

Libraries

Reusing Code

- All but the simplest programs reuse code already written
 - system code
 - `#use <conio>` → C0's input/output functions
 - simple code you wrote in the past
 - `pixel-int.c0` → Encoding and manipulation of pixels as ints
 - complex code somebody else wrote
 - `pixel.o0` → Encoding and manipulation of pixels
 - `jquery.js` → Javascript utilities to build dynamic web pages
- Why?
 - Writing correct code is hard and time-consuming!
- These are **libraries**
 - They separate out code used across many applications from the applications themselves

Abstraction

- Libraries promote **abstraction**
 - Focus on **what** the library code does
 - e.g., print an integer to terminal using `printint`
 - not on **how** it does it
 - the many minute steps to turn a integer into terminal output
- Abstraction has lots of **benefits**
 - Hide inessential details
 - writing code is hard enough without also having to know how `printint` works
 - Make code more manageable
 - if we find a bug in `printint`, there is a single place where to fix it
 - Allow for transparent improvements
 - if we find a better way of printing, update the library not the applications

Computer science is all about abstraction!

What's a Library Anyway?

1. The interface

- Lists the functionalities the library exports and how to use them

```
void printint(int i);
```

Everything we need to use this functionality:

- name of the function
- number and type of arguments
- output type
- contracts

2. The implementation

- The code that implements them

```
void printint(int i) {  
    ...  
}
```

Complex low-level code

3. The documentation

- The explanation of what they do
“print i to standard output”

Human readable,
often in a web page or thick manual

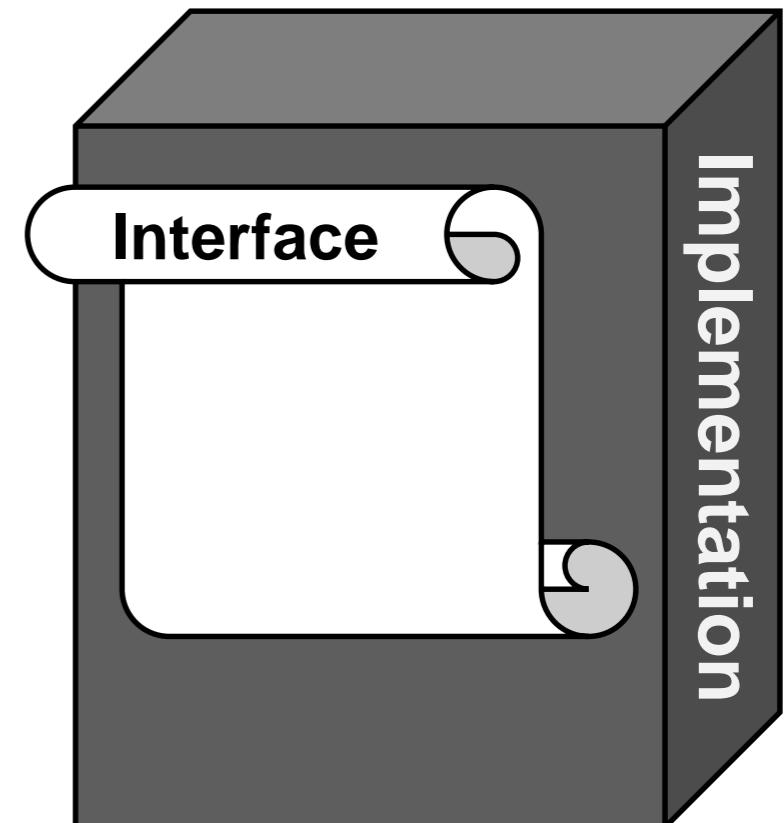
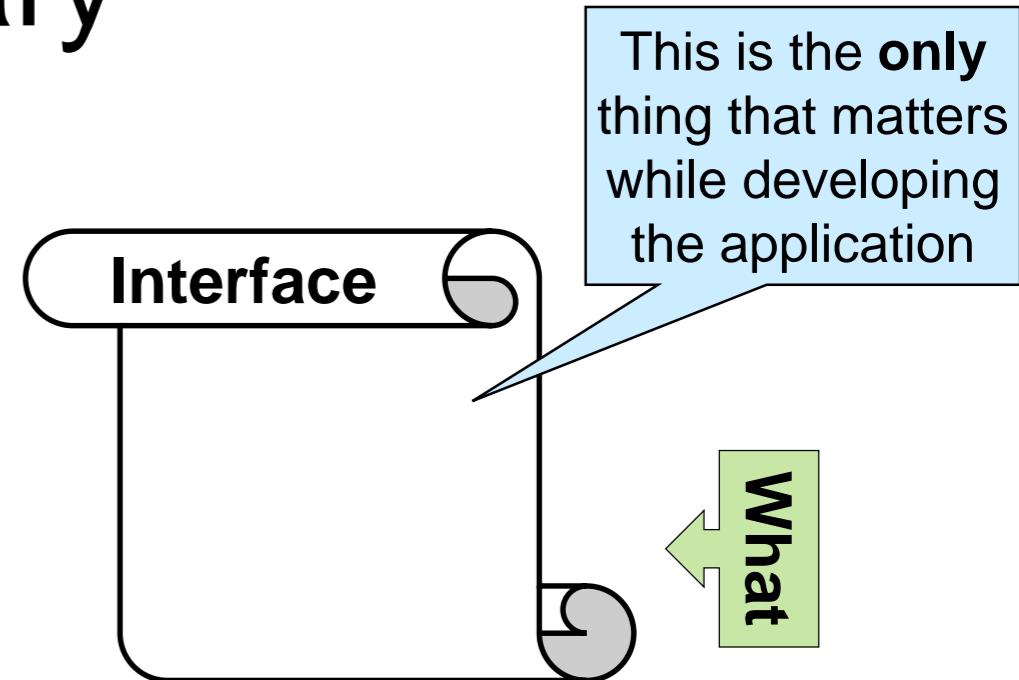
What

How

What

Using a Library

- When **writing** application code, we only use the functionalities listed in the interface
 - *No reliance on implementation*
- When **compiling** the application, we involve an implementation of the library
 - Needed for the application to run
- Implementation is a **black box**



Types of Libraries

- System libraries
 - part of the programming language
#use <conio>
 - No need to load any file to use them

Linux Terminal

```
# cc0 -d my-math-application.c0
```

- User-defined libraries
 - written by users or downloaded from the Internet
pixels.c0 —————— **or pixels.o0**
 - must be compiled with the application

