15-251

Some AWESOME

~~Great Theoretical~~ Ideas

~~in Computer Science~~

about ~~Generating Functions~~

Probability

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Some AWESOME

~~Great Theoretical~~ Ideas

~~in Computer Science~~

about ~~Generating Functions~~

~~Probability~~

Infinity

15-251

What Little Susie Should've Said to Little Johnny

Ideas from the course

- Induction
- Numbers
- Finite Counting and Probability
- A hint of the infinite
  - Infinite row of dominoes
  - Infinite sums (Generating functions!!)
  - Infinite choice trees, and infinite probability
  - Infinite tapes

The Ideal Computer:
no bound on amount of memory
no bound on amount of time

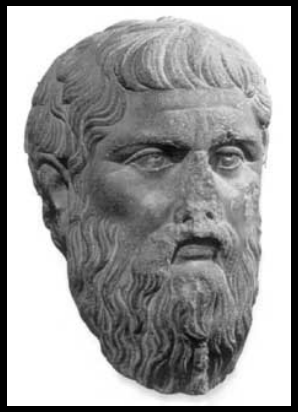
Ideal Computer is defined as a computer with infinite RAM.

You can run a Java program and never have any overflow, or out of memory errors.

Infinite RAM Model

Platonic Version:

One memory location for each natural number 0, 1, 2, …



Aristotelian Version:

Whenever you run out of memory, the computer contacts the factory. A maintenance person is flown by helicopter and attaches 1000 Gig of RAM and all programs resume their computations, as if they had never been interrupted.

Here's a program

```
System.out.print("0.");
for(int i=0;true;i++)
{
   System.out.print( getDigit(i) );
}
```

Here's a program

```
int getDigit(int i)
{
   return 3;
}
```

Here's a program

```
int getDigit(int i)
{
   return i%10;
}
```

Here's a program

Can we do:

Pi?

e?

Any real?

$$\frac{1}{\pi} = 12\sum_{k=0}^{\infty} \frac{(-1)^k (6k)!(13591409 + 545140134k)}{(3k)!(k!)^3 640320^{3k+3/2}}$$

Chudnovsky brothers

←

An Ideal Computer

It can be programmed to print out:

2: 2.00000000000000000000...

1/3: 0.33333333333333333333...

ϕ: 1.61803398874989484820 45...

e: 2.71828182845904523536 02...

π: 3.14159265358979323846264...

Printing Out An Infinite Sequence..

A program P prints out the infinite sequence

$s_0, s_1, s_2, \ldots, s_k, \ldots$

if when P is executed on an ideal computer, it outputs a sequence of symbols such that

-The k^{th} symbol that it outputs is s_k

-For every k, P eventually outputs the k^{th} symbol. I.e., the delay between symbol k and symbol k+1 is not infinite.

Computable Real Numbers

A real number R is computable if there is a program that prints out the decimal representation of R from left to right.

Thus, each digit of R will eventually be output.

Are all real numbers computable?

Describable Numbers

A real number R is describable if it can be denoted unambiguously by a finite piece of English text.

2: "Two."

π: "The area of a circle of radius one."

Are all real numbers describable?

Is every computable real number, also a describable real number?

And what about the other way?

Computable R: some program outputs R
Describable R: some sentence denotes R

Computable $\Rightarrow$ describable

Theorem:

Every computable real is also describable

Computable ⇒ describable

Theorem:

Every computable real is also describable

Proof:

Let R be a computable real that is output by a program P. The following is an unambiguous description of R:

"The real number output by the following program:" P

MORAL: A computer program can be viewed as a description of its output.

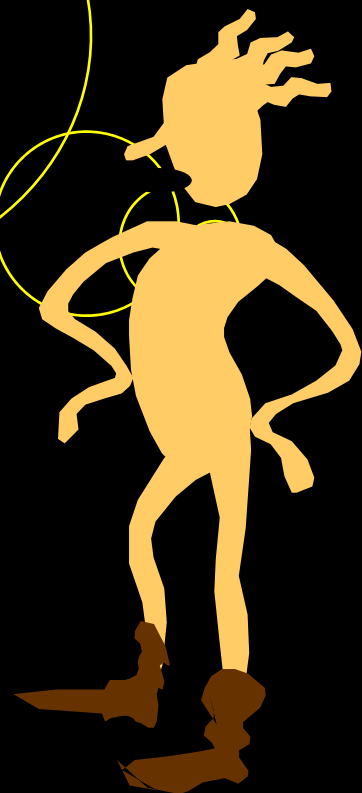
Syntax: The text of the program

Semantics: The real number output by P

Are all reals describable?
Are all reals computable?

We saw that
computable $\Rightarrow$
describable,

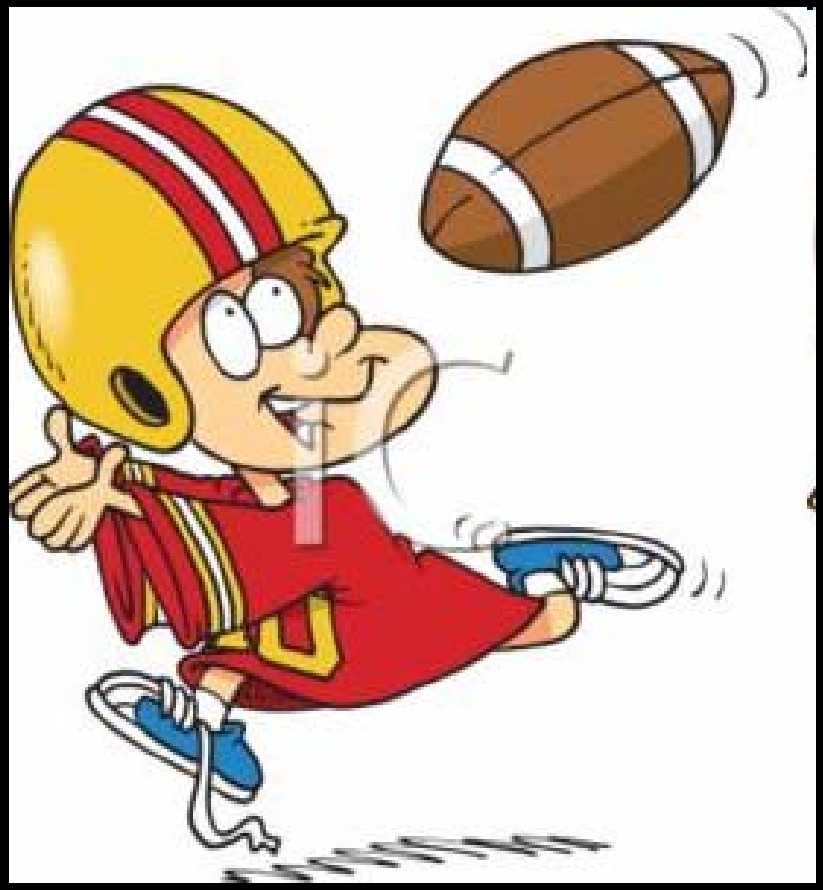
but do we also have
describable $\Rightarrow$
computable?



