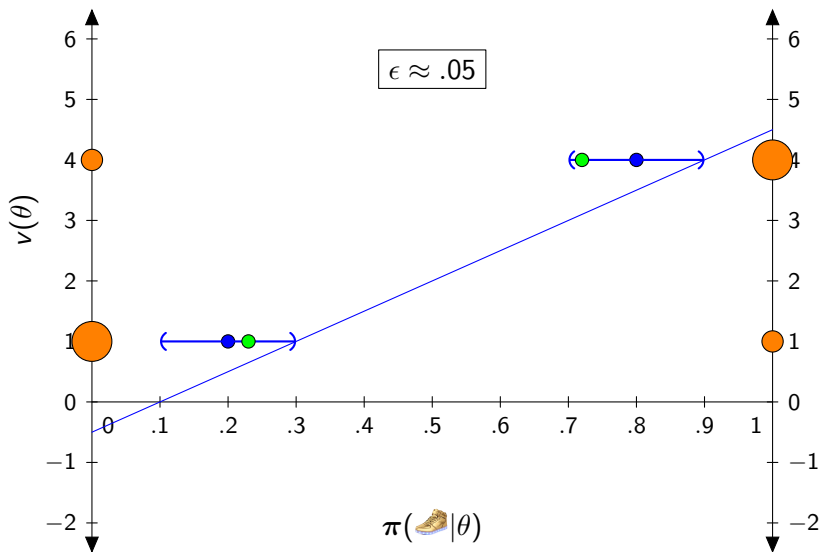
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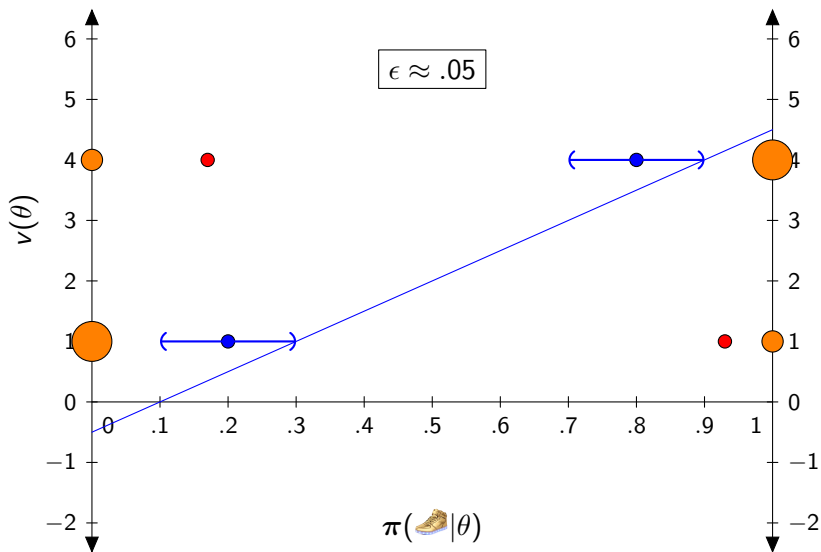
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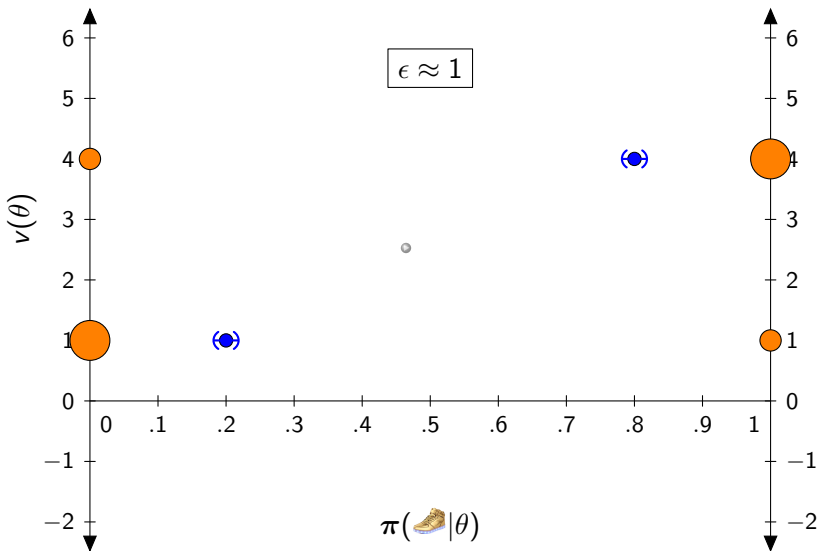
- All results and intuition for robust mechanism design carries over to restricted  $\epsilon$ -robust mechanism design

Parameterized Bayesian IC and IR with  $\epsilon$ 

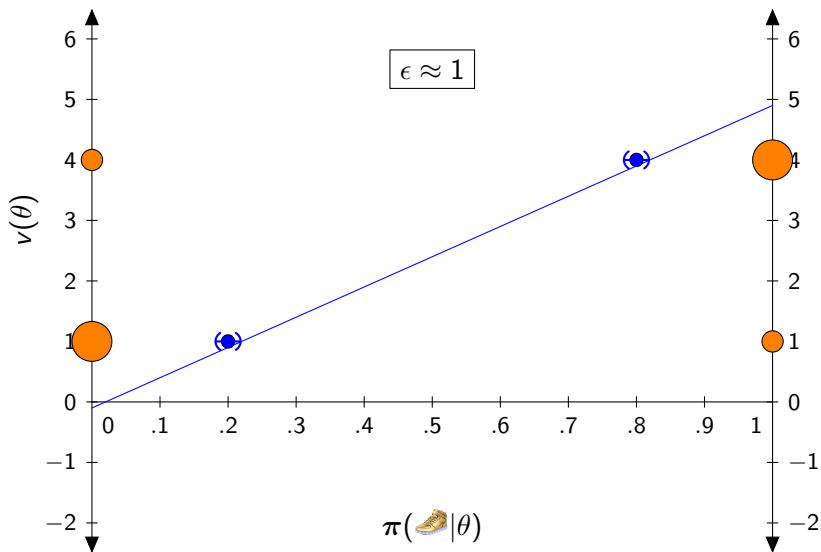
Parameterized Bayesian IC and IR with  $\epsilon$ 

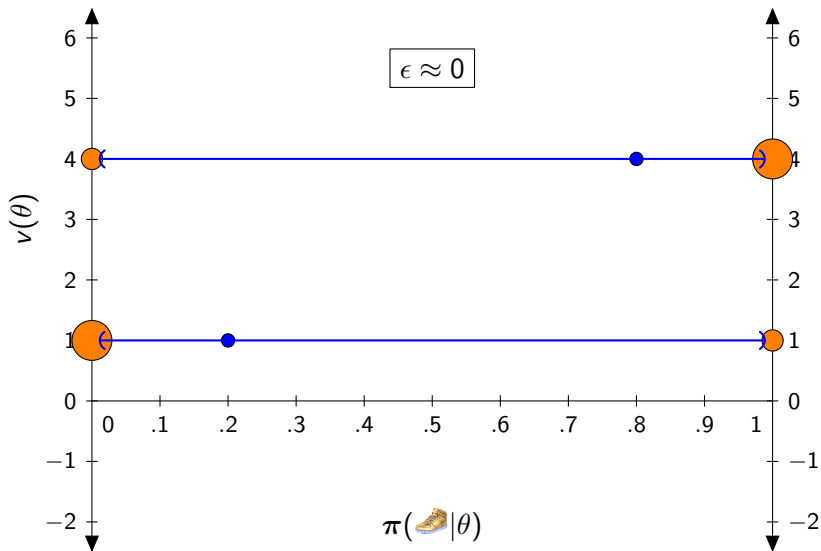
Parameterized Bayesian IC and IR with  $\epsilon$ 

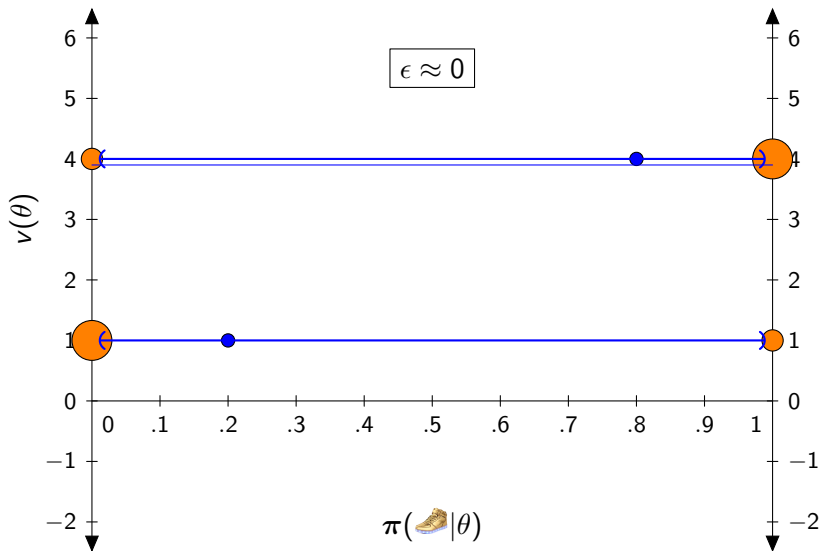
Parameterized Bayesian IC and IR with  $\epsilon$ 

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# Revenue Guarantee for $\epsilon$ -Robust Mechanism Design

*How do  $\epsilon$ -robust mechanisms perform?*

- By our impossibility result, may perform arbitrarily badly
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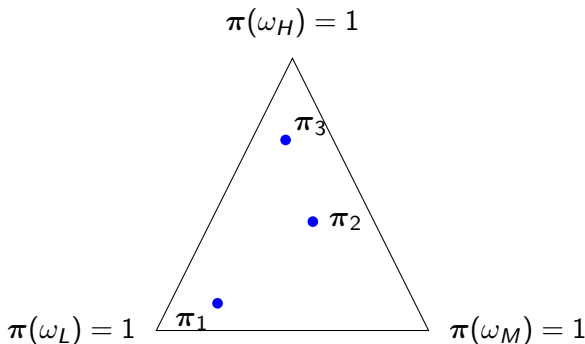
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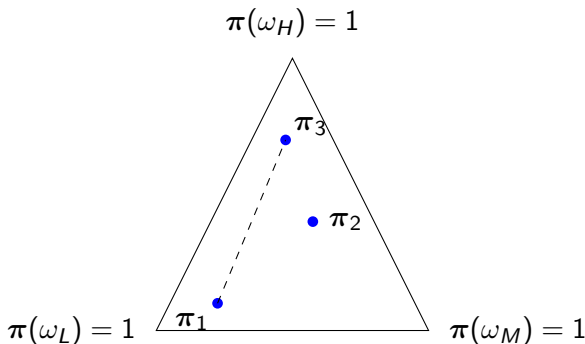
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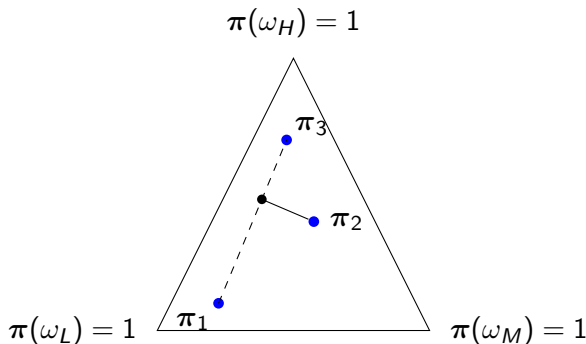




