

15-150

Fall 2024

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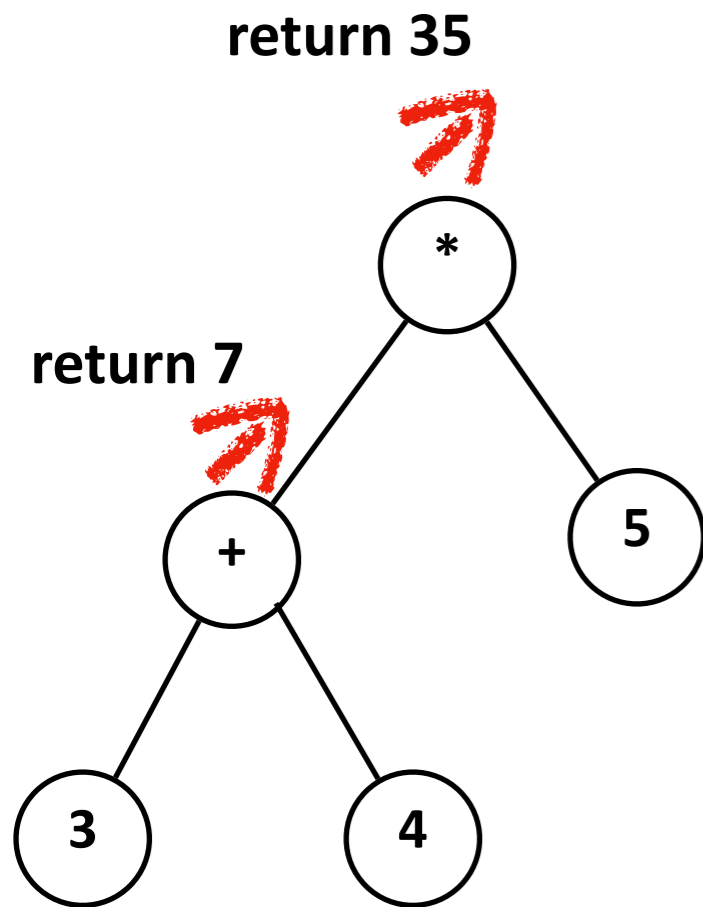
LECTURE 12

