

# Kidney exchanges

(largely follows Abraham, Blum, Sandholm 2007 paper)

Vincent Conitzer

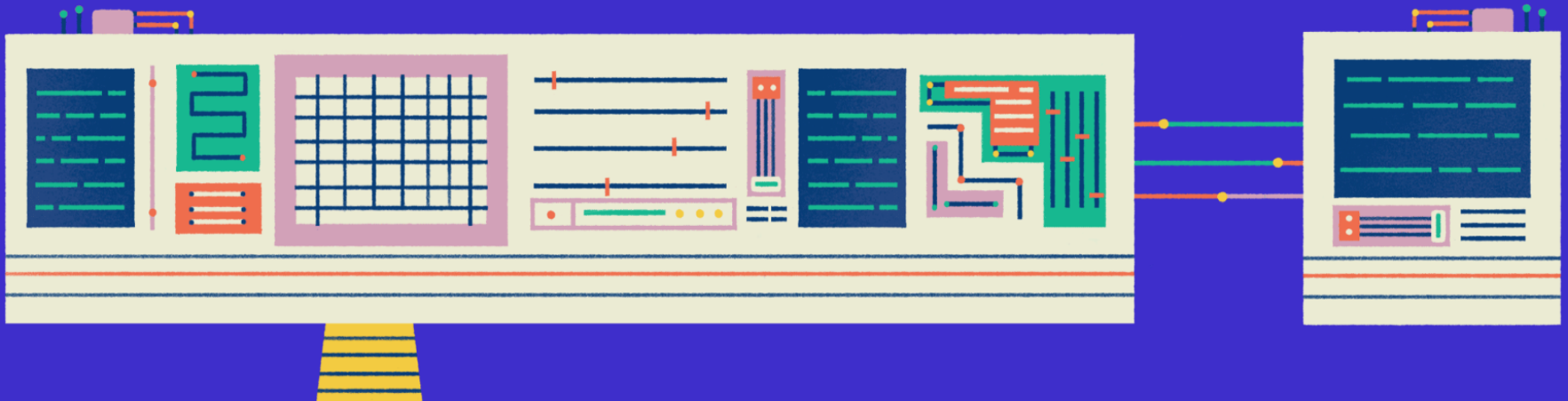
## Prescription AI

This series explores the promise of AI to personalize, democratize, and advance medicine—and the dangers of letting machines make decisions.

THE BOTPERATING TABLE

# How AI changed organ donation in the US

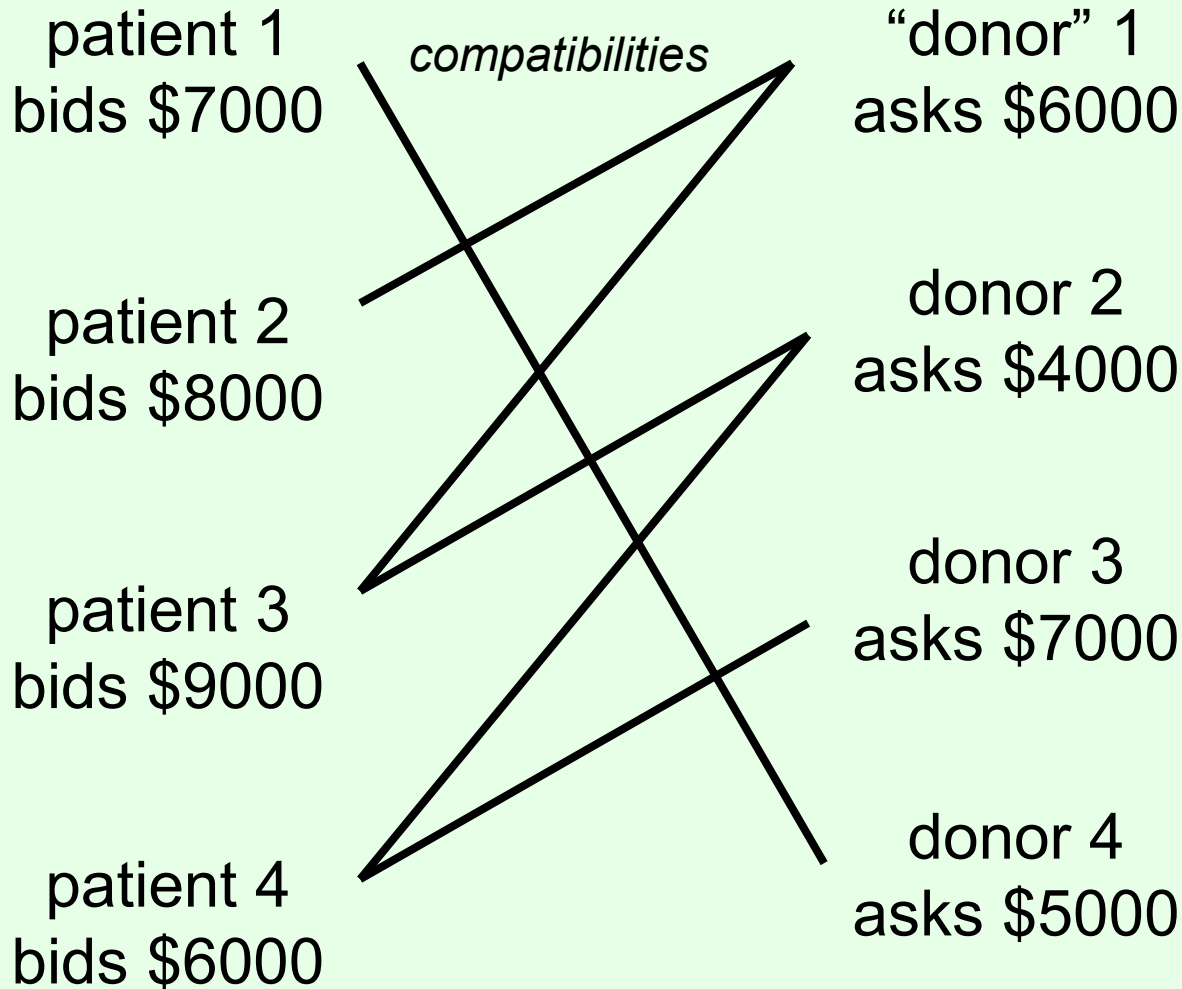
By Corinne Purtill • September 10, 2018



# Kidney transplants

- **Kidneys** filter waste from blood
- Kidney failure results in death in months
- **Dialysis**: regularly get blood filtered in hospital using external machine
  - Low quality of life
- Preferred option: kidney transplant
  - Cadaver kidneys
  - Donation from live person (better)
- Must be compatible
- Shortage of kidneys...