Also called an
API

**Application
Programming
Interface**

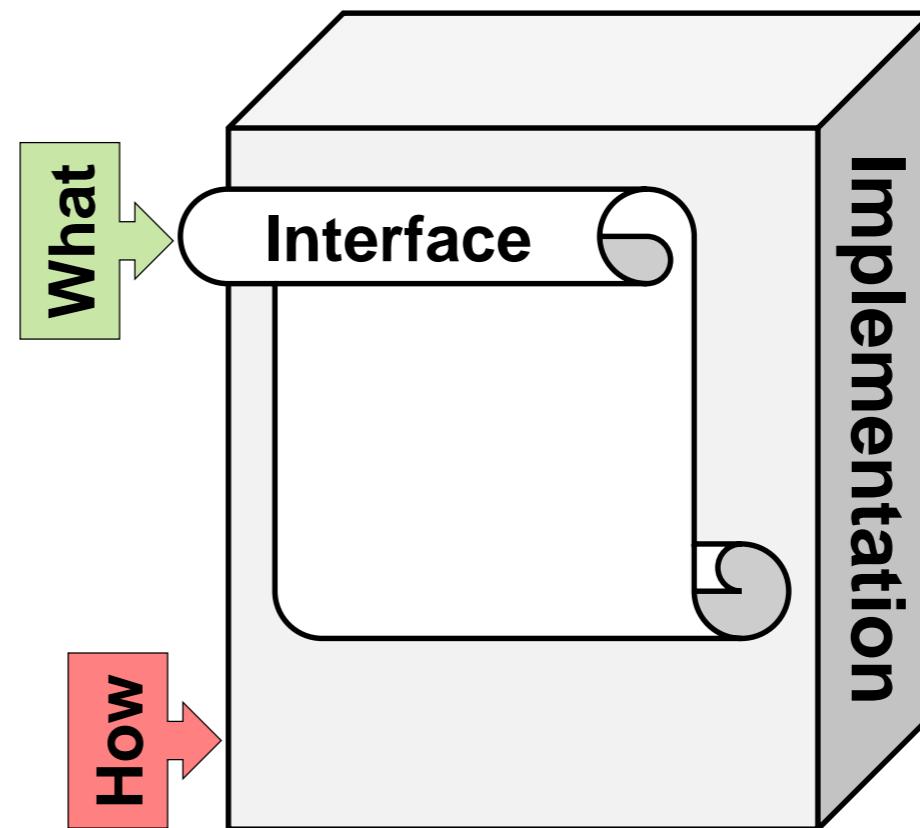
Linux Terminal

```
# cc0 -d pixels.c0 my-image-application.c0
```

Writing a Library

- When writing a library, we need to
 - decide on the interface
 - implement every functionality exported by the interface
 - Fill the black box
 - write lots of documentation

- In this class, we will be writing some of the system libraries that are native in other languages



Abstract Data Types

- A library that defines a **new type** and the ways to use it

E.g.,
pixels

- Defines the type `pixel_t` of pixels
 - The **only** way we shall refer to pixel in application
- Defines functions that manipulate pixels

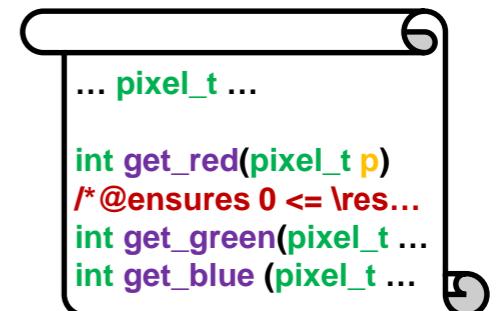
```
int get_red(pixel_t p)
/*@ensures 0 <= \result && \result < 256; @*/
int get_green(pixel_t p) ...
int get_blue(pixel_t p) ...
pixel_t make_pixel(int red, int green, int blue) ...
```

- The **only** operations we shall use to manipulate pixels
 - Except for functions we write using them

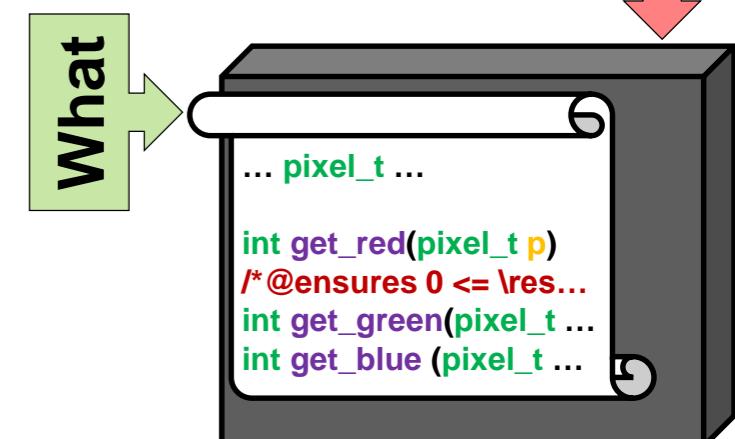
- ADT's promote a **very strong form of abstraction**

- If the client only uses the interface, we can use *any correct implementation* and the application will work the same!

This is the pixel interface



How



Self-Sorting Arrays

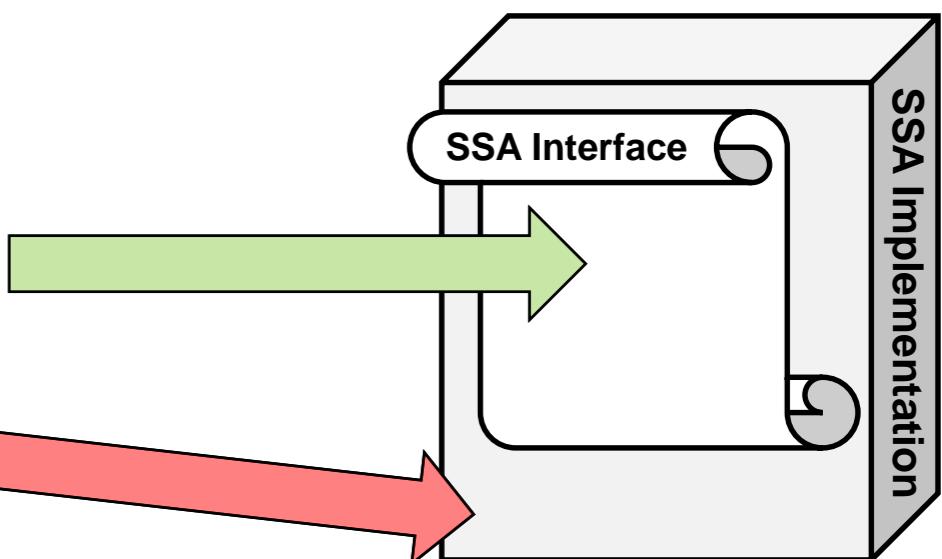
Writing Libraries

- In this course, we will be writing many libraries
- Case study to gain familiarity
 - and define important concepts

Self-sorting arrays (SSA)

- a toy data structure that works just like arrays of strings but
 - has a function that reports the length of the SSA
 - guarantees that its elements are sorted

- What we need to do:
 - A. Define the interface of the SSA library
 - B. Implement it



SSA Interface

```
// typedef _____ ssa_t;
```

Interface Contents

1. A **type** for self-sorting arrays

ssa_t

- SSA's are a data structure
- We need a type to refer to them in code
 - define variables that can hold an SSA
 - define functions that manipulate them, ...

