Bits, Bytes, and Integers – Part 2

15-213: Introduction to Computer Systems
3rd Lecture, May 19, 2022

Instructors:

Zack Weinberg

Autolab/Canvas/Piazza accounts

- **Everyone who is enrolled should have an account**
- **We add students to Autolab and Piazza every morning**
- **If you are on waitlist, just keep on hanging in there**
  - Send all waitlist questions to Amy Weis alweis@andrew.cmu.edu
  - Not to the instructors

- **If you were recently added to the course, read https://piazza.com/class/l2qjkx0vcikb7?cid=8 to learn how to get started on the labs**

Reminder about labs 0 and 1

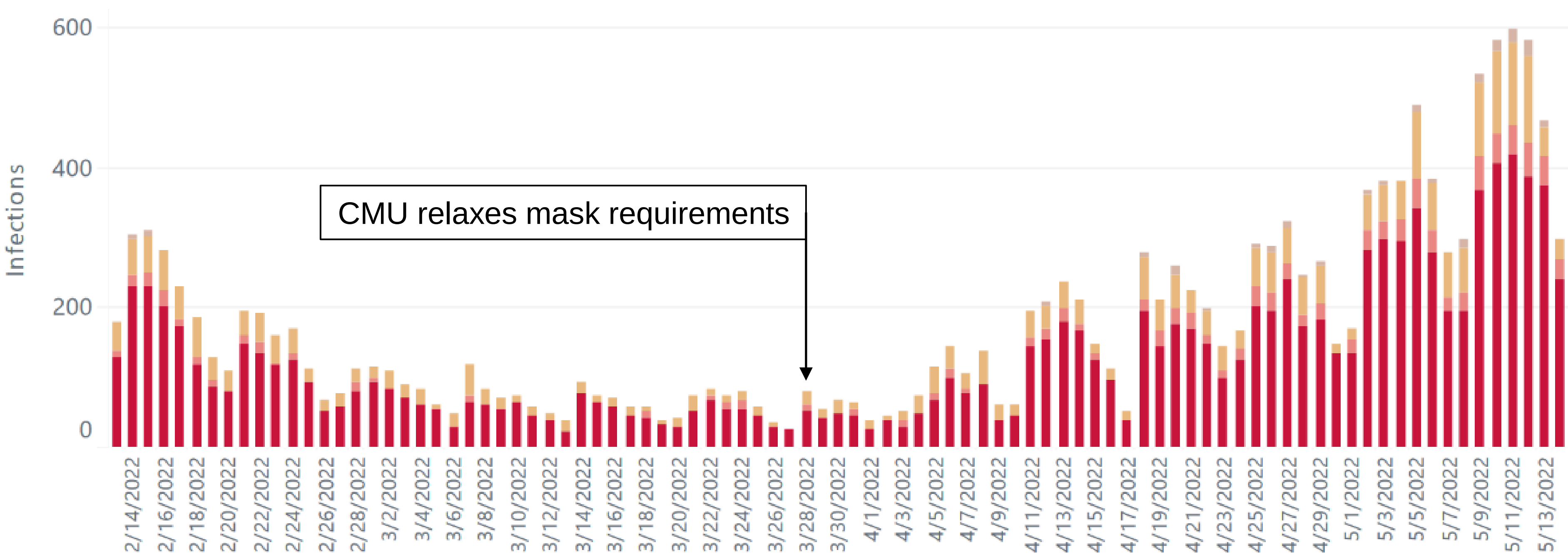
- **Lab 0 is due Monday the 23rd**
  - No late days, no grace days
  - Email instructors if you need an extension
  - It’s supposed to be easy—if it takes you more than a couple hours’ effort, you may not be prepared for this course
- **Lab 1 (data lab) comes out tomorrow (the 20th)**
  - Will be due June 3 (two weeks)
  - However, lab 2 (bomb lab) comes out the 27th and is *also* due on June 3
  - Start early!