Questions we will answer in this (and next) lecture…

Little Susie

Little Johnny



Little Susie and Little Johnny

Little Susie: I hate you
Little Johnny: I hate you more
Little Susie: I hate you times a zillion
Little Johnny: I hate you times infinity
Little Susie: I hate you times infinity

Plus one!

Susie's mistake

Infinity: N.

Infinity plus one : N U {cupcake}

Can we establish a bijection between N and N U {cupcake}?

Susie's mistake

Sure!

f : N -> N U {cupcake}

f(x) = cupcake if x=0

f(x) = x-1 if x>0

0 → cupcake

1 → 0

2 → 1

3 → 2

4 → 3

5 → 4

… …

Correspondence Principle

If two finite sets can be placed into 1-1 onto correspondence, then they have the same size.

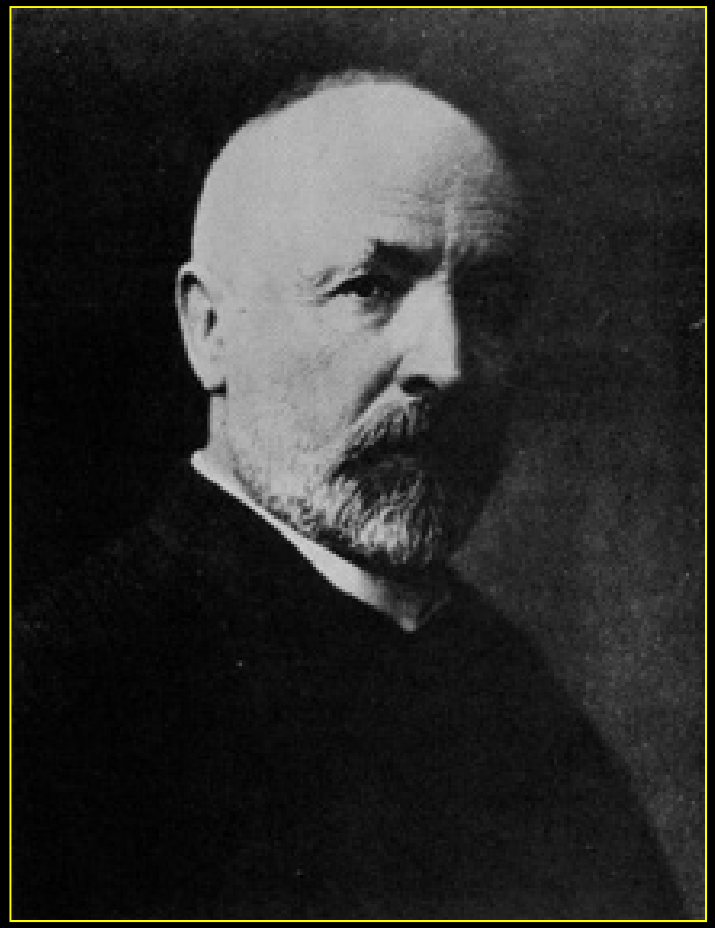
That is, if there exists a bijection between them.

Correspondence Definition

In fact, we can use the correspondence as the definition:

Two finite sets are defined to have the same size if and only if they can be placed into 1-1 onto correspondence.

Georg Cantor (1845-1918)



Cantor's Definition (1874)

Two sets are defined to have the same size if and only if they can be placed into 1-1 onto correspondence.

Cantor's Definition (1874)

Two sets are defined to have
the same cardinality if and only if
they can be placed into
1-1 onto correspondence.

Therefore, N and N U {cupcake} have the same cardinality.

Do N and E have the same cardinality?

N = { 0, 1, 2, 3, 4, 5, 6, 7, ... }

E = { 0, 2, 4, 6, 8, 10, 12, ... }

The even, natural numbers.

E and N do not have the same cardinality! E is a proper subset of N with not one element left over, but an INFINITE amount!

E and N do have the same cardinality!

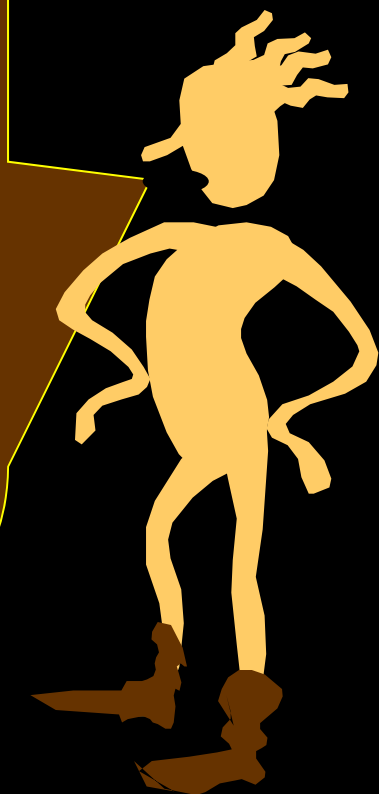
N = 0, 1, 2, 3, 4, 5, …
E = 0, 2, 4, 6, 8,10, …

f(x) = 2x is 1-1 onto.

Lesson:

Cantor's definition only requires that ***some*** 1-1 correspondence between the two sets is onto, not that **all** 1-1 correspondences are onto.

This distinction never arises when the sets are **finite**.