Programming with Continuations

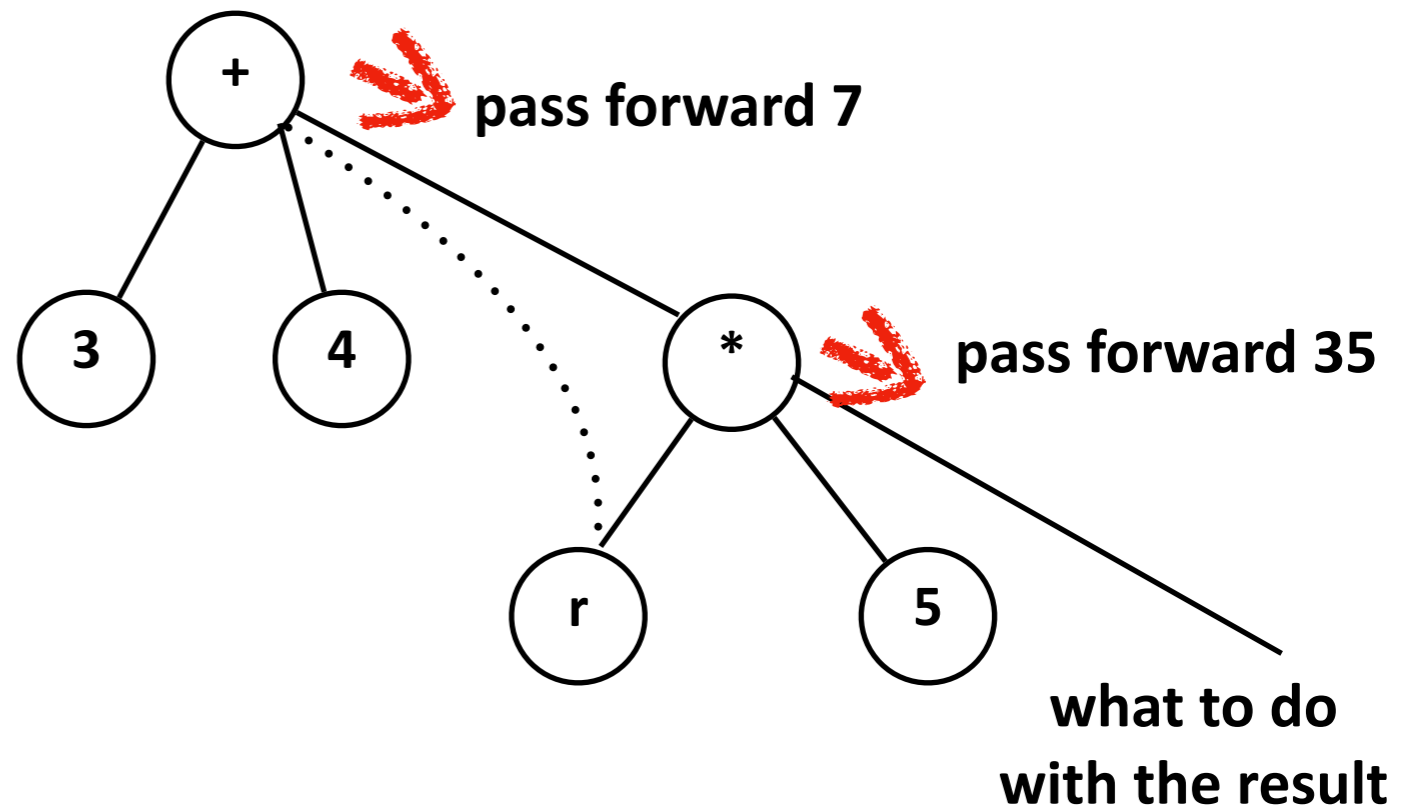
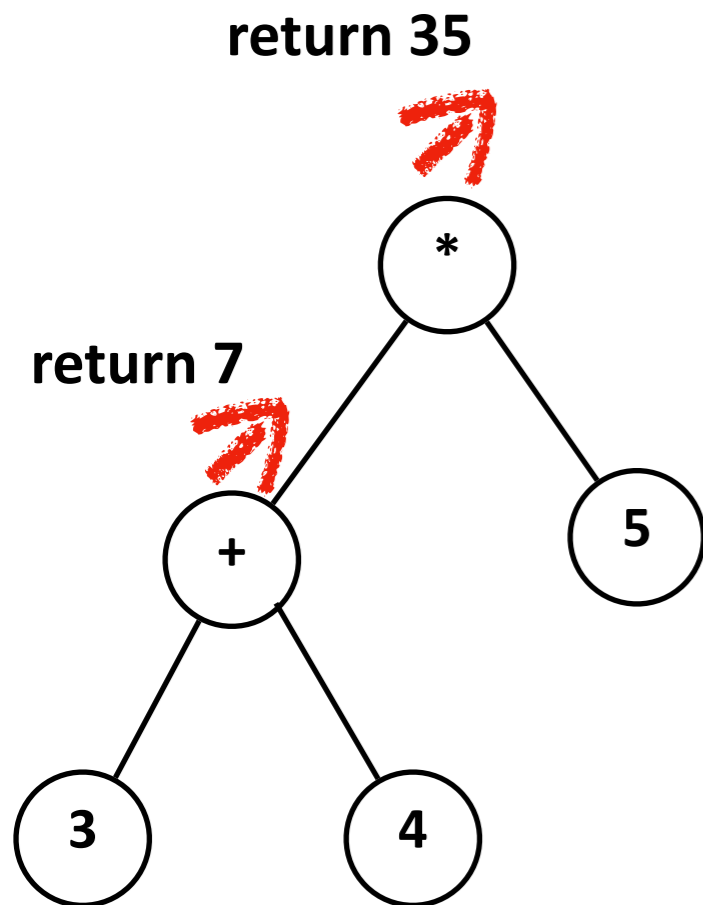
Intuition

- A continuation is a functional argument that controls flow of expression evaluation. It is often used to abstract away the “rest of the computation”.
- They act like “functional accumulators” in tail-recursive functions.
- Useful in backtracking search

Evaluating $(3+4)*5$



Evaluating $(3+4)*5$



Simple cps functions

```
fun add (x,y,k) = k (x + y)
```



continuation

```
fun mult (x,y,k) = k (x * y)
```

We could write them in curried form as well.

Alternatively

```
fun add (x,y) k = k (x + y)
```

```
fun mult (x,y) k = k (x * y)
```

Type of add

```
fun add (x, y, k) = k (x + y)
```

- `add : int * int * (int -> 'a) -> 'a`

the overall return type is the return
type of the continuation

Using a cps function

```
fun add (x,y,k) = k (x + y)
```



continuation

```
add (3, 4, fn r => r)
```

```
=> [3/x,4/y,(fn r => r)/k] k (x + y)
```

```
=> (fn r => r) (3 + 4)
```

```
=> (fn r => r) (7)
```

```
=> [...7/r] r
```

