

## 1 Distributed Problems Beyond Low Rank Approximation

In the previous part of the lecture, we have covered the distributed low rank approximation problem. The problem is optimally solved by the BWZ protocol. However, in general, distributed problems with a similar flavor are not as well understood. We will list some of the open problems below.

**Rank computation.** There are  $s$  servers, each with a real matrix  $A_i \in \mathbb{R}^{n \times n}$ , and the goal is to compute the rank of the sum of the matrices  $\sum_{i=1}^s A_i$ , minimizing the communication cost.

**Linear programming.** There are  $s$  servers, each with a real matrix  $A_i \in \mathbb{R}^{n \times d}$  and a vector  $b_i \in \mathbb{R}^n$ .  $A_i$  and  $b_i$  stand for a set of linear constraints. The goal is to solve the linear program  $\min\{c^T x : A_i x \leq b_i, i = 1, \dots, s\}$ , minimizing the communication cost.

**Maximum matching.** There are  $s$  servers, each with a set of edges  $E_i$  in a graph  $G = (V, E)$ . The goal is to compute the maximum matching when  $E = \bigcup_{i=1}^s E_i$ , minimizing the communication cost.

## 2 Robust Regression

In the following section, we will introduce the robust regression problem. That is, given a matrix  $A$  and a vector  $b$ , we want to find a vector  $x$  that minimizes the objective function

$$\|Ax - b\|_1 = \sum_{i=1}^n |\langle A_i, x \rangle - b_i|, \text{ where } A_i \text{ is the } i\text{-th row of } A.$$

### 2.1 The Linear Programming Solution

We do not have a good closed-form solution for the robust regression problem. However, we can still solve the problem exactly by formulating it as a linear program. The linear program is as follows:

$$\begin{array}{l|l} \text{Minimize} & (1, 1, 1, \dots, 1) \cdot (\alpha^+ + \alpha^-) \\ \text{Subject to} & Ax + \alpha^+ - \alpha^- = b \\ & \alpha^+, \alpha^- \in \mathbb{R}_{\geq 0}^d \\ & x \in \mathbb{R}^d \end{array}$$

One key observation is for each  $i$ , at least one of  $\alpha_i^+$  and  $\alpha_i^-$  is zero. This is because if  $\alpha_i^+ > 0$  and  $\alpha_i^- > 0$ , we can decrease both  $\alpha_i^+$  and  $\alpha_i^-$  by the same amount and the objective function will decrease. Therefore, the optimal solution will have the objective function equal to  $\|Ax - b\|_1$ .

Thus, we can see that the problem can be solved by in  $\text{poly}(nd)$  time using linear programming.

## 2.2 Well-Conditioned Bases

The poly( $nd$ ) time complexity of the linear programming solution is not satisfactory in the context of big data. To develop a faster algorithm, we first introduce the concept of well-conditioned bases.

Recall that in the  $\ell_2$  regression problem, we can use the SVD of  $A$  to decompose  $A = UW$ , where  $U$  has orthonormal columns, so that  $\|Ux\|_2 = \|x\|_2$  for all  $x \in \mathbb{R}^d$ . For the  $\ell_1$  regression problem, we would like to have a similar property. That is, we would like to have a matrix  $U$  such that

$$A = UW \text{ and } \|Ux\|_1 \approx \|x\|_1 \text{ for all } x \in \mathbb{R}^d.$$

To do this, we first write  $A = QW$ , where  $Q$  is a matrix with full column rank. Note that the number of columns  $k$  of  $Q$  is at most  $d$ . Below, we define the  $(Q, 1)$ -norm of a vector  $x$ .

**Definition 2.1.** Let  $Q \in \mathbb{R}^{n \times k}$  be a matrix with full column rank. For any  $x \in \mathbb{R}^k$ , we define

$$\|x\|_{Q,1} = \|Qx\|_1.$$

**Lemma 2.2.**  $\|\cdot\|_{Q,1}$  is a norm on  $\mathbb{R}^k$ .

**Proof.** To show that  $\|\cdot\|_{Q,1}$  is a norm, we need to show that it satisfies the three properties.

For any  $x, y \in \mathbb{R}^k$ , we have

$$\|x + y\|_{Q,1} = \|Q(x + y)\|_1 = \|Qx + Qy\|_1 \leq \|Qx\|_1 + \|Qy\|_1 = \|x\|_{Q,1} + \|y\|_{Q,1}.$$

For any  $x \in \mathbb{R}^k$  and  $c \in \mathbb{R}$ , we have

$$\|cx\|_{Q,1} = \|Q(cx)\|_1 = \|cQx\|_1 = |c| \|Qx\|_1 = |c| \|x\|_{Q,1}.$$

Finally, for any  $x \in \mathbb{R}^k$ , we have  $\|x\|_{Q,1} = \|Qx\|_1 = 0$  if and only if  $Qx = 0$  if and only if  $x = 0$ .

Therefore,  $\|\cdot\|_{Q,1}$  is a norm on  $\mathbb{R}^k$ . ■

For the ease of notation, we assume  $k = d$  in the following discussion.

Let  $C = \{z \in \mathbb{R}^d \mid \|z\|_{Q,1} \leq 1\}$  be the unit ball of the  $(Q, 1)$ -norm. By the triangle inequality and absolute homogeneity of the  $(Q, 1)$ -norm, we know that  $C$  is a convex set. Moreover,  $C$  is symmetric about the origin. For this symmetric convex body, we can apply the Lowner-John ellipsoid theorem.