Reminder about in-person attendance

- This room is full.
- In-person classes are primarily for undergraduate students who are enrolled in 15-***2***13.
- If I told you in email to enroll in 513 for a reason *other than* "you're a graduate student," you are also entitled to be in this room.
- If you're enrolled in 15-***5***13 because you are a graduate student, however, you may attend *only if there is space*.

Reminder about masks

COVID-19 infections, Allegheny County, last three months



- The pandemic is not over
- I think it would be a really good idea if we all continued to wear masks in this classroom

Today: Bits, Bytes, and Integers

- Representing information as bits
- Bit-level manipulations
- Integers
  - **Representation: unsigned and signed; negation**
  - Conversion, casting
  - Extension, truncation, shifting
  - Addition, multiplication
- Representations in memory, pointers, strings

Encoding "Integers"

Unsigned

Given a bit vector x, w bits long…

$$\mathrm{B2U}(x) = \sum_{i=0}^{w-1} x_i \cdot 2^i$$

Signed (twos complement)

$$\mathrm{B2T}(x) = -x_{w-1} \cdot 2^{w-1} + \sum_{i=0}^{w-2} x_i \cdot 2^i$$

Sign Bit

Examples (w = 5)

-16	8	4	2	1
0	1	0	1	0

$$0 + 8 + 0 + 2 + 0 = 10$$

16	8	4	2	1
1	0	1	1	0
-16	8	4	2	1

$$16 + 8 + 0 + 2 + 0 = 26$$

$$-16 + 8 + 0 + 2 + 0 = -10$$

Negation: Complement & Increment

- **Negate through complement and increase**

  `~x + 1 == -x`

- **Why?**
  - `-x + x       == 0` *(by definition)*
  - `~x + x       == 1111…111 == -1`
  - `~x + x + 1 == 0`
  - `(~x+1) + x == 0`
  - `~x+1        == -x`

	x	1 0 0 1 1 1 0 1
+	~x	0 1 1 0 0 0 1 0
	−1	1 1 1 1 1 1 1 1

Example: x = 15213

	Decimal	Hex	Binary
`x`	15213	3B 6D	00111011 01101101
`~x`	-15214	C4 92	11000100 10010010
`~x+1`	-15213	C4 93	11000100 10010011
`y`	-15213	C4 93	11000100 10010011

