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### Linear programming, integer linear programming, mixed integer linear programming

# Example linear program We make reproductions of two paintings



maximize 3x + 2ysubject to  $4x + 2y \leq 16$  $x + 2y \le 8$  $x + y \leq 5$  $x \ge 0$  $y \ge 0$ 

- Painting 1 sells for \$30, painting 2 sells for \$20
- Painting 1 requires 4 units of blue, 1 green, 1 red
- Painting 2 requires 2 blue, 2 green, 1 red
- We have 16 units blue, 8 green, 5 red

#### Solving the linear program graphically



#### **Proving optimality**

maximize 3x + 2ysubject to  $4x + 2y \leq 16$  $x + 2y \le 8$  $x + y \leq 5$  $x \ge 0$  $y \ge 0$ 

Recall: optimal solution: x=3, y=2Solution value = 9+4 = 13

How do we prove this is optimal (without the picture)?

#### Proving optimality...

maximize 3x + 2ysubject to  $4x + 2y \leq 16$  $x + 2y \le 8$  $x + y \leq 5$  $x \ge 0$  $y \ge 0$ 

We can rewrite the blue constraint as  $2x + y \leq 8$ If we add the red constraint  $x + y \leq 5$ we get  $3x + 2y \leq 13$ Matching upper bound! (Really, we added .5 times the blue constraint to 1 times the red constraint)

#### Linear combinations of constraints $b(4x + 2y \le 16) +$ maximize 3x + 2y $g(x + 2y \le 8) +$ $r(x + y \le 5) =$ subject to (4b + g + r)x + $4x + 2y \leq 16$ $(2b + 2g + r)y \leq$ $x + 2y \le 8$ 16b + 8g + 5r4b + g + r must $x + y \leq 5$ be at least 3 $x \ge 0$ 2b + 2g + r must $y \ge 0$ be at least 2 Given this, minimize 16b +

8g + 5r

Using LP for getting the best upper bound on an LP maximize 3x + 2y *minimize* 16b + 8g + 5r subject to subject to  $4x + 2y \leq 16$  $4b + g + r \ge 3$  $x + 2y \le 8$  $2b + 2g + r \ge 2$  $x + y \le 5$  $b \ge 0$  $x \ge 0$ g ≥ 0  $y \ge 0$  $r \ge 0$ 

the dual of the original program

Duality theorem: any linear program has the same optimal value as its dual!

#### Modified LP

maximize 3x + 2ysubject to  $4x + 2y \le 15$  $x + 2y \le 8$  $x + y \leq 5$  $x \ge 0$  $y \ge 0$ 

Optimal solution: x = 2.5, y = 2.5Solution value = 7.5 + 5 = 12.5

Half paintings?

#### Integer (linear) program



#### Mixed integer (linear) program



#### Solving linear/integer programs

- Linear programs can be solved efficiently
  - Simplex, ellipsoid, interior point methods...
- (Mixed) integer programs are NP-hard to solve
  - Quite easy to model many standard NP-complete problems as integer programs (try it!)
  - Search type algorithms such as branch and bound
- Standard packages for solving these

   GNU Linear Programming Kit, CPLEX, Gurobi, …
- LP relaxation of (M)IP: remove integrality constraints
  - Gives upper bound on MIP (~admissible heuristic)

#### Exercise in modeling: knapsack-type problem

- We arrive in a room full of precious objects
- Can carry only 30kg out of the room
- Can carry only 20 liters out of the room
- Want to maximize our total value
- Unit of object A: 16kg, 3 liters, sells for \$11
   There are 3 units available
- Unit of object B: 4kg, 4 liters, sells for \$4
   There are 4 units available
- Unit of object C: 6kg, 3 liters, sells for \$9

   Only 1 unit available
- What should we take?

## The *general* version of this knapsack IP

maximize  $\sum_{j} p_{j} x_{j}$ subject to  $\sum_{j} w_{j} x_{j} \leq W$  $\sum_{j} v_{j} x_{j} \leq V$ (for all *j*)  $x_{j} \leq a_{j}$ (for all *j*)  $x_{j} \geq 0$ ,  $x_{j}$  integer Exercise in modeling: cell phones (set cover)

- We want to have a working phone in every continent (besides Antarctica)
- ... but we want to have as few phones as possible
- Phone A works in NA, SA, Af
- Phone B works in E, Af, As
- Phone C works in NA, Au, E
- Phone D works in SA, As, E
- Phone E works in Af, As, Au
- Phone F works in NA, E

#### Exercise in modeling: hot-dog stands

- We have two hot-dog stands to be placed in somewhere along the beach
- We know where the people that like hot dogs are, how far they are willing to walk
- Where do we put our stands to maximize #hot dogs sold? (price is fixed)

