# 15-441 Computer Networks

Lecture 1

Dave Eckhardt
Peter Steenkiste

(some slides courtesy of Hui Zhang)

# Outline

- First things first
- Administrative overview
- Course non-goals
- Course goals
- Key problems
- Network performance concepts

# **First Things First**

#### Please read Chapter 1 of the text

- To read ahead: most of Chapter 2, but ok to skip for now:
  - 2.5, Reliable Transmission
  - 2.8, Wireless
  - Also, don't kill yourself on 2.4, Error Detection

### Please remind me to let you stretch

I haven't taught an 80-minute class in 1.5 years

## People

#### Professors

- Peter Steenkiste (www.cs/~prs)
- Dave Eckhardt (www.cs/~davide)

#### Teaching assistants

- Mike Cui
- Josh Hailpern
- David Murray
- [watch this space]

#### Course secretary

Barbara Grandillo, Wean Hall 8018

### **Information Sources**

#### Watch the course web page

- http://www.cs.cmu.edu/~441
- We expect you to read the syllabus!
- Handouts, readings, ...

### We expect you to read course bboards

- Official announcements
  - academic.cs.15-441.announce
- Questions/answers
  - academic.cs.15-441

# **Information Sources**

#### Textbook

Peterson and Davie, Computer Networks: A Systems Approach,
 3<sup>rd</sup> Edition, Morgan Kaufmann, 2003

### **Information Sources**

- ~30 lectures
- ~3 paper homeworks
- 1-2 lab homeworks
  - Illustrate networking concepts
- Mid-term and final exams
- 1 programming assignment
  - How to use a network
- 2 programming projects
  - How to build a network

# Grading

Homeworks	15%
Three projects	45%
Midterm exam	15%
Final exam	25%

#### Deadline means deadline

Deadline is 11:59 pm on the specified date

## **Policy on Collaboration**

### Working together is important

- Discuss course material in general terms
- Talk over tough debugging problems
- Parts of the course must be done individually
  - Homeworks, midterm, final, 1<sup>st</sup> programming assignment
- Projects are done by two-student teams
  - Learn how to collaborate
  - But each student must understand the entire project!
- Web page has the details

# Course Non-goals

#### Learn how to configure a Cisco router

- That requires a class all by itself
- Cisco teaches those classes
- Our perspective will be broader

### Become "Internet Experts"

- The Internet will be our frequent motivating example
- Our perspective will be broader

# Why not an "Internet class"?

Is there anything other than the Internet?

Philosophy final exam question:

Define "Universe". Give two examples.

## Why not an "Internet class"?

Is there anything other than the Internet?

Philosophy final exam question

Define "Universe". Give two examples.

- Yes Internet in the 1800's!
- Yes The secret network?
- [Yes What's next?]

### Internet in the 1800's!?!

### Tom Standage, The Victorian Internet

- Telegraph!
- Continent-spanning systems
- Digital transmission of information
- Crypto, cracking
- Nerds
- Attacks on the moral fiber of society
- On-line dating (even an on-line wedding!)

### Distributed message routing despite link outages

- Lines cut by armies in wartime!
- Many problems; solutions eerily similar

### **The Secret Network**

- With the Internet, who needs the phone system?
  - It's a "new era", etc.
- One small detail...
  - From inception, Internet has been a phone system <u>application!</u>
  - To connect two nodes, just ask your telco for a "circuit"...
    - ...somehow there's always copper/fiber waiting for you...
    - ...somehow when it breaks it gets fixed fast...
    - ...somehow your circuit can terminate anywhere...
  - Somehow?

### **Course Goals**

- Think "the network way"
  - Distributed coordination is hard, let's go shopping
- Learn how computer networks work
  - Problems, approaches, protocols, software
- Learn how to write network applications
- Hands-on understanding of network internals
  - Build a simple network in software

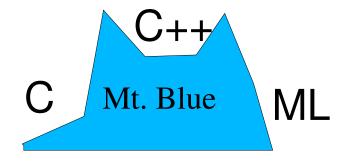
# **Selected Key Problems**

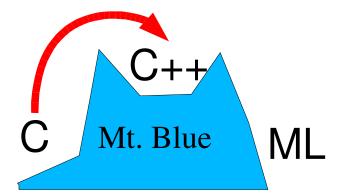
- Two Generals
- Group Membership
- Scaling