Cantor's Definition (1874)

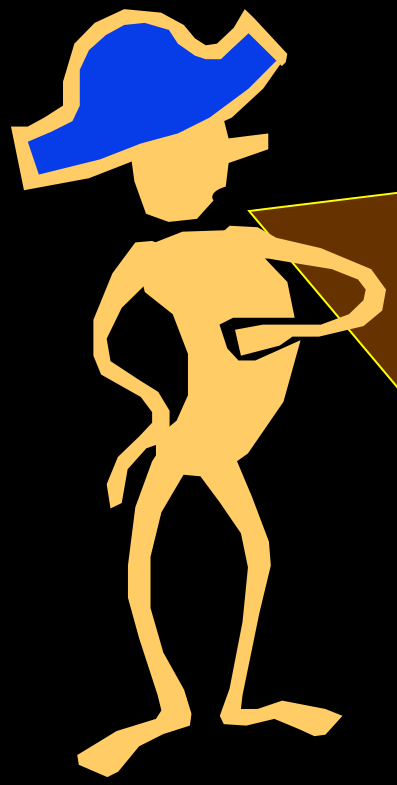
Two sets are defined to have
the same size if and only if they can be
placed into 1-1 onto correspondence.

You just have to get used to this slight subtlety in order to argue about infinite sets!

Do N and Z have the same cardinality?

$N = \{ 0, 1, 2, 3, 4, 5, 6, 7, \ldots \}$

$Z = \{ \ldots, -2, -1, 0, 1, 2, 3, \ldots \}$



No way! Z is infinite in two ways: from 0 to positive infinity and from 0 to negative infinity.

Therefore, there are far more integers than naturals.

Actually, no!

N and Z do have the same cardinality!

N = 0, 1, 2, 3, 4, 5, 6 …
Z = 0, 1, -1, 2, -2, 3, -3, ….

$f(x) = \lceil x/2 \rceil$ if x is odd
$-x/2$ if x is even

Transitivity Lemma

Do E and Z have the same cardinality?

Transitivity Lemma

Lemma: If

$f: A \to B$ is 1-1 onto, and

$g: B \to C$ is 1-1 onto.

Then $h(x) = g(f(x))$ defines a function

$h: A \to C$ that is 1-1 onto

Hence, N, E, and Z all have the same cardinality.

Do N and Q have the same cardinality?

$N = \{ 0, 1, 2, 3, 4, 5, 6, 7, ... \}$

Q = The Rational Numbers

No way!

The rationals are dense: between any two there is a third. You can't list them one by one without leaving out an infinite number of them.

Don’t jump to conclusions!

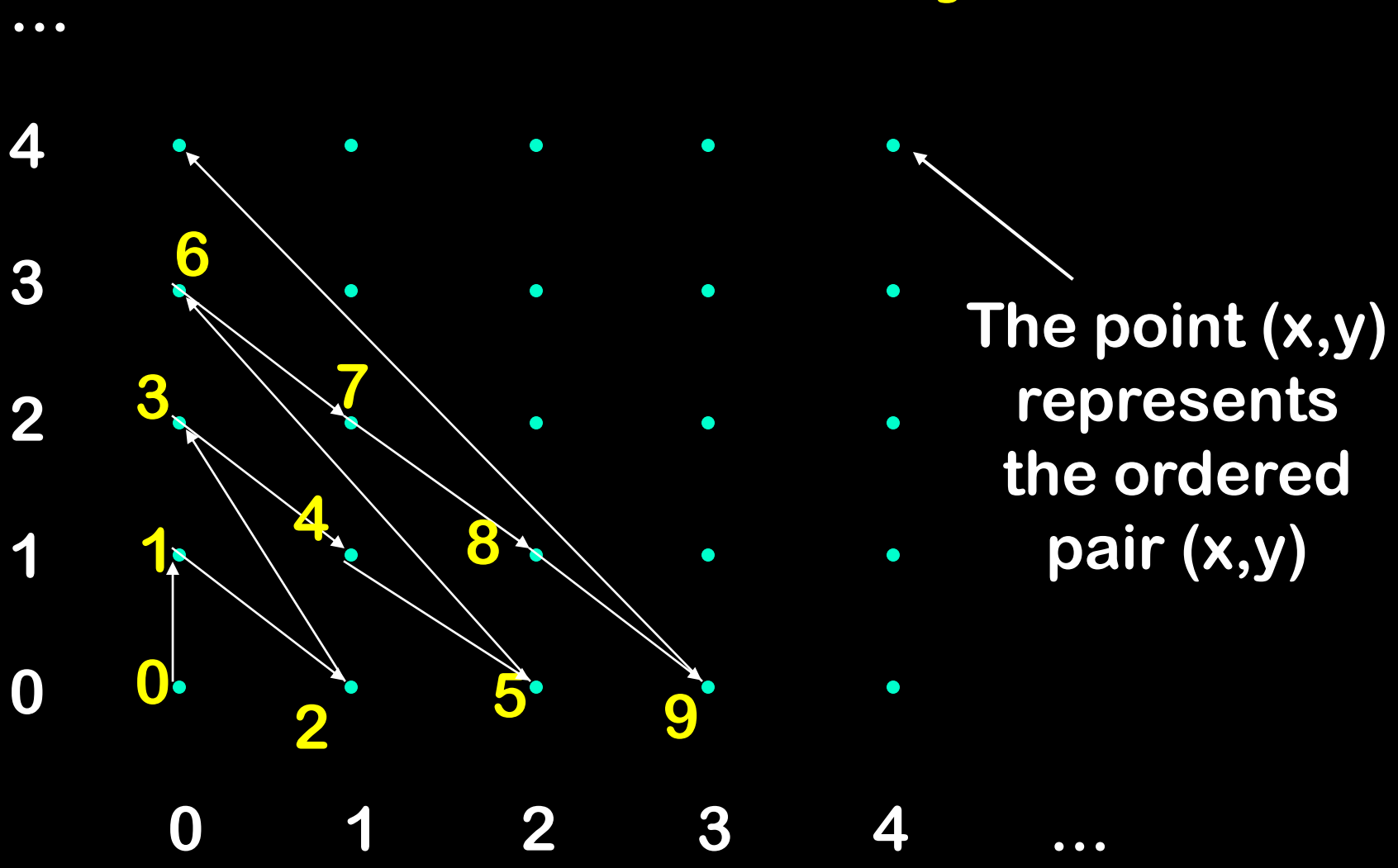
There is a clever way to list the rationals, one at a time, without missing a single one!

