

15-745 Lecture 2b

Dataflow Analysis
Basic Blocks
Related Optimizations
Clarifications and more

Copyright © Seth Copen Goldstein 2005,
Tim Callahan 2007

Lecture 2 1
15-745 © Seth Copen Goldstein 2005

An Improvement: Basic Blocks

- No need to compute this one stmt at a time
- For straight line code:
 - $In[s1; s2] = in[s1]$
 - $Out[s1; s2] = out[s2]$
- Can we combine the gen and kill sets into one set per BB?

	Gen	kill
1: $i \leftarrow 1$	1	8, 4
2: $j \leftarrow 2$	2	
3: $k \leftarrow 3 + i$	3	11
4: $i \leftarrow j$	4	1, 8
5: $m \leftarrow i + k$	5	

- $Gen[BB] = \{2, 3, 4, 5\}$
- $Kill[BB] = \{1, 8, 11\}$
- $Gen[s1; s2] =$
- $Kill[s1; s2] =$

Lecture 2 2
15-745 © Seth Copen Goldstein 2005

An Improvement: Basic Blocks

- No need to compute this one stmt at a time
- For straight line code:
 - $In[s1; s2] = in[s1]$
 - $Out[s1; s2] = out[s2]$
- Can we combine the gen and kill sets into one set per BB?

	Gen	kill
1: $i \leftarrow 1$	1	8, 4
2: $j \leftarrow 2$	2	
3: $k \leftarrow 3 + i$	3	11
4: $i \leftarrow j$	4	1, 8
5: $m \leftarrow i + k$	5	

- $Gen[BB] = \{2, 3, 4, 5\}$
- $Kill[BB] = \{1, 8, 11\}$
- $Gen[s1; s2] = (Gen[s1] - Kill[s2]) + Gen[s2]$
- $Kill[s1; s2] = (Kill[s1] - Gen[s2]) + Kill[s2]$

Lecture 2 3
15-745 © Seth Copen Goldstein 2005

Def-use chains are valuable too

- Def-use chain: for each definition of var x , a list of all uses of that definition
- Computed from ~~liveness analysis~~ *upward-exposed uses*, a backward dataflow problem
- Def-use is **NOT** symmetric to use-def?

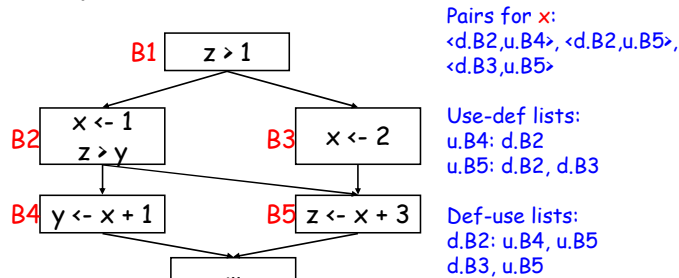
```

graph TD
    A["z > 1"] --> B["x <- 1  
z > y"]
    A --> C["x <- 2"]
    B --> D["y <- x + 1"]
    B --> E["z <- x + 3"]
    C --> E
    D --> F["..."]
    E --> F
    
```

Lecture 2 4
15-745 © Seth Copen Goldstein 2005

Def-use chains are valuable too

- Def-use is **IS** symmetric to use-def.
- If d is in u's UD chain, then u is in d's DU chain.
- Can compute DU chains from UD chains - same basic info, a set of $\langle u, d \rangle$ pairs, just packaged differently.



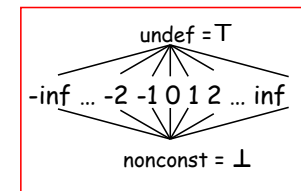
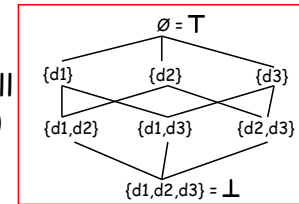
Lecture 2

15-745 © Seth Copen Goldstein 2005

5

What the ... is a Lattice?

- Represents values: for one item, or vector of all (often boolean to powerset)
- Has a defined top and bot
- According to ASU:
 - Top is least info: $\text{top} \wedge X = X$
 - Bot is end: $\text{bot} \wedge X = \text{bot}$
 - Init in[b] with top, out[b] with $F_b(\text{top})$.



Lecture 2

15-745 © Seth Copen Goldstein 2005

6

Dragon versus Muchnick

- From now on, follow Dragon:
- Meet function = confluence for any problem
- Always start at top of lattice, meaning least information
- Meet $(x \wedge y)$ is the closest point where x and y run into each other going down the lattice; might be x, y, or other

(Muchnick seems to allow either direction on lattice, so either starting point, and defines join in addition to meet...)