Complement & Increment Examples

x = 0

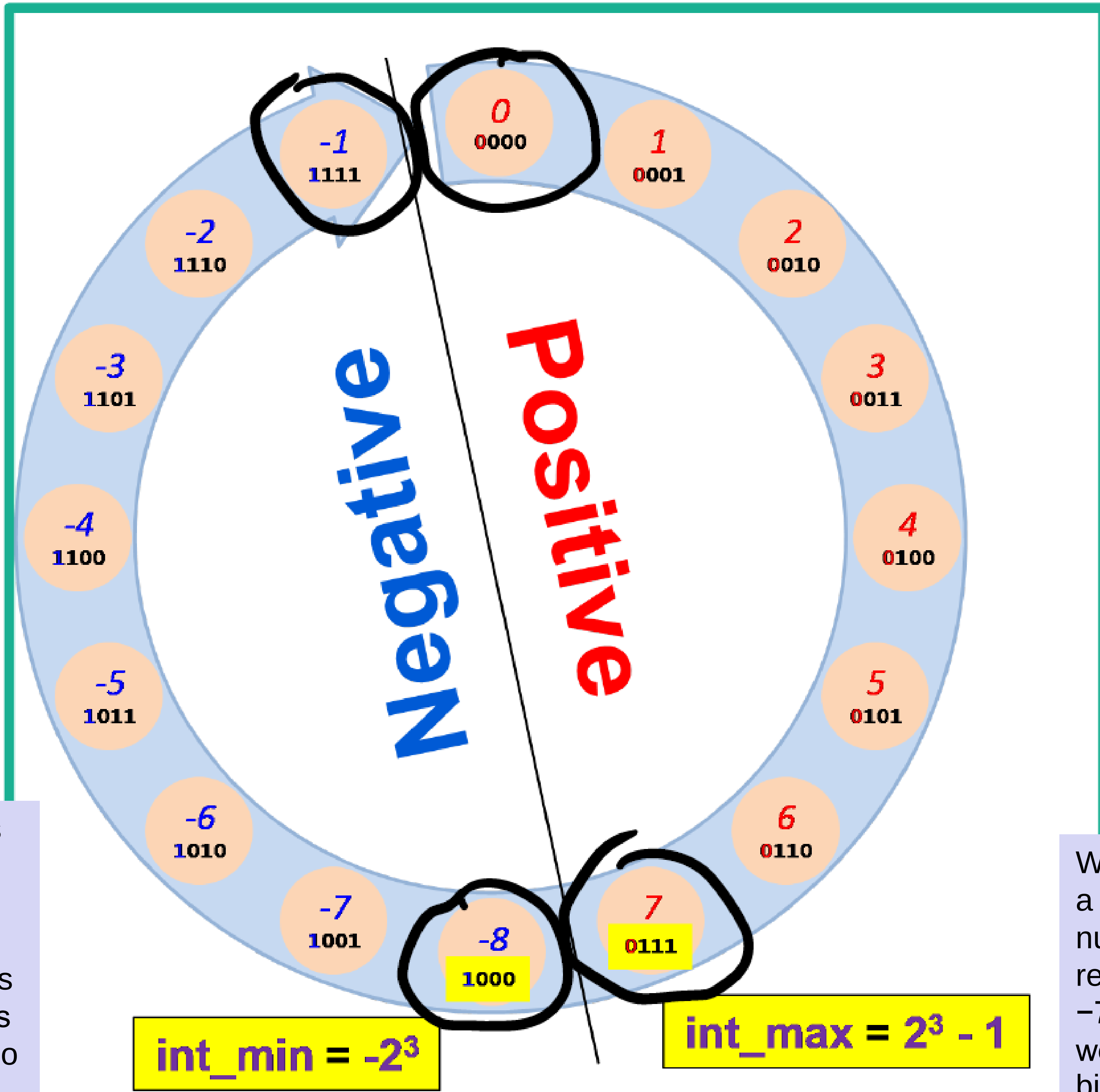
	Decimal	Hex	Binary
0	**0**	00 00	00000000 00000000
~0	**-1**	FF FF	11111111 11111111
~0+1	**0**	00 00	00000000 00000000

x = TMin

	Decimal	Hex	Binary
x	**-32768**	80 00	10000000 00000000
~x	**32767**	7F FF	01111111 11111111
~x+1	**-32768**	80 00	10000000 00000000

Oops!
It's still
negative!

Eight negative values:
–1, – 2, …, – 8



Eight *non*-negative values:
0, 1, …, 7

Mathematicians would prefer it if a 4-bit signed number could represent values −8…8, but that's $2^4 + 1$ values, so they won't all fit.

What if we made a 4-bit signed number only represent values −7…7? Then we wouldn't be using bit pattern `1000`…

Activity: https://www.cs.cmu.edu/afs/cs/academic/class/15213-m22/www/activities/213_lecture3.pdf

Do “model 0” and “model −1”, then stop.

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- Representing information as bits
- **Bit-level manipulations**
- Integers
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Boolean Algebra

- **Developed by George Boole in 19th Century**
  - Algebraic representation of logic
    - Encode "True" as 1 and "False" as 0

And

- **A&B = 1 when both A=1 and B=1**

&	0	1
0	0	0
1	0	1

Or

- **A|B = 1 when either A=1 or B=1**

|	0	1
0	0	1
1	1	1

Not

- **~A = 1 when A=0**

~	
0	1
1	0

Exclusive-Or (Xor)

- **A^B = 1 when either A=1 or B=1, but not both**

^	0	1
0	0	1
1	1	0

General Boolean Algebras

- **Operate on Bit Vectors**
  - Operations applied bitwise

```
  01101001     01101001     01101001
& 01010101   | 01010101   ^ 01010101   ~ 01010101
  01000001     01111101     00111100     10101010
```