First, let's warm up with another interesting example: N can be paired with NxN

Theorem: N and NxN have the same cardinality

Theorem: N and NxN have the same cardinality

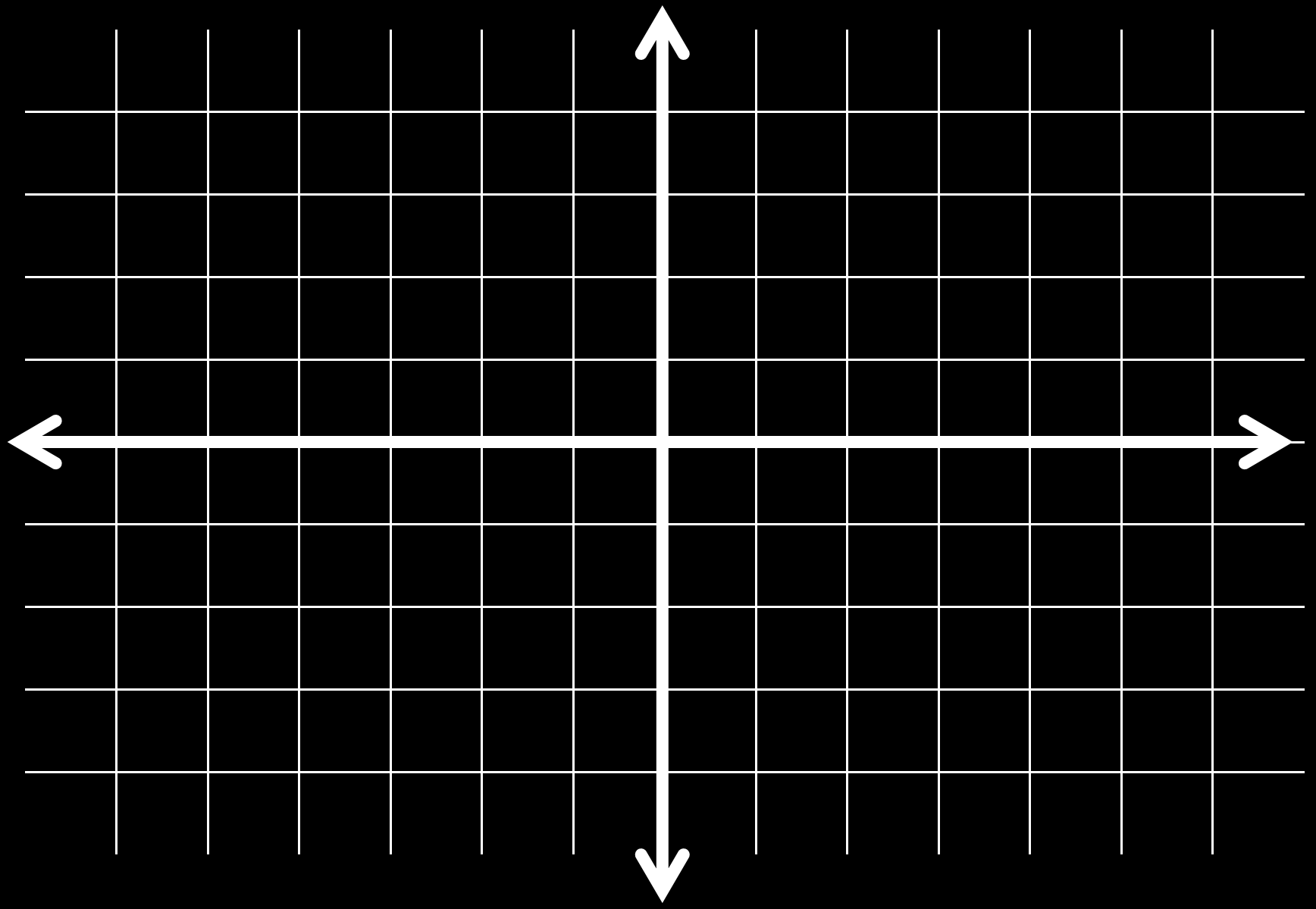
Theorem: N and NxN have the same cardinality



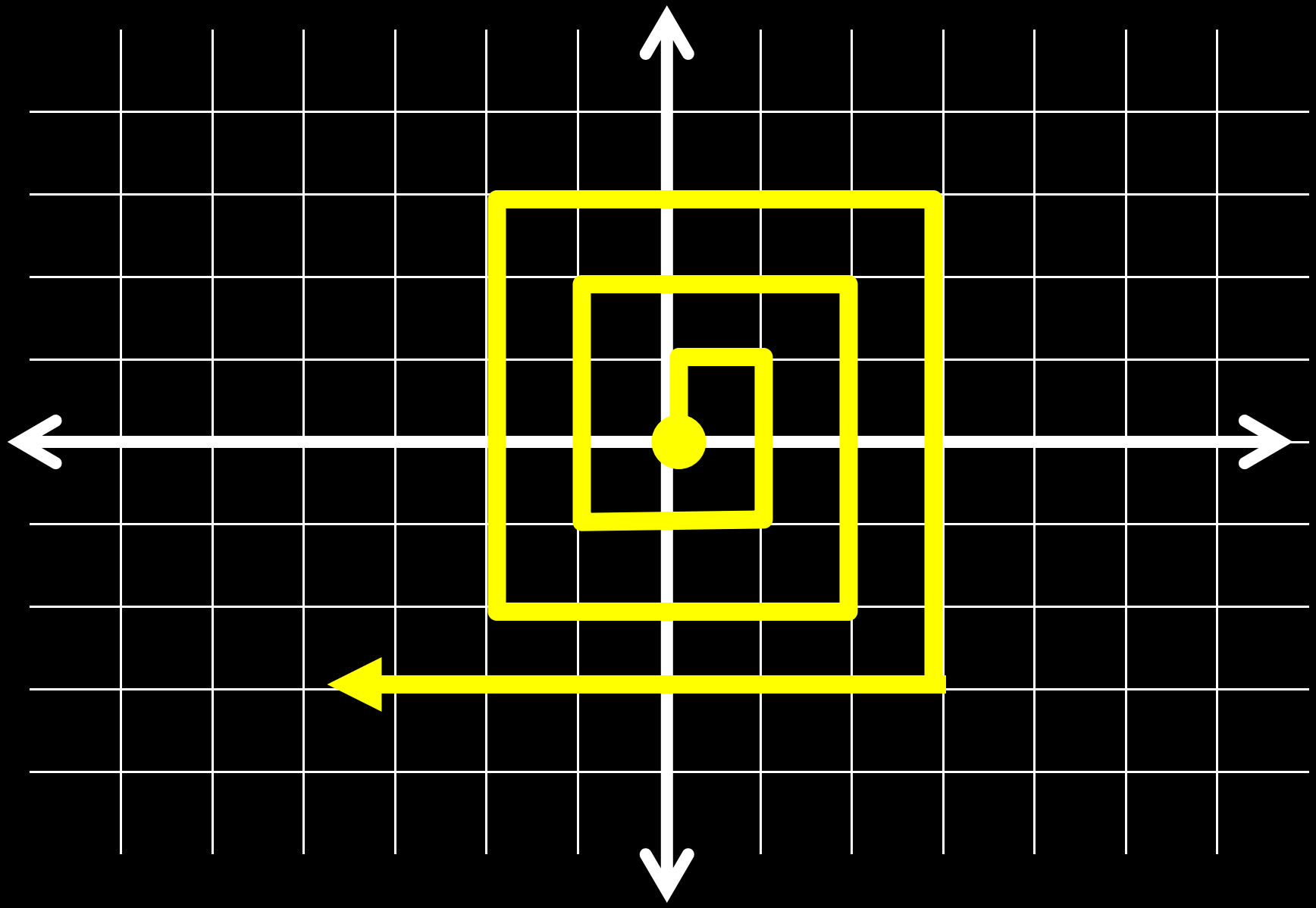
Defining 1-1 onto f: N -> NxN

```
int sum;
for (sum = 0;true;sum++) {
  //generate all pairs with this sum
  for (x = 0;x<=sum;x++) {
    y = sum-x;
        System.out.println(x+“ ”+ y);
  }
}
```