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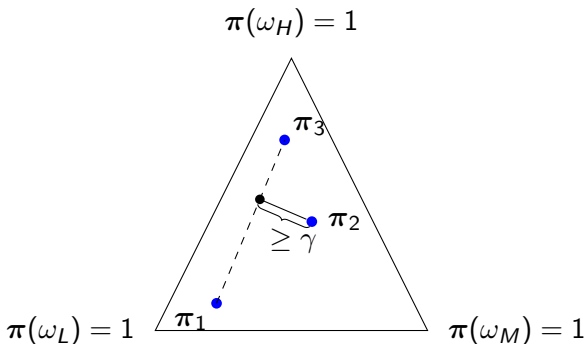
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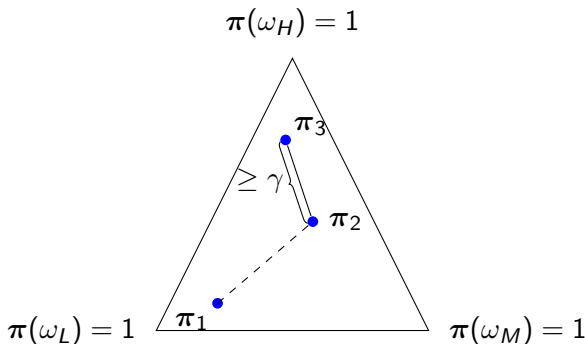
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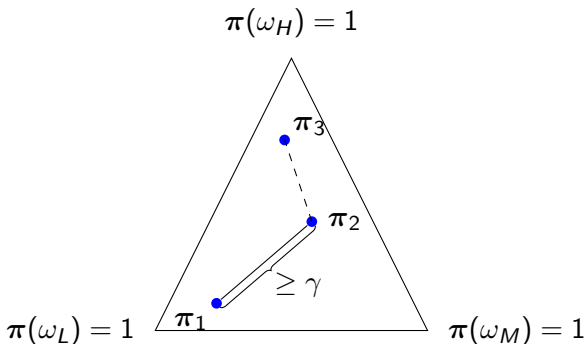
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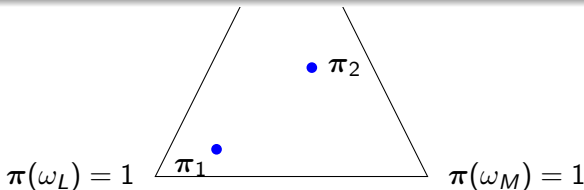
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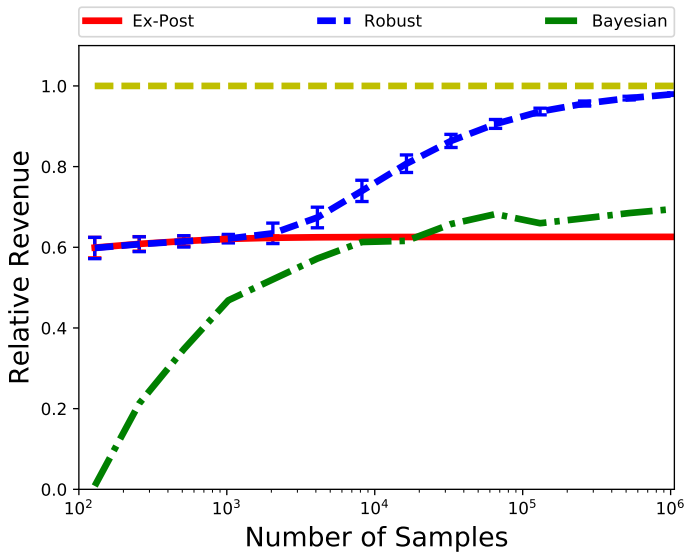
## Sample Complexity of $\epsilon$ -Robust Mechanism Design

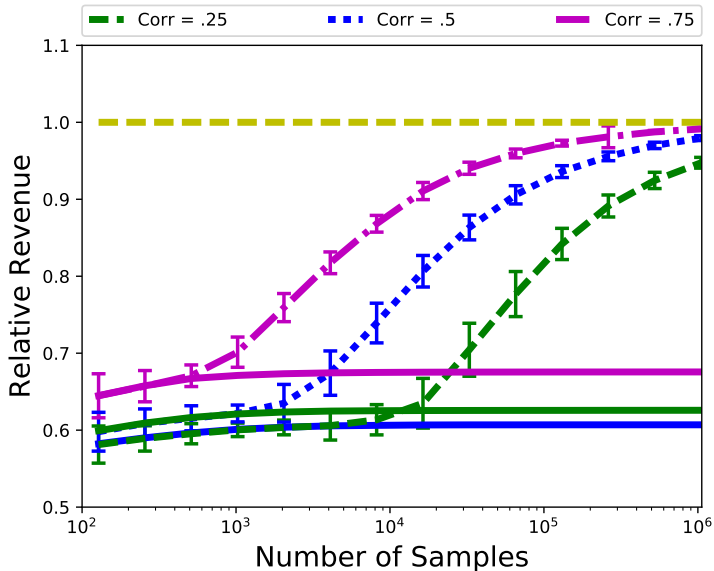
The **sample complexity** for constructing an  $\epsilon$ -robust mechanism that is an **additive  $k$ -approximation** to the optimal mechanism is  $O(\text{poly}(\frac{1}{k}, \frac{1}{\gamma}, |\Theta|, |\Omega|, \mathbf{v}(|\Theta|)))$ . [▶ Proof](#)



# Simulations

- True distribution is discretized bivariate normal distribution
- Sample from the true distribution  $N$  times
- Use Bayesian methods to estimate the distribution
- Calculate empirical confidence intervals for elements of the distribution
- Parameters unless otherwise specified:
  - Correlation = .5
  - $\epsilon = .05$
  - $\Theta = \{1, 2, \dots, 10\}$
  - $|\Omega| = 10$
  - $v(\theta) = \theta$



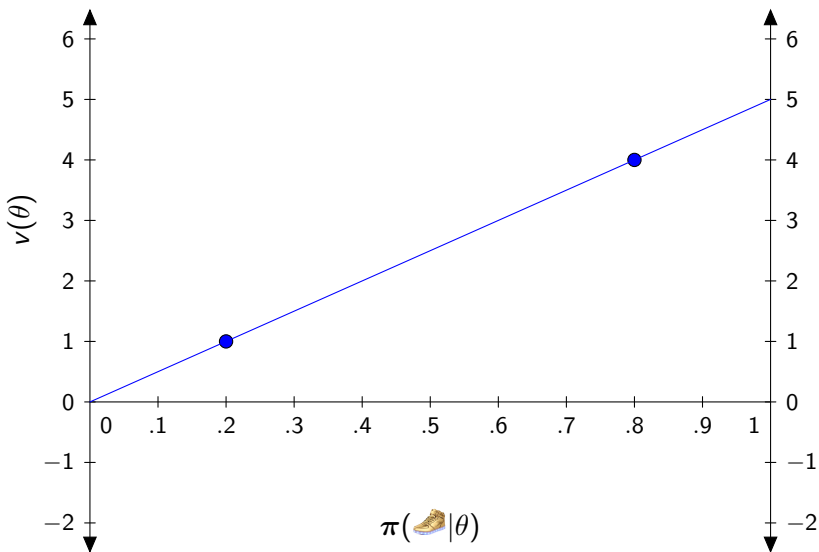




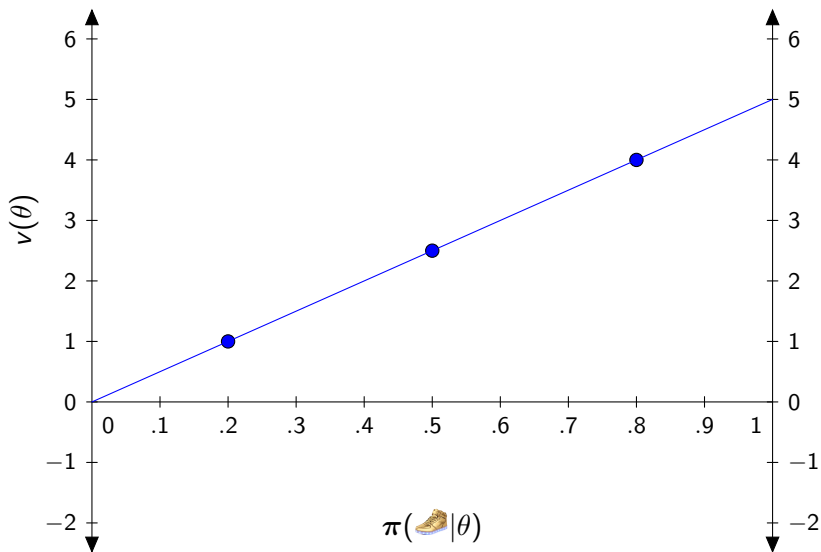
## Related Work for Robust Mechanism Design

- Robust mechanisms (Bergemann and Morris 2005)
- Unknown Correlated Distributions (Lopomo, Rigotti, and Shannon 2009, Fu, Haghpanah, Hartline, and Kleinberg 2014)
- Automated Mechanism Design (Conitzer and Sandholm 2002, 2004; Guo and Conitzer 2010; Sandholm and Likhodedov 2015)
- Robust Optimization (Bertsimas and Sim 2004; Aghassi and Bertsimas 2006)
- Learning Bidder Distributions (Elkind 2007, Blume et. al. 2015, Morgenstern and Roughgarden 2015)

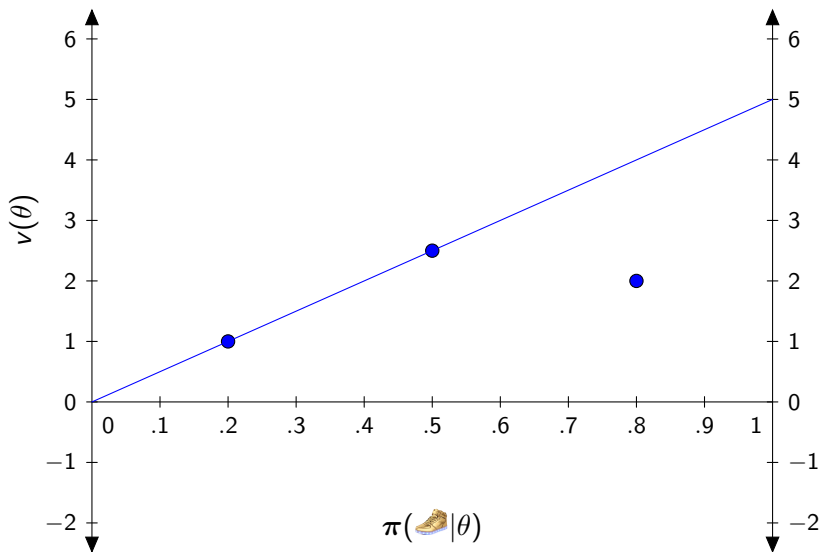
# Full Revenue Extraction with an Ex-Post IC Mechanism



