Lecture 10

Lazy Code Motion

I. Forms of redundancy (quick review)
  - global common subexpression elimination
  - loop invariant code motion
  - partial redundancy

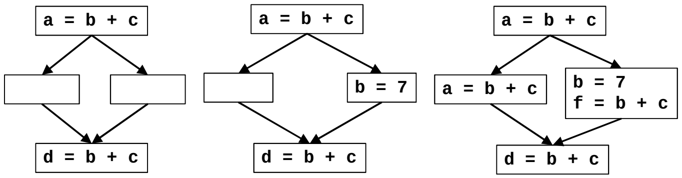
II. Lazy Code Motion Algorithm
  - Mathematical concept: a cut set
  - Basic technique (anticipation)
  - 3 more passes to refine algorithm

Reading: Chapter 9.5

Overview

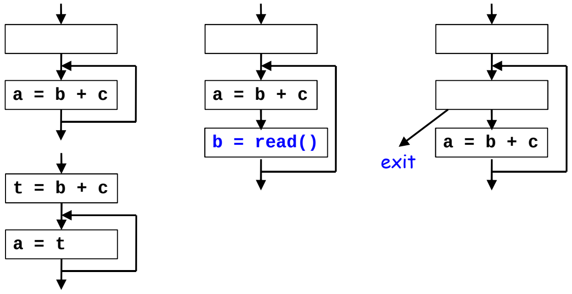
- **Eliminates many forms of redundancy in one fell swoop**
- **Originally formulated as 1 bi-directional analysis**
- **Lazy code motion algorithm**
  - formulated as 4 separate uni-directional passes
    - backward, forward, forward, backward

I. Common Subexpression Elimination



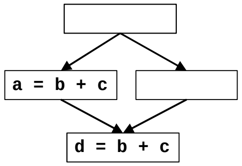
- ***A common expression may have different values on different paths!***
- ***On every path reaching p,***
  - expression b+c has been computed
  - b, c not overwritten after the expression

Loop Invariant Code Motion



- *Given an expression (b+c) inside a loop,*
  - *does the value of b+c change inside the loop?*
  - *is the code executed at least once?*

Partial Redundancy

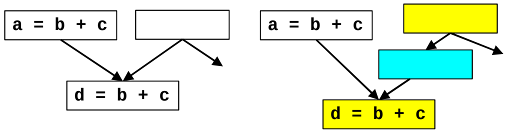


- Can we place calculations of b+c
  such that no path re-executes the same expression
- Partial Redundancy Elimination (PRE)
  - subsumes:
    - global common subexpression (full redundancy)
    - loop invariant code motion (partial redundancy for loops)

II. Lazy Code Motion

- **Key observation:**
  - A bi-directional (!) data flow problem can be replaced with several unidirectional data flow problems → much easier
  - Better result as well!

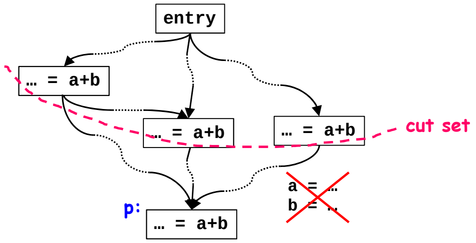
Preparing the Flow Graph



- **Definition:** Critical edges
  - source basic block has multiple successors
  - destination basic block has multiple predecessors
- **Modify the flow graph: (treat every statement as a basic block)**
  - To keep algorithm simple: restrict placement of instructions to the beginning of a basic block
  - Add a basic block for every edge that leads to a basic block with multiple predecessors (not just on critical edges)

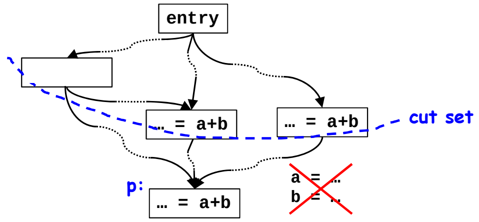
Full Redundancy: A Cut Set in a Graph

Key mathematical concept



- **Full redundancy at p: expression a+b redundant on all paths**
  - a cut set: nodes that separate entry from p
  - a cut set contains calculation of a+b
  - a, b, not redefined

