Lecture 9

Induction Variables and Strength Reduction

I. Overview of optimization

II. Algorithm to find induction variables

Example

```
FOR i = 0 to 100
   A[i] = 0;
```

```
    i = 0
L2: IF i>=100 GOTO L1
    t1 = 4 * i
    t2 = &A + t1
    *t2 = 0
    i = i+1
    GOTO L2
L1:
```

Definitions

- A **basic induction variable** is
  - a variable X whose only definitions within the loop are assignments of the form:

    $X = X+c$ or $X = X-c$,

    where **c** is either a constant or a loop-invariant variable.

- An **induction variable** is
  - a basic induction variable, or
  - a variable defined once within the loop, whose value is a linear function of some basic induction variable at the time of the definition:
    $A = c_1 * B + c_2$

- The **FAMILY of a basic induction variable B** is
  - the set of induction variables A such that each time A is assigned in the loop, the value of A is a linear function of B.

Optimizations

1. **Strength reduction:**
   - Let A be an induction variable in family of basic induction variable B ($A = c_1 * B + c_2$)
     - Create new variable: A'
     - Initialization in preheader: $A' = c_1 * B + c_2$;
     - Track value of B: add after B=B+x: $A' = A' + x * c_1$;
     - Replace assignment to A: A=A'

Optimizations (continued)

2. **Optimizing non-basic induction variables**
   - copy propagation
   - dead code elimination
3. **Optimizing basic induction variables**
   - Eliminate basic induction variables used only for
     - calculating other induction variables and loop tests
   - Algorithm:
     - Select an induction variable A in the family of B, preferably with simple constants ($A = c_1 * B + c_2$).
     - Replace a comparison such as

       `if B > X goto L1`

       with

       `if (A' > c1 * X + c2) goto L1` (assuming c_1 is positive)
     - if B is live at any exit from the loop, recompute it from A'
       - After the exit, $B = (A' - c_2) / c_1$

II. Basic Induction Variables

- **A BASIC induction variable in a loop L**
  - a variable X whose only definitions within L are assignments of the form X = X+c or X = X-c, where c is either a constant or a loop-invariant variable.
- **Algorithm: can be detected by scanning L**
- Example:

```
k = 0;
for (i = 0; i < n; i++) {
   k = k + 3;
   … = m;
   if (x < y)
      k = k + 4;
   if (a < b)
      m = 2 * k;
   k = k – 2;
   … = m;
```

Each iteration may execute a different number of increments/decrements!!

Strength Reduction Algorithm

- **Key idea**:
  - For each induction variable A, ($A = c_1*B+c_2$ at time of definition)
    - variable A' holds expression c_1*B+c_2 at all times
    - replace definition of A with A=A' only when executed

- **Result**:
  - Program is correct
  - Definition of A does not need to refer to B

Finding Induction Variable Families

- **Let B be a basic induction variable**
  - Find all induction variables A in family of B:
    - $A = c_1 * B + c_2$
      (where B refers to the value of B at time of definition)
- **Conditions:**
  - If A has a single assignment in the loop L, and assignment is one of:

```
A = B * c
A = c * B
A = B / c    (assuming A is real)
A = B + c
A = c + B
A = B - c
A = c - B
```

  - OR, ... (next page)

Finding Induction Variable Families (continued)

- Let D be an induction variable in the family of B (D = c_1* B + c_2)
  - If A has a single assignment in the loop L, and assignment is one of:

```
A = D * c
A = c * D
A = D / c    (assuming A is real)
A = D + c
A = c + D
A = D - c
A = c - D
```

  - No definition of D outside L reaches the assignment to A
  - Between the lone point of assignment to D in L and the assignment to A, there are no definitions of B

Summary

- **Precise definitions of induction variables**
- **Systematic identification of induction variables**
- **Strength reduction**
- **Clean up:**
  - eliminating basic induction variables
    - used in other induction variable calculations
    - replacement of loop tests
  - eliminating other induction variables
    - standard optimizations