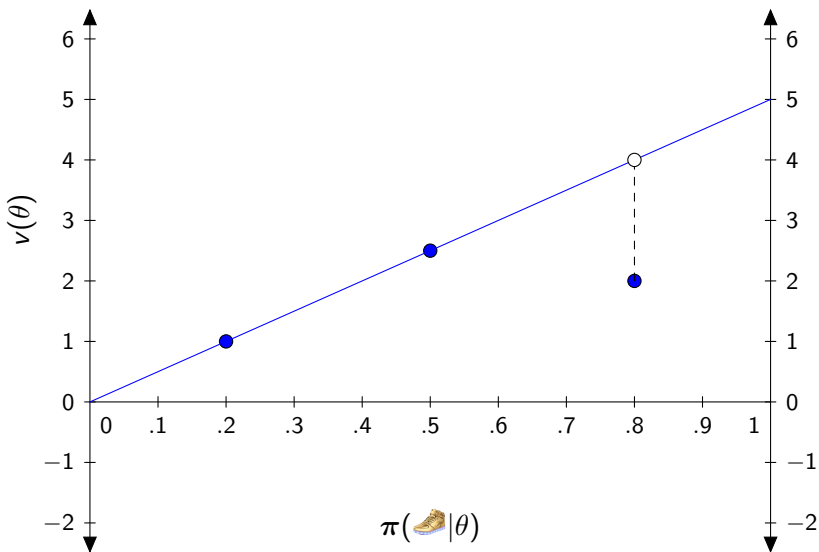
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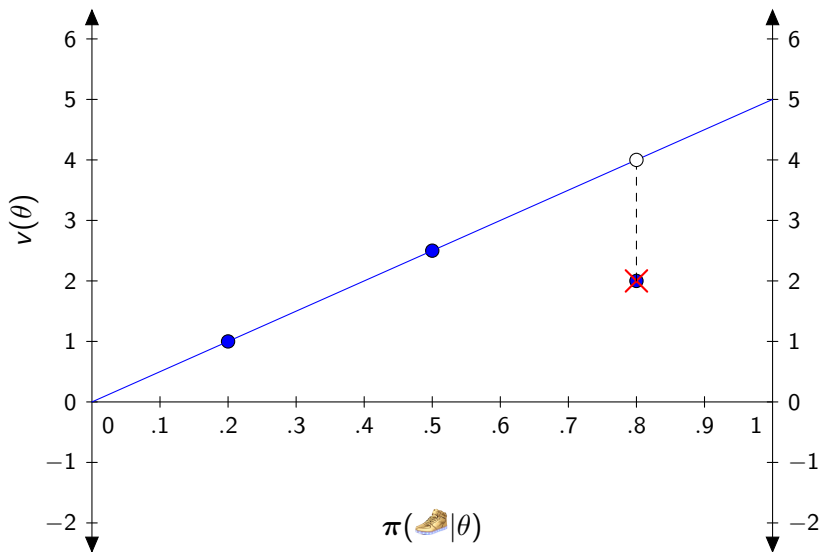
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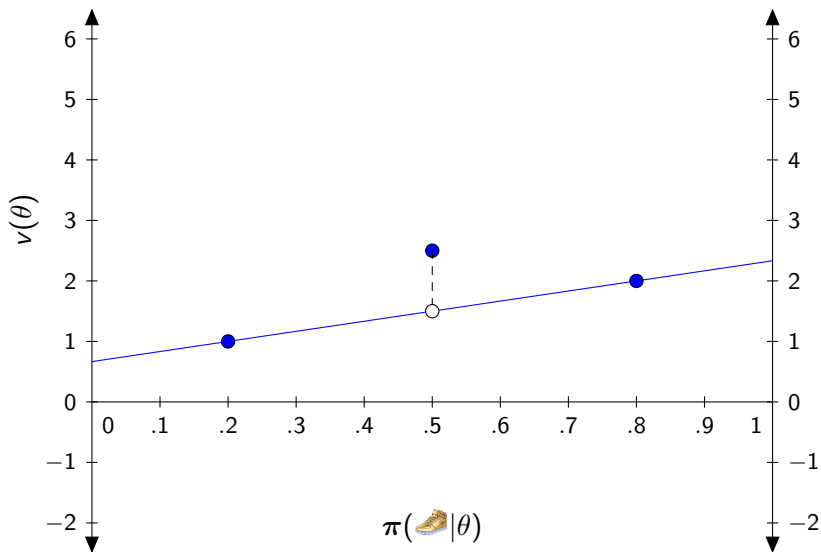
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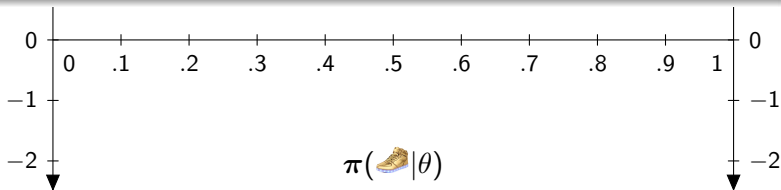
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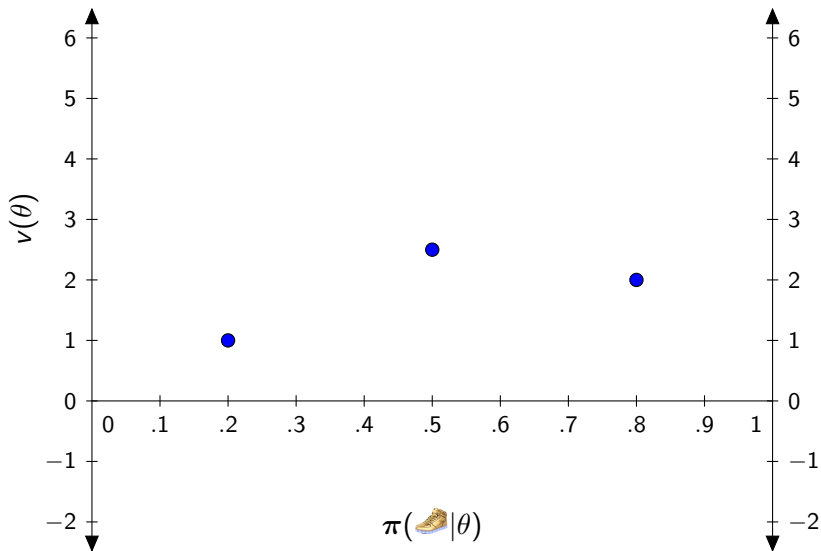
Theorem: Full Revenue Extraction with an Ex-Post IC Mechanism (Albert, Conitzer, Lopomo 2016)

*For a given  $(\pi, \Theta, \Omega)$ , full surplus extraction is possible for a Cremer-McLean mechanism if and only if there exists a linear function  $G : \mathbb{R}^{|\Omega|} \rightarrow \mathbb{R}$  such that  $G(\pi(\bullet|\theta)) = v(\theta)$ .*

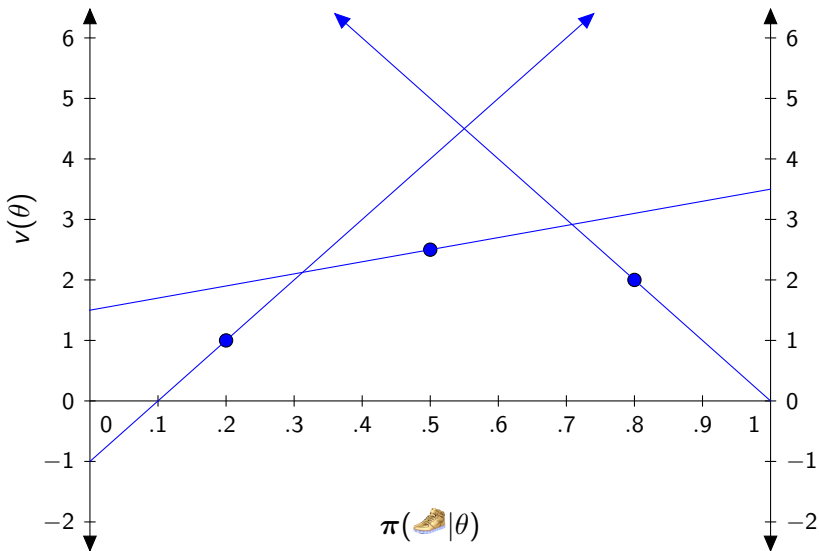




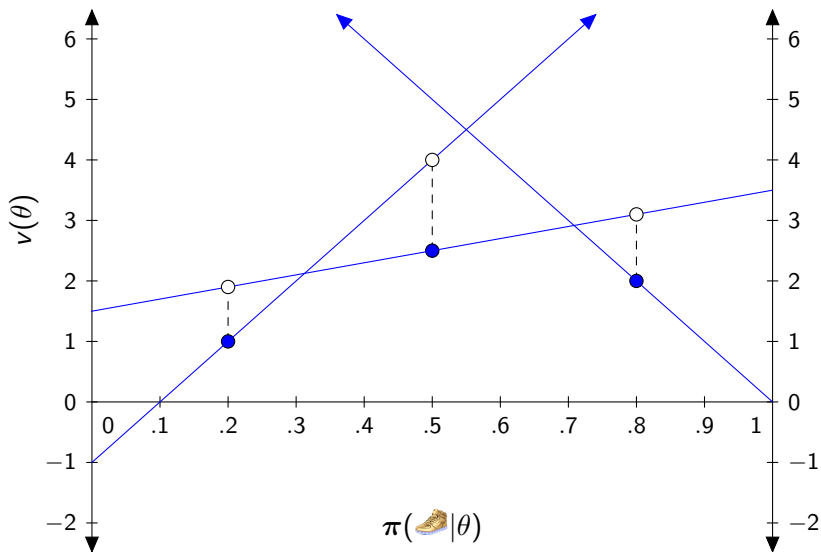
# Full Revenue Extraction with a Bayesian Mechanism



