


Automated mechanism design

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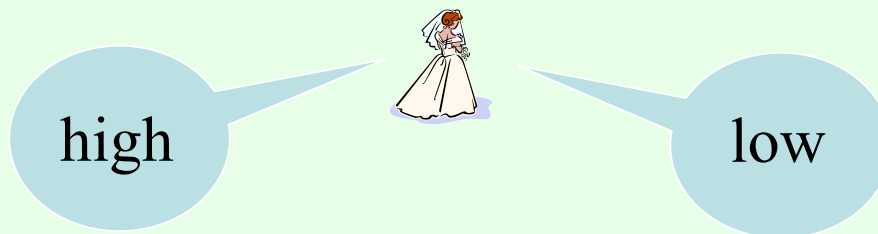
General vs. specific mechanisms





- Mechanisms such as Clarke (VCG) mechanism are very **general**...
- ... but will instantiate to something **specific** in any specific setting
 - This is what we care about

Example: Divorce arbitration

- Outcomes: 
- Each agent is of *high* type w.p. .2 and *low* type w.p. .8
 - Preferences of *high* type:
 - $u(\text{get the painting}) = 11,000$
 - $u(\text{museum}) = 6,000$
 - $u(\text{other gets the painting}) = 1,000$
 - $u(\text{burn}) = 0$
 - Preferences of *low* type:
 - $u(\text{get the painting}) = 1,200$
 - $u(\text{museum}) = 1,100$
 - $u(\text{other gets the painting}) = 1,000$
 - $u(\text{burn}) = 0$

Clarke (VCG) mechanism



 Both pay 5,000	 Husband pays 200
 Wife pays 200	 Both pay 100

Expected sum of divorcees' utilities = 5,136

“Manual” mechanism design has yielded

- some **positive results**:
 - “Mechanism x achieves properties P in any setting that belongs to class C ”
- some **impossibility results**:
 - “There is no mechanism that achieves properties P for all settings in class C ”

Difficulties with manual mechanism design

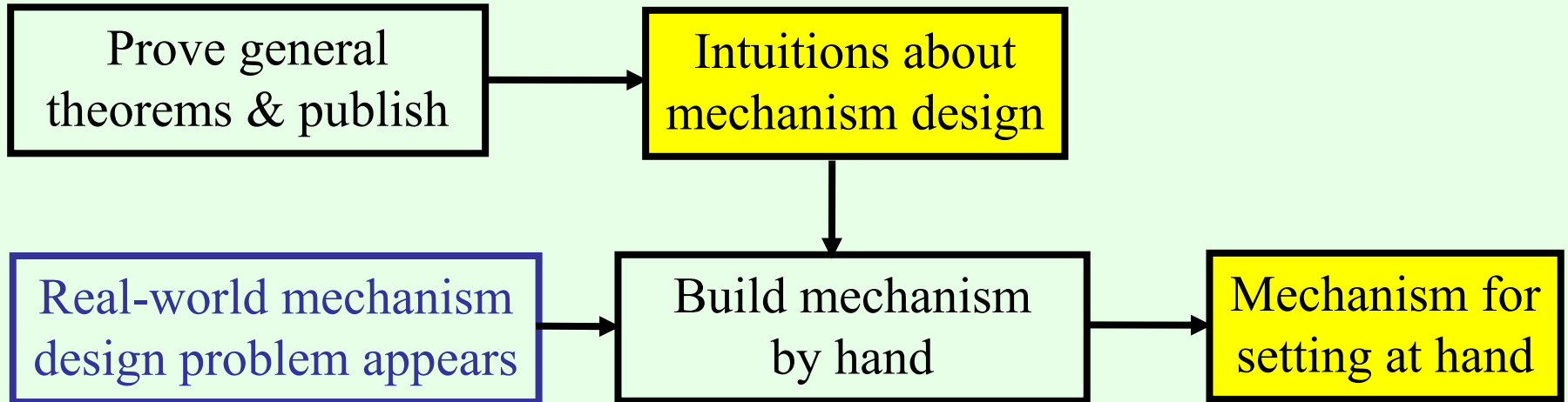
- Design problem instance comes along
 - Set of outcomes, agents, set of possible types for each agent, prior over types, ...
- What if **no** canonical mechanism covers this instance?
 - Unusual objective, or payments not possible, or ...
 - Impossibility results may exist for the general class of settings
 - But instance may have additional structure (restricted preferences or prior) so good mechanisms exist (but unknown)
- What if a canonical mechanism **does** cover the setting?
 - Can we use instance's structure to get higher objective value?
 - Can we get stronger nonmanipulability/participation properties?
- Manual design for every instance is prohibitively slow