Onto the Rationals!



The point at x,y represents x/y



The point at x,y represents x/y

Hold it!

You've included both 1,1 and 2,2 –
They correspond to the same rational.

Also, 0/0, 1/0, 2/0, … are not rational!

Hold it!

0	0/0	7	-1/0	14	2/0
1	0/1	8	-1/1	15	2/-1
2	1/1	9	-1/2	16	2/-2
3	1/0	10	0/2	17	1/-2
4	1/-1	11	1/2	18	0/-2
5	0/-1	12	2/2	19	-1/-2
6	-1/-1	13	2/1	20	-2/-2

Hold it!

0	0/0	7	-1/0	14	2/0
1	0/1	8	-1/1	15	2/-1
2	1/1	9	-1/2	16	2/-2
3	1/0	10	0/2	17	1/-2
4	1/-1	11	1/2	18	0/-2
5	0/-1	12	2/2	19	-1/-2
6	-1/-1	13	2/1	20	-2/-2

Hold it!

0		7		14	
1	0/1	8		15	2/-1
2	1/1	9	-1/2	16	
3		10		17	
4	1/-1	11	1/2	18	
5		12		19	
6		13	2/1	20	

Hold it!

0	0/1
1	1/1
2	1/-1
3	-1/2
4	1/2
5	2/1
6	2/-1

Hold it!

It's okay. We can just skip those. So instead of assigning 0 to 0/0 we will assign it to 0/1, and so on.

This way, we'll use all the naturals and we'll hit all the rationals without duplication.

Cantor-Bernstein-Schroeder

If there exists an injection from A to B and an injection from B to A, then there exists a bijection between B and A.

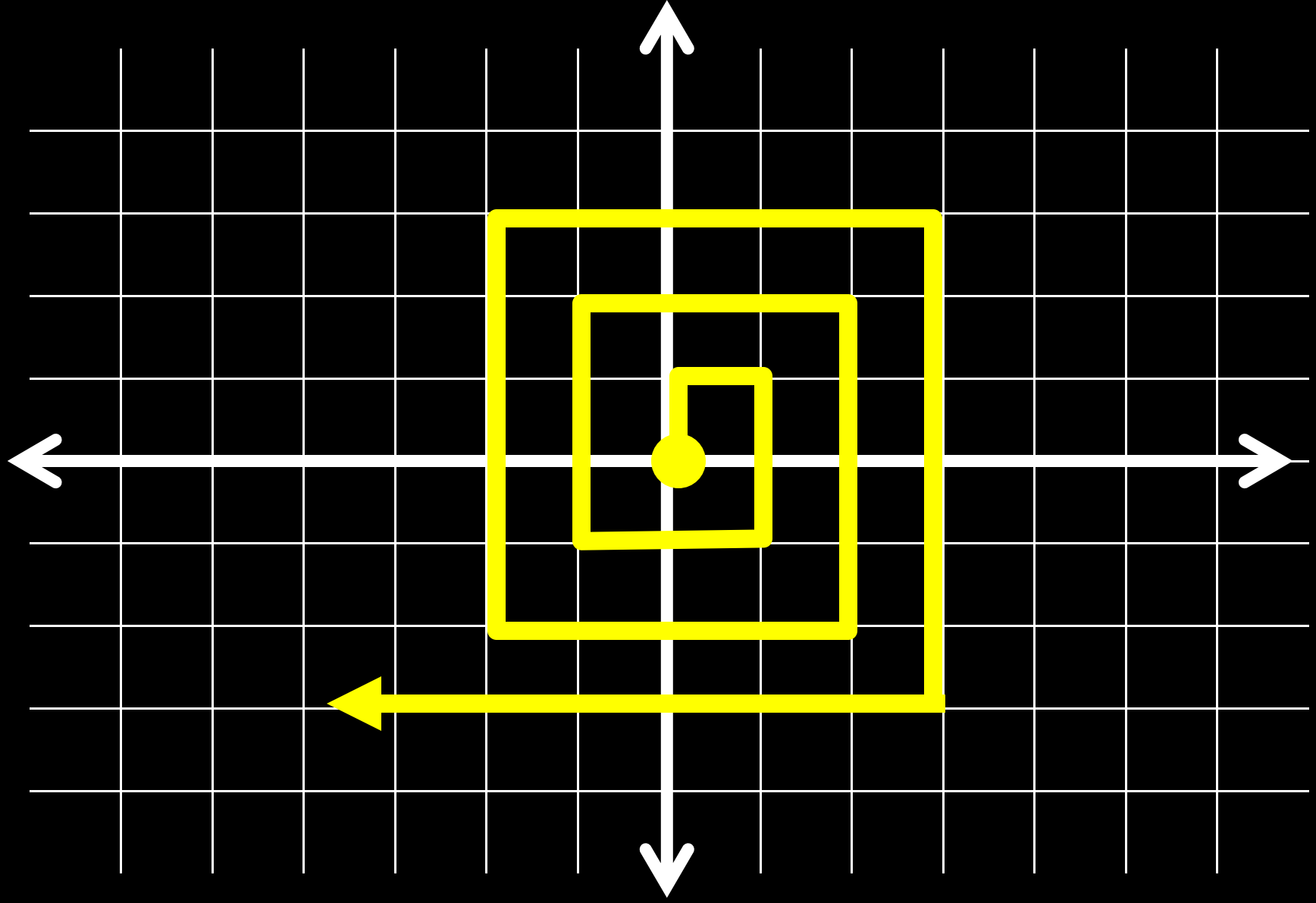
Easy to prove for finite sets, trickier for infinite sets.