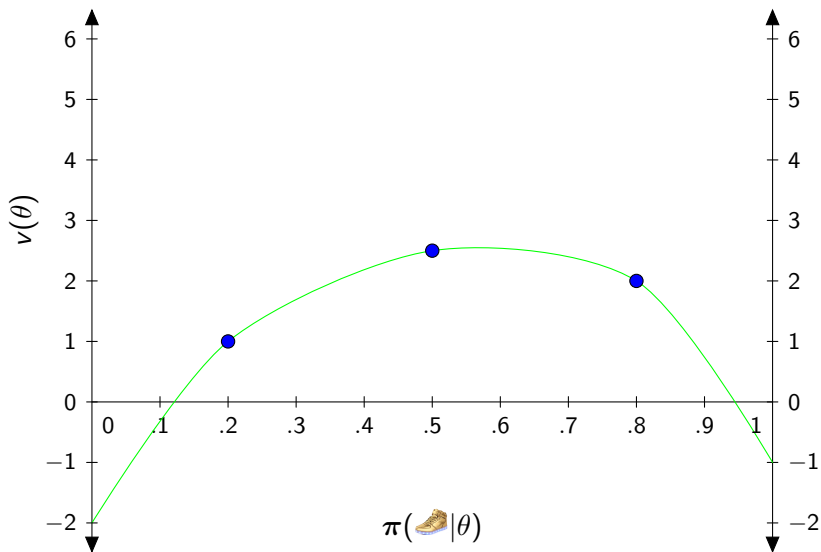
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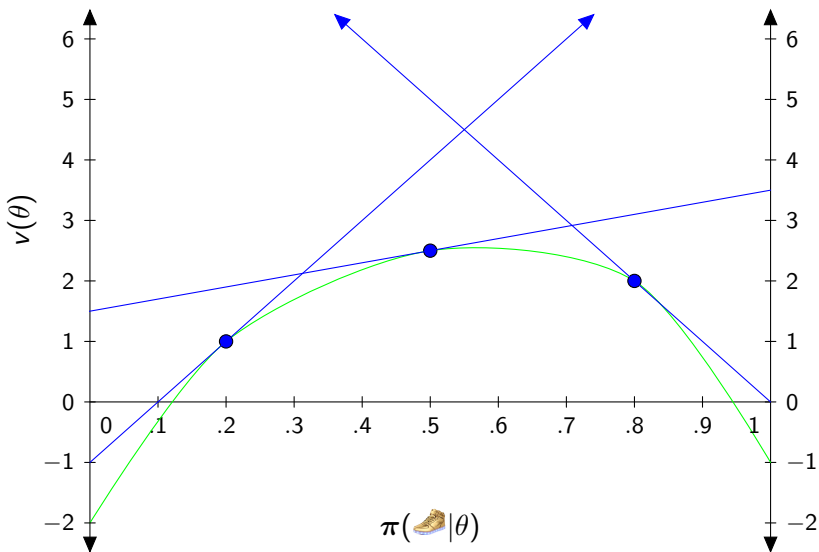
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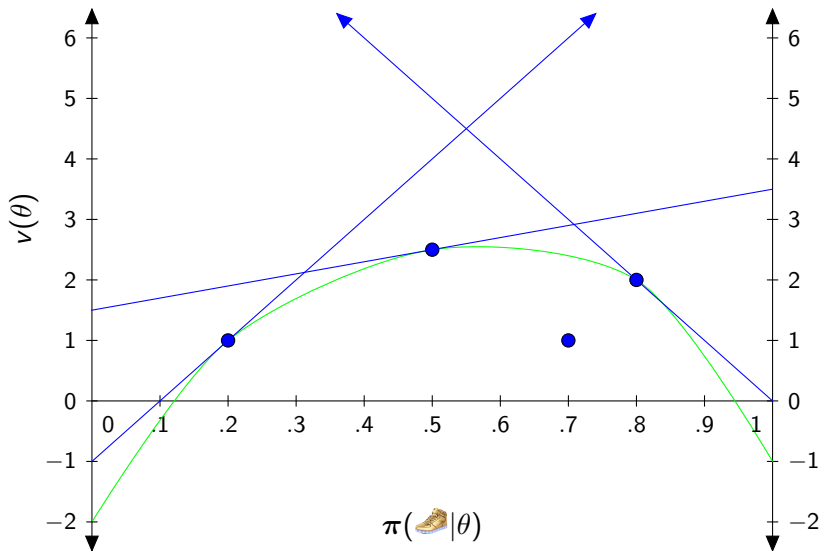
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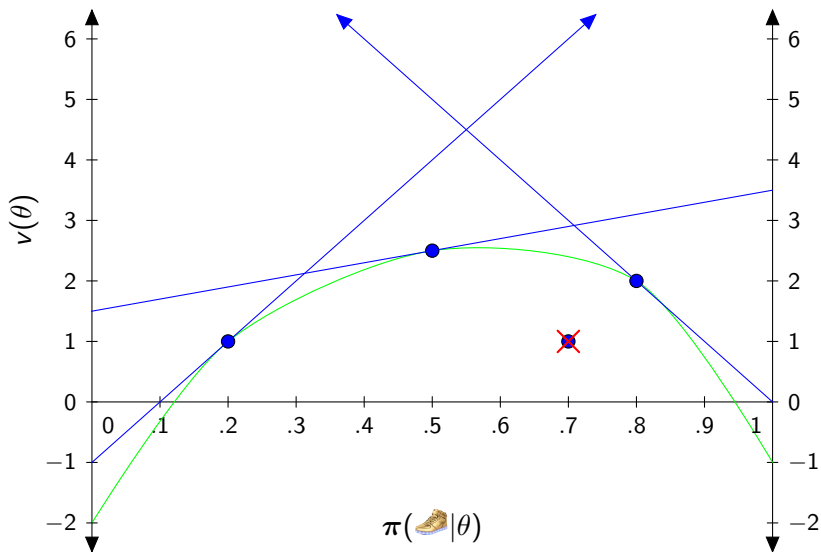
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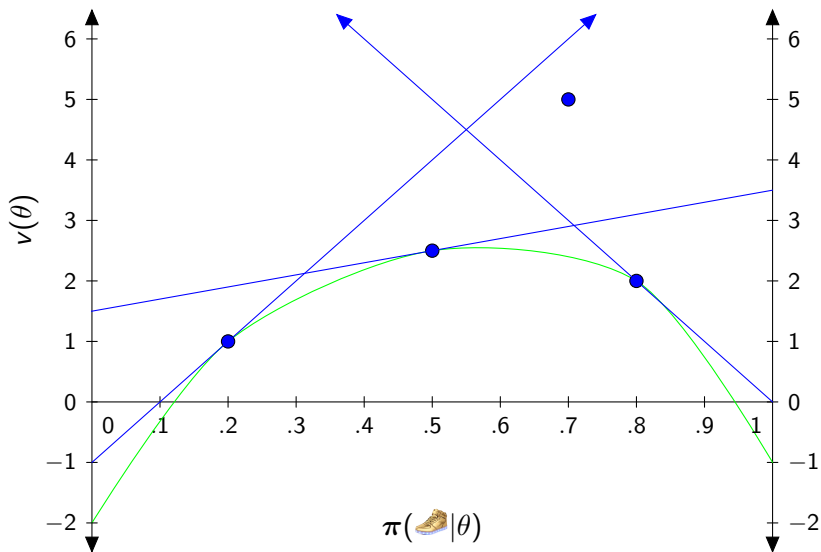
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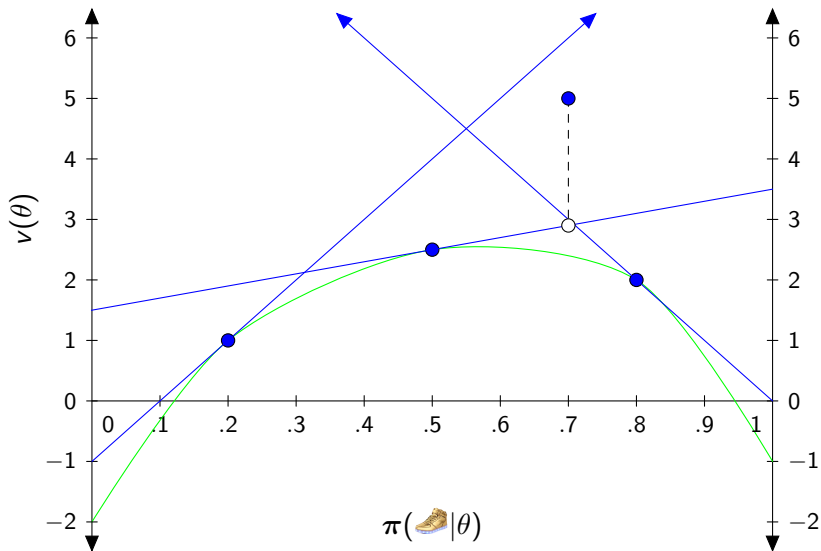


# Full Revenue Extraction with a Bayesian Mechanism

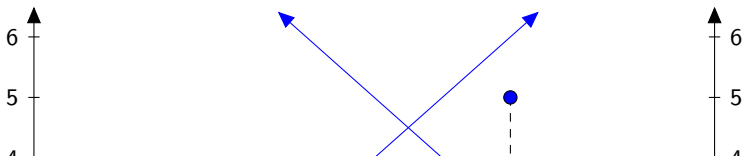




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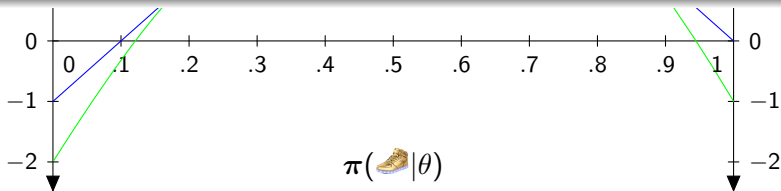


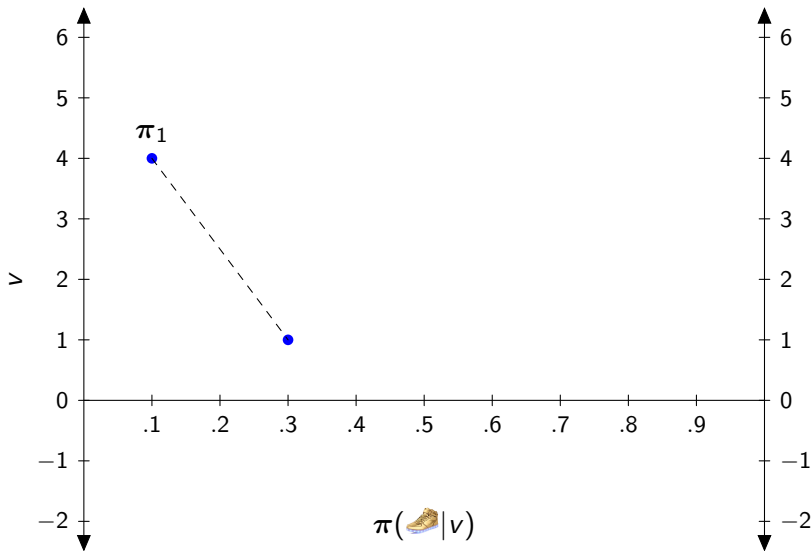
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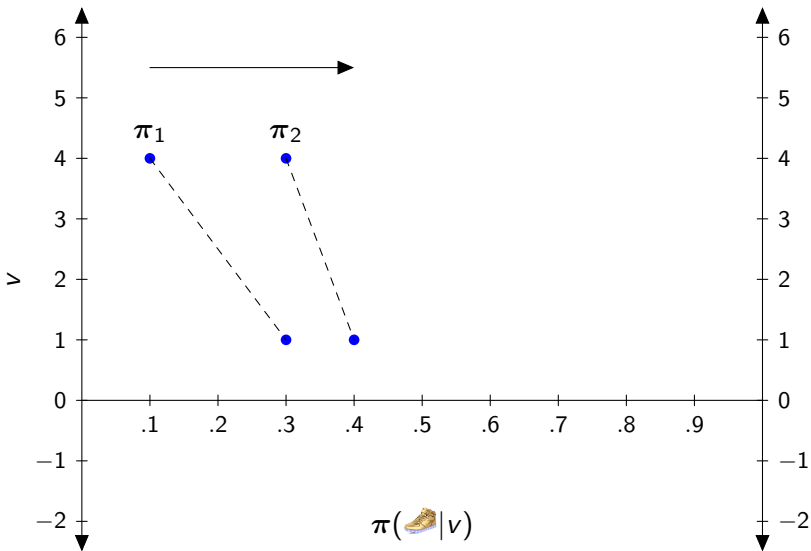


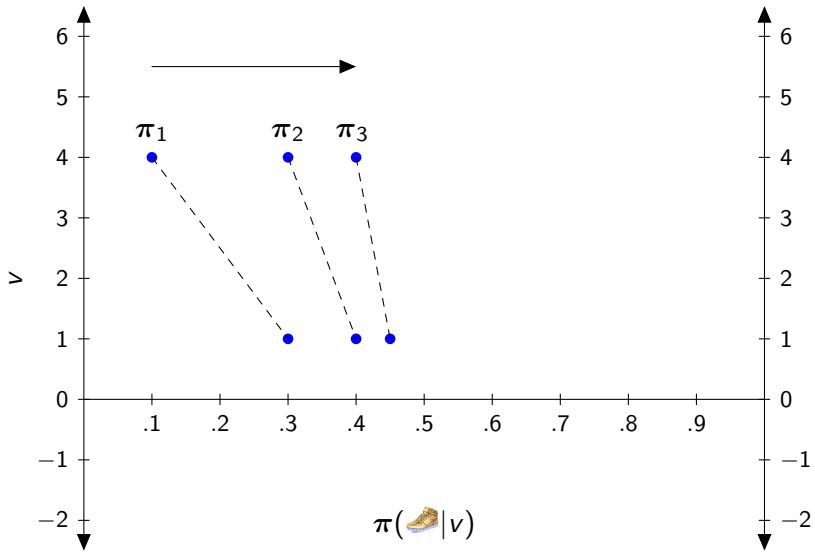
Theorem: Full Surplus Extraction with a Bayesian Mechanism  
(Albert, Conitzer, Lopomo 2016)