Automated mechanism design (AMD)

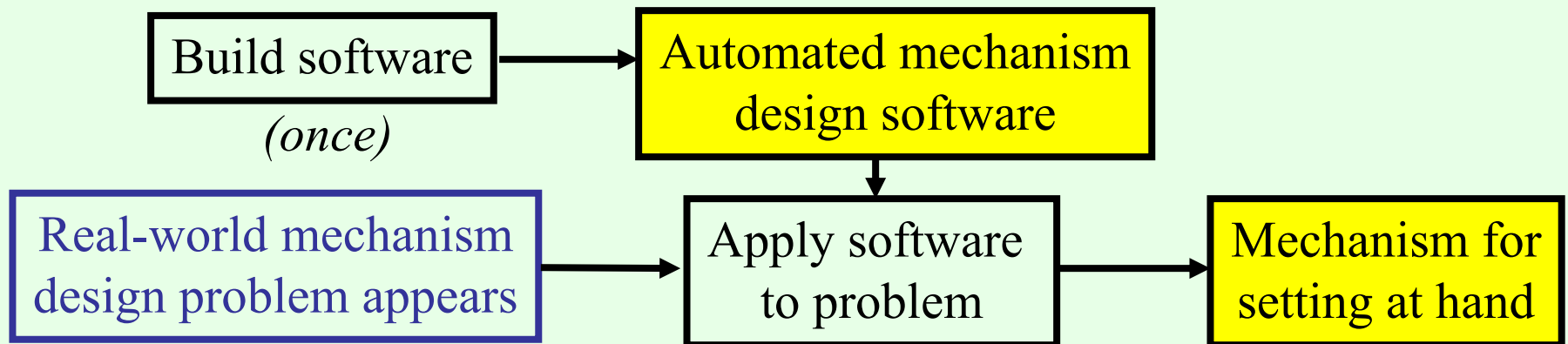
- Idea: Solve mechanism design **as optimization problem** automatically
- Create a mechanism **for the specific setting at hand** rather than a class of settings
- Advantages:
 - Can lead to greater value of designer's objective than known mechanisms
 - Sometimes circumvents economic impossibility results & always minimizes the pain implied by them
 - Can be used in new settings & for unusual objectives
 - Can yield stronger incentive compatibility & participation properties
 - Shifts the burden of design from human to machine