Convention: types
exported by a
library end in **_t**

● We do **not** want the client to learn the details of this type

What →

- This type is **abstract** for the client: *just a name*
- We define it as a **pseudo-typedef**

A commented-out **typedef**
with underscores

```
// typedef _____ ssa_t;
```

Another convention

How →

- The implementation will contain the actual definition of **ssa_t**
 - **Concrete** type of SSA's

```
// typedef _____ ssa_t;  
  
// Operations
```

Interface Contents

2. The **operations** provided by the library to manipulate SSA's

- What should these be?
 - SSA's are just fancy arrays
 - We will need SSA versions of the standard operations on arrays
 - create a new array
 - read a value from an array index
 - replace the value at an array index
 - *Plus a function that returns the length*
 - not just in contracts, but in regular code

Interface Contents

2. The **operations** provided by the library to manipulate SSA's

- Creating a new SSA

Newly
created
SSA

```
ssa_t ssa_new(int size); // akin to alloc_array(string, size)
```

Number of elements

This is a **function prototype**:
a function definition without a body

- Reading the value at an index of an SSA

```
string ssa_get(ssa_t A, int i); // akin to ... A[i] ...
```

- Replacing the value at an index of an SSA

```
void ssa_set(ssa_t A, int i, string x); // akin to A[i] = x
```

➤ unlike regular arrays, this may rearrange the contents of A to keep it sorted

- Returning the length of an SSA

```
int ssa_len(ssa_t A); // akin to \length(A) but better
```

➤ unlike regular array, this can be used anywhere in code

Recall that our SSAs
contain strings
(we'll learn later how to generalize)

```
// typedef _____ ssa_t;

int ssa_len(ssa_t A);

ssa_t ssa_new(int size) ;

string ssa_get(ssa_t A, int i);

void ssa_set(ssa_t A, int i, string x);
```

Interface Contents

3. The **contracts** of each operation

- The client needs to be able to write safe code
 - Provide arguments that satisfy the preconditions of each function
 - Use the functions' postconditions to reason about follow-up code
- Reading the value at an index of an SSA
 - Same contracts as native A[i]
 - /*@requires 0 <= i && i < \length(A); @*/
 - So:


```
string ssa_get(ssa_t A, int i) // akin to ... A[i] ...
/*@requires 0 <= i && i < \length(A); @*/ ;
```

Can this be right?

SSA Interface

```
// typedef _____ ssa_t;  
  
int ssa_len(ssa_t A);  
  
ssa_t ssa_new(int size);  
  
string ssa_get(ssa_t A, int i);  
  
void ssa_set(ssa_t A, int i, string x);
```

Interface Contents

3. The **contracts** of each operation

```
string ssa_get(ssa_t A, int i) // akin to ... A[i] ...  
/*@requires 0 <= i && i < \length(A); @*/ ;
```

Can this be right?

- **\length** is defined only for C0 arrays
 - But SSAs are not C0 arrays
 - length can be retrieved
 - elements stay sorted
 - should be manipulated only with operations in the SSA interface
- We can however use **ssa_len**

```
string ssa_get(ssa_t A, int i) // akin to ... A[i] ...  
/*@requires 0 <= i && i < ssa_len(A); @*/ ;
```

Interface Contents

3. The **contracts** of each operation

- With `ssa_len`, we can give a meaningful precondition to `ssa_get`
- and to `ssa_set`
- and a postcondition to `ssa_new`

```
// typedef _____ssa_t;

int ssa_len(ssa_t A)
/*@ensures \result >= 0; @*/ ;

ssa_t ssa_new(int size)
/*@requires 0 <= size; @*/
/*@ensures ssa_len(\result) == size; @*/ ;

string ssa_get(ssa_t A, int i)
/*@requires 0 <= i && i < ssa_len(A); @*/ ;

void ssa_set(ssa_t A, int i, string x)
/*@requires 0 <= i && i < ssa_len(A); @*/ ;
```

Interface Contents

SSA Interface

```
// typedef _____ *ssa_t;  
  
int ssa_len(ssa_t A) /*@requires A != NULL; */ /*@ensures \result >= 0; */  
  
ssa_t ssa_new(int size) /*@requires 0 <= size; */ /*@ensures \result != NULL; */ /*@ensures ssa_len(\result) == size; */  
  
string ssa_get(ssa_t A, int i) /*@requires A != NULL; */ /*@requires 0 <= i && i < ssa_len(A); */  
  
void ssa_set(ssa_t A, int i, string x) /*@requires A != NULL; */ /*@requires 0 <= i && i < ssa_len(A); */
```

- But what kind of type can `ssa_t` be?
 - An array? ✗
 - No way to get the length of an array in C0
 - An `int`, `bool` or `char`? ✗
 - No way to represent arbitrarily many strings
 - A `string`? ✗
 - Let's not go there ...
 - A struct? ✗
 - Structs cannot be passed as function arguments directly

- Then, `ssa_t` must be a **pointer**

- Update the pseudo-`typedef` to reflect this
- Disallow `NULL` as a valid `ssa_t`

- Every operation that takes an `ssa_t` has a `NULL`-check as a precondition
- Every operation that returns an `ssa_t` has a `NULL`-check as a postcondition

We never use `NULL` for an empty data structure

Interface Contents

SSA Interface

```
// typedef _____ * ssa_t;

int ssa_len(ssa_t A)
/*@requires A != NULL; @*/
/*@ensures \result >= 0; @*/
;

ssa_t ssa_new(int size)
/*@requires 0 <= size; @*/
/*@ensures \result != NULL; @*/
/*@ensures ssa_len(\result) == size; @*/
;

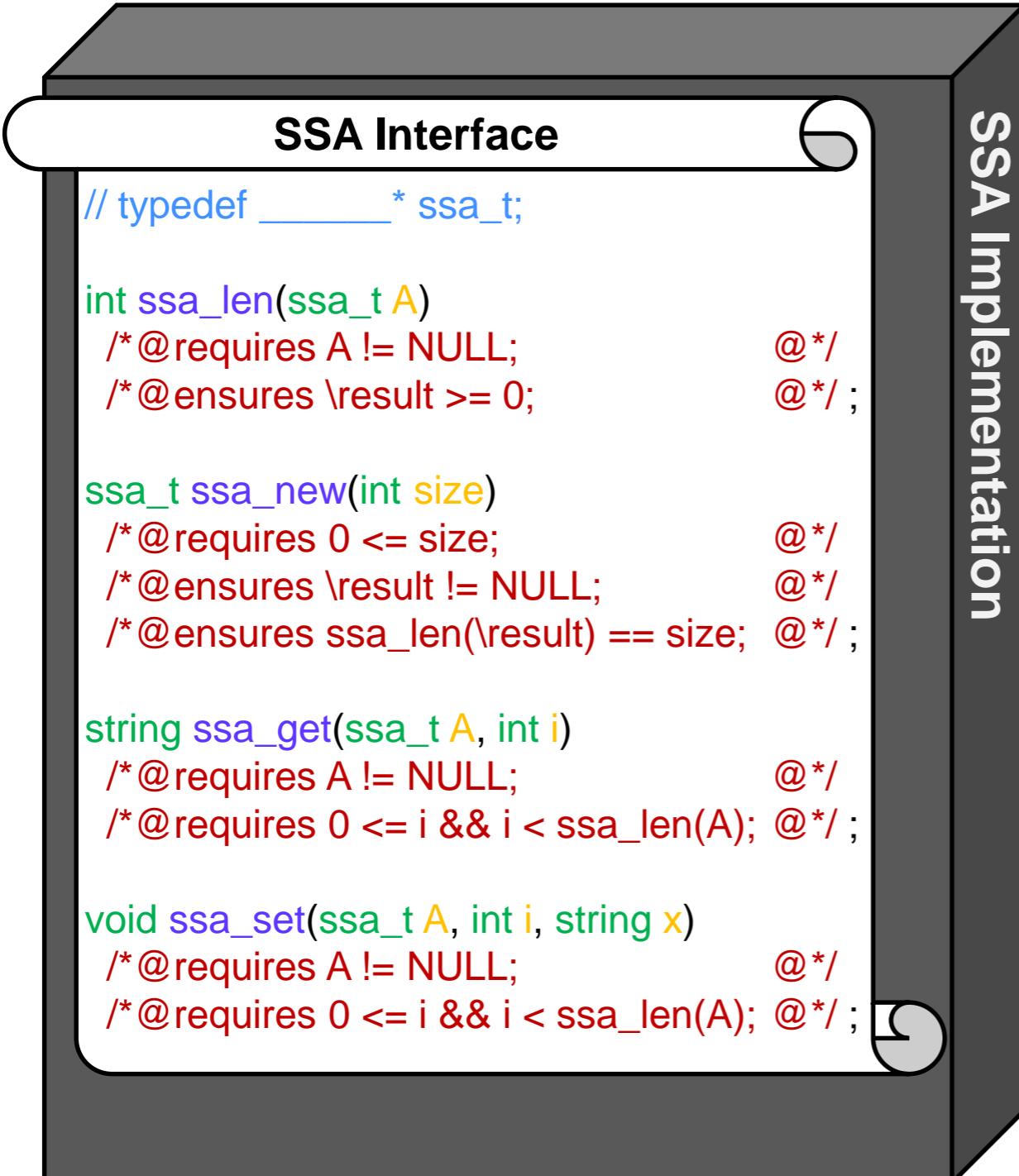
string ssa_get(ssa_t A, int i)
/*@requires A != NULL; @*/
/*@requires 0 <= i && i < ssa_len(A); @*/
;

void ssa_set(ssa_t A, int i, string x)
/*@requires A != NULL; @*/
/*@requires 0 <= i && i < ssa_len(A); @*/
;
```