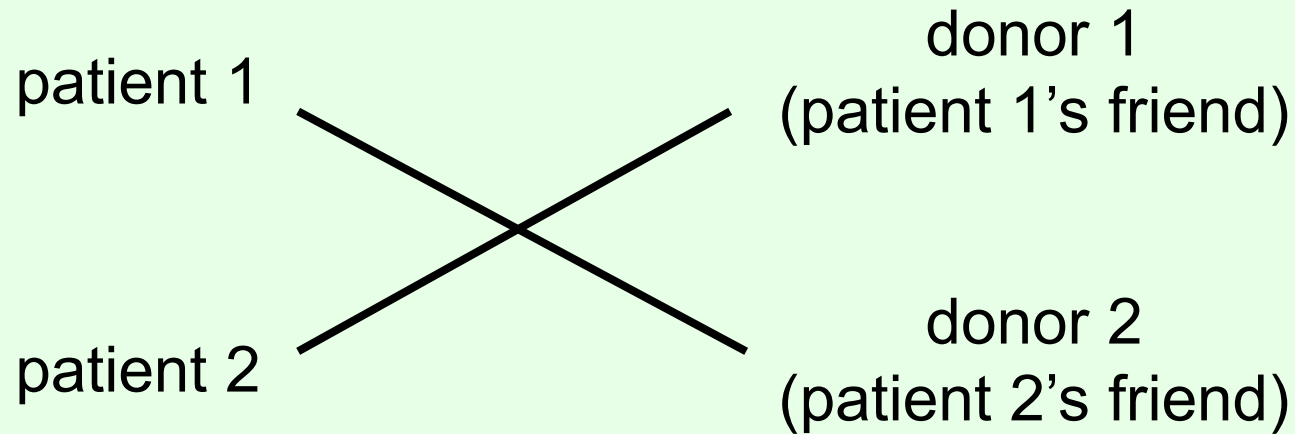
# An imaginary kidney exchange with money



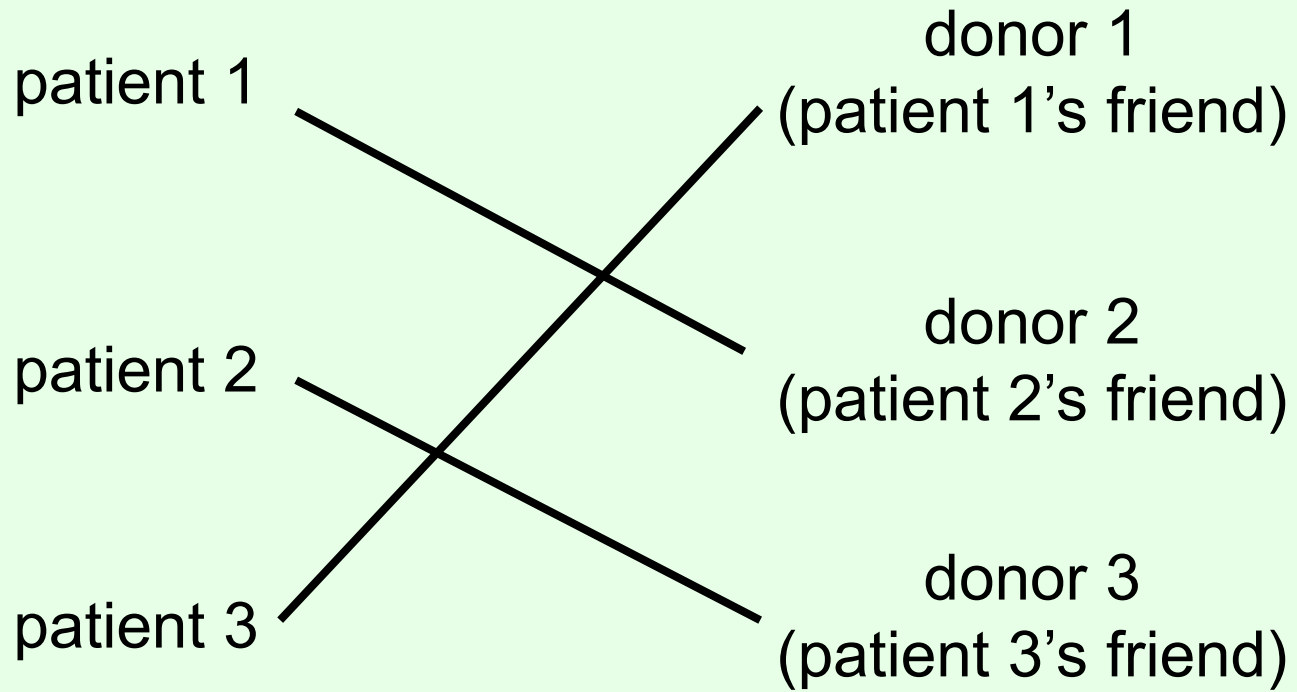
# Selling kidneys is illegal!

- Large international black market
  - Desperate people on both ends...
- What can we do legally?