Classical vs. automated mechanism design

Classical



Automated



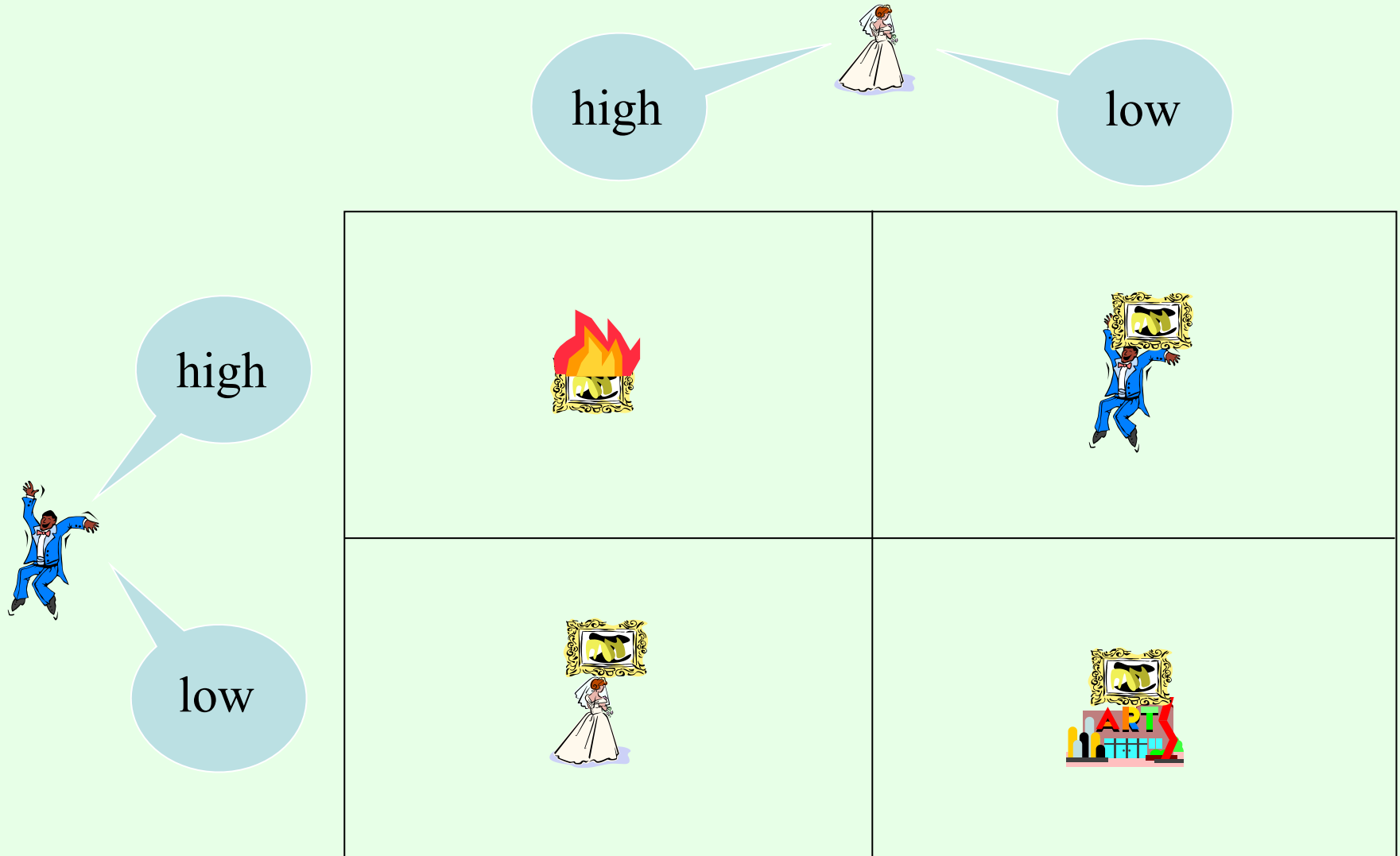
Input

- Instance is given by
 - Set of possible *outcomes*
 - Set of *agents*
 - For each agent
 - set of possible *types*
 - *probability distribution* over these types
 - *Objective function*
 - Gives a value for each outcome for each combination of agents' types
 - E.g. social welfare, payment maximization
 - Restrictions on the mechanism
 - Are payments allowed?
 - Is randomization over outcomes allowed?
 - What versions of incentive compatibility (IC) & individual rationality (IR) are used?