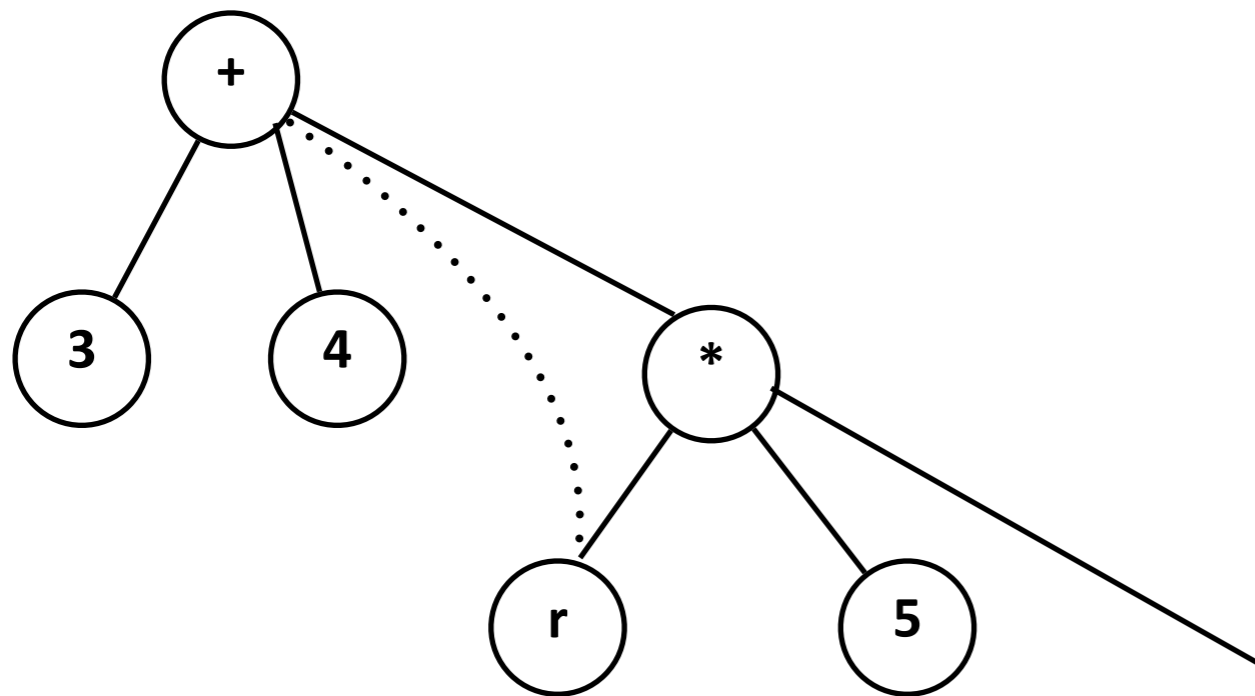
```
=> 7
```



```
fun add (x,y,k) = k (x + y)
```

```
fun mult (x,y,k) = k (x * y)
```

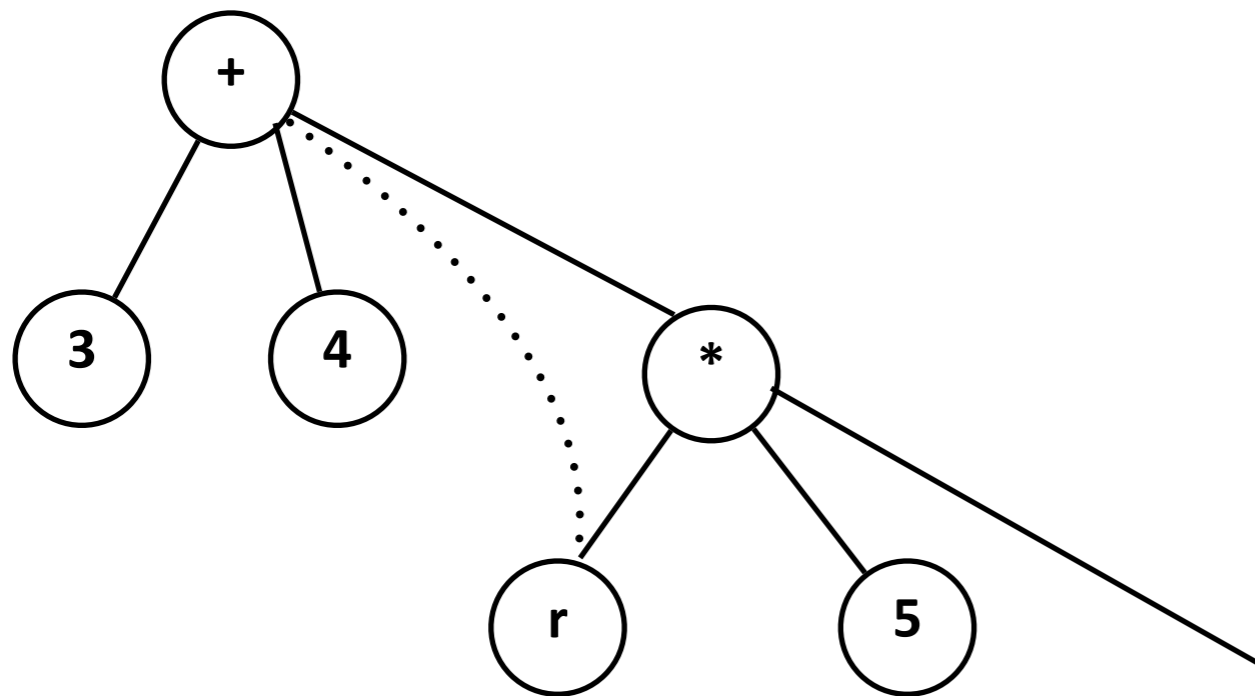
cps function to compute $(3+4)*5$



```
fun add (x,y,k) = k (x + y)
```

```
fun mult (x,y,k) = k (x * y)
```