What

Client Application

Using a library



- The client only knows **what** the library does
 - the library interface
 - the library documentation
- The client does not know **how** it does it
 - treat the implementation as a black box
 - even if its code is available
 - it may change!

Searching an SSA

- Client code that uses binary search to check if a value is in an SSA
 - This is OK because SSAs are sorted!

```
bool is_in(string x, ssa_t A, int n)
//@requires A != NULL;
//@requires n == ssa_len(A);
{
  int lo = 0;
  int hi = n;
  while (lo < hi)
    //@loop_invariant 0 <= lo && lo <= hi && hi <= n;
  {
    int mid = lo + (hi - lo) / 2;
    //@assert lo <= mid && mid < hi;
    string a = ssa_get(A, mid);
    int cmp = string_compare(a, x);
    if (cmp == 0) return true;
    if (cmp < 0) {
      lo = mid + 1;
    } else { //@assert cmp > 0;
      hi = mid;
    }
  }
  return false;
}
```

Precondition of
ssa_len and ssa_get

```
// typedef _____* ssa_t;

int ssa_len(ssa_t A)
/*@requires A != NULL;
/*@ensures \result >= 0;
@*/;

sssa_t ssa_new(int size)
/*@requires 0 <= size;
/*@ensures \result != NULL;
/*@ensures ssa_len(\result) == size; @*/;

string ssa_get(ssa_t A, int i)
/*@requires A != NULL;
/*@requires 0 <= i && i < ssa_len(A); @*/;

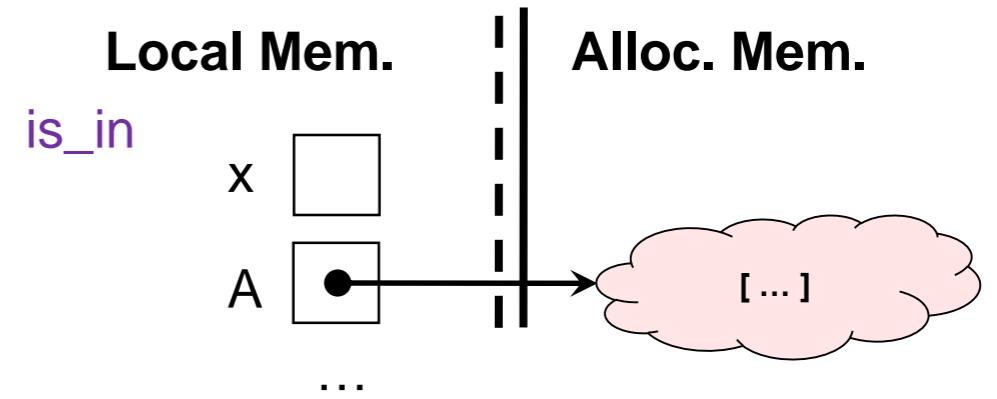
void ssa_set(ssa_t A, int i, string x)
/*@requires A != NULL;
/*@requires 0 <= i && i < ssa_len(A); @*/;
```

- All array operations are replaced with functions from the SSA interface
 - returns <0 if a “less than” x, 0 if equal, >0 otherwise
- **Safety** is supported by loop invariant and assertion
- For **correctness**, we would need to implement SSA versions of **gt_seg** and **lt_seg**

Searching an SSA

- Client view of memory

```
bool is_in(string x, ssa_t A, int n)
//@requires n == ssa_len(A);
{
    int lo = 0;
    int hi = n;
    while (lo < hi)
        //@loop_invariant 0 <= lo && lo <= hi && hi <= n;
    {
        int mid = lo + (hi - lo) / 2;
        //@assert lo <= mid && mid < hi;
        string a = ssa_get(A, mid);
        int cmp = string_compare(a, x);
        if (cmp == 0) return true;
        if (cmp < 0) {
            lo = mid + 1;
        } else { //@assert cmp > 0;
            hi = mid;
        }
    }
    return false;
}
```