Cantor-Bernstein-Schroeder

Injection from N to Q

$f(x) = x$

Injection from Q to N



The point at x,y represents x/y

Injection from Q to N

0		7		14	
1	0/1	8		15	2/-1
2	1/1	9	-1/2	16	
3		10		17	
4	1/-1	11	1/2	18	
5		12		19	
6		13	2/1	20	

Injection from Q to N

0	7	14
1←0/1	8	15←2/-1
2←1/1	9 ← -1/2	16
3	10	17
4←1/-1	11←1/2	18
5	12	19
6	13←2/1	20

Cantor-Bernstein-Schroeder

Injection from Q to N

Just eliminate the invalid matchings. We match all the valid rationals and never duplicate a natural.

While this misses some naturals, it's still an injection.

Countable Sets

We call a set **countable** if it can be placed into 1-1 onto correspondence with the natural numbers N.

Hence

N, E, Q, and Z are all countable

Do N and R have the same cardinality?

$N = \{ 0, 1, 2, 3, 4, 5, 6, 7, \ldots \}$

R = The Real Numbers

No way!

You will run out of natural numbers long before you match up every real.

Now hang on a minute!

You can't be sure that there isn't some clever correspondence that you haven't thought of yet.

I am sure!
Cantor proved it.

To do this, he invented a very important technique called
"Diagonalization"

Theorem: The set of reals between 0 and 1 is not countable.

Proof: (by contradiction)

Suppose $R_{[0,1]}$ is countable.

Let f be a 1-1 onto function from N to $R_{[0,1]}$.

Make a list L as follows:

0: decimal expansion of f(0)
1: decimal expansion of f(1)
…
k: decimal expansion of f(k)
…

Theorem: The set of reals between 0 and 1 is not countable.

Proof: (by contradiction)

Suppose $R_{[0,1]}$ is countable.

Let f be a 1-1 onto function from N to $R_{[0,1]}$.

Make a list L as follows:

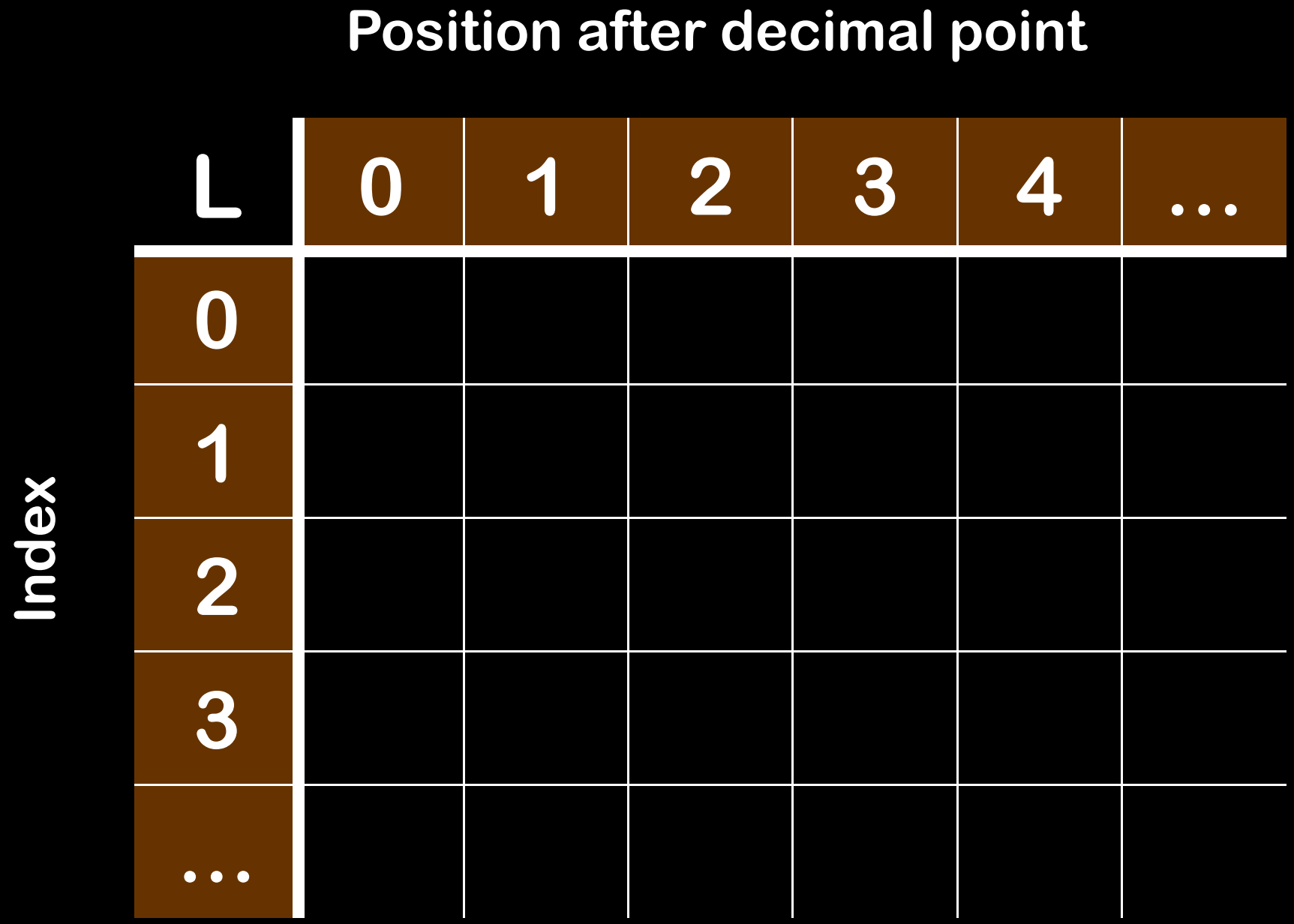
0: 0.33333333333333333…

1: 0.31415926565783959 3…

…

k: 0.235094385543905834…

…


Position after decimal point
L
0
1
2
3
4
...
Index
0
1
2
3
...