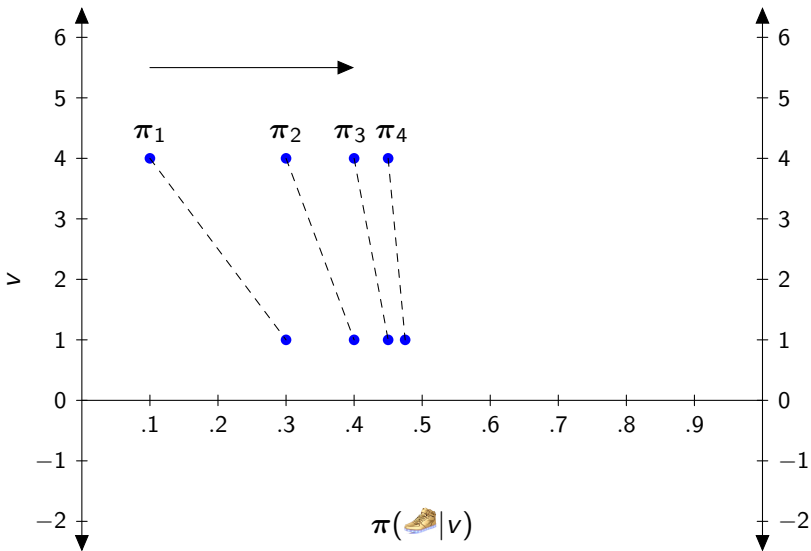
For a given  $(\pi, \Theta, \Omega)$ , full surplus extraction is possible for a Bayesian mechanism if and only if there exists a concave function  $G : \mathbb{R}^{|\Omega|} \rightarrow \mathbb{R}$  such that  $G(\pi(\bullet|\theta)) = v(\theta)$ . [Return](#)

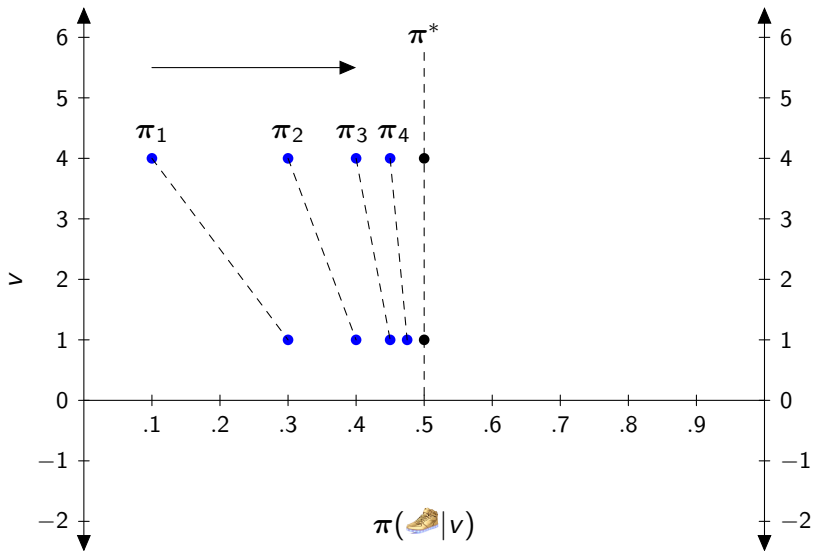


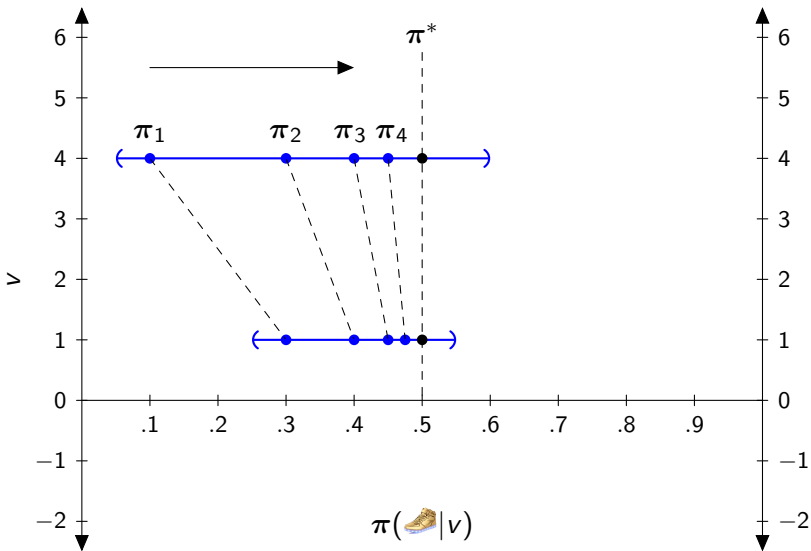




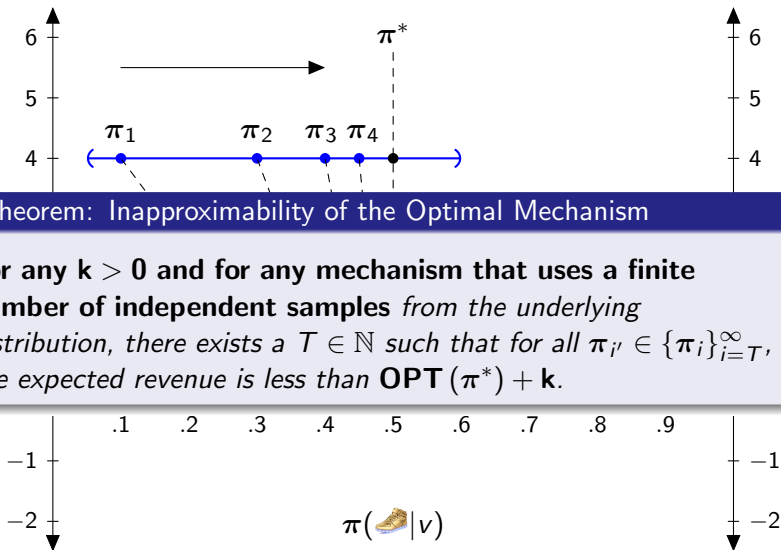






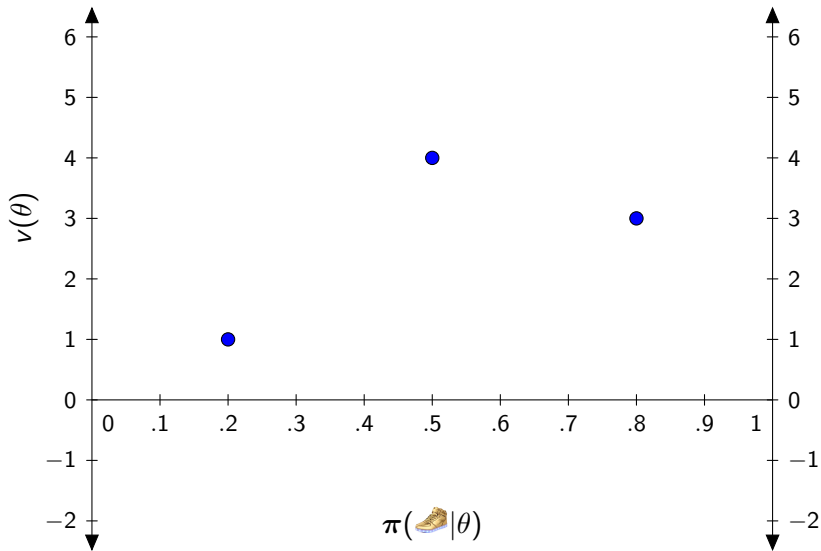


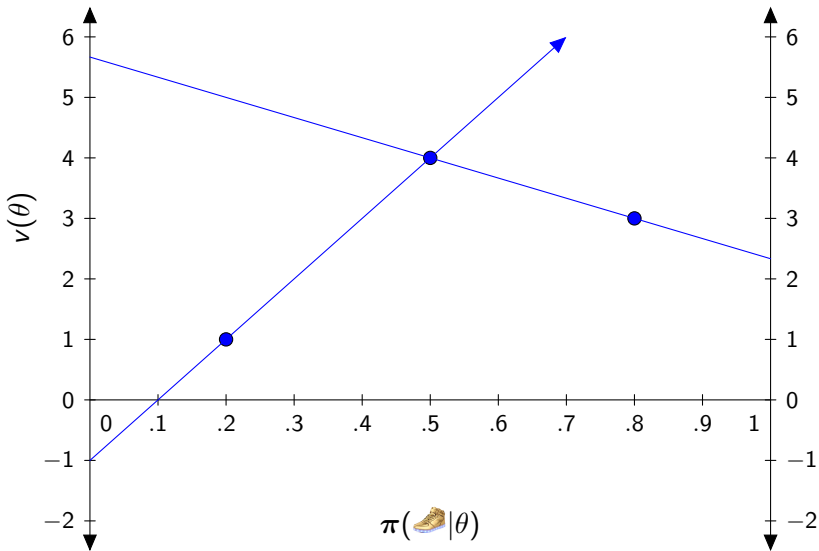


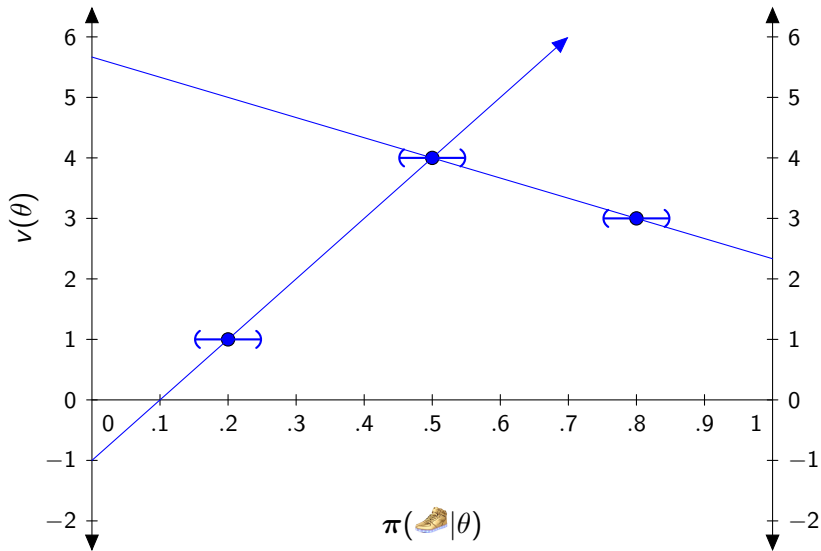


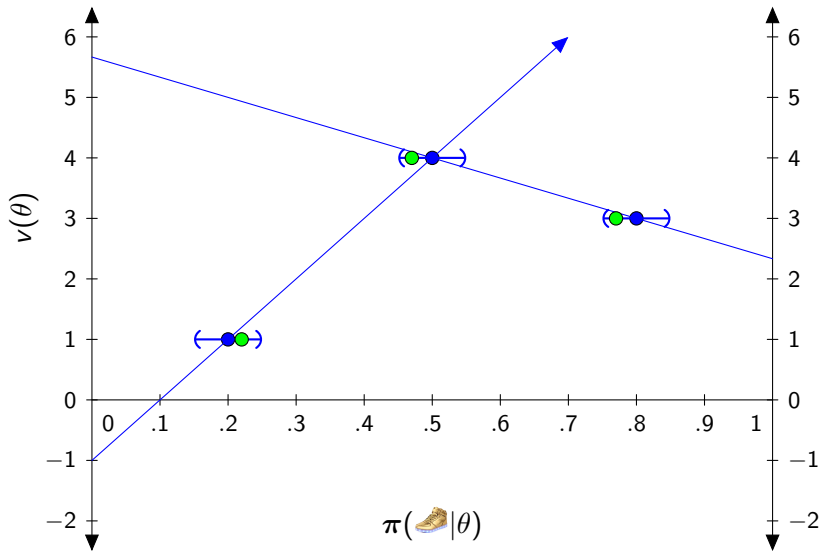
### Theorem: Inapproximability of the Optimal Mechanism