cps function to compute $(3+4)*5$



```
add (3, 4, fn r => mult (r,5, _____ ))
```

```
fun add (x,y,k) = k (x + y)
```

```
fun mult (x,y,k) = k (x * y)
```

cps function to compute $(3+4)*5$

```
add (3, 4, fn r => mult (r,5, fn x=>x))
```

What if we wanted to return the result as a string?

```
fun add (x,y,k) = k (x + y)
```

```
fun mult (x,y,k) = k (x * y)
```

cps function to compute $(3+4)*5$

```
add (3, 4, fn r => mult(r,5,Int.toString))
```

```
==> (fn r => mult (r, 5, Int.toString))(3 + 4)
```

```
==> [7/r] mult (r, 5, Int.toString)
```

```
==> [7/r] Int.toString(r*5)
```

```
==> Int.toString(35) ==> "35"
```

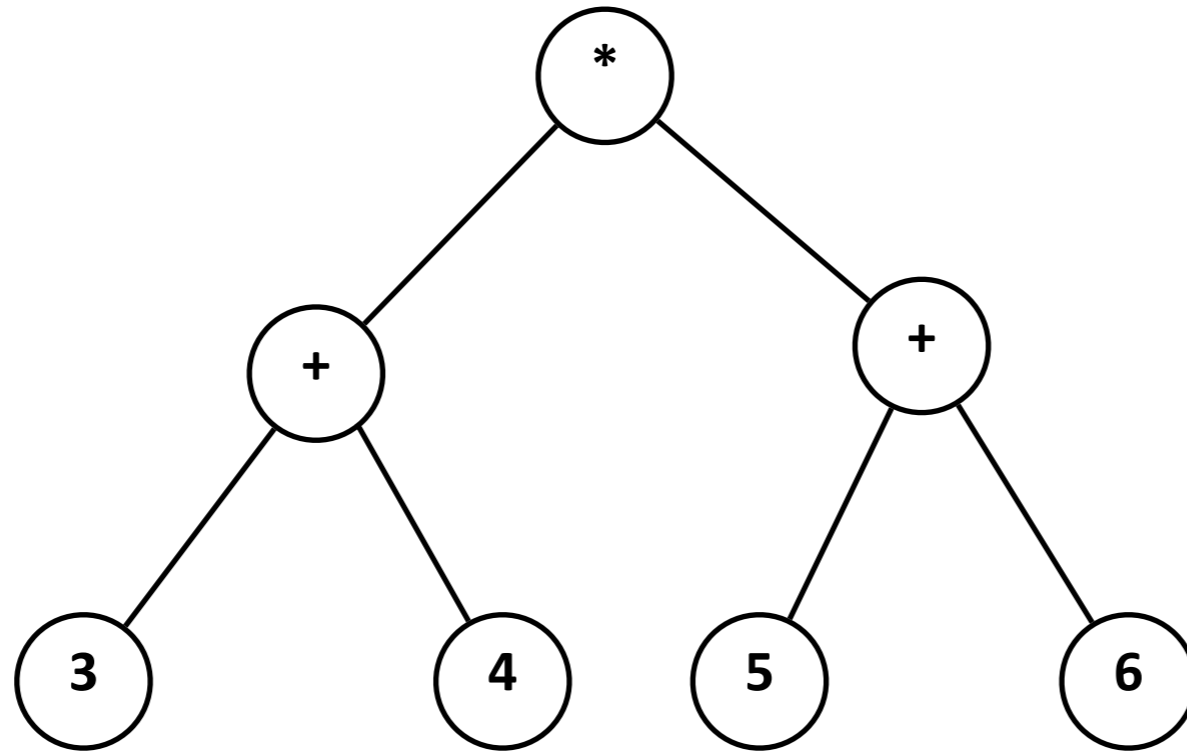
```
fun add (x,y,k) = k (x + y)
```

```
fun mult (x,y,k) = k (x * y)
```

```
(3 + 4) * (5 + 6)
```

```
add (3,4, fn r1 =>  
      add (5, 6, fn r2 =>  
            mult (r1, r2, fn r3 => r3)))
```

$$(3 + 4) * (5 + 6)$$



add 3 4 r1

add 5 6 r2

mult r1 r2 r3

return r3

add 3 and 4 and put the result in **r1** ...