**Theorem 2.3.** Let  $C$  be a symmetric convex body in  $\mathbb{R}^d$ . Then there exists an ellipsoid  $E$  such that

$$E \subseteq C \subseteq \sqrt{d}E.$$

Algebraically, there exists some  $G \in \mathbb{R}^{d \times d}$  such that  $E = \{z \in \mathbb{R}^d \mid z^\top G^\top G z \leq 1\}$ .

**Corollary 2.4.** There exists a matrix  $G \in \mathbb{R}^{d \times d}$  such that for all  $z \in \mathbb{R}^d$ , we have

$$\sqrt{z^\top G^\top G z} \leq \|z\|_{Q,1} \leq \sqrt{d} \cdot \sqrt{z^\top G^\top G z}.$$

Define  $U = QG^{-1}$ . We will show that  $U$  is a well-conditioned basis such that  $\|Ux\|_1 \approx \|x\|_1$  for all  $x \in \mathbb{R}^d$ . Let  $z = G^{-1}x$ . Then, we can compute the followings:

$$\begin{aligned}\|Ux\|_1 &= \|QG^{-1}x\|_1 = \|Qz\|_1 = \|z\|_{Q,1} \cdot \\ z^\top G^\top Gz &= x^\top (G^{-1})^\top G^\top G(G^{-1})x = x^\top x = \|x\|_2^2.\end{aligned}$$

Combining the above two equations with Corollary 2.4, we have

$$\|x\|_2 \leq \|Ux\|_1 \leq \sqrt{d} \cdot \|x\|_2.$$

To conclude the discussion, we will use the following relationship between the  $\ell_1$  and  $\ell_2$  norms.

**Fact 2.5.** *Let  $x \in \mathbb{R}^d$ . Then,  $\|x\|_2 \leq \|x\|_1 \leq \sqrt{d} \cdot \|x\|_2$ .*

Therefore, we can see that

$$\frac{\|x\|_1}{\sqrt{d}} \leq \|x\|_2 \leq \|Ux\|_1 \leq \sqrt{d} \|x\|_2 \leq \sqrt{d} \|x\|_1.$$

Here, we have shown that  $U$  preserves the  $\ell_1$  norm up to a factor of  $\sqrt{d}$ .

### 2.3 Net Argument

Previously, in the  $\ell_2$  regression problem, given a sketching matrix  $S$ , we used a net argument to show that the sketching matrix  $S$  preserves the  $\ell_2$  norm for a finite set of vectors, and in general, each vector is close to, and thus can be approximated by some vector in the net. We would like to obtain a similar result for the  $\ell_1$  regression problem in this section.

**Net for the Ball.** Let  $B = \{x \in \mathbb{R}^d \mid \|x\|_1 = 1\}$  be the unit ball of the  $\ell_1$  norm. We would like to find a  $\gamma$ -net  $N$  for  $B$  with respect to the  $\ell_1$  norm. That is, for any  $x \in B$ , there exists  $y \in N$  such that  $\|x - y\|_1 \leq \gamma$ . To do this, we could construct the net  $N$  greedily:

*While there is a point  $x \in B$  of distance larger than  $\gamma$  from all points in  $N$ , add  $x$  to  $N$ .*

Consider each point  $x \in N$ , if we construct a  $\ell_1$  ball of radius  $\gamma/2$  around  $x$ , then the balls are all disjoint by the triangle inequality. Moreover, the balls are all contained in the  $\ell_1$  ball of radius  $1 + \gamma/2$  around the origin. Therefore, we can apply a volume argument to bound the size of the net  $N$ . Note that the ratio of volume of  $d$ -dimensional  $\ell_1$  ball of radius  $1 + \gamma/2$  to that of  $\ell_1$  ball of radius  $\gamma/2$  is  $(1 + \gamma/2)^d / (\gamma/2)^d$ . Therefore, we can conclude that

$$|N| \leq \left( \frac{1 + \gamma/2}{\gamma/2} \right)^d.$$

**Net for the Subspace.** After constructing the net for the ball, we would like to lift the net to a net for the subspace. Let  $U$  be a well-conditioned basis such that

$$\|x\|_1 \leq \|Ux\|_1 \leq d \|x\|_1.$$

Unlike the  $\ell_2$  norm, the  $\ell_1$  norm is not preserved by the linear transformation  $U$ . Therefore, we need a denser net for the ball to ensure the performance of the net for the subspace. Let  $N$  be a  $(\gamma/d)$ -net for the ball  $B$  with respect to the  $\ell_1$  norm, and let  $M = \{Ux \mid x \in N\}$ . From the previous paragraph, we already know that  $|M| \leq (1 + \gamma/(2d))^d / (\gamma/(2d))^d = \left(\frac{d}{\gamma}\right)^{O(d)}$ .

We conclude with the following lemma.

**Lemma 2.6.** *For every  $x \in B$ , there exists  $y \in M$  such that  $\|Ux - y\|_1 \leq \gamma$ .*

**Proof.** Let  $x' \in N$  be the closest point to  $x$ , then  $\|x - x'\|_1 \leq \gamma/d$ . Therefore,

$$\|Ux - Ux'\|_1 \leq d \|x - x'\|_1 \leq \gamma.$$

Setting  $y = Ux'$ , we have  $\|Ux - y\|_1 \leq \gamma$ . ■