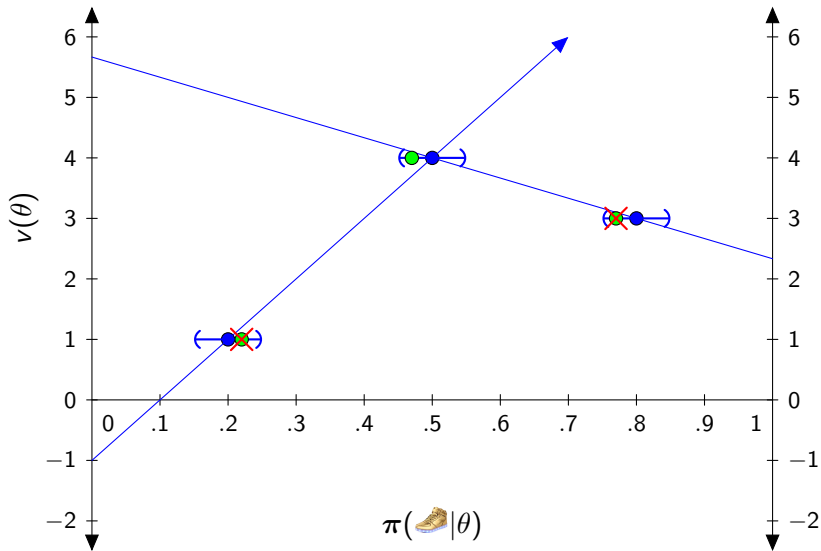
For any  $k > 0$  and for any mechanism that uses a finite number of independent samples from the underlying distribution, there exists a  $T \in \mathbb{N}$  such that for all  $\pi_{i'} \in \{\pi_i\}_{i=1}^{\infty}, T$ , the expected revenue is less than  $\text{OPT}(\pi^*) + k$ .

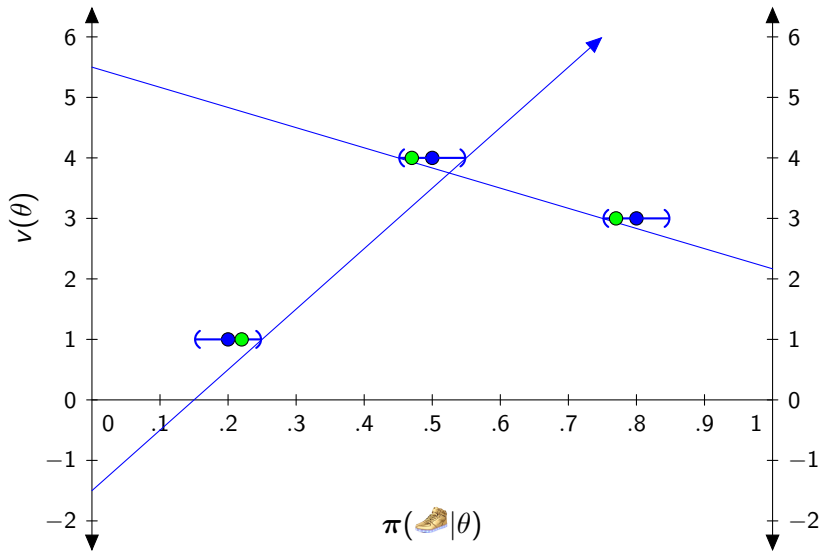


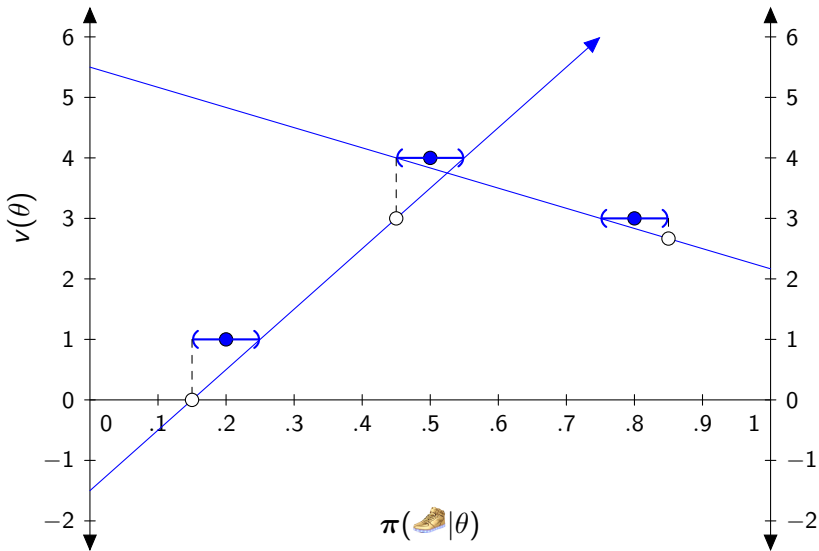




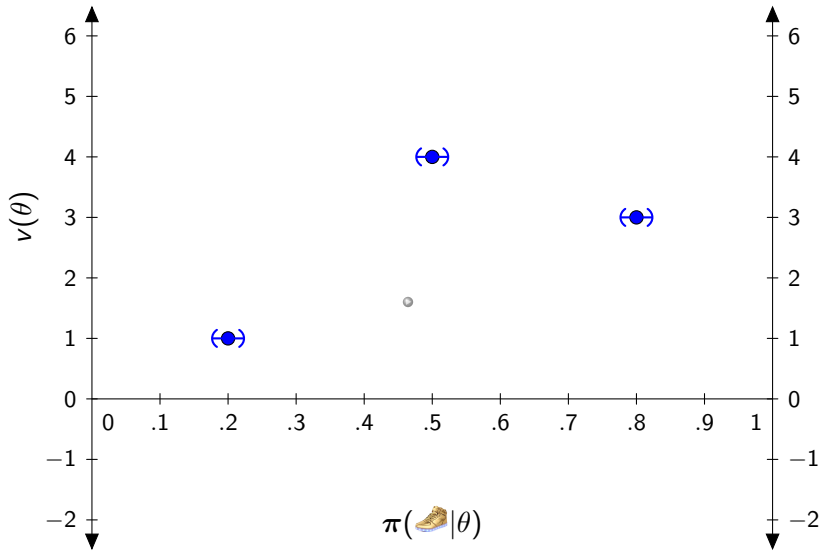


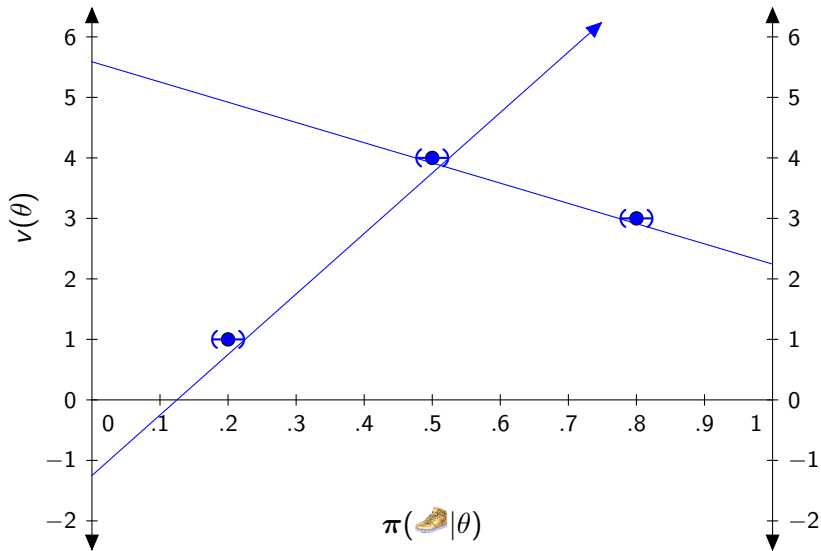


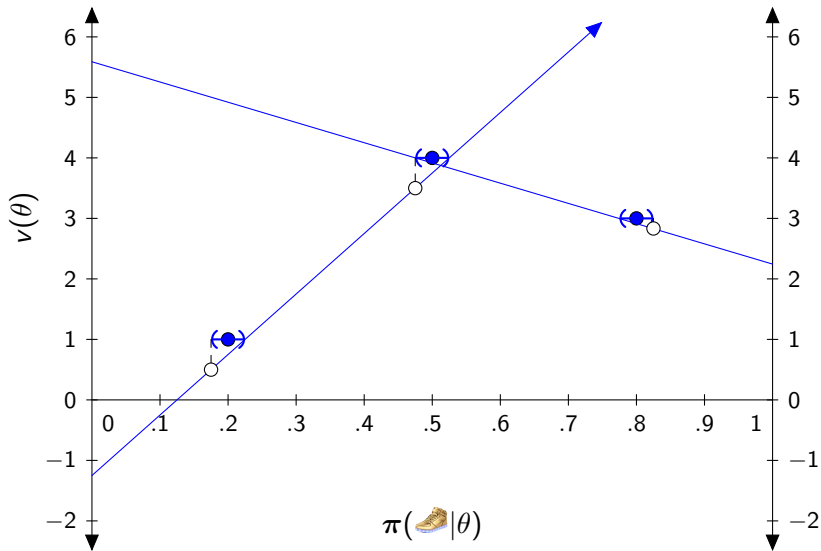


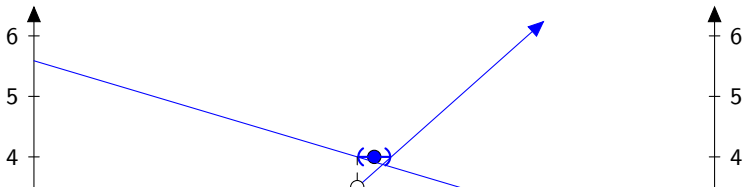






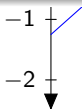




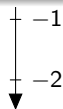


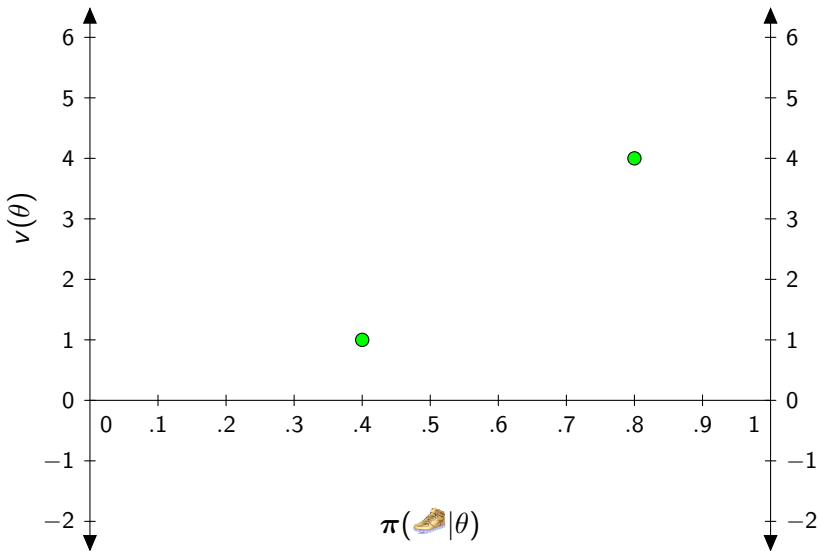
### Theorem: Sufficient Correlation Implies Near Optimal Revenue