makes control flow and intermediate results explicit

```
(* sum : int list -> int
   REQUIRES: true
   ENSURES:  sum L returns the sum of all the elements in L
*)
```

```
  fun sum ([ ] : int list) : int = 0
    | sum (x::xs) = x + sum xs
```

```
(* Using tail-recursion: *)
```

```
(* tsum : int list * int -> int
```

```
   REQUIRES: true
```

```
   ENSURES:  tsum (L, acc)  $\cong$  (sum L) + acc
```

```
*)
```

```
  fun tsum ([ ] : int list, acc : int) : int = acc
    | tsum (x::xs, acc) = tsum(xs, x + acc)
```

```
(* Using foldr: *)
```

```
  val Lsum = foldr (op +) 0
```

```
(* sum : int list -> int
   REQUIRES: true
   ENSURES:  sum L returns the sum of all the elements in L
*)
```

```
  fun sum ([ ] : int list) : int = 0
    | sum (x::xs) = x + sum xs
```

```
(* Using tail-recursion: *)
(* tsum : int list * int -> int
   REQUIRES: true
   ENSURES:  tsum (L, acc)  $\cong$  (sum L) + acc
*)
```

```
  fun tsum ([ ] : int list, acc : int) : int = acc
    | tsum (x::xs, acc) = tsum(xs, x + acc)
```

```
(* Using continuation-passing style: *)
(* csum : int list -> (int -> 'a) -> 'a
   REQUIRES: k is total
   ENSURES:  csum L k  $\cong$  k (sum L)
*)
```



```
(* sum : int list -> int
   REQUIRES: true
   ENSURES:  sum L returns the sum of all the elements in L
*)
```

```
  fun sum ([ ] : int list) : int = 0
    | sum (x::xs) = x + sum xs
```

```
(* Using tail-recursion: *)
(* tsum : int list * int -> int
   REQUIRES: true
   ENSURES:  tsum (L, acc)  $\cong$  (sum L) + acc
*)
```

```
  fun tsum ([ ] : int list, acc : int) : int = acc
    | tsum (x::xs, acc) = tsum(xs, x + acc)
```

```
(* Using continuation-passing style: *)
(* csum : int list -> (int -> 'a) -> 'a
   REQUIRES: k is total
   ENSURES:  csum L k  $\cong$  k (sum L)
*)
```

```
  fun csum ([ ] : int list) (k: int -> 'a) : 'a = _____
    | csum (x::xs) k = _____
```



```
(* sum : int list -> int
   REQUIRES: true
   ENSURES:  sum L returns the sum of all the elements in L
*)
```

```
  fun sum ([ ] : int list) : int = 0
    | sum (x::xs) = x + sum xs
```

```
(* Using tail-recursion: *)
(* tsum : int list * int -> int
   REQUIRES: true
   ENSURES:  tsum (L, acc)  $\cong$  (sum L) + acc
*)
```

```
  fun tsum ([ ] : int list, acc : int) : int = acc
    | tsum (x::xs, acc) = tsum(xs, x + acc)
```

```
(* Using continuation-passing style: *)
(* csum : int list -> (int -> 'a) -> 'a
   REQUIRES: k is total
   ENSURES:  csum L k  $\cong$  k (sum L)
*)
```

```
  fun csum ([ ] : int list) (k: int -> 'a) : 'a = k(0)
    | csum (x::xs) k = csum xs (fn s => k(x + s))
```

```
fun csum ([ ] : int list) (k: int -> 'a) : 'a = k(0)
  | csum (x::xs) k = csum xs (fn s => k(x + s))
```

csum [2,3] (fn s => s)

≡ csum [3] (fn s' => (fn s => s)(2 + s'))

≡ csum [] (fn s'' => (fn s' => (fn s => s)(2 + s')) (3 + s''))

≡ (fn s'' => (fn s' => (fn s => s) (2 + s')) (3 + s'')) 0

≡ (fn s' => (fn s => s) (2 + s')) (3 + 0)

≡ (fn s => s) (2 + 3)

≡ 5

Comparing tree contents to list contents

```
datatype tree = Empty | Node of tree * int * tree
```

```
(* inorder : tree * int list --> int list
```

```
  REQUIRES: true
```

```
  ENSURES: inorder(T, acc)  $\cong$  L @ acc, where L consists of the  
           elements of T as encountered in an in-order traversal of T.
```

```
*)
```

```
fun inorder(Empty, acc) = acc
```

```
  | inorder(Node(left, n, right), acc) =  
    inorder(left, n::inorder(right, acc))
```