Output

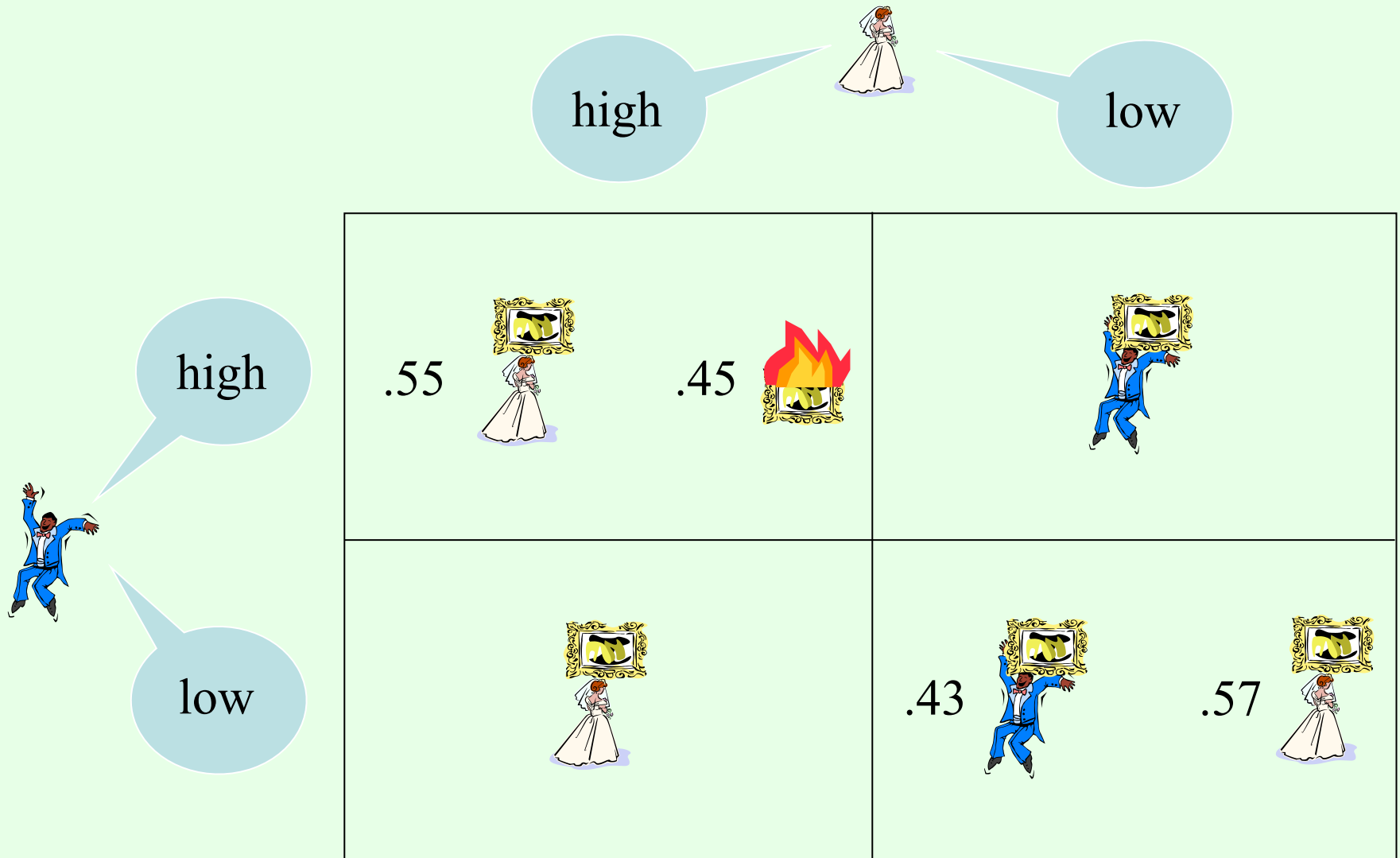
- *Mechanism*
 - A mechanism maps combinations of agents' revealed types to outcomes
 - *Randomized mechanism* maps to probability distributions over outcomes
 - Also specifies payments by agents (if payments allowed)
- ... which
 - satisfies the IR and IC constraints
 - maximizes the expectation of the objective function

Optimal BNE incentive compatible deterministic mechanism without payments for maximizing sum of divorcees' utilities



Expected sum of divorcees' utilities = 5,248

Optimal BNE incentive compatible *randomized* mechanism without payments for maximizing sum of divorcees' utilities