Lecture 2

15-745 © Seth Copen Goldstein 2005

7

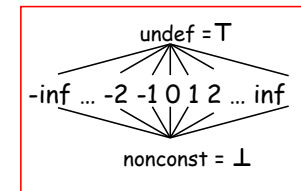
Better Constant Propagation

- What about:


```

x <- 1
if (y > z)
    x <- 1
a <- x
            
```
- Use a better lattice
- Meet:

$a \wedge \text{top}$	$\rightarrow a$
$a \wedge \text{bot}$	$\rightarrow \text{bot}$
$c \wedge c$	$\rightarrow c$
$c \wedge d$ (if $c \neq d$)	$\rightarrow \text{bot}$
- Init all vars to: ~~bot~~ or top?



Lecture 2

15-745 © Seth Copen Goldstein 2005

8

Available Expressions

- $X+Y$ is "available" at statement S if
 - $x+y$ is computed along every path from the start to S AND
 - neither x nor y is modified after the last evaluation of $x+y$

$a \leftarrow b+c$

$b \leftarrow a-d$

$c \leftarrow b+c$

$d \leftarrow a-d$

Lecture 2

15-745 © Seth Copen Goldstein 2005

9

Computing Available Expressions

- Forward or backward?
- Values?
- Lattice?
- $gen[b] =$
- $kill[b] =$
- $in[b] =$
- $out[b] =$
- initialization?

Lecture 2

15-745 © Seth Copen Goldstein 2005

10

Computing Available Expressions

- Forward
- Values: all expressions
- Lattice: available, not-avail
- $gen[b] =$ if b evals expr e and doesn't define variables used in e
- $kill[b] =$ if b assigns to x , then all exprs using x are killed.
- $out[b] = in[b] - kill[b] \cup gen[b]$
- $in[b] =$ what to do at a join point?
- initialization?

Lecture 2

15-745 © Seth Copen Goldstein 2005

11

Computing Available Expressions

- Forward
- Values: all expressions
- Lattice: available, not-avail
- $gen[b] =$ if b evals expr e and doesn't define variables used in e
- $kill[b] =$ if b assigns to x , exprs(x) are killed
 $out[b] = (in[b] - kill[b]) \cup gen[b]$
- $in[b] =$ An expr is avail only if avail on ALL edges, so: $in[b] = \cap$ over all $p \in pred(b)$, $out[p]$
- Initialization
 - All nodes except entry are set to ALL avail
 - Entry is set to NONE avail

Lecture 2

15-745 © Seth Copen Goldstein 2005

12

Constructing Gen & Kill

Stmt s	gen[s]	kill[s]
t ← x op y	{x op y}-kill[s]	{exprs containing t}
t ← M[a]	{M[a]}-kill[s]	
M[a] ← b		
f(a, ...)		{M[x] for all x}
t ← f(a,...)		

Lecture 2

15-745 © Seth Copen Goldstein 2005

13

Constructing Gen & Kill

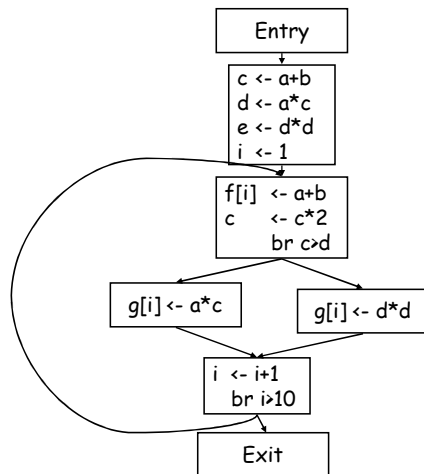
Stmt s	gen[s]	kill[s]
t ← x op y	{x op y}-kill[s]	{exprs containing t}
t ← M[a]	{M[a]}-kill[s]	{exprs containing t}
M[a] ← b	{}	{for all x, M[x]}
f(a, ...)	{}	{for all x, M[x]}
t ← f(a,...)	{}	{exprs containing t for all x, M[x]}