Partial Redundancy: Completing a Cut Set



- **Partial redundancy at p: redundant on some but not all paths**
  - Add operations to create a cut set containing a+b
  - Note: Moving operations up can eliminate redundancy
- **Constraint on placement: no wasted operation**
  - a+b is "anticipated" at B if its value computed at B will be used along ALL subsequent paths
  - a, b not redefined, no branches that lead to exit without use
- **Range where a+b is anticipated → Choice**

Pass 1: Anticipated Expressions

This pass does most of the heavy lifting in eliminating redundancy

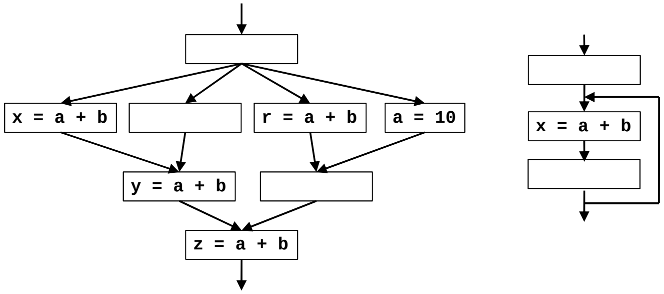
- **Backward pass: Anticipated expressions**
  Anticipated[b].in: Set of expressions anticipated at the entry of b
  - An expression is anticipated if its value computed at point p will be used along ALL subsequent paths

	Anticipated Expressions
Domain	Sets of expressions
Direction	backward
Transfer Function	$f_b(x) = EUse_b \cup (x - EKill_b)$ EUse: used exp, EKill: exp killed
$\wedge$	$\cap$
Boundary	in[exit] = ∅
Initialization	in[b] = {all expressions}

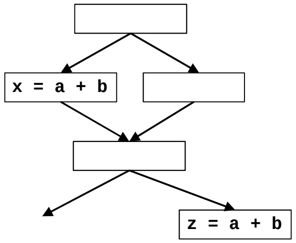
- **First approximation:**
  - place operations at the frontier of anticipation
    (boundary between not anticipated and anticipated)

Examples (1)

See the algorithm in action

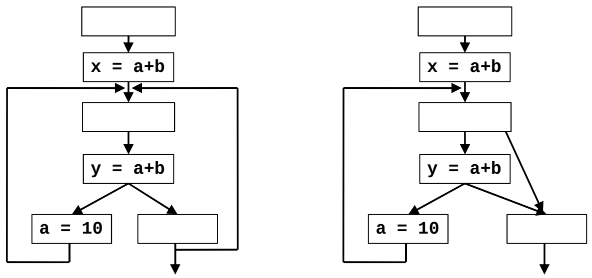


Examples (2)



- Cannot eliminate all redundancy

Examples (3)

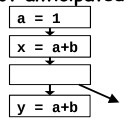


- Do you know how the algorithm works without simulating it?

Pass 2: Place As Early As Possible

There is still some redundancy left!

- **First approximation: frontier between "not anticipated" & "anticipated"**
- **Complication: anticipation may oscillate**



- Pretend we calculate expression e whenever it is anticipated
- e will be **available at p** if e has been "anticipated but not subsequently killed" on all paths reaching p

	Available Expressions
Domain	Sets of expressions
Direction	forward
Transfer Function	$f_b(x) = (\text{Anticipated}[b].\text{in} \cup x) - \text{EKill}_b$
$\wedge$	$\cap$
Boundary	out[entry] = ∅
Initialization	out[b] = {all expressions}

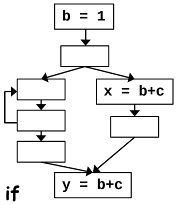
Early Placement

- **earliest(b)**
  - set of expressions added to block b under early placement
- **Place expression at the earliest point anticipated and not already available**
  - earliest(b) = anticipated[b].in - available[b].in
- **Algorithm**
  - For all basic block b, if x+y ∈ earliest[b]
    - at beginning of b:
      create a new variable t
      t = x+y,
      replace every original x+y by t

Pass 3: Lazy Code Motion

Let's be lazy without introducing redundancy.

