

## 15-859(B) Machine Learning Theory

### Lecture 1: intro, basic models and issues

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### Admin

- Course web page. Textbook covers about 1/2 of course material.
- 6 hwk assignments. Exercises/problems.
- Small project: explore a theoretical question, try some experiments, or read a paper and explain the idea. Short writeup and possibly presentation. Small groups ok.
- Take-home exam (worth roughly 2 hwks).
- "volunteers" for hwk grading.

OK, let's get to it...

### Machine learning can be used to...

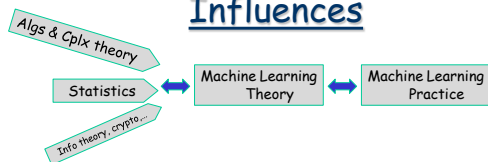
- recognize speech, faces,
- play games, steer cars,
- adapt programs to users,
- classify documents, protein sequences,...

#### Goals of machine learning theory:

develop and analyze models to understand:

- what kinds of tasks we can hope to learn, and from what kind of data,
- what types of guarantees might we hope to achieve,
- other common issues that arise.

### Influences



#### Goals of machine learning theory:

develop and analyze models to understand:

- what kinds of tasks we can hope to learn, and from what kind of data,
- what types of guarantees might we hope to achieve,
- other common issues that arise.

### A typical setting

- Imagine you want a computer program to help you decide which email messages are **spam** and which are important.
- Might represent each message by  $n$  features. (e.g., return address, keywords, spelling, etc.)
- Take sample  $S$  of data, labeled according to whether they were/weren't **spam**.
- Goal of algorithm is to use data seen so far produce good prediction rule (a "hypothesis")  $h(x)$  for future data.

### The concept learning setting

E.g.,

| \$\$ | meds | Mr. | bad spelling | known-sender | spam? |
|------|------|-----|--------------|--------------|-------|
| Y    | N    | Y   | Y            | N            | Y     |
| N    | N    | N   | Y            | Y            | N     |
| N    | Y    | N   | N            | N            | Y     |
| Y    | N    | N   | N            | Y            | N     |
| N    | N    | Y   | N            | Y            | N     |
| Y    | N    | N   | Y            | N            | Y     |
| N    | N    | Y   | N            | N            | N     |
| N    | Y    | N   | Y            | N            | Y     |

Given data, some reasonable rules might be:

- Predict **SPAM** if  $\neg$ known AND ( $$$$$  OR meds)
- Predict **SPAM** if  $$$$ +$  meds - known  $>$  0.

•...

## Big questions

- (A) How might we automatically generate rules that do well on observed data?  
[algorithm design]
- (B) What kind of confidence do we have that they will do well in the future?  
[confidence bound / sample complexity]

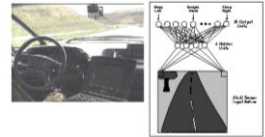
for a given learning alg. how much data do we need...

## Power of basic paradigm

Many problems solved by converting to basic "concept learning from structured data" setting.

- E.g., document classification
  - convert to bag-of-words
  - Linear separators do well
- E.g., driving a car

- convert image into features.
- Use neural net with several outputs.



## Natural formalization (PAC)

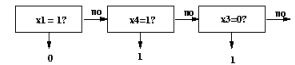
Email msg Spam or not?

- We are given sample  $S = \{(x,y)\}$ .
  - View labels  $y$  as being produced by some target function  $f$ .
- Alg does optimization over  $S$  to produce some hypothesis (prediction rule)  $h$ .
- Assume  $S$  is a random sample from some probability distribution  $D$ . Goal is for  $h$  to do well on new examples also from  $D$ .

I.e.,  $\Pr_D[h(x) \neq f(x)] < \epsilon$ .

err(h)

## Example of analysis: Decision Lists



Say we suspect there might be a good prediction rule of this form.

- Design an efficient algorithm  $A$  that will find a consistent DL if one exists.
- Show that if  $S$  is of reasonable size, then  $\Pr[\text{exists consistent DL } h \text{ with } \text{err}(h) > \epsilon] < \delta$ .
- This means that  $A$  is a good algorithm to use if  $f$  is, in fact, a DL.

If  $S$  is of reasonable size, then  $A$  produces a hypothesis that is Probably Approximately Correct.

