

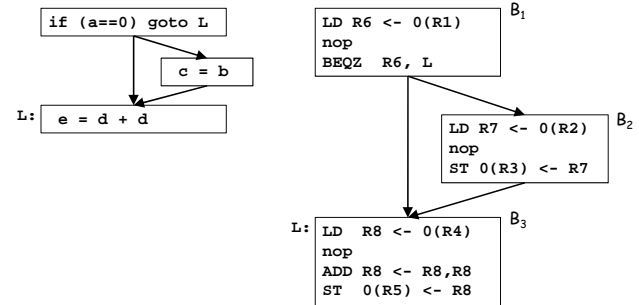
Lecture 18

Global Scheduling

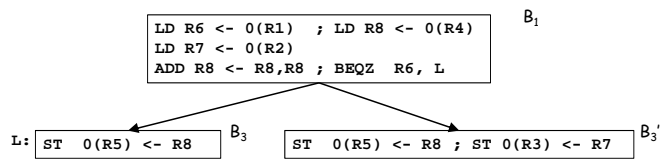
Reading: Chapter 10.4

Introduction to Global Scheduling

Assume each clock can execute 2 operations of any kind.



Result of Code Scheduling



Terminology

Control equivalence:

- Two operations o_1 and o_2 are *control equivalent* if o_1 is executed if and only if o_2 is executed.

Control dependence:

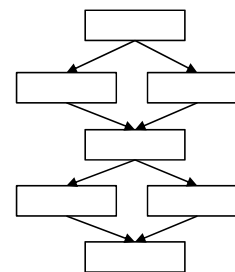
- An op o_2 is *control dependent* on op o_1 if the execution of o_2 depends on the outcome of o_1 .

Speculation:

- An operation o is *speculatively* executed if it is executed before all the operations it depends on (control-wise) have been executed.

Requirements:

- does not raise an exception
- satisfies data dependences



Code Motions

Goal: Shorten execution time **probabilistically**

Moving instructions up:

- Move instruction to a cut set (from entry)
- **Speculation**: even when not anticipated.

Moving instructions down:

- Move instruction to a cut set (from exit)
- May execute extra instruction
- Can duplicate code

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A Note on Data Dependences

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General-Purpose Applications

- Lots of data dependences
- Key performance factor: **memory latencies**
- **Move memory fetches up**
 - Speculative memory fetches can be expensive
- **Control-intensive: get execution profile**
 - **Static estimation**
 - Innermost loops are frequently executed
 - back edges are likely to be taken
 - Edges that branch to exit and exception routines are not likely to be taken
 - **Dynamic profiling**
 - Instrument code and **measure** using representative data

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A Basic Global Scheduling Algorithm

- Schedule innermost loops first
- Only upward code motion
- No creation of copies
- Only one level of speculation

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Program Representation

- A **region** in a control flow graph is:
 - a set of **basic blocks** and all the **edges** connecting these blocks,
 - such that control from outside the region **must enter through a single entry block**.
- A **function** is represented as a **hierarchy of regions**
 - The whole control flow graph is a region
 - Each natural loop in the flow graph is a region
 - Natural loops are hierarchically nested
- **Schedule regions from inner to outer**
 - treat inner loop as a black box unit
 - can **schedule around it but not into it**
 - **ignore all the loop back edges** → get an acyclic graph

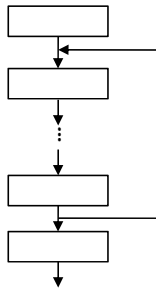
Algorithm

```
Compute data dependencies;  
For each region from inner to outer {  
  For each basic block B in prioritized topological order {  
    CandBlocks = ControlEquiv{B} ∪  
      Dominated-Successors{ControlEquiv{B}};  
    CandInsts = ready operations in CandBlocks;  
    For (t = 0, 1, ... until all operations from B are scheduled) {  
      For (n in CandInst in priority order) {  
        if (n has no resource conflicts at time t) {  
          S(n) = < B, t >  
          Update resource commitments  
          Update data dependencies  
        }  
      }  
      Update CandInsts;  
    }  
  }  
}
```

Priority functions: non-speculative before speculative

Extensions

- Prepass before scheduling: **loop unrolling**
- Especially important to move operation up loop back edges



Summary

- **Global scheduling**
 - Legal code motions
 - Heuristics