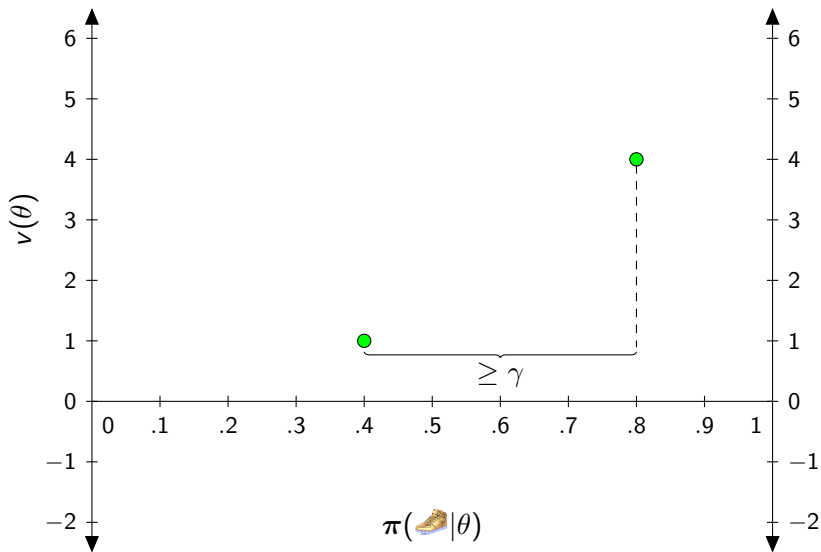
For any distribution  $\pi^*$  that satisfies the FSE condition and given any positive constant  $k > 0$ , **there exists**  $\epsilon > 0$  and a mechanism such that for all distributions,  $\pi'$ , for which for all  $\theta \in \Theta$ ,  $\|\pi^*(\cdot|\theta) - \pi'(\cdot|\theta)\| < \epsilon$ , the revenue generated by the mechanism is greater than or equal to  $\text{OPT}(\pi^*) - k$ . [▶ Return](#)