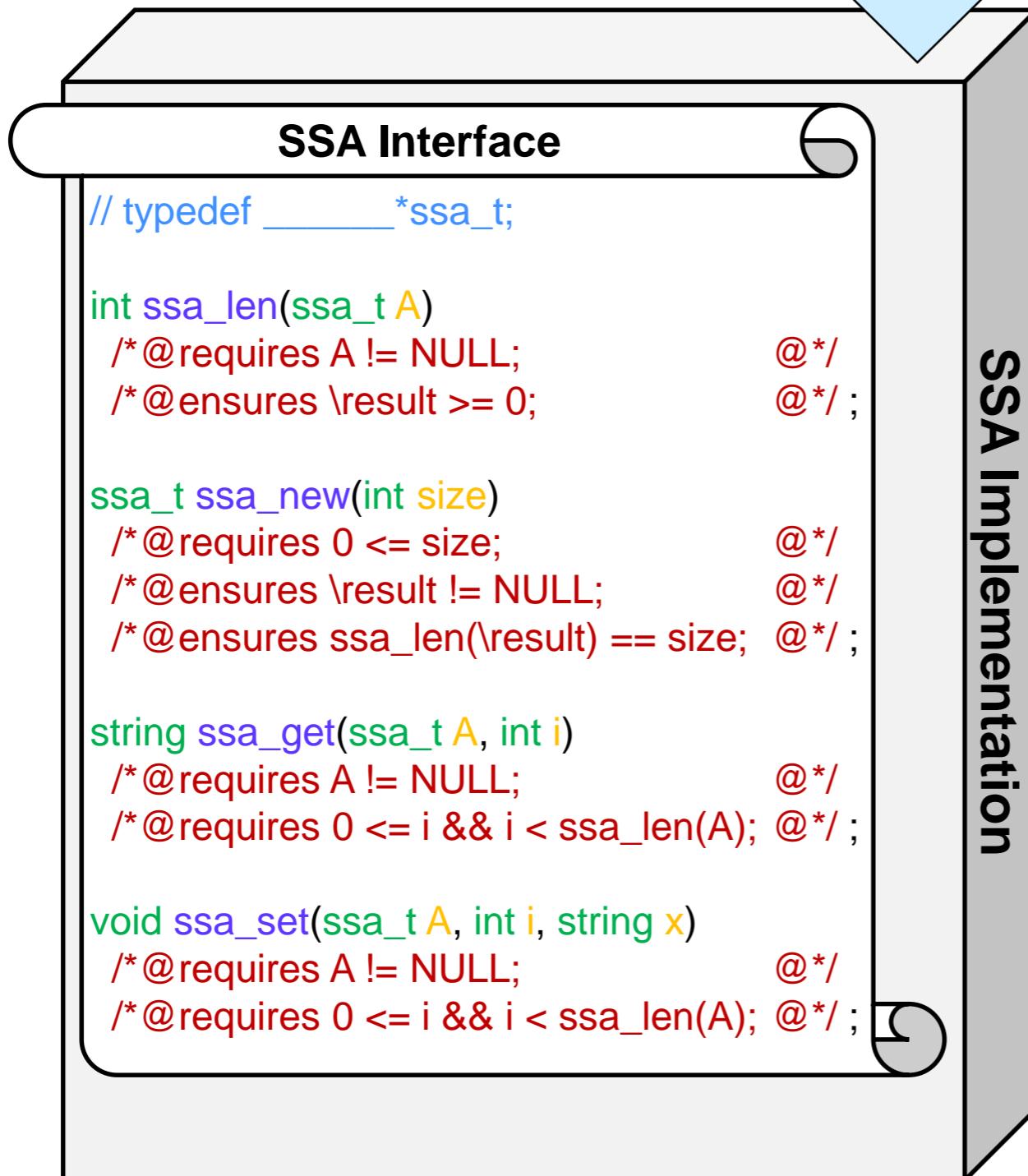


- The client has no knowledge of how `A` is represented in memory

SSA Implementation

Implementing SSAs

Now we've got
to fill the box



- Define the type `ssa_t`
 - **Concrete** type
- Write code for every function
- Make sure it is safe and correct

How

```
// typedef _____ * ssa_t;
// ...
```

Concrete Type

- Store elements in a C0 array, but keep track of the length
 - Package them together in a struct

This is the **concrete** implementation type

- Define an internal nickname for it
 - So that the code is succinct and readable
 - It's convenient that it **not** be a pointer
- Define the abstract type exported to the client
 - This is what connects the concrete implementation type with the exported abstract type

Internal nickname

```
// Implementation-side type
struct ssa_header {
    string[] data; // sorted
    int length; // = \length(data)
};
typedef struct ssa_header ssa;

// ... rest of implementation

// Client type
typedef ssa* ssa_t;
```

SSA Implementation

Abstract client type

Client vs. Implementation View

An SSA containing "a", "b" and "c"

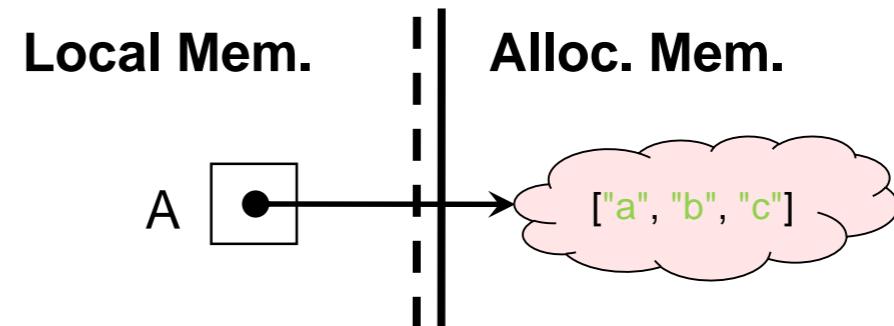
```
// Implementation-side type
struct ssa_header {
    string[] data; // sorted
    int length;   // = \length(data)
};
typedef struct ssa_header ssa;

// ... rest of implementation

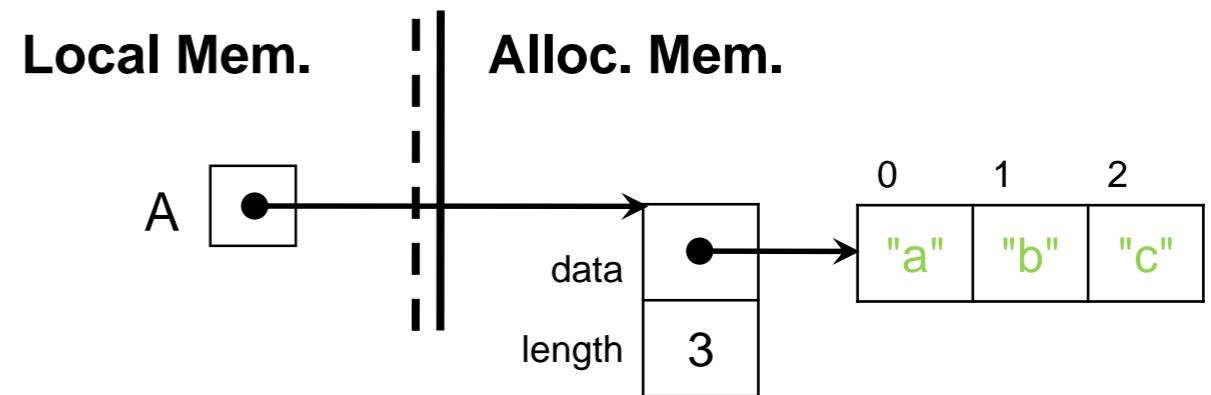
// Client type
typedef ssa* ssa_t;
```

SSA Implementation

- Client view



- Implementation view



SSA Interface

```
string ssa_get(ssa_t A, int i)
/*@requires A != NULL;
/*@requires 0 <= i && i < ssa_len(A); @*/;
```

Implementing ssa_get

```
struct ssa_header {
    string[] data; // sorted
    int length;   // = \length(data)
};
typedef struct ssa_header ssa;

string ssa_get(ssa* A, int i)
//@requires A != NULL;
//@requires 0 <= i && i < ssa_len(A);
{
    return A->data[i];
}

// ... rest of implementation
```

SSA Implementation

- Simply return the i -th element of the underlying array
$$\text{return } A\text{-}>\text{data}[i]$$
- Is this safe? We need to check
 - $A \neq \text{NULL}$
 - By 1st precondition
 - $i \leq 0$
 - By 2nd precondition (first conjunct)
 - $i < \text{\length}(A\text{-}>\text{data})$
 - We know that $i < \text{ssa_len}(A)$
 - but we don't know how $\text{ssa_len}(A)$ and $\text{\length}(A\text{-}>\text{data})$ are related
 - ***Not supported!***



Let's also write `ssa_len`

```

struct ssa_header {
    string[] data; // sorted
    int length;   // = \length(data)
};
typedef struct ssa_header ssa;

int ssa_len(ssa* A)
//@requires A != NULL;
//@ensures \result >= 0;
{
    return A->length;
}

string ssa_get(ssa* A, int i)
//@requires A != NULL;
//@requires 0 <= i && i < ssa_len(A);
{
    return A->data[i];
}

// ... rest of implementation

