

AP CS A and AB: New/Experienced A Tall Order?

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Schedule

- **8:30 – 10:00**
 - Intro, resources, what's new
- **10:15 - 12:00**
 - Design
 - Inheritance, Interfaces & Abstract Classes
- **12:45 – 3:30**
 - Collections -> Analysis -> Big O
 - Reading sample questions and their rubrics

Intro

- A v. AB?
- Years teaching?
- Java experience?
- The cards (years teaching A/AB, 2 things)
- Materials (AB q/ref, A2 & AB1 samples, role play)

Resources

- **AP Central**
 - **Course descriptions**
 - **Java subset**
 - **Sample syllabi**
- **AP list-serve**
- **JETT workshops**
 - **8/6 & 7 at Florida International**

What's new?

- **2-D arrays to AB**
- **Well, Java...**
 - **No reference parameters**
 - **Collections**
- **But also...**
 - **Design**
 - **Analysis**
 - **Priority Queues**

Design

- **Interfaces – proscribe a set of behaviors**
 - methods are public by default
 - NO implementation can be provided
 - NO instance variables can be declared
 - cannot, therefore, be instantiated
 - what does that mean?
 - examples: stack and queue
- **Classes realize interfaces by implementing all methods**

Design...

- **vs. Abstract classes**
 - some methods implemented
 - others designated as abstract (which forces the class to be abstract) – these must be implemented by class(es) that extend this abstract class (ultimately)
 - can have instance variables

Design...

- Let's look at A2
- Let's look at some other classes

Collections

- Lists, Sets, Maps
- Raises the level of design (and discourse)
- But, then, what is (are) the issue(s)?

Collections...

- What do they do?

- List

- ordered (positionally) – a sequence
- duplicates?

- Set

- unordered collection
- duplicates?

- Map

- Maps (unique) keys to values

Collections...

- How do they do it?
 - Analyze their algorithms
 - But what does that mean?

Algorithm Analysis (kudos to ola)

- How do you measure performance?
 - It's faster! It's more elegant! It's safer! It's cooler!
- Use mathematics to analyze the *algorithm* (performance as function of input size)
- Implementation is another matter
 - cache, compiler optimizations, OS, memory,...

What do we need?

- **Need empirical tests and mathematical tools**
 - **Compare by running**
 - 30 seconds vs. 3 seconds,
 - 5 hours vs. 2 minutes
 - Two weeks to implement code
- **We need a vocabulary to discuss tradeoffs**

What is big-Oh about?

- Intuition: avoid details when they don't matter, and they don't matter when input size (N) is big enough
 - For polynomials, use only leading term, ignore coefficients: linear, quadratic

$$y = 3x$$

$$y = x^2$$

$$y = 6x - 2$$

$$y = x^2 - 6x + 9$$

$$y = 15x + 44$$

$$y = 3x^2 + 4x$$

O-notation, family of functions

- **first family is $O(n)$, the second is $O(n^2)$**
 - **Intuition: family of curves, same shape**
 - **More formally: $O(f(n))$ is an *upper-bound*, when n is large enough the expression $cf(n)$ is larger**
 - **Intuition: linear function: double input, double time, quadratic function: double input, quadruple the time**

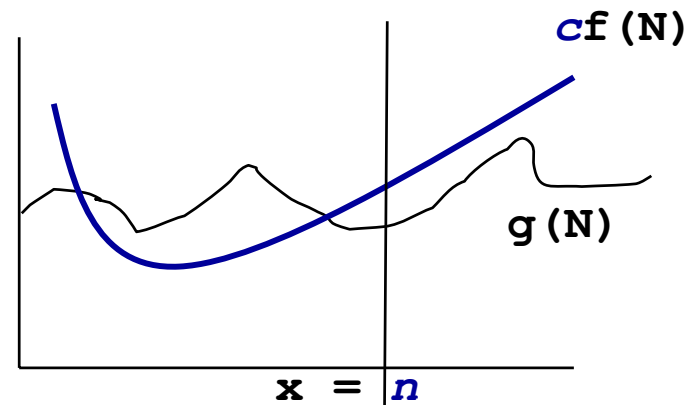
Reasoning about algorithms

- **We have an $O(n)$ algorithm,**
 - For 5,000 elements takes 3.2 seconds
 - For 10,000 elements takes 6.4 seconds
 - For 15,000 elements takes?

- **We have an $O(n^2)$ algorithm**
 - For 5,000 elements takes 2.4 seconds
 - For 10,000 elements takes 9.6 seconds
 - For 15,000 elements takes ...?

More formal definition

- O-notation is an upper-bound, this means that N is $O(N)$, but it is also $O(N^2)$; we try to provide *tight* bounds. Formally:
 - A function $g(N)$ is $O(f(N))$ if there exist constants c and n such that $g(N) < cf(N)$ for all $N > n$



Other definitions

- $g(n)$ is $O(f(n))$ if $\lim g(n)/f(n) = c$ as $n \rightarrow \infty$
- Informally, think of O as " \leq "
- Similarly, there's a notation for lower bounds...

Big-Oh calculations from code

- Search for element in array:
 - What is complexity (using O-notation)?
 - If array doubles, what happens to time?

```
for(int i=0; i < a.length; i++) {  
    if (a[i].equals(target)) return true;  
}  
return false;
```

- Best case? Average case? Worst case?

Measures of complexity

- **Worst case**
 - Good upper-bound on behavior
 - Never get worse than this

- **Average case**
 - What does average mean?
 - Averaged over all inputs? Assuming uniformly distributed random data?

Some helpful mathematics

- $1 + 2 + 3 + 4 + \dots + N$
 - $N(N+1)/2 = N^2/2 + N/2$ is $O(N^2)$
- $N + N + N + \dots + N$ (total of N times)
 - $N * N = N^2$ which is $O(N^2)$
- $1 + 2 + 4 + \dots + 2^N$
 - $2^{N+1} - 1 = 2 \times 2^N - 1$ which is $O(2^N)$

The usual suspects

- $O(1)$
- $O(\log n)$
- $O(n)$
- $O(n \log n)$
- $O(n^2)$
- $O(n^3)$
- ...
- $O(2^n)$

Multiplying and adding big-Oh

- **Suppose we do a linear search then we do another one**
 - **What is the complexity?**
 - **If we do 100 linear searches?**
 - **If we do n searches on an array of size n ?**

Multiplying and adding

- **Binary search followed by linear search?**
 - **What are big-Oh complexities? Sum?**
 - **50 binary searches? N searches?**

- **What is the number of elements in the list (1,2,2,3,3,3)?**
 - **What about (1,2,2, ..., n,n,...,n)?**
 - **How can we reason about this?**

Analysis for AP collections

- ListQueue
- ArrayStack (where's the top?)

- Lists, Maps, Sets

Priority Queues

- Idea
- Implementation(s)

Language details

- **No reference parameters**
 - All passes are by value (primitives/objects)
- **How do you change something**
 - Through “modifier” methods
 - Through returning a reference to a modified object ($x = \text{“changed } x\text{”}$)