- **All of the Properties of Boolean Algebra Apply**

Example: Representing & Manipulating Sets

Representation

- Width w bit vector represents subsets of {0, …, w–1}
- $a_j = 1$ if $j \in A$
  - 01101001 { 0, 3, 5, 6 }
  - *76543210*

  - 01010101 { 0, 2, 4, 6 }
  - *76543210*

Operations

&	Intersection	01000001	{ 0, 6 }
|	Union	01111101	{ 0, 2, 3, 4, 5, 6 }
^	Symmetric difference	00111100	{ 2, 3, 4, 5 }
~	Complement	10101010	{ 1, 3, 5, 7 }

Bit-Level Operations in C

Operations &, |, ~, ^ Available in C

- Apply to any "integral" data type
  - long, int, short, char, unsigned
- View arguments as bit vectors
- Arguments applied bit-wise

Examples (Char data type)

- ~0x41 →
- ~0x00 →
- 0x69 & 0x55 →
- 0x69 | 0x55 →

Hex	Decimal	Binary
0	0	0000
1	1	0001
2	2	0010
3	3	0011
4	4	0100
5	5	0101
6	6	0110
7	7	0111
8	8	1000
9	9	1001
A	10	1010
B	11	1011
C	12	1100
D	13	1101
E	14	1110
F	15	1111

Bit-Level Operations in C

- **Operations &, |, ~, ^ Available in C**
  - Apply to any “integral” data type
    - long, int, short, char, unsigned
  - View arguments as bit vectors
  - Arguments applied bit-wise
- **Examples (Char data type)**
  - ~0x41 → 0xBE
    - $\sim 0100\ 0001_2 \rightarrow 1011\ 1110_2$
  - ~0x00 → 0xFF
    - $\sim 0000\ 0000_2 \rightarrow 1111\ 1111_2$
  - 0x69 & 0x55 → 0x41
    - $0110\ 1001_2\ \&\ 0101\ 0101_2 \rightarrow 0100\ 0001_2$
  - 0x69 | 0x55 → 0x7D
    - $0110\ 1001_2\ |\ 0101\ 0101_2 \rightarrow 0111\ 1101_2$

Hex	Decimal	Binary
0	0	0000
1	1	0001
2	2	0010
3	3	0011
4	4	0100
5	5	0101
6	6	0110
7	7	0111
8	8	1000
9	9	1001
A	10	1010
B	11	1011
C	12	1100
D	13	1101
E	14	1110
F	15	1111

Contrast: Logic Operations in C

- **Contrast to Bit-Level Operators**
  - **Logic Operations &, ||, !**
    - View 0 as "Fals
    - Anything nonze
    - Alway
    - Early

Watch out for && vs. & (and || vs. |)… one of the more common oopsies in C programming

- **Example**
  - !0x41 →
  - !0x00 →
  - !!0x41→ 0x01
  - 0x69 && 0x55 → 0x01
  - 0x69 || 0x55 → 0x01
  - p && *p (avoids null pointer access)

Activity: https://www.cs.cmu.edu/afs/cs/academic/class/15213-m22/www/activities/213_lecture3.pdf