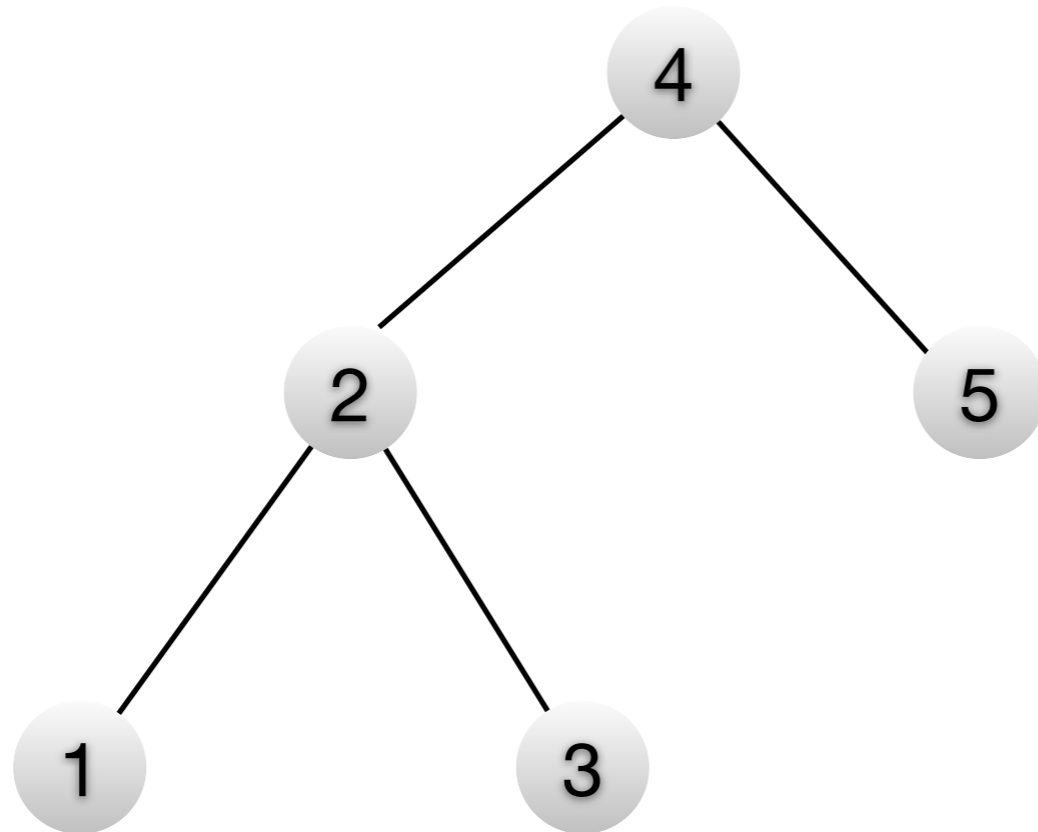
```
(* treematch : tree -> int list -> bool
```

```
  REQUIRES: k is total
```

```
  ENSURES: treematch T L returns true if L consists of the  
           elements of T as encountered in an in-order traversal of T,  
           and returns false otherwise.
```

```
*)
```

Suppose T is bound to the following tree:



```
val true = treematch T [1,2,3,4,5]
```

```
val false = treematch T [1,4,5,2]
```

```
fun inorder(Empty, acc) = acc
  | inorder(Node(left, n, right), acc) =
      inorder(left, n::inorder(right, acc))
```

```
(* same : int list * int list -> bool *)
```

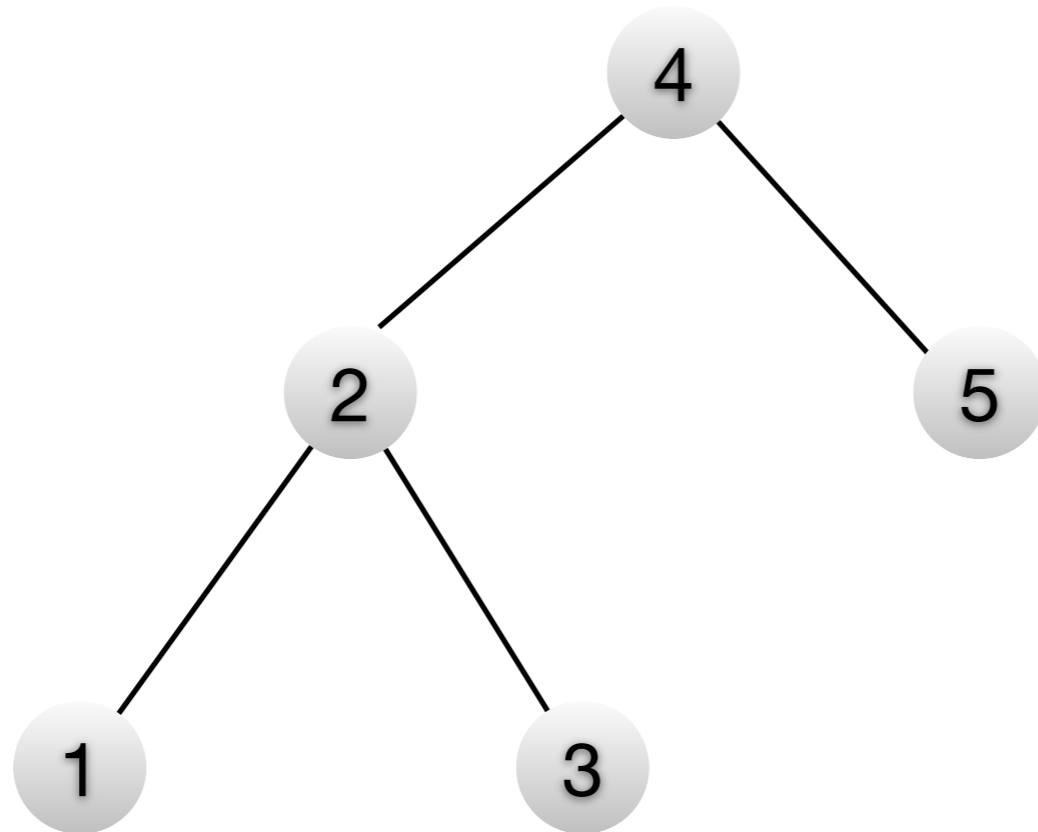
```
fun same ([ ], [ ]) = true
  | same (x::xs, y::ys) = (x = y) andalso same(xs, ys)
  | same _ = false
```

```
(* treematch : tree -> int list -> bool
   REQUIRES: k is total
   ENSURES: treematch T L returns true if L consists of the
             elements of T as encountered in an in-order
             traversal of T, and returns false otherwise.
```

```
*)
```

```
fun treematch T L = same(inorder(T, nil) ,L))
```