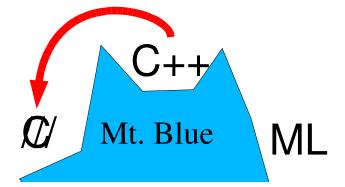
- Problem (formulated by Jim Gray?)
  - Two cooperating armies
    - Each size 2X
    - Separated by...
  - One opposing army
    - Size 3X
- Idealized "combat" (think: Diplomacy, Risk, ...)
  - 4X vs. 3X: probable win
  - 2X vs. 3X: certain loss

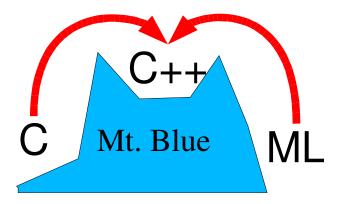
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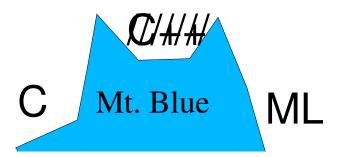
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  - Armies can communicate via messenger
- Protocol 0
  - C: "Attack at dawn!"
  - What if C's messenger is captured by C++?

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#### Protocol 1

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- ML: "Ok!"
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#### Protocol 1

- C: "Attack at dawn! Ok?"
- ML: "Ok!"
- What if ML's messenger is captured?

### Seemingly-trivial coordination is impossible!

## **Group Membership**

#### Group of nodes on a network

Require distributed election of a "leader"

#### Sample solution

- "Dstributed election" algorithm chooses among group members
- If a node enters or leaves during election, re-start algorithm

#### Results

- Works great for 10 nodes
- Fails horribly for 1,000,000 nodes
  - If inter-node-join time approximates election time...
  - Election process never completes

### **Group Membership**

- Problem: "group membership" is undefined
  - By the time you can compute it, it's changed
- Lots of algorithms will run into trouble
  - "To acquire a node number, add one to the largest current node number" – oops!
- Key network functions must face this environment
  - Routing, naming

## Scaling

#### "DOD Standard Internet Protocol"

- RFC 760, 1980: Addresses are fixed length of four octets (32 bits). An address begins with a one octet network number, followed by a three octet local address. This three octet field is called the "rest" field.
- Result: 254 networks (surely enough!)

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- ~16k "Class B" networks of ~64k hosts (CMU)
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- ~4m "Class C" networks of ~255 hosts
- Then "subnets", then "CIDR"
- "Surely enough" evaporates pretty fast!

# **Network Performance Concepts**

#### Throughput

- "How many things per unit time?"
- Mb/s = megabits per second
- KB/s = kilobytes per second

#### Latency

- "How long until my message arrives?"
- ms = millisecond (10<sup>-6</sup>), µs = microsecond (10<sup>-9</sup>)

#### Reciprocal "in theory"

bits/second = (1/(seconds/bit))

#### Relationship much more complex

### **Hen Performance**

#### Old riddle

"If a hen and a half lays an egg and a half in a day and a half, how long does it take to get a dozen eggs?"

#### Egg Latency

How long does it take for one hen to lay one egg?

#### Henhouse throughput (eggs per day)

- Increases with number of hens
- Does <u>not</u> mean you can build henhouse, get first egg in 1 hour
- What is minimum time to 12 eggs?

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### **Two Problems**

#### Small problem

- Mega != Mega
- Computer people: Megabyte (MB) = 2<sup>20</sup> bytes
- Network people: Megabit (Mb) = 106 bits
- It's 8.4 seconds, not 8

### **Two Problems**

#### Small problem

- Mega != Mega
- Computer people: MegaByte (MB) = 2<sup>20</sup> bytes
- Network people: Megabit (Mb) = 10<sup>6</sup> bits
- It's 8.4 seconds, not 8

#### Big problem

- I forgot to tell you... Mike lives on the Moon
  - (Extra credit: What is Mike's last name?)
- It takes radio waves 1.3 seconds to get there

### Message latency = sum of

- Propagation delay (distance/lightspeed)
- Transmission time (size/throughput)
- Queue delay (ignore for now)

