

RECITATION 4 – THE STACK

15-213-m12

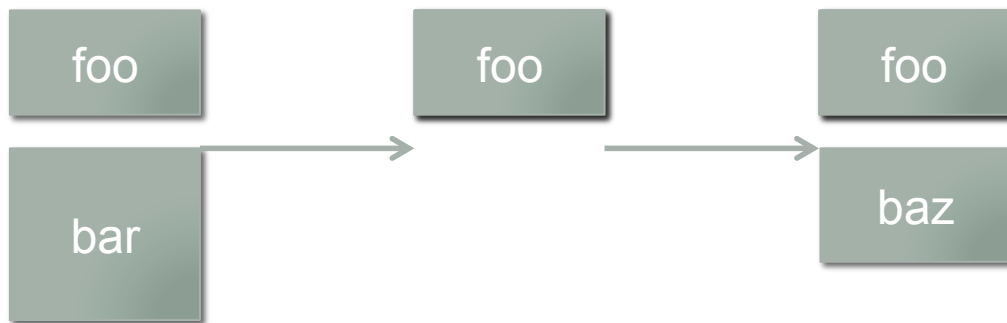
Rick Benua

The Stack

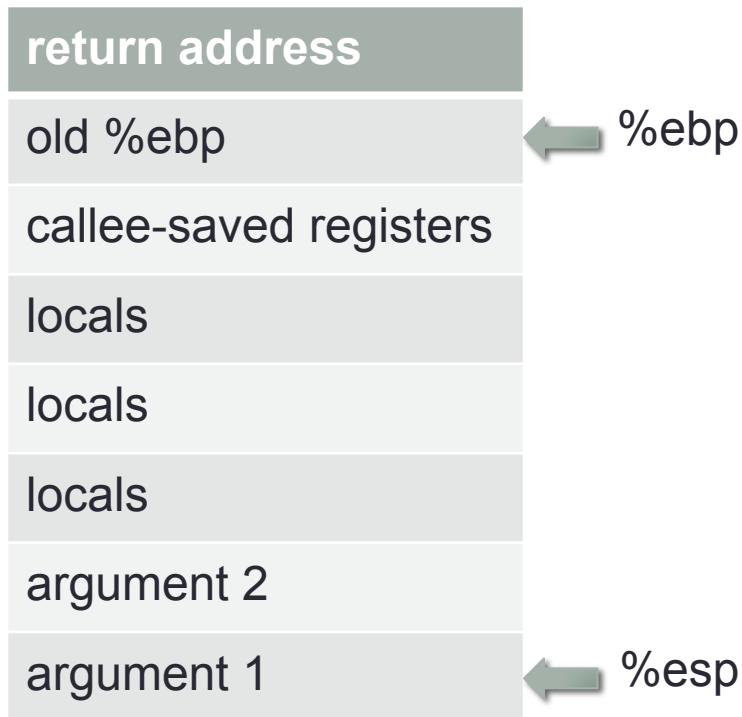
- Region of memory dedicated to local variables and arguments for **all** functions currently being executed
- Maintained using registers `%esp` and `%ebp` (on IA32)
 - `%esp` points to the top of the stack (**actually the lowest address**)
 - `%ebp` points to the base of the current “frame” – section of data associated with current function
- Modern compilers don't need `%ebp` for this
 - Omitted by default on x86-64
 - `%rbp` can be another GPR
 - Can pass compiler flags to omit it on IA32

The Stack

- Memory on the stack can be accessed without checks
 - Callee reaches into caller's frame to find arguments
 - Caller may pass a pointer into its stack frame to callee
 - (read as input, or write result, or both!)
- Callee may **NOT** return pointers into its stack
 - Stack space is "freed" upon return
 - Reused for next function call

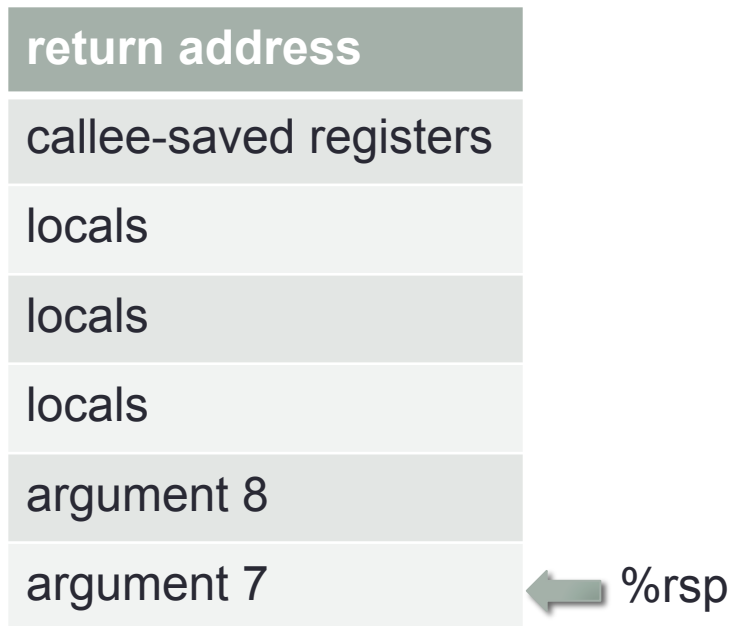


Anatomy of a Stack Frame – IA32



- Just before calling a function
- arguments to next call pushed on stack in reverse order