Expected sum of divorcees' utilities = 5,510

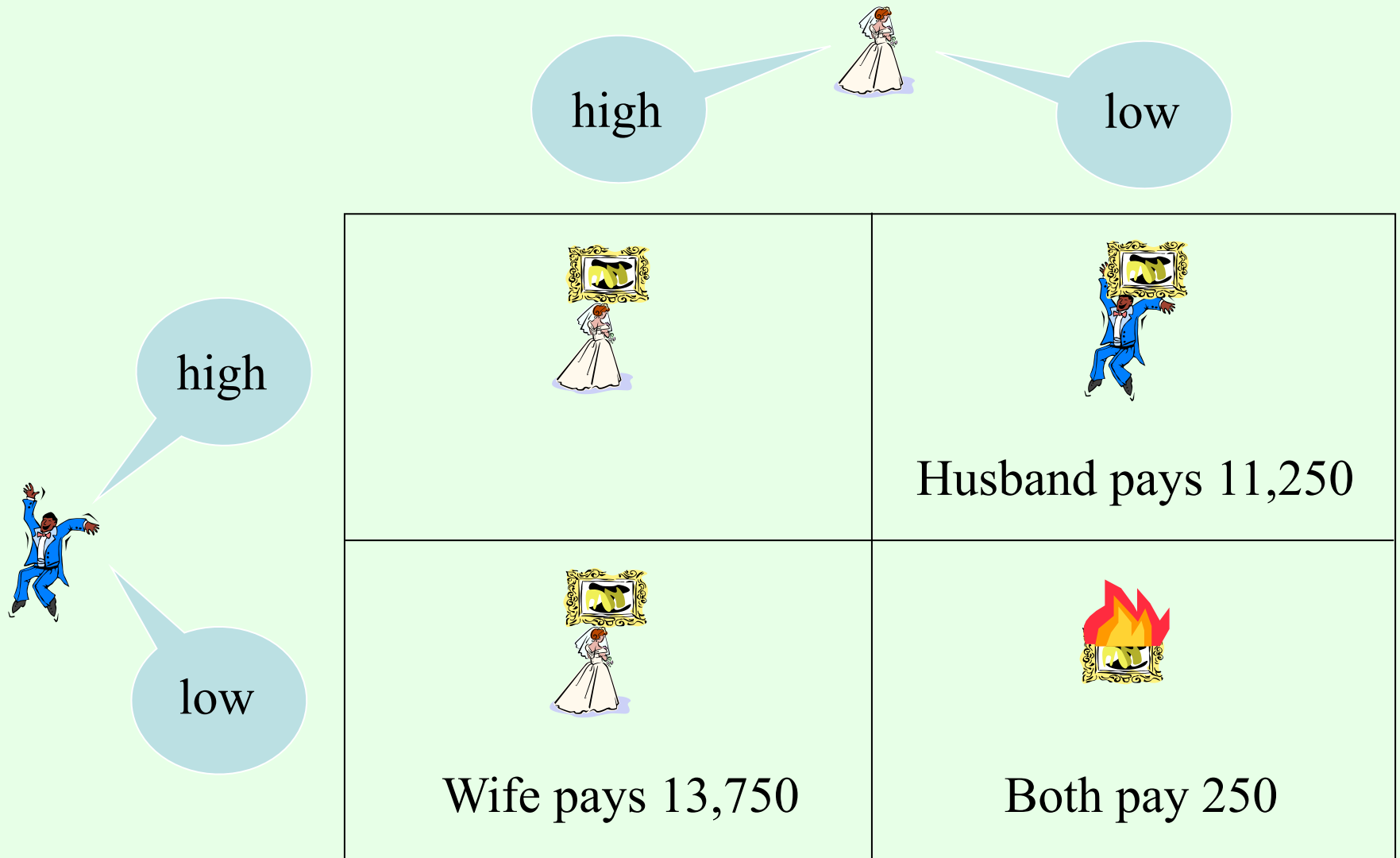
Optimal BNE incentive compatible randomized mechanism *with payments* for maximizing sum of divorcees' utilities



 Wife pays 1,000	

Expected sum of divorcees' utilities = 5,688

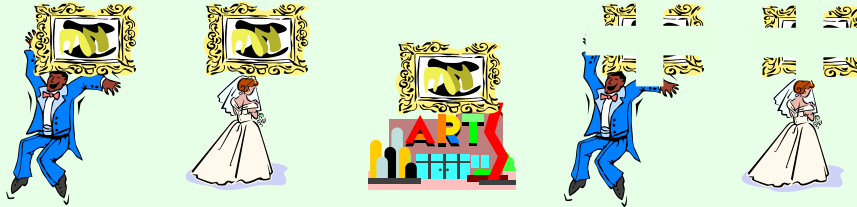
Optimal BNE incentive compatible randomized mechanism with payments for *maximizing arbitrator's revenue*



Expected sum of divorcees' utilities = 0

Arbitrator expects 4,320









Modified divorce arbitration example



- Outcomes:
- Each agent is of *high* type with probability 0.2 and of *low* type with probability 0.8
 - Preferences of *high* type:
 - $u(\text{get the painting}) = 100$
 - $u(\text{other gets the painting}) = 0$
 - $u(\text{museum}) = 40$
 - $u(\text{get the pieces}) = -9$
 - $u(\text{other gets the pieces}) = -10$
 - Preferences of *low* type:
 - $u(\text{get the painting}) = 2$
 - $u(\text{other gets the painting}) = 0$
 - $u(\text{museum}) = 1.5$
 - $u(\text{get the pieces}) = -9$
 - $u(\text{other gets the pieces}) = -10$

Optimal *dominant-strategies* incentive compatible randomized mechanism for maximizing expected sum of utilities



<p>high</p> <p>.47  .4  .13 </p>	<p>.96  .04 </p>
<p>.96  .04 </p>	<p></p>

How do we set up the optimization?