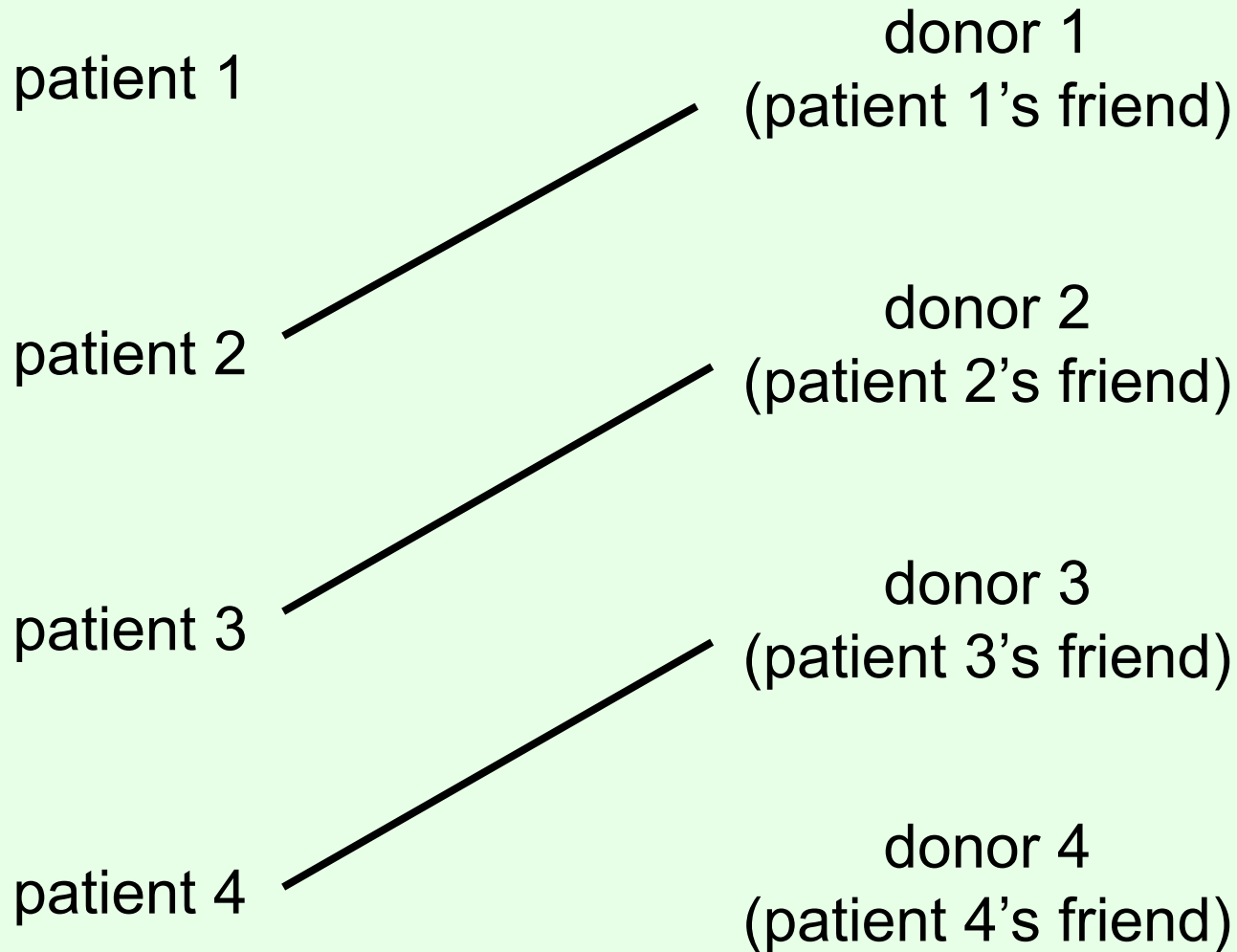
# Kidney exchange



# Kidney exchange (3-cycle)

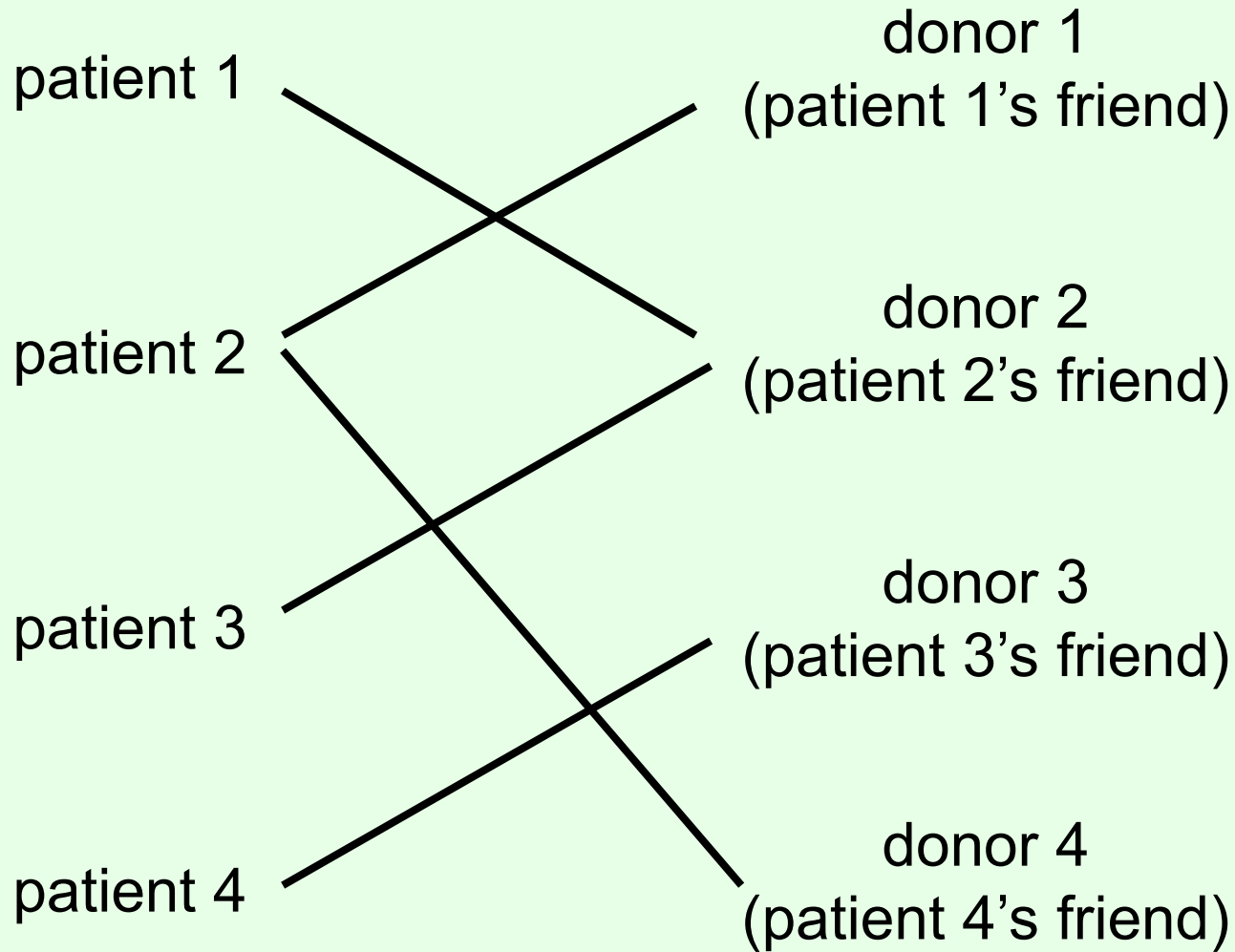