Lecture 2

15-745 © Seth Copen Goldstein 2005

14

Example

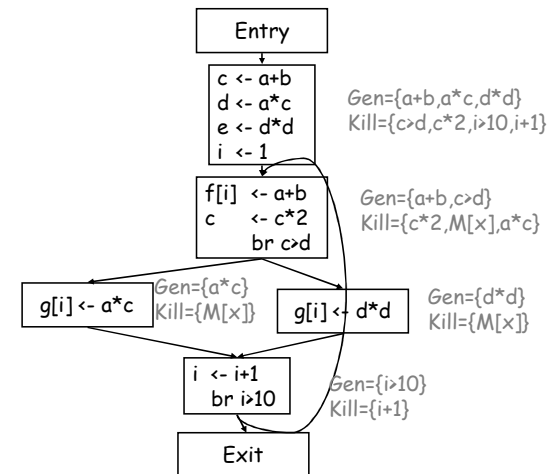


Lecture 2

15-745 © Seth Copen Goldstein 2005

15

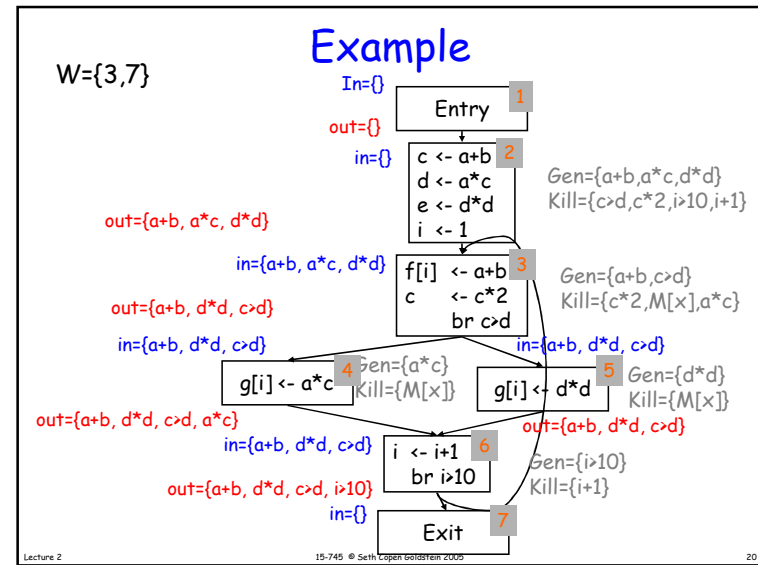
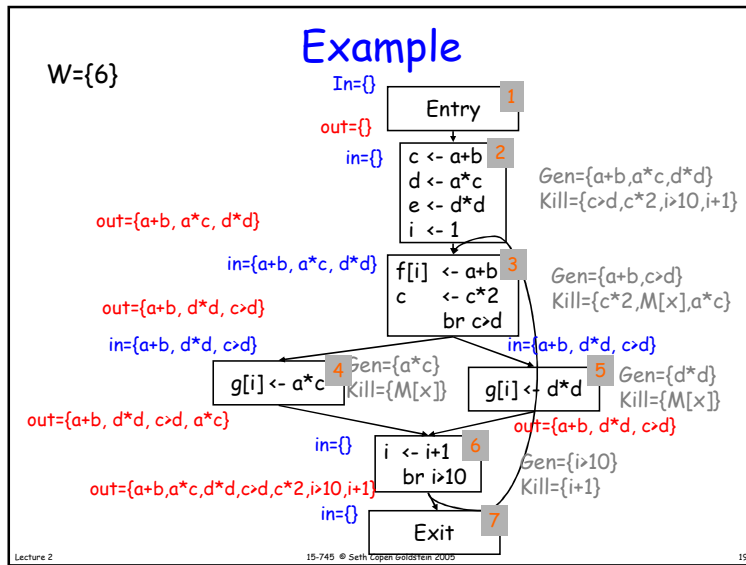
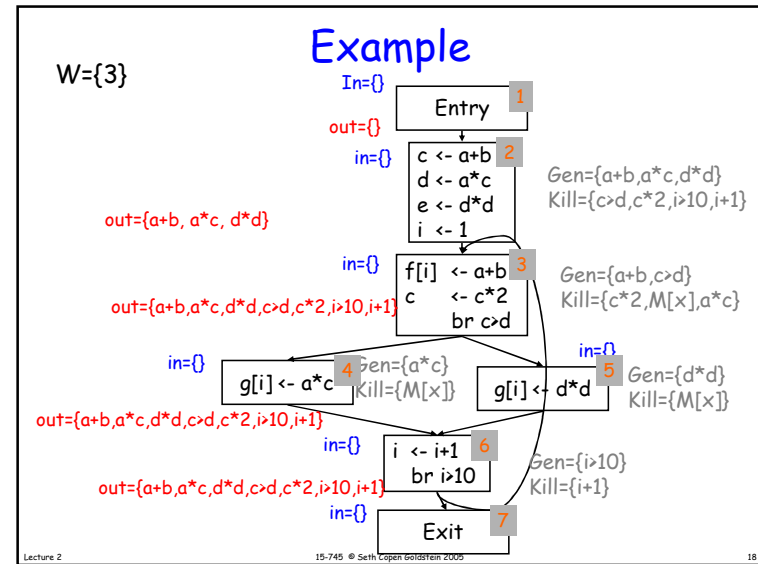
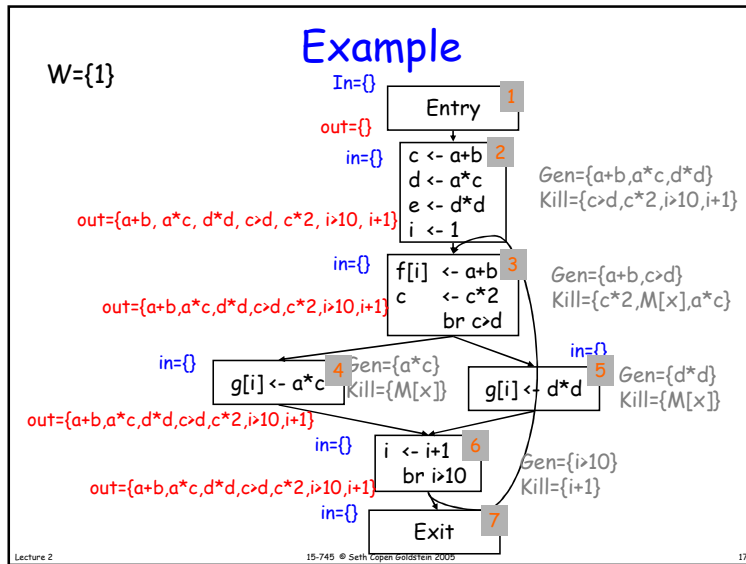
Example