```

SSA Implementation

- Simply return the length field
- Is this safe? We need to check
 - **A != NULL**
 - By precondition
- Does this help us with `ssa_get`?
 - No useful postcondition
 - Peeking at the code of `ssa_len` would be operational reasoning!



```

int ssa_len(ssa_t A)
/*@requires A != NULL;
/*@ensures \result >= 0;
@*/;
@*/;

string ssa_get(ssa_t A, int i)
/*@requires A != NULL;
/*@requires 0 <= i && i < ssa_len(A); @*/
@*/;
@*/;

```

Let's also write `ssa_len`

```
struct ssa_header {
    string[] data; // sorted
    int length;   // = \length(data)
};
typedef struct ssa_header ssa;
```

```
int ssa_len(ssa* A)
//@requires A != NULL;
//@ensures \result >= 0;
//@ensures \result == \length(A->data);
```

```
{  
    return A->length;  
}
```

```
string ssa_get(ssa* A, int i)
//@requires A != NULL;
//@requires 0 <= i && i < ssa_len(A);
{
    return A->data[i];
}
```

```
// ... rest of implementation
```

SSA Implementation

```
int ssa_len(ssa_t A)
/*@requires A != NULL;
/*@ensures \result >= 0;
@*/;

string ssa_get(ssa_t A, int i)
/*@requires A != NULL;
/*@requires 0 <= i && i < ssa_len(A); @*/;
```

- Add a *useful* postcondition
 $\text{\result} == \text{\length}(\text{A}->\text{data})$
- Is this safe? We need to check
 - $\text{A} \neq \text{NULL}$ ✓
 - By precondition
- Is `ssa_len` correct?
 - No relation between $\text{A}->\text{length}$ and $\text{\length}(\text{A}->\text{data})$ ✗

Innocent mistake: define `ssa_new` as

```
ssa* ssa_new(int size) {
    ssa* A = alloc(ssa);
    A->length = size;
    A->data = alloc_array(string, size+1);
    return A;
}
```

Let's also write `ssa_len`

```
struct ssa_header {  
    string[] data; // sorted  
    int length;   // = \length(data)  
};  
typedef struct ssa_header ssa;
```

```
int ssa_len(ssa* A)  
//@requires A != NULL;  
//@requires A->length == \length(A->data);  
//@ensures \result >= 0;  
//@ensures \result == \length(A->data);  
{  
    return A->length;  
}
```

```
string ssa_get(ssa* A, int i)  
//@requires A != NULL;  
//@requires 0 <= i && i < ssa_len(A);  
{  
    return A->data[i];  
}  
// ... rest of implementation
```

SSA Implementation

SSA Interface

```
int ssa_len(ssa_t A)  
/*@requires A != NULL;  
/*@ensures \result >= 0;  
 @*/  
  
string ssa_get(ssa_t A, int i)  
/*@requires A != NULL;  
/*@requires 0 <= i && i < ssa_len(A); @*/;
```

- Add it as a precondition
 $A->\text{length} == \text{\length}(A->\text{data})$
- Is this safe? We need to check
 - $A \neq \text{NULL}$
 - By precondition
- Is `ssa_len` correct?
 - $A->\text{length} == \text{\length}(A->\text{data})$
 - By new precondition
 - $\text{\result} == A->\text{length}$
 - By code
 - $\text{\result} == \text{\length}(A->\text{data})$
 - By previous two



```

int ssa_len(ssa_t A)
/*@requires A != NULL;
/*@ensures \result >= 0;
@*/
@*/;

string ssa_get(ssa_t A, int i)
/*@requires A != NULL;
@*/
/*@requires 0 <= i && i < ssa_len(A); @*/
@*/;
```

Back to ssa_get

```

struct ssa_header {
    string[] data; // sorted
    int length;    // = \length(data)
};
typedef struct ssa_header ssa;
```

```

int ssa_len(ssa* A)
//@requires A != NULL;
//@requires A->length == \length(A->data);
//@ensures \result >= 0;
//@ensures \result == \length(A->data);
{
    return A->length;
}
```

```

string ssa_get(ssa* A, int i)
//@requires A != NULL;
//@requires 0 <= i && i < ssa_len(A);
{
    return A->data[i];
}

// ... rest of implementation
```

SSA Implementation

- Is the code for `ssa_get` safe?
 - The new postcondition of `ssa_len` takes care of the remaining safety check
 - $i < \length(A->data)$ ✓
 - But `ssa_len` has a new precondition
 - $A->length == \length(A->data)$
 - we need to have a reason for why it is satisfied
 - but we don't
 - ***Not supported!*** ✗

Back to ssa_get

```
struct ssa_header {
    string[] data; // sorted
    int length;   // = \length(data)
};

typedef struct ssa_header ssa;
```

```
int ssa_len(ssa* A)
//@requires A != NULL;
//@requires A->length == \length(A->data);
//@ensures \result >= 0;
//@ensures \result == \length(A->data);
{
    return A->length;
}
```

```
string ssa_get(ssa* A, int i)
//@requires A != NULL;
//@requires A->length == \length(A->data);
//@requires 0 <= i && i < ssa_len(A);
{
    return A->data[i];
}
```

// ... rest of implementation

SSA Implementation

- Is the code for ssa_get safe?
 - Add

A->length == \length(A->data)

as a precondition to ssa_get to support the safety of ssa_len

➤ A->length == \length(A->data)

□ By new precondition



```
int ssa_len(ssa_t A)
/*@requires A != NULL;
/*@ensures \result >= 0;
@*/;
@*/;
```

```
string ssa_get(ssa_t A, int i)
/*@requires A != NULL;
/*@requires 0 <= i && i < ssa_len(A); @*/;
@*/;
```

Representation Invariants

Where are we?

```
struct ssa_header {
    string[] data; // sorted
    int length;   // = \length(data)
};
typedef struct ssa_header ssa;
```

```
int ssa_len(ssa* A)
//@requires A != NULL;
//@requires A->length == \length(A->data);
//@ensures \result >= 0;
//@ensures \result == \length(A->data);
{
    return A->length;
}
```

```
string ssa_get(ssa* A, int i)
//@requires A != NULL;
//@requires A->length == \length(A->data);
//@requires 0 <= i && i < ssa_len(A);
{
    return A->data[i];
}
```

// ... rest of implementation

SSA Implementation

```
int ssa_len(ssa_t A)
/*@requires A != NULL;
/*@ensures \result >= 0;
@*/
@*/;

string ssa_get(ssa_t A, int i)
/*@requires A != NULL;
/*@requires 0 <= i && i < ssa_len(A); @*/
@*/;
```

- All our code is safe
- Both functions have preconditions

A != NULL

A->length == \length(A->data)

- `ssa_set` will need them too
- and `ssa_new` will have them as postconditions

Where are we?

```
struct ssa_header {
    string[] data; // sorted
    int length;   // = \length(data)
};
typedef struct ssa_header ssa;
```

```
int ssa_len(ssa* A)
//@requires A != NULL;
//@requires A->length == \length(A->data);
//@ensures \result >= 0;
//@ensures \result == \length(A->data);
{
    return A->length;
}
```

```
string ssa_get(ssa* A, int i)
//@requires A != NULL;
//@requires A->length == \length(A->data);
//@requires 0 <= i && i < ssa_len(A);
{
    return A->data[i];
}
```

// ... rest of implementation

SSA Implementation

```
int ssa_len(ssa_t A)
/*@requires A != NULL;
/*@ensures \result >= 0;
@*/
@*/;

string ssa_get(ssa_t A, int i)
/*@requires A != NULL;
@*/
/*@requires 0 <= i && i < ssa_len(A); @*/
@*/;
```