## How can we find a consistent DL?

| $x_1$ | $x_2$ | $x_3$ | $x_4$ | $x_5$ | label |
|-------|-------|-------|-------|-------|-------|
| 1     | 0     | 0     | 1     | 1     | +     |
| 0     | 1     | 1     | 0     | 0     | -     |
| 1     | 1     | 1     | 0     | 0     | +     |
| 0     | 0     | 0     | 1     | 0     | -     |
| 1     | 1     | 0     | 1     | 1     | +     |
| 1     | 0     | 0     | 0     | 1     | -     |

if ( $x_1=0$ ) then -, else  
 if ( $x_2=1$ ) then +, else  
 if ( $x_4=1$ ) then +, else -

## Decision List algorithm

- Start with empty list.
- Find if-then rule consistent with data.  
(and satisfied by at least one example)
- Put rule at bottom of list so far, and cross off examples covered. Repeat until no examples remain.

If this fails, then:

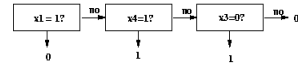
- No DL consistent with remaining data.
- So, no DL consistent with original data.

OK, fine. Now why should we expect it to do well on future data?

## Confidence/sample-complexity

- Consider some DL  $h$  with  $\text{err}(h) > \epsilon$ , that we're worried might fool us.
  - Chance that  $h$  is consistent with  $S$  is at most  $(1-\epsilon)^{|S|}$ .
  - Let  $|H|$  = number of DLs over  $n$  Boolean features.  $|H| < n!4^n$ . (for each feature there are 4 possible rules, and no feature will appear more than once)
- So,  $\Pr[\text{some DL } h \text{ with } \text{err}(h) > \epsilon \text{ is consistent}] \leq |H|(1-\epsilon)^{|S|} \leq |H|e^{-\epsilon|S|}$ .
- This is  $< \delta$  for  $|S| > (1/\epsilon)[\ln(|H|) + \ln(1/\delta)]$  or about  $(1/\epsilon)[n \ln n + \ln(1/\delta)]$

## Example of analysis: Decision Lists



Say we suspect there might be a good prediction rule of this form.

- DONE** Design an efficient algorithm  $A$  that will find a consistent DL if one exists.
- DONE** Show that if  $|S|$  is of reasonable size, then  $\Pr[\text{exists consistent DL } h \text{ with } \text{err}(h) > \epsilon] < \delta$ .
- So, if  $f$  is in fact a DL, then whp  $A$ 's hypothesis will be approximately correct. "PAC model"

## PAC model more formally:

- We are given sample  $S = \{(x,y)\}$ .
  - Assume  $x$ 's come from some fixed probability distribution  $D$  over instance space.
  - View labels  $y$  as being produced by some target function  $f$ .
- Alg does optimization over  $S$  to produce some hypothesis (prediction rule)  $h$ . Goal is for  $h$  to do well on new examples also from  $D$ . I.e.,  $\Pr_D[h(x) \neq f(x)] < \epsilon$ .

Algorithm PAC-learns a class of functions  $C$  if:

- For any given  $\epsilon > 0, \delta > 0$ , any target  $f \in C$ , any dist.  $D$ , the algorithm produces  $h$  of  $\text{err}(h) < \epsilon$  with prob. at least  $1-\delta$ .
- Running time and sample sizes polynomial in relevant parameters:  $1/\epsilon, 1/\delta, n$  (size of examples),  $\text{size}(f)$ .
- Require  $h$  to be poly-time evaluable. Learning is called "proper" if  $h \in C$ . Can also talk about "learning  $C$  by  $H$ ".

We just gave an alg to PAC-learn decision lists.

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Some notes:

- Can either view alg as requesting examples (button/oracle model) or just as function of  $S$ , with guarantee if  $S$  is suff. lg.
- "size(f)" term comes in when you are looking at classes where some fns could take  $> \text{poly}(n)$  bits to write down. (e.g., decision trees, DNF formulas)

## Confidence/sample-complexity

- What's great is there was nothing special about DLs in our argument.
- All we said was: "if there are not *too* many rules to choose from, then it's unlikely one will have fooled us just by chance."
- And in particular, the number of examples needs to only be proportional to  $\log(|C|)$ . (notice big difference between  $|C|$  and  $\log(|C|)$ )

## Occam's razor

William of Occam (~1320 AD):

"entities should not be multiplied unnecessarily" (in Latin)

Which we interpret as: "in general, prefer simpler explanations".

Why? Is this a good policy? What if we have different notions of what's simpler?