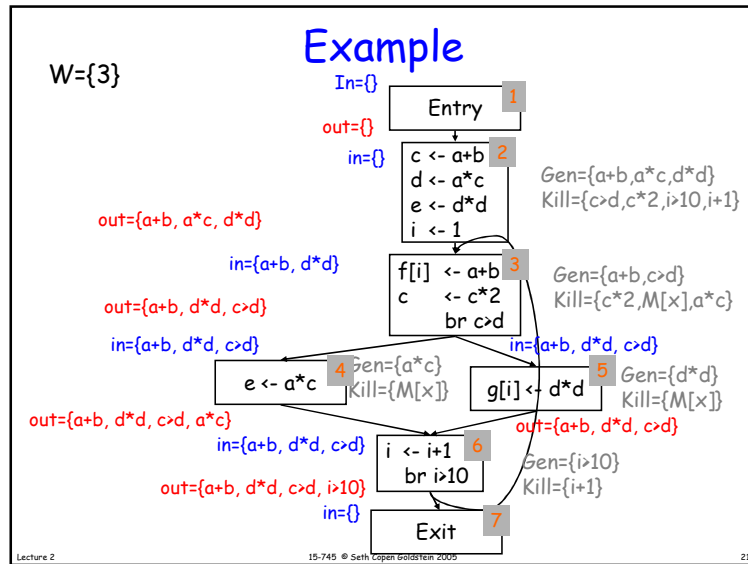


Lecture 2

15-745 © Seth Copen Goldstein 2005

16

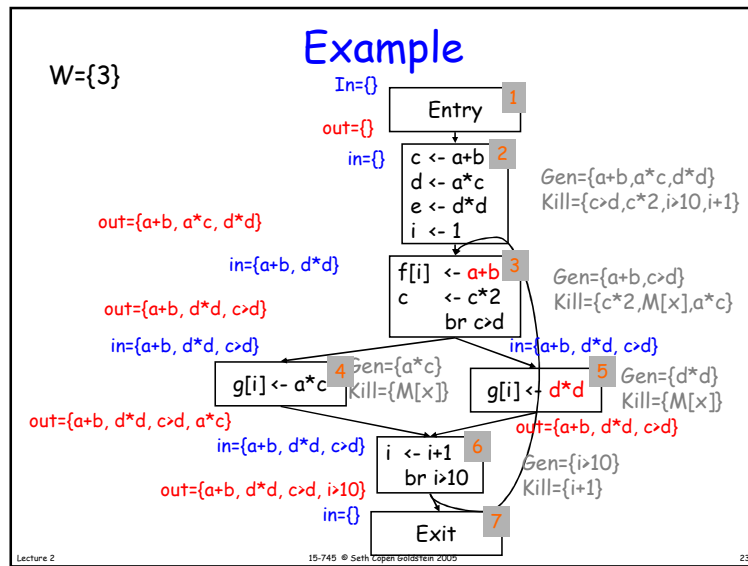




CSE

- Calculate Available expressions
- For every stmt in program
 - If expression, $x \text{ op } y$, is available {
 - Compute reaching expressions for $x \text{ op } y$ at this stmt
 - foreach stmt in RE of the form $t \leftarrow x \text{ op } y$
 - rewrite at: $t' \leftarrow x \text{ op } y$
 - $t \leftarrow t'$
- }
 - replace $x \text{ op } y$ in stmt with t'

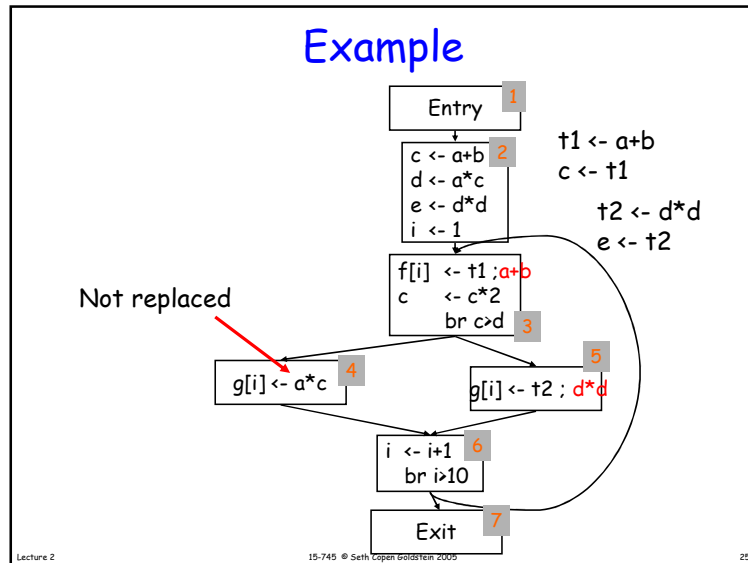
Lecture 2 15-745 © Seth Copen Goldstein 2005 22



Calculating RE

- Could be dataflow problem, but not needed enough, so ...
- To find RE for $x \text{ op } y$ at stmt S
 - traverse cfg backward from S until
 - reach $t \leftarrow x + y$ (& put into RE)
 - reach definition of x or y

Lecture 2 15-745 © Seth Copen Goldstein 2005 24



Dataflow Summary

	Union	intersection
Forward	Reaching defs	Available exprs
Backward	Live variables	

Later in course we look at bidirectional dataflow

Lecture 2 15-745 © Seth Copen Goldstein 2005 26

- ### Dataflow Framework
- Lattice
 - Universe of values