# Another example

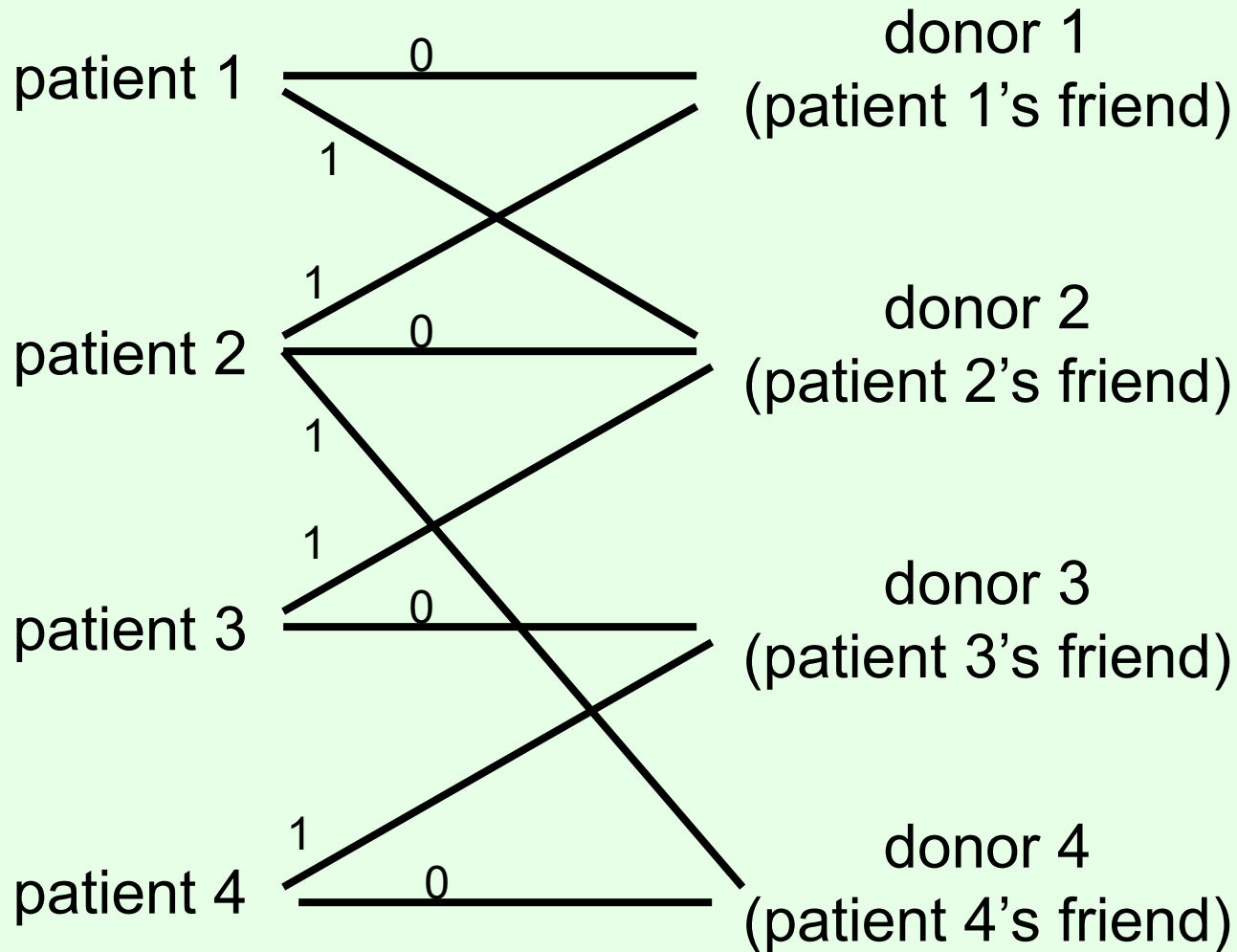




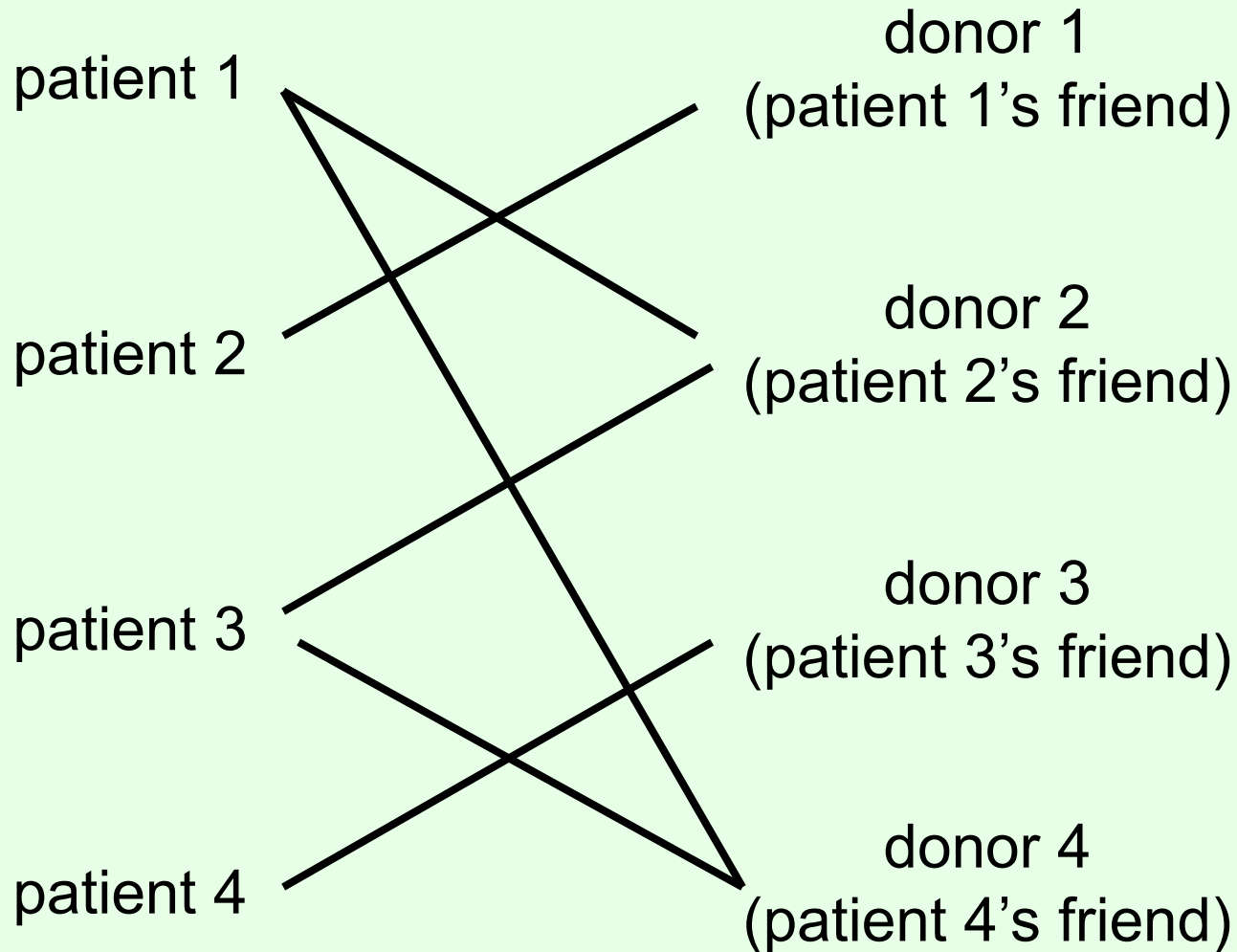
# More complex example



# Solving kidney exchange as