- **Delay creating redundancy to reduce register pressure**

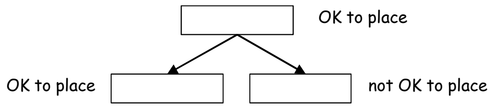


- **An expression e is postponable at a program point p if**
  - all paths leading to p have seen the earliest placement of e but not a subsequent use

	Postponable Expressions
Domain	Sets of expressions
Direction	forward
Transfer Function	$f_b(x) = (\text{earliest}[b] \cup x) - \text{EUse}_b$
$\wedge$	$\cap$
Boundary	out[entry] = ∅
Initialization	out[b] = {all expressions}

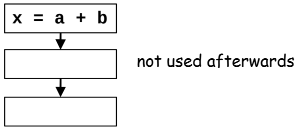
Latest: frontier at the end of "postponable" cut set

- latest[b] = (earliest[b] ∪ postponable.in[b]) ∩
  (EUse$_b$ ∪ ¬($\bigcap_{s \in succ[b]}$(earliest[s] ∪ postponable.in[s])))
    - OK to place expression: earliest or postponable
    - Need to place at b if either
        - used in b, or
        - not OK to place in one of its successors
- Works because of pre-processing step (an empty block was introduced to an edge if the destination has multiple predecessors)
    - if b has a successor that cannot accept postponement, b has only one successor
    - The following does not exist:



Pass 4: Cleaning Up

Finally… this is easy, it is like liveness



- **Eliminate temporary variable assignments unused beyond current block**
- **Compute: Used.out[b]: sets of used (live) expressions at exit of b.**

	Used Expressions
Domain	Sets of expressions
Direction	backward
Transfer Function	$f_b(x)$ = (EUse[b] ∪ x) - latest[b]
∧	∪
Boundary	in[exit] = ∅
Initialization	in[b] = ∅

Code Transformation

- For all basic blocks b,

  if (x+y) ∈ (latest[b] ∩ used.out[b])

  at beginning of b:

  add new t = x+y

  replace every original x+y by t

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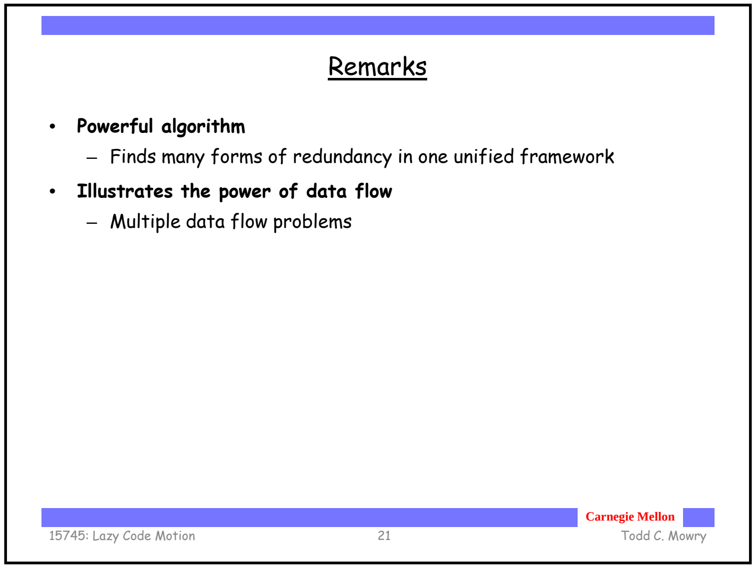
4 Passes for Partial Redundancy Elimination

- ***Heavy lifting*: Cannot introduce operations not executed originally**
    - Pass 1 (backward): Anticipation: range of code motion
    - Placing operations at the frontier of anticipation gets most of the redundancy
- ***Squeezing the last drop of redundancy*:**
  An anticipation frontier may cover a subsequent frontier
    - Pass 2 (forward): Availability
    - Earliest: anticipated, but not yet available
- ***Push the cut set out -- as late as possible***
  To minimize register lifetimes
    - Pass 3 (forward): Postponability: move it down provided it does not create redundancy
    - Latest: where it is used or the frontier of postponability
- ***Cleaning up***
    - Pass 4: **Remove temporary assignment**



Remarks

- **Powerful algorithm**
  - Finds many forms of redundancy in one unified framework
- **Illustrates the power of data flow**
  - Multiple data flow problems