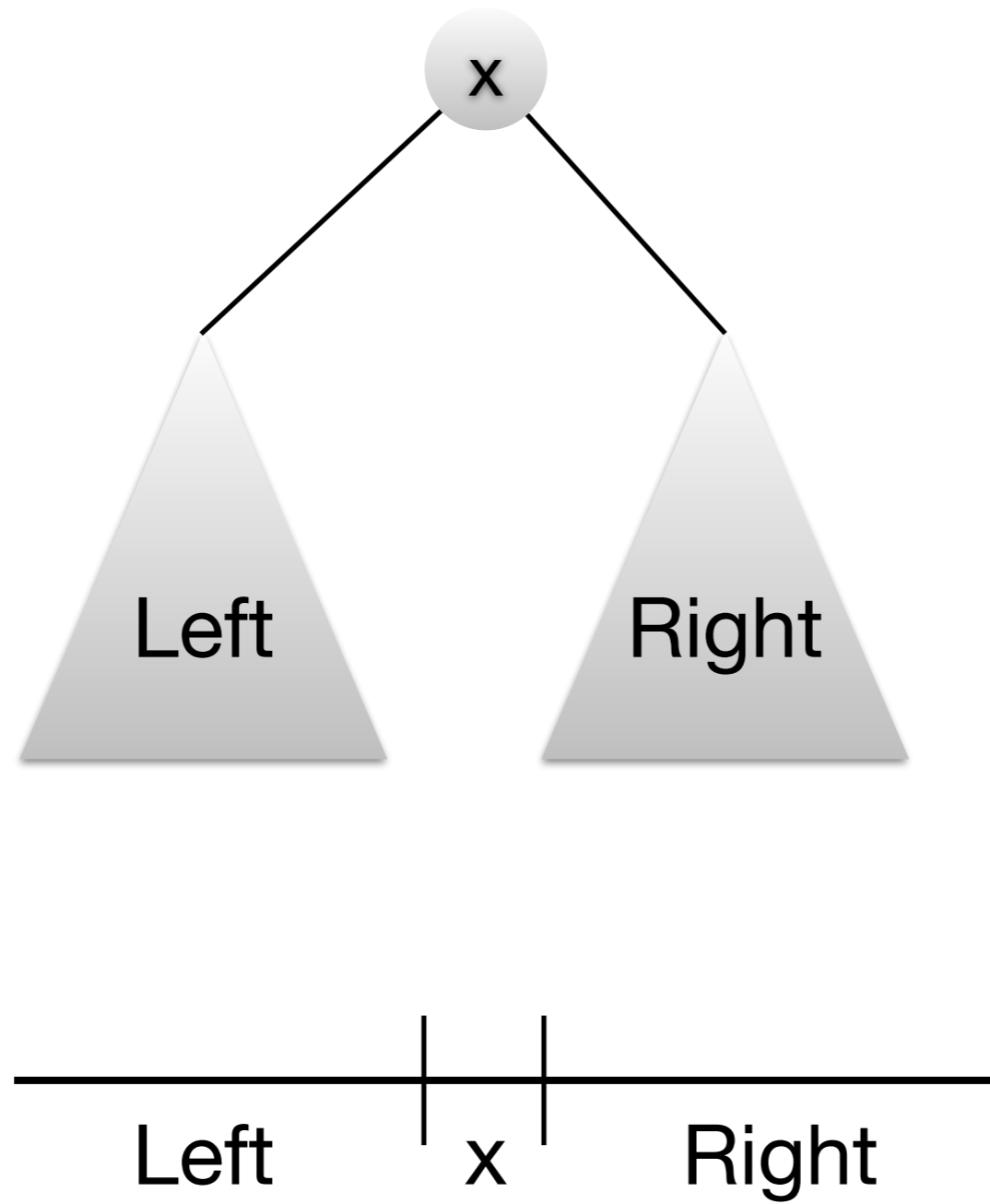
Suppose T is bound to the following tree:



Could we make treematch be faster?

```
val false = treematch T [1,4,5,2]
```


prefix : tree -> int list -> (int list -> bool) -> bool



Matching tree contents to prefix of a list

```
(* prefix : tree -> int list -> (int list -> bool) -> bool
```

```
  REQUIRES: k is total
```

```
  ENSURES: prefix T L k ==> true, if L ≅ L1 @ L2, such that  
           the inorder traversal of T is equal to L1, and k(L2) ≅ true.  
           false, otherwise.
```

*)

```
fun prefix (Empty) L k = k(L)
```

```
  | prefix (Node(left, n, right)) L k = prefix left L _____
```

```
(* treematch' : tree -> int list -> bool
```

```
  REQUIRES: true
```

```
  ENSURES: treematch' T L returns true if L consists of the  
           elements of T as encountered in an in-order traversal of T,  
           and returns false otherwise.
```

*)

```
fun treematch' T L = prefix T L List.null
```

Matching tree contents to prefix of a list

```
(* prefix : tree -> int list -> (int list -> bool) -> bool
```

```
  REQUIRES: k is total
```

```
  ENSURES: prefix T L k ==> true, if L ≅ L1 @ L2, such that  
           the inorder traversal of T is equal to L1, and k(L2) ≅ true.  
           false, otherwise.
```

```
*)
```

```
fun prefix (Empty) L k = k(L)
```

```
  | prefix (Node(left, n, right)) L k =
```

```
    prefix left L (fn [] => false
```

```
                  |(y::ys) => (n=y) andalso (prefix right ys k))
```

```
(* treematch' : tree -> int list -> bool
```

```
  REQUIRES: true
```

```
  ENSURES: treematch' T L returns true if L consists of the  
           elements of T as encountered in an in-order traversal of T,  
           and returns false otherwise.
```

```
*)
```

```
fun treematch' T L = prefix T L List.null
```

Search problems

- Implemented using 2 continuations:
 - success (what to do if search is successful)
 - failure (takes unit as argument, implements backtracking)

search : ('a -> bool) -> 'a tree -> ('a -> 'b) -> (unit -> 'b) -> 'b



success continuation



failure continuation

```
(* search : ('a -> bool) -> 'a tree -> ('a -> 'b) -> (unit -> 'b) -> 'b
  REQUIRES: p, sc, and f are total.
  ENSURES:  search p T sc fc ≡ sc(x), if p(x) ≡ true for some x in T
           ≡ fc(), otherwise
           (if more than one x satisfies p(x) ≡ true, then use the first
            encountered in a pre-order traversal of T).
```

*)

```
fun search p Empty sc fc = fc()
  | search p (Node(left, x, right)) sc fc =
      if p(x) then sc(x)
      else
          search p left sc (fn () => search p right sc fc)
```

```
(* findeven : int tree -> string
  REQUIRES: true
  ENSURES:  findeven(T) returns the string representation of the
           first even integer found in a pre-order traversal of T,
           if there is such an integer. Otherwise, findeven(T)
           returns "none found".
```

*)

```
fun findeven T = search (fn n => n mod 2 = 0) T Int.toString (fn () => "none found")
```

```
(* search : ('a -> bool) -> 'a tree -> ('a -> 'b) -> (unit -> 'b) -> 'b
  REQUIRES: p, sc, and f are total.
  ENSURES:  search p T sc fc ≡ sc(x), if p(x) ≡ true for some x in T
           ≡ fc(), otherwise
           (if more than one x satisfies p(x) ≡ true, then use the first
            encountered in a pre-order traversal of T).
```

*)

```
fun search p Empty sc fc = fc()
  | search p (Node(left, x, right)) sc fc =
      if p(x) then sc(x)
      else
          search p left sc (fn () => search p right sc fc)
```

```
(* findeven : int tree -> string
  REQUIRES: true
  ENSURES:  findeven(T) returns SOME x where x is the
            first even integer found in a pre-order traversal of T,
            if there is such an integer. Otherwise, findeven(T)
            returns NONE.
```

*)

```
fun findeven T = search (fn n => n mod 2 = 0) T SOME (fn () => NONE)
```