maximum weighted bipartite matching



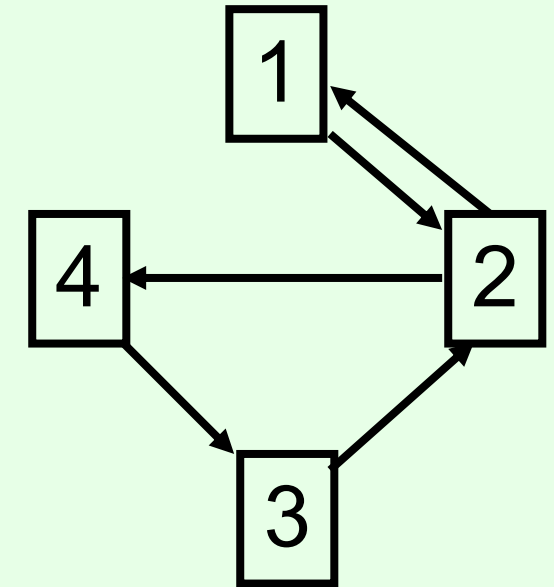
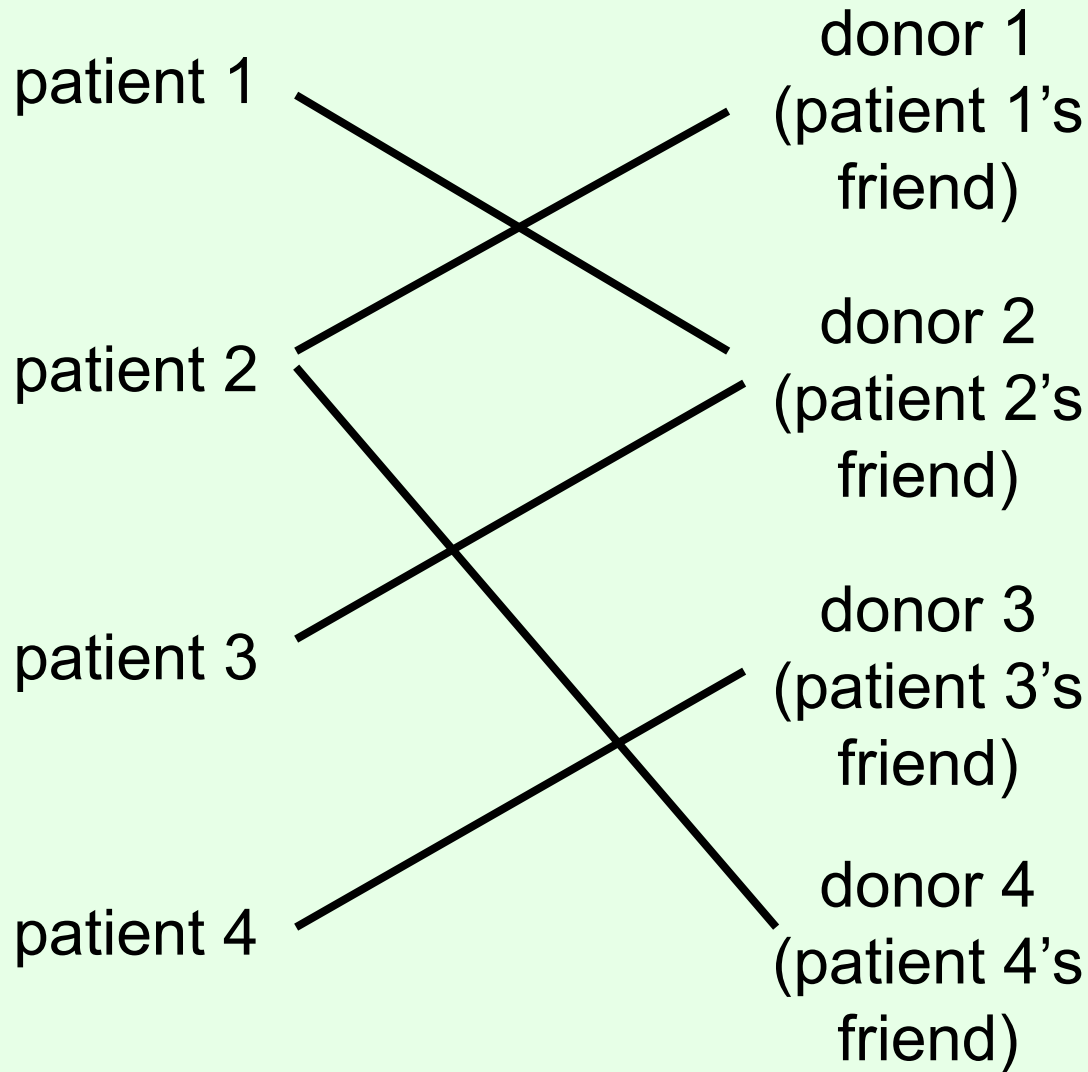
# Which solution is better?