Position after decimal point

Index \ L	0	1	2	3	4	…
0	3	3	3	3	3	3
1	3	1	4	1	5	9
2	1	2	4	8	1	2
3	4	1	2	2	6	8
…						

digits along
the diagonal

L	0	1	2	3	4	…
0	d_0					
1		d_1				
2			d_2			
3				d_3		
…					…	

L	0	1	2	3	4
0	d_0				
1		d_1			
2			d_2		
3				d_3	
…					…

Define the following real number

$\text{Confuse}_L = . \; C_0 \quad C_1 \quad C_2 \quad C_3 \quad C_4 \quad C_5 \; \ldots$

L	0	1	2	3	4
0	d_0				
1		d_1			
2			d_2		
3				d_3	
…					…

Define the following real number

$Confuse_L = . C_0 \quad C_1 \quad C_2 \quad C_3 \quad C_4 \quad C_5 \ \ldots$

$$C_k = \begin{cases} 5, \text{ if } d_k = 6 \\ 6, \text{ otherwise} \end{cases}$$

L	0	1	2	3	4
0	$C_0 \neq d_0$	C_1	C_2	C_3	C_4
1		d_1			
2			d_2		
3				d_3	
…					…

…

$$C_k = \begin{cases} 5, \text{ if } d_k = 6 \\ 6, \text{ otherwise} \end{cases}$$

L	0	1	2	3	4
0	d_0				
1	C_0	$C_1 \neq d_1$	C_2	C_3	C_4
2			d_2		
3				d_3	
…					…

$$C_k = \begin{cases} 5, \text{ if } d_k = 6 \\ 6, \text{ otherwise} \end{cases}$$

…

L	0	1	2	3	4
0	d_0				
1		d_1			
2	C_0	C_1	$C_2 \neq d_2$	C_3	C_4
3				d_3	
…					…

$$C_k = \begin{cases} 5, \text{ if } d_k = 6 \\ 6, \text{ otherwise} \end{cases}$$

Diagonalized!

By design, Confuse_L can't be on the list L!

Confuse_L differs from the k^{th} element on the list L in the k^{th} position.

This contradicts the assumption that the list L is complete; i.e., that the map $f: N \text{ to } R_{[0,1]}$ is onto.

The set of reals is uncountable!
(Even the reals between 0 and 1.)

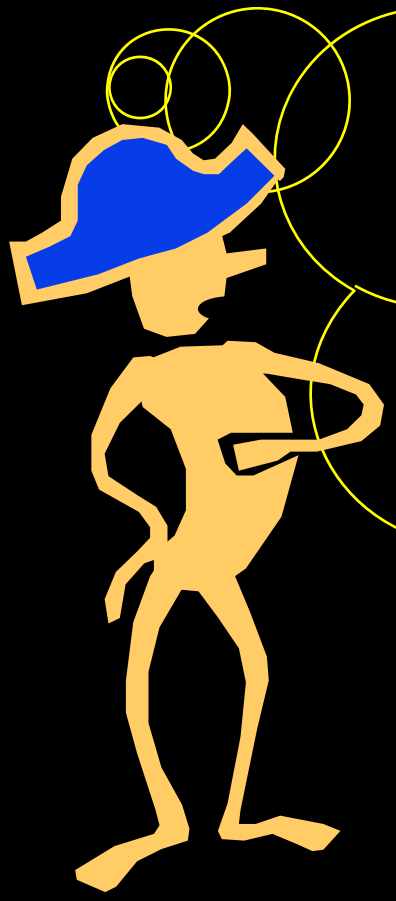
An aside: you **can** set up a correspondence between R and $R_{[0,1]}$.

Hold it!
Why can't the same argument be used to show that the set of rationals Q is uncountable?

The argument is the same for Q until the punchline.

However, since CONFUSE_L is not necessarily rational, there is no contradiction from the fact that it is missing from the list L.

Back to the questions we were asking earlier



Are all reals describable?
Are all reals computable?

We saw that
computable $\Rightarrow$
describable,

but do we also have
describable $\Rightarrow$
computable?

Standard Notation

Σ = Any finite alphabet

Example: {a,b,c,d,e,…,z}

Σ^* = All finite strings of symbols from Σ including the empty string ε

Theorem: Every infinite subset S of Σ^* is countable

Proof:

Sort S by first by length and then alphabetically.

Map the first word to 0, the second to 1, and so on….

Stringing Symbols Together

Σ = The symbols on a standard keyboard

For example:

The set of all possible Java programs is a subset of Σ^*

The set of all possible finite pieces of English text is a subset of Σ^*

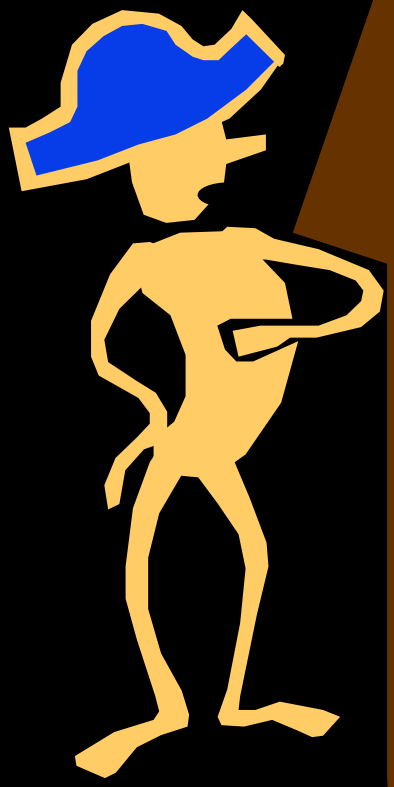
Thus:

The set of all possible Java programs is countable.

The set of all possible finite length pieces of English text is countable.

