

Lecture 8

Induction Variables and Strength Reduction

- I. Overview of optimization
- II. Algorithm to find induction variables

Example

```
FOR i = 0 to 100
  A[i] = 0;

  i = 0
L2: IF i>=100 GOTO L1
  t1 = 4 * i
  t2 = &A + t1
  *t2 = 0
  i = i+1
  GOTO L2
L1:
```

Definitions

- A **basic induction variable** is
 - a variable X whose only definitions within the loop are assignments of the form:
 $X = X+c$ or $X = X-c$,
where c is either a **constant** or a **loop-invariant variable**.
- An **induction variable** is
 - a basic induction variable, or
 - a variable defined once within the loop, whose value is a linear function of some basic induction variable at the time of the definition:
 $A = c_1 * B + c_2$
- The **FAMILY of a basic induction variable B** is
 - the set of induction variables A such that each time A is assigned in the loop, the value of A is a linear function of B .

Optimizations

- 1. **Strength reduction:**
 - Let A be an induction variable in family of basic induction variable B ($A = c_1 * B + c_2$)
 - Create new variable: A'
 - Initialization in preheader: $A' = c_1 * B + c_2$;
 - Track value of B :
add after $B=B+x$: $A' = A' + x * c_1$;
 - Replace assignment to A : $A = A'$

Optimizations (continued)

2. Optimizing non-basic induction variables

- copy propagation
- dead code elimination

3. Optimizing basic induction variables

- Eliminate basic induction variables used only for
 - calculating other induction variables and loop tests
- Algorithm:
 - Select an induction variable A in the family of B , preferably with simple constants ($A = c_1 * B + c_2$).
 - Replace a comparison such as
`if B > X goto L1`
with
`if (A' > c1 * X + c2) goto L1` (assuming c_1 is positive)
 - if B is live at any exit from the loop, recompute it from A'
 - After the exit, $B = (A' - c_2) / c_1$

II. Basic Induction Variables

• A BASIC induction variable in a loop L

- a variable X whose only definitions within L are assignments of the form $X = X+c$ or $X = X-c$, where c is either a constant or a loop-invariant variable.

• Algorithm: can be detected by scanning L

• Example:

```
k = 0;
for (i = 0; i < n; i++) {
    k = k + 3;
    ... = m;
    if (x < y)
        k = k + 4;
    if (a < b)
        m = 2 * k;
    k = k - 2;
    ... = m;
```

Each iteration may execute a different number of increments/decrements!!

Strength Reduction Algorithm

• Key idea:

- For each induction variable A , ($A = c_1*B+c_2$ at time of definition)
 - variable A' holds expression c_1*B+c_2 at all times
 - replace definition of A with $A=A'$ only when executed

• Result:

- Program is correct
- Definition of A does not need to refer to B

Finding Induction Variable Families

• Let B be a basic induction variable

- Find all induction variables A in family of B :
 - $A = c_1 * B + c_2$
(where B refers to the value of B at time of definition)

• Conditions:

- If A has a single assignment in the loop L , and assignment is one of:

```
A = B * c
A = c * B
A = B / c (assuming A is real)
A = B + c
A = c + B
A = B - c
A = c - B
```

- OR, ... (next page)

Finding Induction Variable Families (continued)

– Let D be an induction variable in the family of B ($D = c_1 * B + c_2$)

- If A has a single assignment in the loop L , and assignment is one of:

$$\begin{aligned} A &= D * c \\ A &= c * D \\ A &= D / c \quad (\text{assuming } A \text{ is real}) \\ A &= D + c \\ A &= c + D \\ A &= D - c \\ A &= c - D \end{aligned}$$

- No definition of D outside L reaches the assignment to A
- Between the lone point of assignment to D in L and the assignment to A , there are no definitions of B

Summary

- **Precise definitions of induction variables**
- **Systematic identification of induction variables**
- **Strength reduction**
- **Clean up:**
 - eliminating basic induction variables
 - used in other induction variable calculations
 - replacement of loop tests
 - eliminating other induction variables
 - standard optimizations