- They are **fundamental properties** an **ssa*** must obey to be the representation of a valid SSA
 - NULL is not a valid SSA
 - The length field must be equal to the length of the array field data
- These are **invariants** of our representation:
 - Preconditions of every library function that takes an SSA as a parameter
 - Postcondition of every library function that returns or modifies an SSA

Representation Invariants

```
struct ssa_header {
    string[] data; // sorted
    int length;   // = \length(data)
};
typedef struct ssa_header ssa;
```

```
int ssa_len(ssa* A)
//@requires A != NULL;
//@requires A->length == \length(A->data);
//@requires is_sorted(A->data, 0, A->length);
//@ensures \result >= 0;
//@ensures \result == \length(A->data);
{
    return A->length;
}
```

```
string ssa_get(ssa* A, int i)
//@requires A != NULL;
//@requires A->length == \length(A->data);
//@requires is_sorted(A->data, 0, A->length);
//@requires 0 <= i && i < ssa_len(A);
{
    return A->data[i];
}
```

SSA Implementation

● Representation invariants

- Preconditions of every library function that takes an SSA as a parameter
- Postcondition of every library function that returns or modifies an SSA
- Also called **data structure invariants**

● Do **ssa*** have other representation invariants?

- Yes! A->data should be sorted
`is_sorted(A->data, 0, A->length)`

```
int ssa_len(ssa_t A)
/*@requires A != NULL;
/*@ensures \result >= 0;
@*/;
@*/;

string ssa_get(ssa_t A, int i)
/*@requires A != NULL;
/*@requires 0 <= i && i < ssa_len(A); @*/
@*/;
```

Representation Invariants

```
struct ssa_header {
    string[] data; // sorted
    int length;   // = \length(data)
};

typedef struct ssa_header ssa;

bool is_ssa(ssa* A) { ... }

int ssa_len(ssa* A)
// @requires is_ssa(A);
// @ensures \result >= 0;
// @ensures \result == \length(A->data);
{
    return A->length;
}

string ssa_get(ssa* A, int i)
// @requires is_ssa(A);
// @requires 0 <= i && i < ssa_len(A);
{
    return A->data[i];
}

// ... rest of implementation
```

SSA Implementation

SSA Interface

```
int ssa_len(ssa_t A)
/* @requires A != NULL;
/* @ensures \result >= 0;
@*/;
@*/;

string ssa_get(ssa_t A, int i)
/* @requires A != NULL;
@*/
/* @requires 0 <= i && i < ssa_len(A); @*/;
```

- Factor them out into a single function that checks that they are satisfied

is_ssa

- **Representation invariant function**

Convention: name of representation invariant functions start with **is_**

Representation Invariants

```
struct ssa_header {  
    string[] data; // sorted  
    int length;   // = \length(data)  
};  
typedef struct ssa_header ssa;  
  
bool is_ssa(ssa* A) { ... }
```

```
ssa* ssa_new(int size)  
//@requires size >= 0;  
//@ensures is_ssa(\result);  
//@ensures ssa_len(\result) == size;  
{  
    ssa* A = alloc(ssa);  
    A->data = alloc_array(string, size);  
    A->length = size;  
    return A;  
}
```

```
void ssa_set(ssa* A, int i, string x)  
//@requires is_ssa(A);  
//@requires 0 <= i && i < ssa_len(A);  
//@ensures is_ssa(A);  
{ /* left as exercise */ }
```

SSA Implementation

SSA Interface

```
ssa_t ssa_new(int size)  
/*@requires 0 <= size; */  
/*@ensures \result != NULL; */  
/*@ensures ssa_len(\result) == size; */  
  
void ssa_set(ssa_t A, int i, string x)  
/*@requires A != NULL; */  
/*@requires 0 <= i && i < ssa_len(A); */
```

- The remaining functions

- Precondition of every `ssa*` parameter

Defining the internal type `ssa` **not** to be a pointer allows simpler allocations

- Postcondition of

- every returned `ssa*`
- every modified `ssa*` parameter

The representation Invariant Function

```
struct ssa_header {  
    string[] data; // sorted  
    int length;   // = \length(data)  
};  
typedef struct ssa_header ssa;  
  
// Representation invariant  
bool is_ssa (ssa* A) {  
    return A != NULL  
    && A->length == \length(A->data)  
    && is_sorted(A->data, 0, A->length);  
}  
  
// ... rest of implementation
```

SSA Implementation

Let's write it!

- 1st attempt: simply copy the contracts it stands for
 - **Problem:** \length can only be used in contracts

X

The representation Invariant Function

```
struct ssa_header {  
    string[] data; // sorted  
    int length;   // = \length(data)  
};  
typedef struct ssa_header ssa;  
  
// Representation invariant  
bool is_ssa (ssa* A)  
//@requires A->length == \length(A->data);  
{  
    return A != NULL  
        && is_sorted(A->data, 0, A->length);  
}  
  
// ... rest of implementation
```

SSA Implementation

Let's write it!

- 2nd attempt: move that part in the precondition of **is_ssa**
 - **Problem:** this is unsafe!
 - A may be NULL
 - NULL checked only *after* the precondition

✗

The representation Invariant Function

```
struct ssa_header {  
    string[] data; // sorted  
    int length;   // = \length(data)  
};  
typedef struct ssa_header ssa;  
  
// Representation invariant  
bool is_array_expected_length(string[] A, int len) {  
    //@assert \length(A) == len;  
    return true;  
}  
  
bool is_ssa(ssa* A) {  
    return A != NULL  
        && is_array_expected_length(A->data, A->length)  
        && is_sorted(A->data, 0, A->length);  
}  
  
// ... rest of implementation
```

SSA Implementation

Let's write it!

- 3rd attempt: move it into a helper function



The representation Invariant Function

```
struct ssa_header {  
    string[] data; // sorted  
    int length;   // = \length(data)  
};  
typedef struct ssa_header ssa;  
  
// Representation invariant  
bool is_ssa (ssa* A) {  
    if (A == NULL) return false;  
    //@assert A->length == \length(A->data);  
    return is_sorted(A->data, 0, A->length);  
}  
  
// ... rest of implementation
```

SSA Implementation

Let's write it!

- Alternative 3rd attempt: turn it into an **//@assert** in **is_ssa**



Things to Note

- The representation invariant function `is_ssa` is **NOT** part of interface
 - Clients are allowed to manipulate SSA's **only using the interface** functions
 - If the library is correct, `is_ssa(A)` will always return true
 - Client must ensure the safety of library calls according to the interface
 - `A != NULL` only
 - Providing `is_ssa` to clients would encourage them to bypass the interface
 - use `is_ssa` to test if hacks are successful
 - The representation invariant function is an implementation device to ensure the safety and correctness of the library code
 - Used while developing the library
 - Every function that takes an SSA A must have `//@requires is_ssa(A);`
 - Every function that modifies an input SSA A must have `//@ensures is_ssa(A);`
 - Every function that returns an SSA must have `//@ensures is_ssa(\result);`

Things to Note

- The contracts in the interface and the implementation are **different**

Interface	Implementation
<pre>void ssa_set(ssa_t A, int i, string x) /*@requires A != NULL; @*/ /*@requires 0 <= i && i < ssa_len(A); @*/;</pre>	<pre>void ssa_set(ssa* A, int i , string x) //@requires is_ssa(A); //@requires 0 <= i && i < ssa_len(A); //@ensures is_ssa(A);</pre>