# Long cycles are impractical

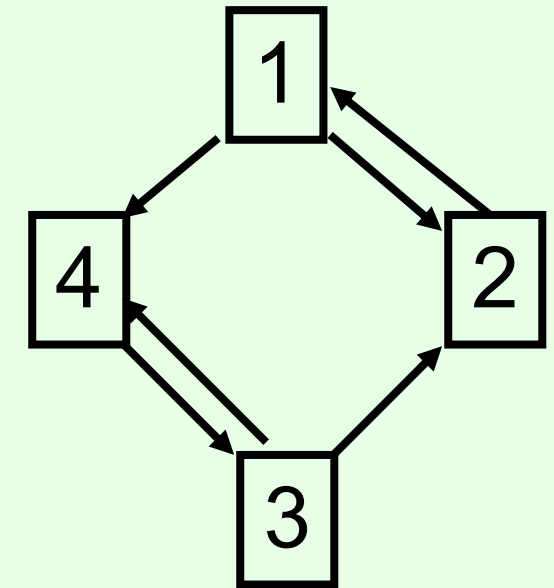
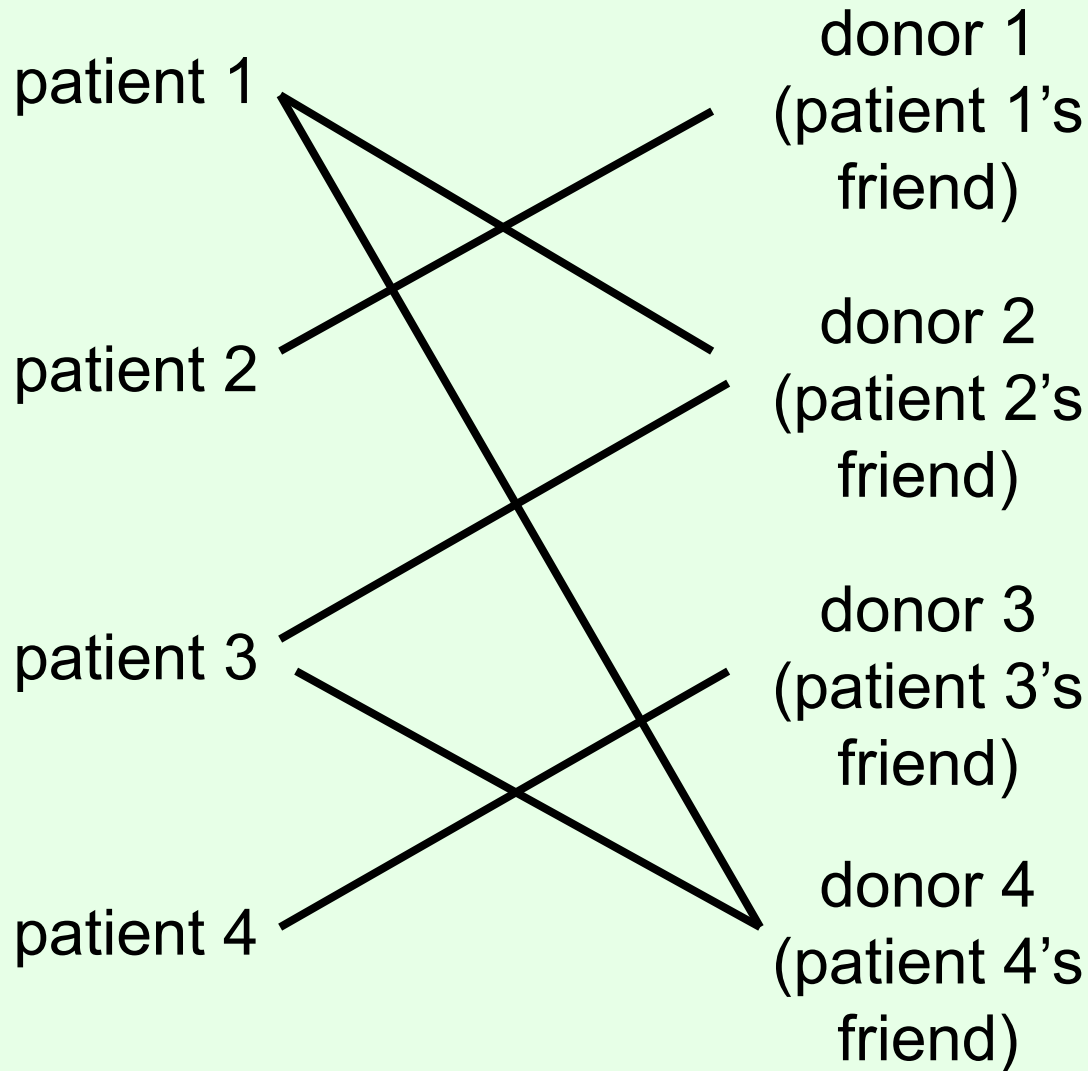
- All patients in a cycle must be operated on simultaneously
  - Otherwise donor can wait for friend to receive kidney, then back out
  - Contracts to donate an organ not binding
- If last-minute test reveals incompatibility, whole thing falls apart
- Require each cycle has length at most  $k$