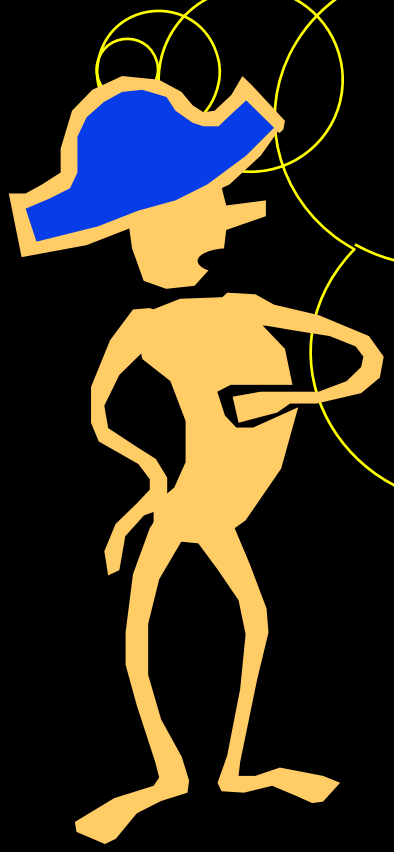
There are countably many Java program and uncountably many reals.

Hence,
Most reals are not computable!



I see!
There are countably many descriptions and uncountably many reals.

Hence:
Most real numbers are not describable!



Are all reals describable? NO

Are all reals computable? NO

We saw that computable $\Rightarrow$ describable,

but do we also have describable $\Rightarrow$ computable?

Is there a real number that can be described, but not computed?

Wait till the next lecture!

We know there are at least 2 infinities.
(the number of naturals, the number of reals.)

Are there more?

Definition: Power Set

The power set of S is the set of all subsets of S.

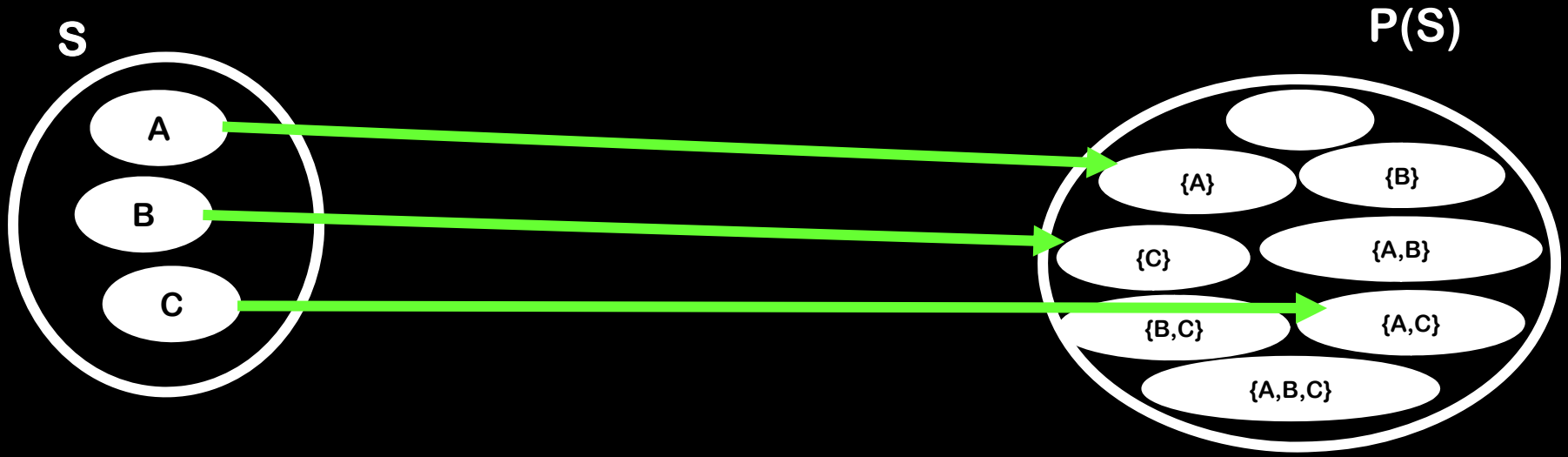
The power set is denoted as P(S).

Proposition:

If S is finite, the power set of S has cardinality $2^{|S|}$

Theorem: S can't be put into bijection with P(S)

Theorem: S can't be put into bijection with P(S)



Suppose f:S → P(S) is a bijection.

Let $CONFUSE_f$ contain all and only those elements that are not in the sets they map to

Since f is onto, exists $y \in S$ such that $f(y) = CONFUSE_f$.

Is y in $CONFUSE_f$?

YES: Definition of $CONFUSE_f$ implies no

NO: Definition of $CONFUSE_f$ implies yes

This proves that there are at least a countable number of infinities.

The first infinity is called:

$\aleph_0$

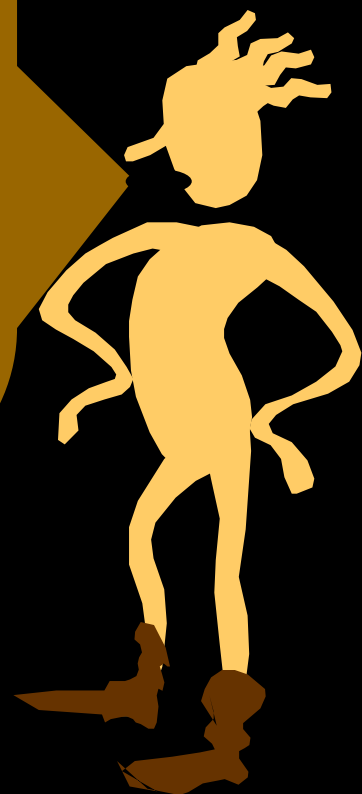
$|\mathbb{N}|, |\mathcal{P}(\mathbb{N})|, |\mathcal{P}(\mathcal{P}(\mathbb{N}))|, \ldots$

Are there any more infinities?

N, P(N), P(P(N)), …

Let S be the union of all of them!

Then S cannot be bijected to any of them!



In fact, the same argument can be used to show that no single infinity is big enough to count the number of infinities!

Cantor wanted to show that there was no infinity between |N| and |P(N)|

Cantor called his conjecture the "Continuum Hypothesis."

However, he was unable to prove it. This helped fuel his depression.

The Continuum Hypothesis can't be proved or disproved from the standard axioms of set theory!

This has been proved!

Little Susie and Little Johnny

What Little Susie should've said to Little Johnny:

Little Johnny: I hate you times infinity

Little Susie: I hate you times 2 to the infinity!

Little Johnny: I hate you times 2 to the 2 to the infinity!

…

Here's What You Need to Know…

Cantor's Definition:
Two sets have the same cardinality if there exists a bijection between them.

|E| = |N| = |Z| = |Q| (and proofs),
Cantor-Bernstein-Schroeder

Proof that there is no bijection between N and R

Countable versus Uncountable

Power sets and their properties