- The implementation contracts are more detailed
 - `is_ssa(A)` checks `A != NULL`
 - The implementation contains more information, so it needs to check more things
 - There is no point having `//@ensures A != NULL` in the interface
 - `ssa_set` is called with a *copy* of the address of `A`
 - when returning, the original has not changed, even if `ssa_set` modified its copy of `A`
 - If original `A` was not `NULL` when calling `ssa_set`, it will not be `NULL` when returning from it

Overall Implementation

```
// Implementation-side type
struct ssa_header {
    string[] data; // sorted
    int length; // = \length(data)
};
```

```
// Representation invariant
bool is_sorted(string* A, int lo, int hi) {
    /* left as exercise */
}
```

```
bool is_ssa(ssa* A) {
    if (A == NULL) return false;
    // @assert A->length == \length(A->data);
    return is_sorted(A->data, 0, A->length);
}
```

```
// Implementation of interface functions
int ssa_len(ssa* A)
// @requires is_ssa(A);
// @ensures \result >= 0;
// @ensures \result == \length(A->data);
{
    return A->length;
}
```

```
ssa* ssa_new(int size)
// @requires size >= 0;
// @ensures is_ssa(\result);
// @ensures ssa_len(\result) == size;
{
    ssa* A = alloc(ssa);
    A->data = alloc_array(string, size);
    A->length = size;
    return A;
}

string ssa_get(ssa* A, int i)
// @requires is_ssa(A);
// @requires 0 <= i && i < ssa_len(A);
{
    return A->data[i];
}

void ssa_set(ssa* A, int i, string x)
// @requires is_ssa(A);
// @requires 0 <= i && i < ssa_len(A);
// @ensures is_ssa(A);
{ /* left as exercise */ }

// Client type
typedef ssa* ssa_t;
```

SSA Implementation

- By **convention**, we put the interface **after** the implementation in the same file

SSA Interface

```
// typedef _____ * ssa_t;

int ssa_len(ssa_t A)
/* @requires A != NULL; */ // @*
/* @ensures \result >= 0; */ // @*

sssa_t ssa_new(int size)
/* @requires 0 <= size; */ // @*
/* @ensures \result != NULL; */ // @*
/* @ensures ssa_len(\result) == size; */ // @*

string ssa_get(ssa_t A, int i)
/* @requires A != NULL; */ // @*
/* @requires 0 <= i && i < ssa_len(A); */ // @*

void ssa_set(ssa_t A, int i, string x)
/* @requires A != NULL; */ // @*
/* @requires 0 <= i && i < ssa_len(A); */ // @*
```

How

What

Structure of a C0 Library File

```
***** IMPLEMENTATION *****/
// Implementation-side type
struct ssa_header {
    ...
};

typedef struct ssa_header ssa;

// Representation invariant
bool is_ssa(ssa* A) {
    ...
}

// Implementation of interface functions
int ssa_len(ssa* A) { ... }

...

// Client type
typedef ssa* ssa_t;

***** LIBRARY INTERFACE *****/
// typedef _____ * ssa_t;

int ssa_len(ssa_t A)
/*@requires A != NULL;
/*@ensures \result >= 0;
    @*/
    @*/;
```

Implementation

- Implementation
 - Concrete type definition
 - Representation invariant function
 - Implementation of interface functions
 - Client type definition

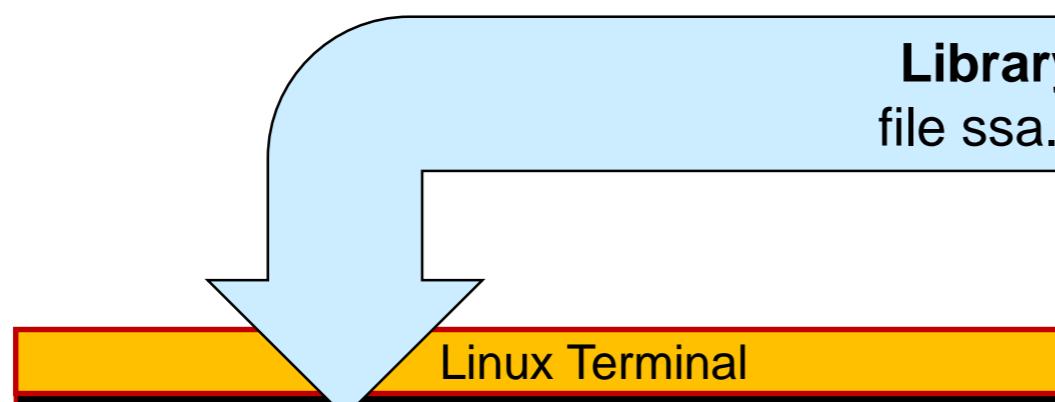
Interface

- Interface
 - Abstract type name
 - Pseudo-**typedef**
 - Prototype of exported functions

We will revisit this

Compiling a Library in a C0 Application

- Library file contains both implementation and interface
- When compiling,
library files come **before** application files
 - The application uses library interface types and functions
 - They need to be defined first
➤ this happens in the library



The diagram shows the "SSA Interface" and "SSA Implementation" side-by-side. The "SSA Interface" section contains declarations for `ssa_t`, `int ssa_len(ssa_t A)`, `string ssa_get(ssa_t A, int i)`, `void ssa_set(ssa_t A, int i, string x)`, and a `typedef` for `ssa_t`. The "SSA Implementation" section contains the actual implementations for these functions, including memory allocation and deallocation, and string manipulation. The code is annotated with various annotations such as `@requires`, `@ensures`, and `/* ... */`.

```
SSA Interface
// typedef _____ * ssa_t;
int ssa_len(ssa_t A)
/*@requires A != NULL;
/*@ensures \result >= 0;
ssat ssa_new(int size)
/*@requires 0 <= size;
/*@ensures \result != NULL;
/*@ensures ssa_len(result) == size;
string ssa_get(ssa_t A, int i)
/*@requires A != NULL;
/*@requires 0 <= i && i < ssa_len(A);
void ssa_set(ssa_t A, int i, string x)
/*@requires A != NULL;
/*@requires 0 <= i && i < ssa_len(A);
/*@ensures \result >= 0;
/*@ensures \result == \length(A->data);
{ return A->length;
}

SSA Implementation
a* ssa_new(int size)
@requires size >= 0;
@ensures is_ssa(result);
@ensures ssa_len(result) == size;
sa* A = alloc(ssa);
A->data = alloc_array(string, size);
A->length = size;
turn A;

string ssa_get(ssa* A, int i)
@requires is_ssa(A);
@requires 0 <= i && i < ssa_len(A);
return A->data[i];

void ssa_set(ssa* A, int i, string x)
@requires is_ssa(A);
@requires 0 <= i && i < ssa_len(A);
@ensures is_ssa(result);
@ensures \result == \length(A->data);
{ /* left as exercise */
}

// Client type
typedef ssa* ssa_t;
```

The diagram shows the "Application file main.c0" containing the code for the `main` function. It includes calls to `ssa_new`, `ssa_set`, and `is_in`.

```
bool is_in(string x, ssa_t A, int n)
//@requires n == ssa_len(A);
{ ... ssa_get(A, mid) ... }

int main() {
    ssa_t A = ssa_new(42);
    ... ssa_set(A, 7, "hello") ...
    ... is_in("hello", A, 42) ...
    return 0;
}
```