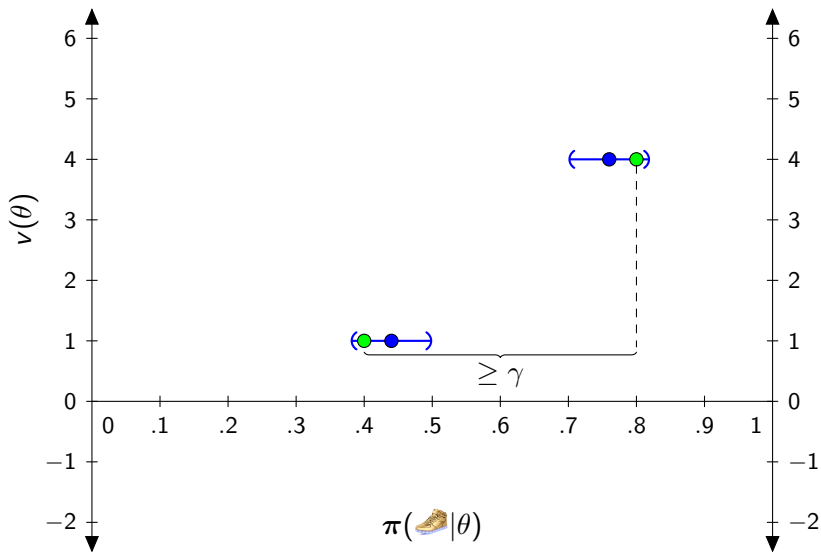


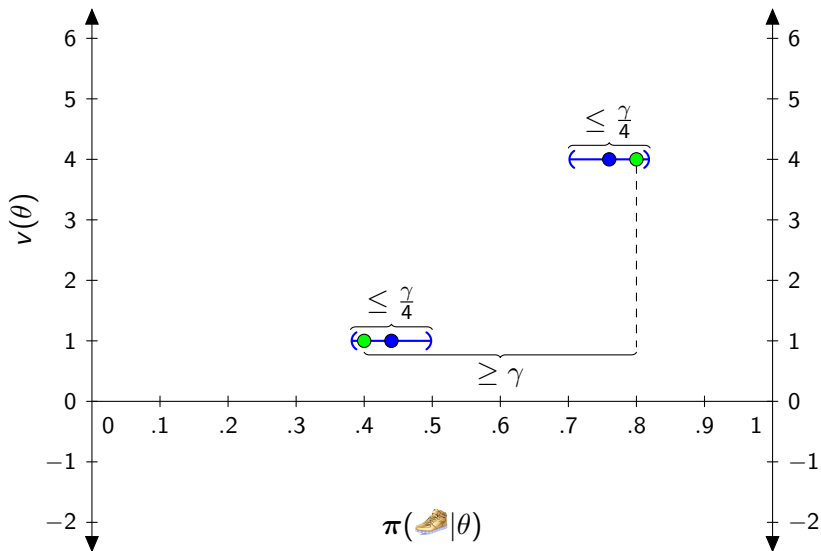
$\pi(\text{shoe}|\theta)$



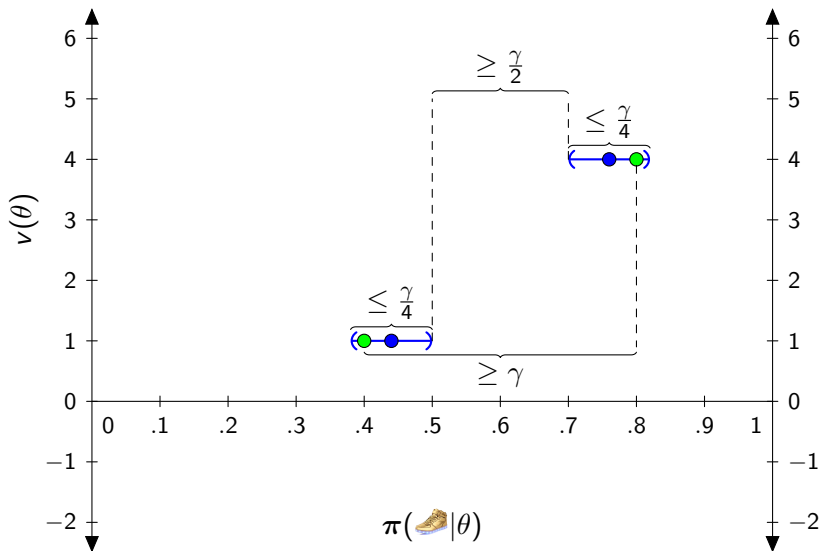


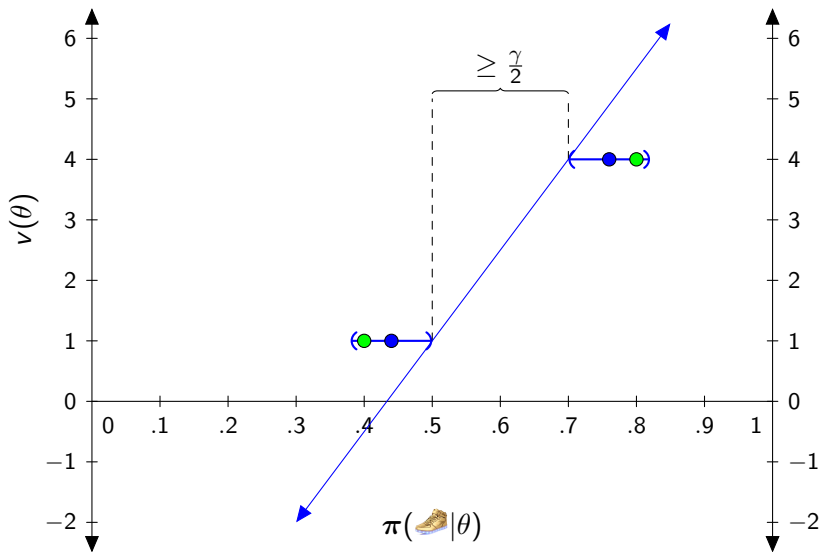


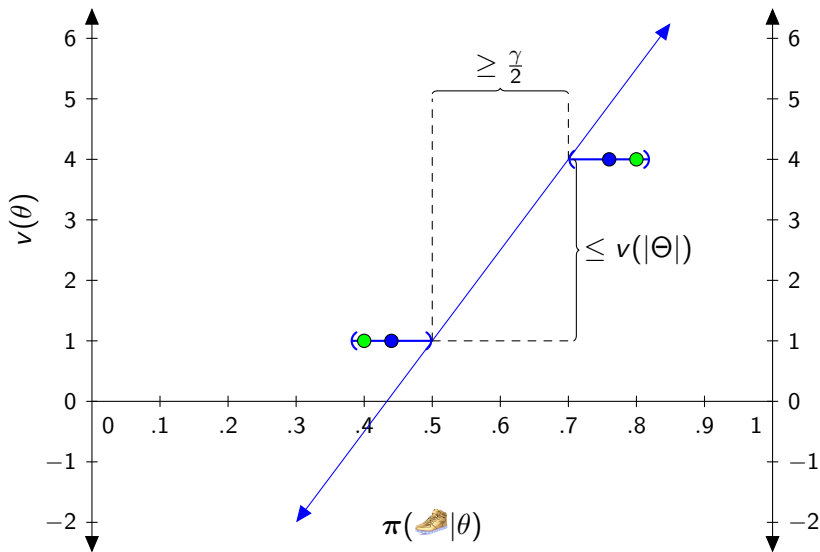


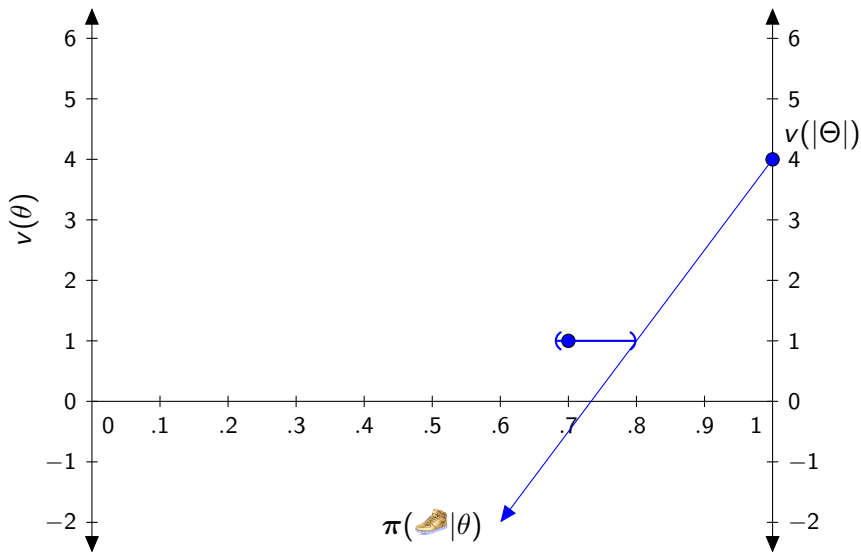


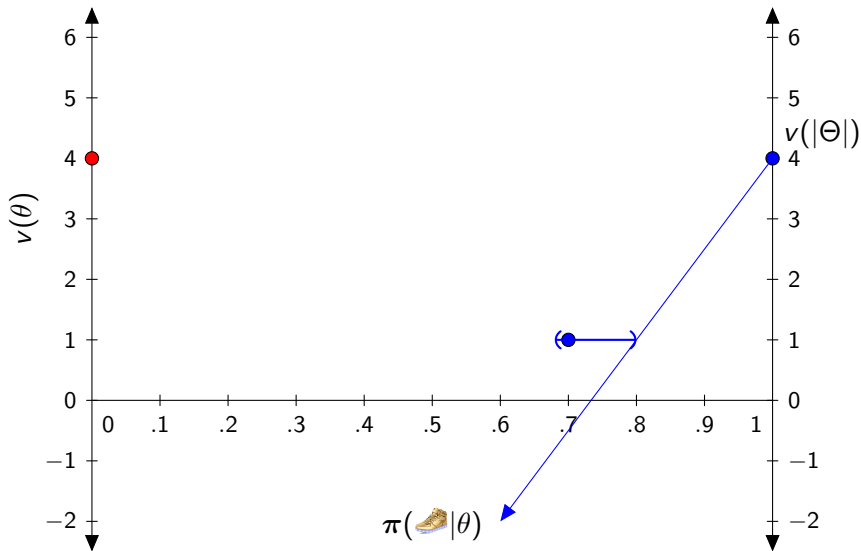


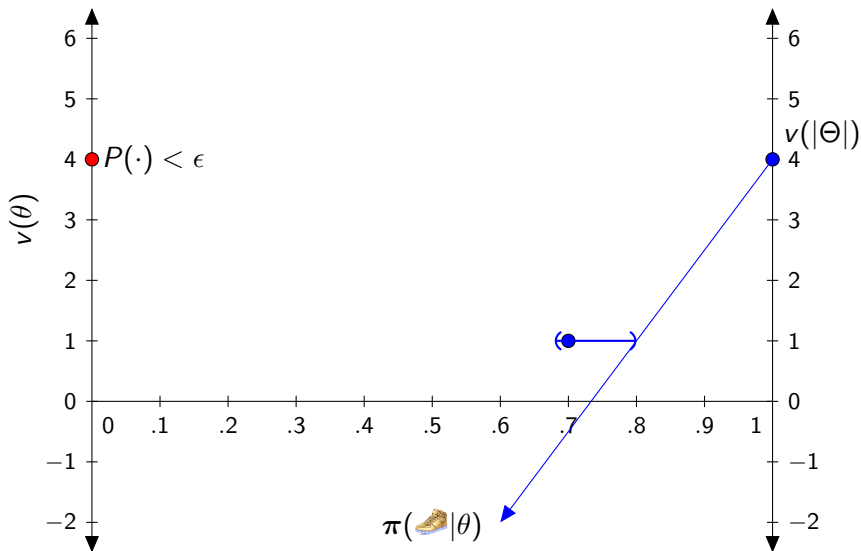


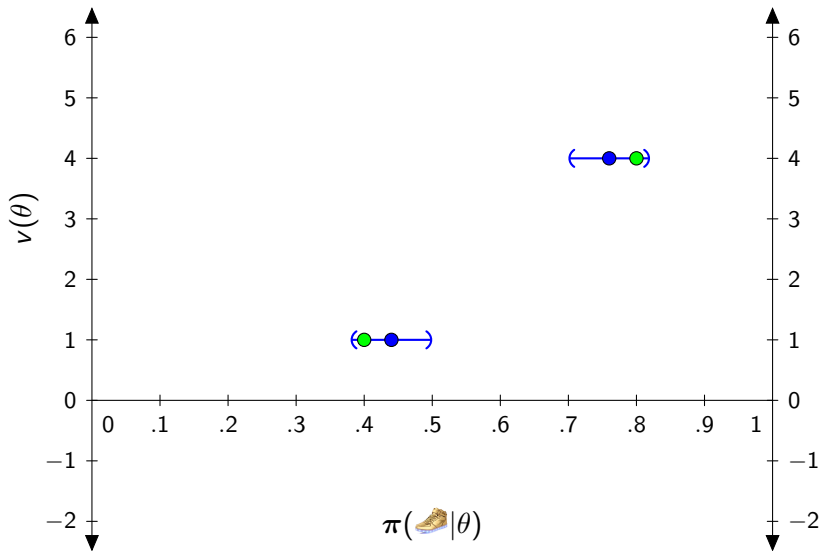


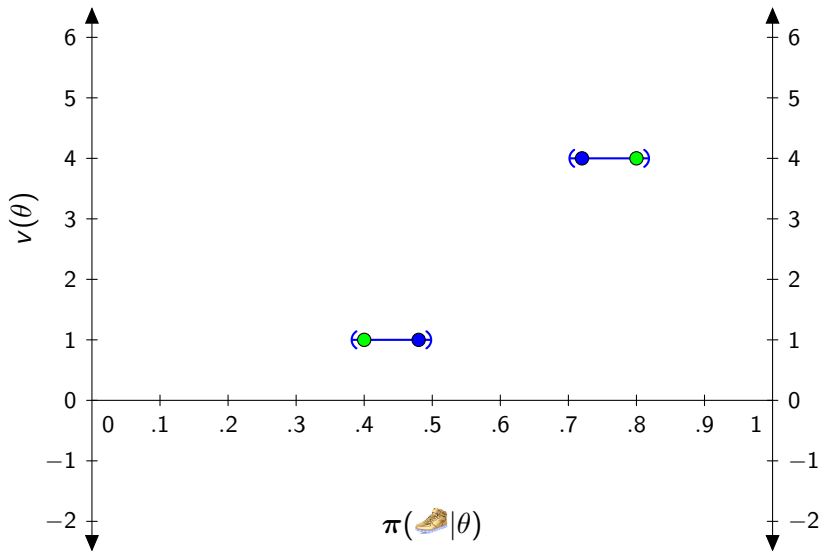




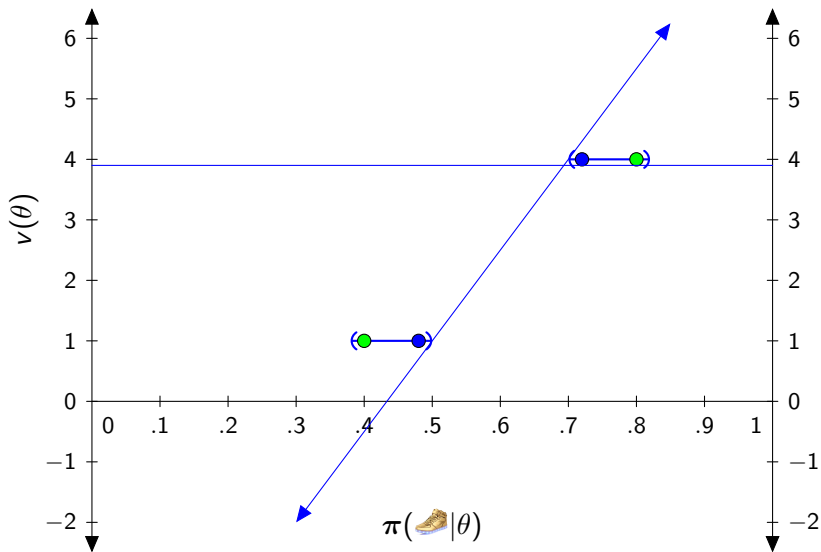


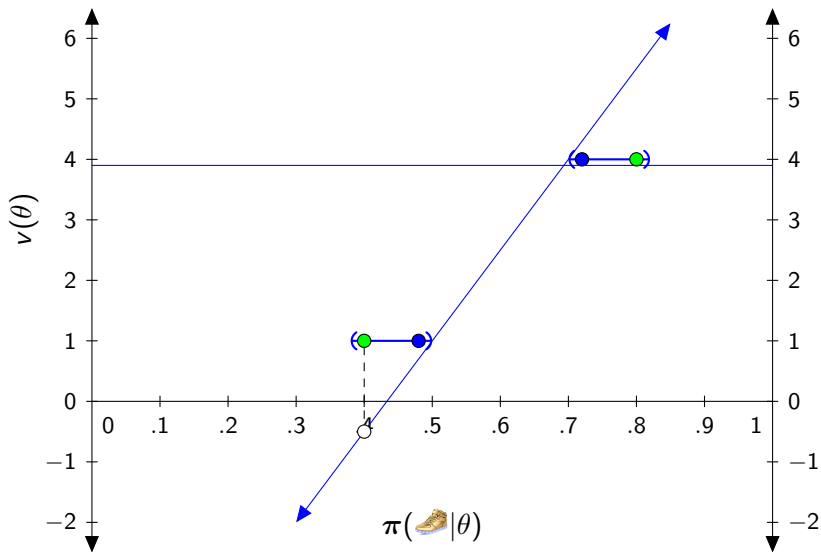


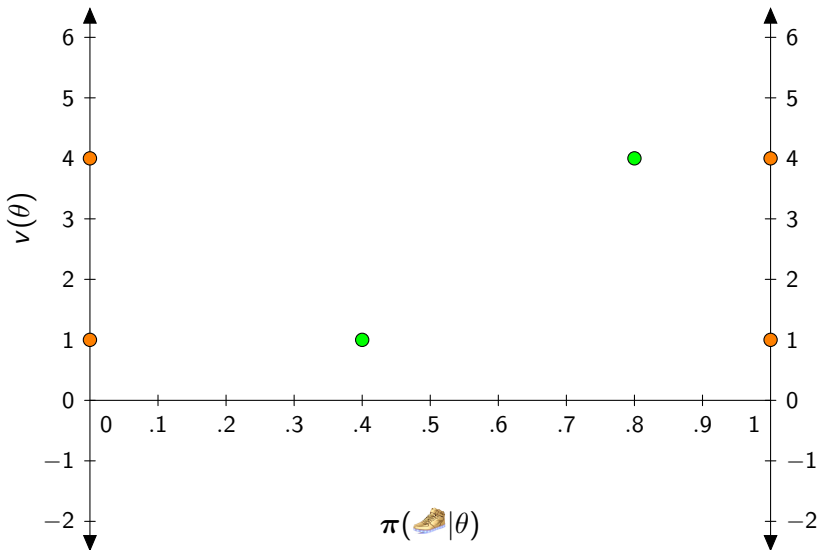


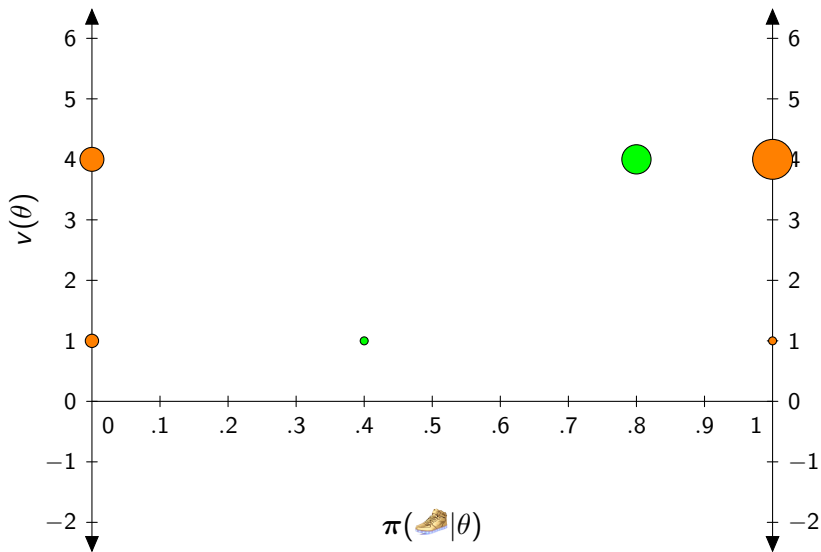


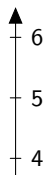
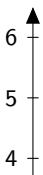






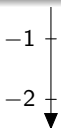






## Sample Complexity of $\epsilon$ -Robust Mechanism Design

The **sample complexity** for constructing an  $\epsilon$ -robust mechanism with an expected revenue that is an **additive  $k$ -approximation** to the expected revenue of the optimal mechanism over a true distribution that is  $\gamma$ -separated and satisfies the FSE condition is  $O(\text{poly}(\frac{1}{k}, \frac{1}{\gamma}, |\Theta|, |\Omega|, \mathbf{v}(|\Theta|)))$ . [▶ Return](#)



$\pi(\text{👟}|\theta)$