# Different representation



edge from  $i$  to  $j$  =  
patient  $i$  wants  
donor  $j$ 's kidney

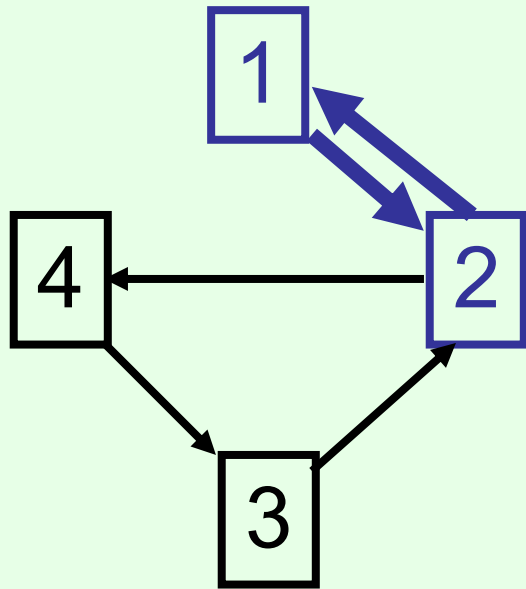
# Different representation



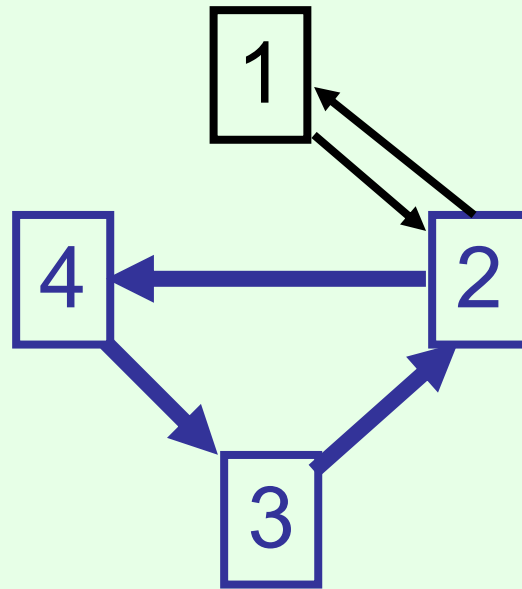
edge from  $i$  to  $j$  =  
patient  $i$  wants  
donor  $j$ 's kidney

# Market clearing problem

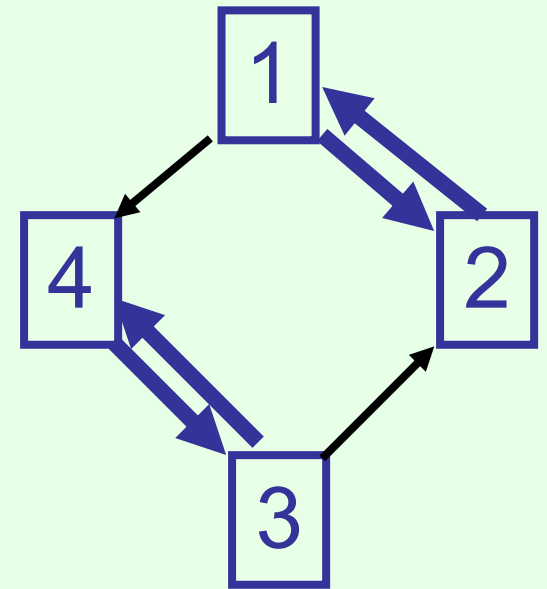
- Try to cover as many vertices as possible with (vertex-)disjoint cycles of length at most  $k$