- Use linear programming
- Variables:
 - $p(o \mid \theta_1, \dots, \theta_n)$ = probability that outcome o is chosen given types $\theta_1, \dots, \theta_n$
 - (maybe) $\pi_i(\theta_1, \dots, \theta_n)$ = i 's payment given types $\theta_1, \dots, \theta_n$
- Strategy-proofness constraints: for all $i, \theta_1, \dots, \theta_n, \theta_i'$:
$$\sum_o p(o \mid \theta_1, \dots, \theta_n) u_i(\theta_i, o) + \pi_i(\theta_1, \dots, \theta_n) \geq$$
$$\sum_o p(o \mid \theta_1, \dots, \theta_i', \dots, \theta_n) u_i(\theta_i, o) + \pi_i(\theta_1, \dots, \theta_i', \dots, \theta_n)$$
- Individual-rationality constraints: for all $i, \theta_1, \dots, \theta_n$:
$$\sum_o p(o \mid \theta_1, \dots, \theta_n) u_i(\theta_i, o) + \pi_i(\theta_1, \dots, \theta_n) \geq 0$$
- Objective (e.g. sum of utilities)
$$\sum_{\theta_1, \dots, \theta_n} p(\theta_1, \dots, \theta_n) \sum_i (\sum_o p(o \mid \theta_1, \dots, \theta_n) u_i(\theta_i, o) + \pi_i(\theta_1, \dots, \theta_n))$$
- Also works for BNE incentive compatibility, ex-interim individual rationality notions, other objectives, etc.
- For deterministic mechanisms, use mixed integer programming (probabilities in $\{0, 1\}$)
 - Typically designing the optimal deterministic mechanism is NP-hard

Computational complexity of automatically designing deterministic mechanisms

- Many different variants
 - **Objective** to maximize: Social welfare/revenue/designer's agenda for outcome
 - **Payments** allowed/not allowed
 - **IR constraint**: ex interim IR/ex post IR/no IR
 - **IC constraint**: Dominant strategies/Bayes-Nash equilibrium
- The above already gives $3 * 2 * 3 * 2 = 36$ variants
- Approach: Prove hardness for the case of only 1 type-reporting agent
 - results imply hardness in more general settings

DSE & BNE incentive compatibility constraints coincide when there is only 1 (reporting) agent

Dominant strategies:

Reporting truthfully is optimal for *any* types the others report

	t_{21}	t_{22}
t_{11}	o_5	o_9
t_{12}	o_3	o_2

$$u_1(t_{11}, o_5) \geq u_1(t_{11}, o_3)$$

AND

$$u_1(t_{11}, o_9) \geq u_1(t_{11}, o_2)$$

Bayes-Nash equilibrium:

Reporting truthfully is optimal *in expectation* over the other agents' (true) types

	t_{21}	t_{22}
t_{11}	o_5	o_9
t_{12}	o_3	o_2

$$P(t_{21})u_1(t_{11}, o_5) +$$

$$P(t_{22})u_1(t_{11}, o_9) \geq$$

$$P(t_{21})u_1(t_{11}, o_3) +$$

$$P(t_{22})u_1(t_{11}, o_2)$$

With only 1 reporting agent, the constraints are the same

	t_{21}
t_{11}	o_5
t_{11}	o_3

$$u_1(t_{11}, o_5) \geq u_1(t_{11}, o_3)$$

is equivalent to

$$P(t_{21})u_1(t_{11}, o_5) \geq P(t_{21})u_1(t_{11}, o_3)$$

Ex post and *ex interim* individual rationality constraints coincide when there is only 1 (reporting) agent

Ex post:

Participating never hurts (for any types of the other agents)

	t_{21}	t_{22}
t_{11}	o_5	o_9
t_{12}	o_3	o_2

$$u_1(t_{11}, o_5) \geq 0$$

AND

$$u_1(t_{11}, o_9) \geq 0$$

Ex interim:

Participating does not hurt *in expectation* over the other agents' (true) types

	t_{21}	t_{22}
t_{11}	o_5	o_9
t_{12}	o_3	o_2

$$P(t_{21})u_1(t_{11}, o_5) + P(t_{22})u_1(t_{11}, o_9) \geq 0$$

With only 1 reporting agent, the constraints are the same

	t_{21}
t_{11}	o_5
t_{11}	o_3

$$u_1(t_{11}, o_5) \geq 0$$

is equivalent to

$$P(t_{21})u_1(t_{11}, o_5) \geq 0$$

How hard is designing an optimal *deterministic* mechanism?

NP-complete (even with 1 reporting agent):	Solvable in polynomial time (for any constant number of agents):
<ol style="list-style-type: none">1. Maximizing social welfare (no payments)2. Designer's own utility over outcomes (no payments)3. General (linear) objective that doesn't regard payments4. Expected revenue	<ol style="list-style-type: none">1. Maximizing social welfare (not regarding the payments) (VCG)

1 and 3 hold even with no IR constraints

AMD can create optimal (expected-revenue maximizing) **combinatorial** auctions

- Instance 1

- 2 items, 2 bidders, 4 types each (LL, LH, HL, HH)
- H=utility 2 for that item, L=utility 1
- But: utility 6 for getting both items if type HH (complementarity)
- Uniform prior over types
- Optimal *ex-interim* IR, BNE mechanism (0 = item is burned):
- Payment rule not shown
- Expected revenue: 3.94 (VCG: 2.69)

- Instance 2

- 2 items, 3 bidders
- Complementarity and substitutability
- Took 5.9 seconds
- Uses randomization

	LL	LH	HL	HH
LL	0,0	0,2	2,0	2,2
LH	0,1	1,2	2,1	2,2
HL	1,0	1,2	2,1	2,2
HH	1,1	1,1	1,1	1,1

Optimal mechanisms for a public good

- AMD can design optimal mechanisms for public goods, **taking money burning into account as a loss**
- Bridge building instance
 - Agent 1: High type (prob .6) values bridge at 10. Low: values at 1
 - Agent 2: High type (prob .4) values bridge at 11. Low: values at 2
 - Bridge costs 6 to build
- Optimal mechanism (*ex-post* IR, BNE):

Outcome rule

	Low	High
Low	Don't build	Build
High	Build	Build

Payment rule

	Low	High
Low	0, 0	0, 6
High	4, 2	.67, 5.33

- There is no general mechanism that achieves budget balance, *ex-post* efficiency, and *ex-post* IR [Myerson-Satterthwaite 83]
- However, for this instance, AMD found such a mechanism

Combinatorial public goods problems

- AMD for interrelated public goods
- Example: building a bridge and/or a boat
 - 2 agents each uniform from types: {None, Bridge, Boat, Either}
 - Type indicates which of the two would be useful to the agent
 - If something is built that is useful to you, you get 2, otherwise 0
 - Boat costs 1 to build, bridge 3
- Optimal mechanism (*ex-post* IR, dominant strategies):

Outcome rule
 ($P(\text{none}), P(\text{boat}),$
 $P(\text{bridge}), P(\text{both}))$

	None	Boat	Bridge	Either
None	(1,0,0,0)	(0,1,0,0)	(1,0,0,0)	(0,1,0,0)
Boat	(.5,.5,0,0)	(0,1,0,0)	(0,.5,0,.5)	(0,1,0,0)
Bridge	(1,0,0,0)	(0,1,0,0)	(0,0,1,0)	(0,0,1,0)
Either	(.5,.5,0,0)	(0,1,0,0)	(0,0,1,0)	(0,1,0,0)

- Again, no money burning, but outcome not always efficient
 - E.g., sometimes nothing is built while boat should have been

Additional & future directions

- **Scalability** is a major concern
 - Can sometimes create **more concise** LP formulations
 - Sometimes, some constraints are implied by others
 - In **restricted domains** faster algorithms sometimes exist
 - Can sometimes make use of partial characterizations of the optimal mechanism
- Automatically generated mechanisms can be **complex/hard to understand**
 - Can we make automatically designed mechanisms more intuitive?
- Using AMD to create **conjectures** about general mechanisms