#### Advanced Languages

Presented by Hormoz Zarnani and Eric Malloy

#### The Design and Implementation of a Certifying Compiler

George C. Necula and Peter Lee

## Compiler Correctness

- •Compiler produces correct output
- • That is, generated assembly code is functionally equivalent to the high-level code

# Traditional Approaches

- • Formal compiler verification
	- –Hand proofs
	- –Mechanical proofs
- •**Testing** 
	- – Automatically generating test patterns and checking validity of corresponding outputs

### Why Traditional Approaches Do Not Work

- • Formal compiler verification
	- –Verify algorithm rather than implementation
	- – Not automatic
		- Requires human intervention and expertise
	- –Must redo proofs if compiler changed
- •**Testing** 
	- –Generated patterns are usually inadequate

### Why Traditional Approaches Do Not Work (cont.)

•Cannot handle optimizing compilers

–Place many restrictions on optimizations

## Better Approach

- •Proving full correctness is too expensive
- • Instead, employ a method that is less expensive and yet gives satisfactory confidence
	- –Check individual compilations

### Touchstone: A Certifying Compiler

• Compiles a strongly typed subset of C into optimized DEC Alpha assembly language

### Structure of Touchstone



Touchstone – Advantages

- •Easy to employ
- • Also can transform conventional compilers to certifying ones
- •Can be applied to any type safe language
- • Places no restriction on optimizations allowed
- • Only VCGen and Proof Checker have to be correct

## Touchstone – Disadvantages

- •Applies to only type-safe languages
- •But C is not type-safe

## Touchstone – Conclusion

- •Does not fully address problem
- Kut is noval work and avealla But is novel work and excellent starting point

### Checking System Rules Using System-Specific Programmer-Written Compiler **Extensions**

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# System Rules Violations

- • What are the some of the ways we might find violations of system rules in a program?
	- –Model Checkers
	- –Theorem Provers
	- –**Testing**
	- –Code Inspections
	- –**Compilers**

## Metal-Level Compilation

- • System specific "meta" semantics
	- –Essentially these are rules for a systems API's
	- – Implemented as runtime extensions to the compiler
	- – Capable of discovering errors in complex code as well as optimization opportunities

### Extensions to the xg++ Compiler

- • xg++ is and extensible compiler based on  $g_{++}$
- • Extensions are written in a high level state machine language called "metal"

## How Metal Extensions Work

- • Metal is compiled using the mcc compiler and dynamically linked into xg++ at compile time
- • Pattern comparisons are done based on xg++'s internal representation

### Metal Example

```
{ #include "linux-includes.h" }
sm check_interrupts {
  // Variables
  // used in patterns
  decl { unsigned } flags;
  // Patterns
  // to specify enable/disable functions.
  pat enable = { sti(); }
             \vert { restore_flags(flags); };
 pat disable = {cli(); };
  // States
  // The first state is the initial state.
  is_enabled: disable ==> is_disabled
     | enable ==> { err("double enable"); }
  is_disabled: enable ==> is_enabled
     | disable ==> { err("double disable"); }
     // Special pattern that matches when the SM
     // hits the end of any path in this state.
     \frac{1}{2} $end_of_path$ ==>
        { err("exiting w/intr disabled!"); }
     ÷
}
```
## What Can Be Checked

- •Assertion side-effects
- • Checking assertions of constant scalar variables
- •Temporal orderings of system calls
- •Memory management
- • Global checking of blocking routines and reference counts

### Memory Management Error Counts



### Blocking with an Interrupt Disabled



## Advantages\Disadvantages

• What advantages and disadvantages does this approach have?