 - Meet operator
 - Basic attributes (e.g., gen, kill)
 - Traversal order
 - Transfer function
- Lecture 2 15-745 © Seth Copen Goldstein 2005 27

- ### Dataflow Framework
- Another formulation: "Meet over Paths" (MOP)
 - To find $in[B]$,
 - Enumerate all paths from entry to B to get P
 - For each path p_x in P, $p_x = \{entry \rightarrow b_i \rightarrow b_j \rightarrow \dots b_n\}$, calculate sum of transfer functions:

$$s_x = F_{b_n}(\dots F_{b_j}(F_{b_i}(out[entry])))$$
 - Then do one big Meet over all the s_x values
 - Not practical; more of theoretical interest...
- Lecture 2 15-745 © Seth Copen Goldstein 2005 28

Dataflow Framework

- Finally...why not put the values on the edges?
 - (Muchnick thinks it's a good idea...)

$in[n] = \bigwedge_{p \in \text{pred}[n]} out[p]$
 $out[n] = F_n(in[n])$

$e[n \rightarrow s] = \bigwedge_{p \in \text{pred}[n]} F_{n,s}(e[p \rightarrow n])$

Lecture 2 15-745 © Seth Copen Goldstein 2005 29

Dataflow : Edges

- Muchnick's example: smart const prop

Lecture 2 15-745 © Seth Copen Goldstein 2005 30

Dataflow : Edges

- My example (again, const prop)

Lecture 2 15-745 © Seth Copen Goldstein 2005 31