#### Message to Mike

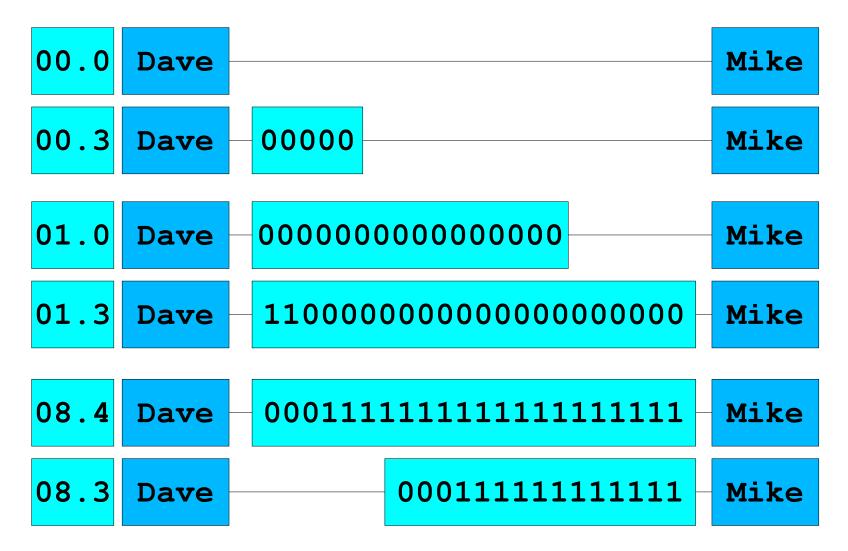
- Propagation delay is 1.3 seconds (one-way)
  - Also known as "link delay"
- Transmission time is 8.4 seconds, total is 9.7 (121% of 8)

### By the way: RTT (round-trip-time)

Time to send a 0-bit message there and back: 2.6 seconds

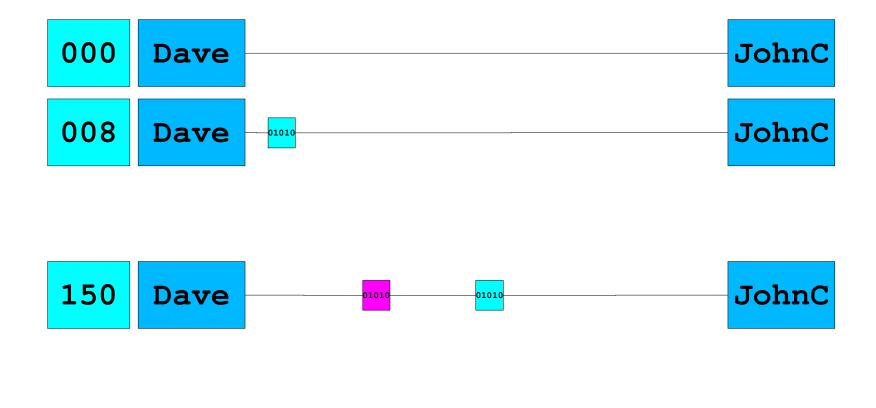
### Propagation delay vs. transmission time

- May vary widely
- Earth-to-Moon is 1.3 seconds (<< 8.4)</li>
  - Delay is a minor compared to transmission time
  - Can transmit part of message, receive back status
    - "Got that part ok" or "Oops, send it again"



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- Earth-to-Mars is 300-1225 seconds (>> 8.4)
  - Delay <u>vastly exceeds</u> transmission time
  - Link holds multiple entire messages



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**JohnC** 

300

42

Dave

#### Bandwidth-delay product

- megabits/second X link-delay
- This many bits are always "in flight" / "queued in link"

#### What if Mike says "Stop!! My buffer is full!"?

- One b-d product of bits are "in flight" to him already
- You will queue <u>another</u> b-d product before you hear his alert!

### Message throughput (≠ link throughput)

- How many messages per second can you send to Mike?
- Depends on b-d product vs. message size
- Depends on message protocol (= waiting protocol) you use

#### See text for more-dignified treatment

No hens, no Loonies, no Martians

#### Things to watch out for

- Is "delay" one-way or round-trip?
- Mega vs. mega, kilo vs. kilo
- Do we mean link latency or message latency?
- Do we mean link throughput or message throughput?

# Things Which Aren't Throughput

#### Bandwidth

- Properly, measured in Hertz
- Difference between max & min frequency of transmission band
- Routinely abused by CS people to mean "throughput"

#### Goodput

- Used to mean "productive throughput"
- Ignore "waste" if part of a message is transmitted multiple times

### **Back to the Internet**

#### Another reason the Internet isn't perfect

- Fatally overoptimized for single-planet case
- Will work to low-Earth orbit
- Efficiency problems talking to the Moon
- Forget about Mars

### InterPlanetary Internet

http://www.ipnsig.org/

# Summary

#### First things first

- Read Chapter 1
- Study socket-programming example as a refresher
- Project 1 (individual) out Wednesday
- Course non-goals, goals
  - "Networking perspective", Internet as running example
- Key problems
  - Distributed coordination; scaling
- Network performance concepts
  - Throughput vs. latency, ...