$k=2$



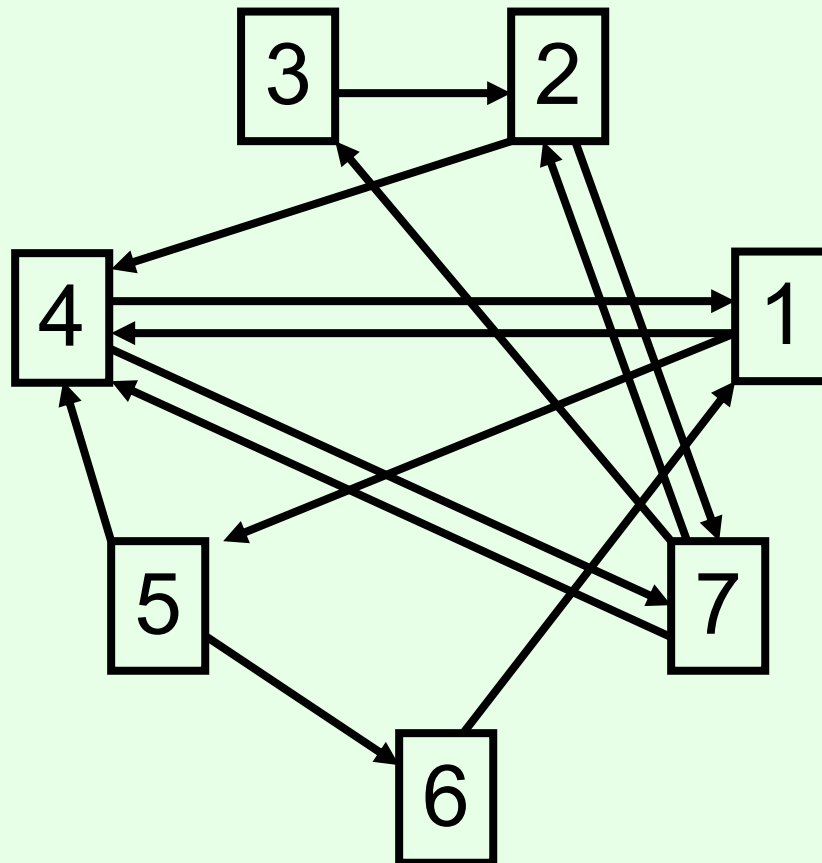
$k=3$



$k=2,3$

# Market clearing problem

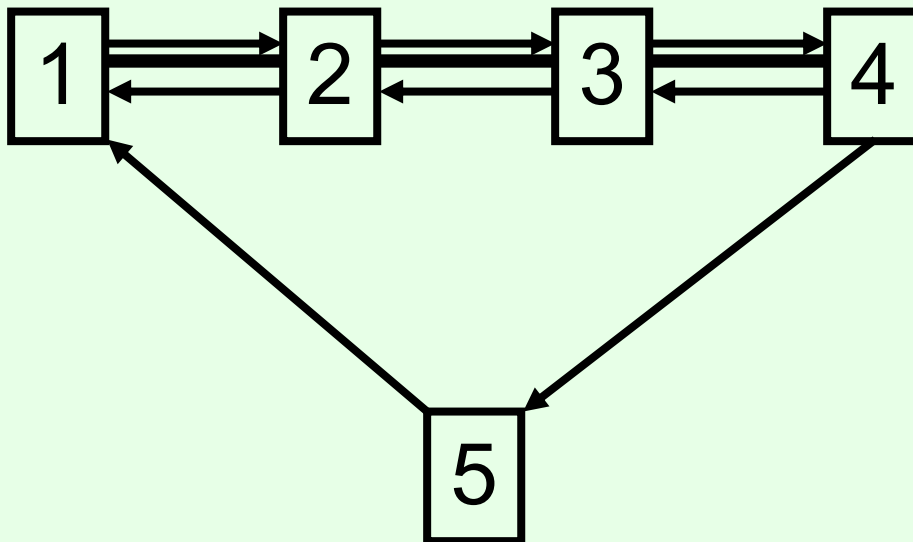
- Try to cover as many vertices as possible with (vertex-)disjoint cycles of length at most  $k$





# Special case: $k=2$

- If edges go in both directions, replace by undirected edge
- Remove other edges



- Maximum matching problem!

# Complexity

- $k = 2$ : in  $P$  by maximum matching
- $k = \text{number of vertices (no constraint)}$ : in  $P$  by maximum weighted bipartite matching
- $k = 3, 4, 5, \dots$ : NP-hard!

# An integer programming formulation

- For each edge from  $i$  to  $j$ , make a binary variable  $x_{ij}$ 
  - 1 if  $i$  gets  $j$ 's kidney, 0 otherwise
- maximize  $\sum_{ij} x_{ij}$
- subject to:
- for every  $i$ :  $\sum_j x_{ij} = \sum_j x_{ji}$ 
  - (number of kidneys received by  $i$  = number of kidneys given by  $i$ )
- for every  $j$ :  $\sum_i x_{ij} \leq 1$ 
  - ( $j$  gives at most 1 kidney)
- for every path  $i_1 i_2 \dots i_k i_{k+1}$  with  $i_1 \neq i_{k+1}$ :  $\sum_{1 \leq j \leq k} x_{i_j i_{j+1}} \leq k-1$ 
  - (no path of length  $k$  that doesn't end up where it started, hence no cycles greater than  $k$ )

# Another integer programming formulation

(turns out better)

- For each cycle  $c$  of length at most  $k$ , make a binary variable  $x_c$ 
  - 1 if all edges on this cycle are used, 0 otherwise
- maximize  $\sum_c |c| x_c$
- subject to:
- for every vertex  $i$ :  $\sum_{c: i \text{ in } c} x_c \leq 1$ 
  - (every vertex in at most one used cycle)

# Program size

- Even for small  $k$ , number of paths/cycles is too large in reasonably large exchanges
- Solution: generate constraints/variables on the fly during solving
  - Constraint/column generation

# Another integer program (not in paper)

- Say an “event” is a set of simultaneous operations
- Denote events by  $t = 1, \dots, T$  (how big should  $T$  be?)
- For each edge from  $i$  to  $j$ , for each  $t$ , make a binary variable  $x_{ijt}$ 
  - 1 if  $i$  gets  $j$ 's kidney in event  $t$ , 0 otherwise
- maximize  $\sum_{i,j,t} x_{ijt}$
- subject to:
- for every  $i$ ,  $t$ :  $\sum_j x_{ijt} = \sum_j x_{jit}$ 
  - (number of kidneys received by  $i$  in event  $t$  = number of kidneys given by  $i$  in event  $t$ )
- for every  $j$ :  $\sum_{i,t} x_{ijt} \leq 1$ 
  - ( $j$  gives at most 1 kidney overall)
- for every  $t$ :  $\sum_{i,j} x_{ijt} \leq k$ 
  - (at most  $k$  operations per event)

# Other applications

- Barter exchanges: agents want to swap items without paying money
- Peerflix (DVDs)
- Read It Swap It (books)
- Intervac (holiday houses)
- National odd shoe exchange
  - People with different foot sizes
  - Amputees

# Modeling

- What assumptions have we implicitly made in modeling a kidney exchange?
- What problems might come up that we haven't thought about?
- What additional aspects could one model to get even better results?