Anatomy of a Stack Frame – x86-64



- No base pointer – compiler uses offset from %rsp to find return value
- Arguments passed in registers, but can spill over onto the stack

Buffer Lab

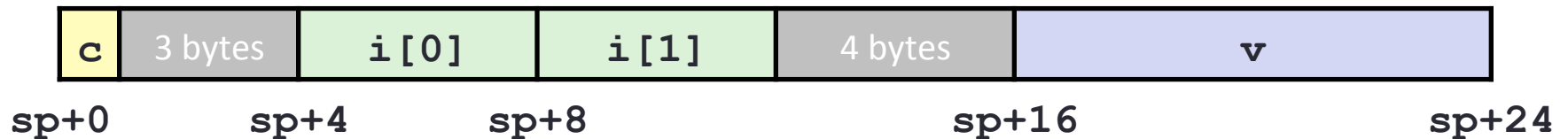
- Out now!
- Due Tuesday
- More examination of programs
 - Create buffer overflow exploits for a known program
- READ THE HANDOUT
 - FOR THE LOVE OF GOD, READ THE ENTIRE HANDOUT
- Series of incrementally more complex exploits

Buffer Overflow

- Common idiom in code: Copy input from user into buffer, then process it
- Copy may not check length of input
 - Part of the point of this lab is to teach you to **not** do that
- Can reach beyond buffer into other parts of stack
- Strings generally written in from low – high addresses
- “up” the stack, including into saved `%ebp` or return address!
- This is very bad.

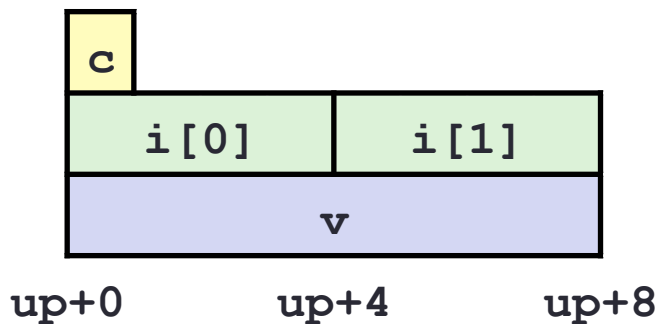
Structures

- Structures combine sets of related values that can be passed around together
- Values not necessarily contiguous in memory
 - Each value must be aligned to its size
 - Entire struct must be aligned to the largest constraint of any member
- Each member is at a constant offset from the beginning of the struct



Unions

- Structures store values “next to” each other
- Unions store values “on top of” each other
- Casting between types does conversion
- Union access does not



```

0x00001f30 <main+0>: push  %ebp
0x00001f31 <main+1>: mov   %esp,%ebp
0x00001f33 <main+3>: sub   $0x10,%esp
0x00001f36 <main+6>: movl  $0x39,-0xc(%ebp)
0x00001f3d <main+13>: movl  $0x0,-0x10(%ebp)
0x00001f44 <main+20>: jmp   0x1f62 <main+50>
0x00001f46 <main+22>: mov   -0xc(%ebp),%eax
0x00001f49 <main+25>: add   $0x1,%eax
0x00001f4c <main+28>: mov   %eax,%ecx
0x00001f4e <main+30>: shr   $0x1f,%ecx
0x00001f51 <main+33>: lea  (%eax,%ecx,1),%eax
0x00001f54 <main+36>: sar   %eax
0x00001f56 <main+38>: mov   %eax,-0xc(%ebp)
0x00001f59 <main+41>: mov   -0x10(%ebp),%eax
0x00001f5c <main+44>: lea  0x1(%eax),%eax
0x00001f5f <main+47>: mov   %eax,-0x10(%ebp)
0x00001f62 <main+50>: mov   -0x10(%ebp),%eax
0x00001f65 <main+53>: cmp   $0x7,%eax
0x00001f68 <main+56>: jle  0x1f46 <main+22>
0x00001f6a <main+58>: mov   -0x10(%ebp),%eax
0x00001f6d <main+61>: cmp   $0x1,%eax
0x00001f70 <main+64>: je   0x1f7b <main+75>
0x00001f72 <main+66>: movl  $0x1,-0x8(%ebp)
0x00001f79 <main+73>: jmp   0x1f82 <main+82>
0x00001f7b <main+75>: movl  $0x0,-0x8(%ebp)
0x00001f82 <main+82>: mov   -0x8(%ebp),%eax
0x00001f85 <main+85>: mov   %eax,-0x4(%ebp)
0x00001f88 <main+88>: mov   -0x4(%ebp),%eax
0x00001f8b <main+91>: add   $0x10,%esp
0x00001f8e <main+94>: pop   %ebp
0x00001f8f <main+95>: ret

```

```

int main(){
    int x = 57;
    int y = 0;
    for(; y < 8; y++){
        x = (x + 1) / 2;
    }
    if(y != 1){
        return 1;
    }
    else{
        return 0;
    }
}

```

```

struct{
  int i;
  char c[3];
  struct s *n;
  double d;
  short s;
} s;

```

0x00	i	i	i	i	c[0]	c[1]	c[2]	--
0x08	n	n	n	n	--	--	--	--
0x10	d	d	d	d	d	d	d	d
0x18	s	s	--	--	--	--	--	--