## Occam's razor (contd)

A computer-science-ish way of looking at it:

- Say "simple" = "short description".
- At most  $2^s$  explanations can be  $< s$  bits long.
- So, if the number of examples satisfies:

Think of as 10x #bits to write down h.

$$|S| > (1/\epsilon)[s \ln(2) + \ln(1/\delta)]$$

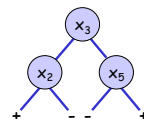
Then it's unlikely a bad simple explanation will fool you just by chance.

## Occam's razor (contd)<sup>2</sup>

Nice interpretation:

- Even if we have different notions of what's simpler (e.g., different representation languages), we can both use Occam's razor.
- Of course, there's no guarantee there will be a short explanation for the data. That depends on your representation.

## Decision trees



- Decision trees over  $\{0,1\}^n$  not known to be PAC-learnable.
- Given any data set  $S$ , it's easy to find a consistent DT if one exists. How?
- Where does the DL argument break down?
- Simple heuristics used in practice (ID3 etc.) don't work for all  $c \in C$  even for uniform  $D$ .
- Would suffice to find the (apx) smallest DT consistent with any dataset  $S$ , but that's NP-hard.

## More examples

Other classes we can PAC-learn: (how?)

- Monomials [conjunctions, AND-functions]
  - $x_1 \wedge x_4 \wedge x_6 \wedge x_9$
- 3-CNF formulas (3-SAT formulas)
- OR-functions, 3-DNF formulas
- k-Decision lists (each if-condition is a conjunction of size k), k is constant.

Given a data set  $S$ , deciding if there is a consistent 2-term DNF formula is NP-complete. Does that mean 2-term DNF is hard to learn?

## More examples

Hard to learn  $C$  by  $C$ , but easy to learn  $C$  by  $H$ , where  $H = \{2\text{-CNF}\}$ .

Given a data set  $S$ , deciding if there is a consistent 2-term DNF formula is NP-complete. Does that mean 2-term DNF is hard to learn?

## If computation-time is no object, then any class is PAC-learnable

- Occam bounds  $\Rightarrow$  any class is learnable if computation time is no object:
  - Let  $s_1=10, \delta_1 = \delta/2$ . For  $i=1,2,\dots$  do:
    - Request  $(1/\delta)[s_i + \ln(1/\delta_i)]$  examples  $S_i$ .
    - Check if there is a function of size at most  $s_i$  consistent with  $S_i$ . If so, output it and halt.
    - $s_{i+1} = 2s_i, \delta_{i+1} = \delta_i/2$ .
  - At most  $\delta_1 + \delta_2 + \dots \leq \delta$  chance of failure.
  - Total data used:  $O((1/\epsilon)[\text{size}(f)+\ln(1/\delta)\ln(\text{size}(f))])$ .

## More about the PAC model

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- Running time and sample sizes polynomial in relevant parameters:  $1/\epsilon, 1/\delta, n, \text{size}(f)$ .
- Require  $h$  to be poly-time evaluable. Learning is called "proper" if  $h \in C$ . Can also talk about "learning  $C$  by  $H$ ".
- What if your alg only worked for  $\delta = \frac{1}{2}$ , what would you do?
- What if it only worked for  $\epsilon = \frac{1}{4}$ , or even  $\epsilon = \frac{1}{2}-1/n$ ? This is called **weak-learning**. Will get back to later.
- **Agnostic learning** model: Don't assume anything about  $f$ . Try to reach  $\text{err}(H) + \epsilon$ .

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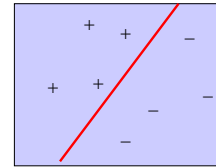
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Drawbacks of model:

- In the real world, labeled examples are much more expensive than running time. **Poly(size(f)) not enough**.
- "Prior knowledge/beliefs" might be not just over form of target but other relations to data.
- Doesn't address other kinds of info (**cheap unlabeled data, pairwise similarity information**).
- Only considers "one shot" learning.

## Extensions we'll get at later:

- Replace  $\log(|H|)$  with "effective number of degrees of freedom".



- There are infinitely many linear separators, but not that many really different ones.
- Other more refined analyses.

## Some open problems

Can one efficiently PAC-learn...

- an intersection of 2 halfspaces? (**2-term DNF trick doesn't work**)
- $C=\{\text{fns with only } O(\log n) \text{ relevant variables}\}$ ? (or even  $O(\log \log n)$  or  $\omega(1)$  relevant variables)? **This is a special case of DTs, DNFs.**
- Monotone DNF over uniform  $D$ ?
- Weak agnostic learning of monomials.