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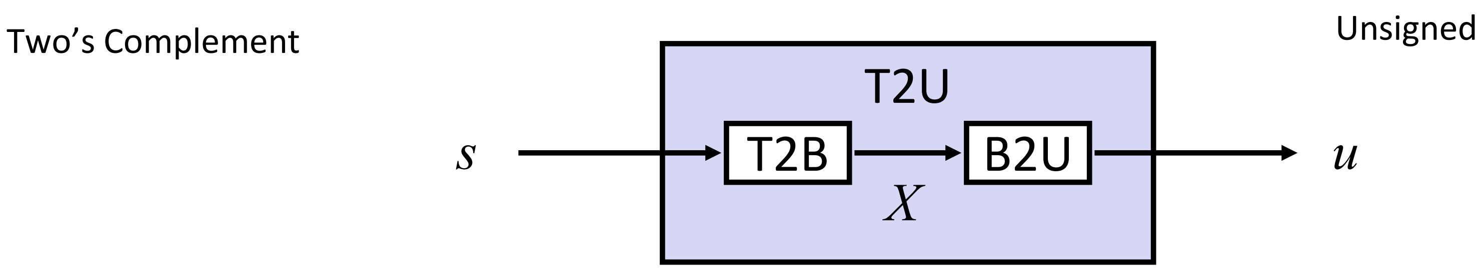
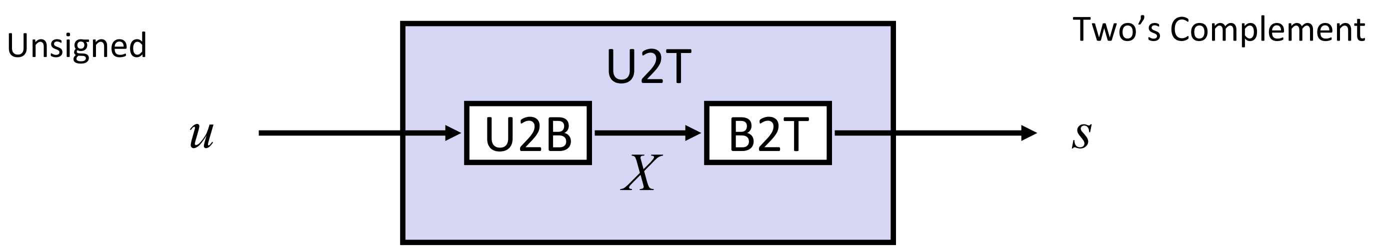
Mapping Between Signed & Unsigned



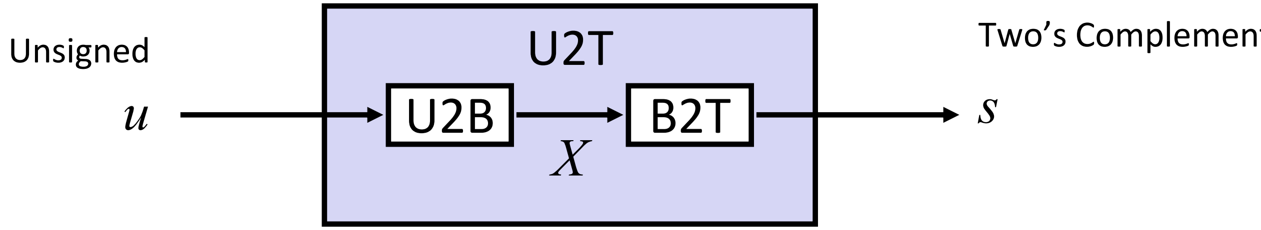
Maintain Same Bit Pattern



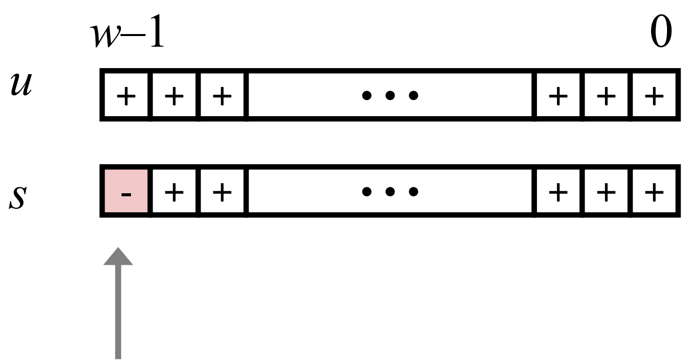
Maintain Same Bit Pattern

- **Mappings between unsigned and two's complement numbers:**
  Keep bit representations and reinterpret

Relation between Signed & Unsigned



Maintain Same Bit Pattern



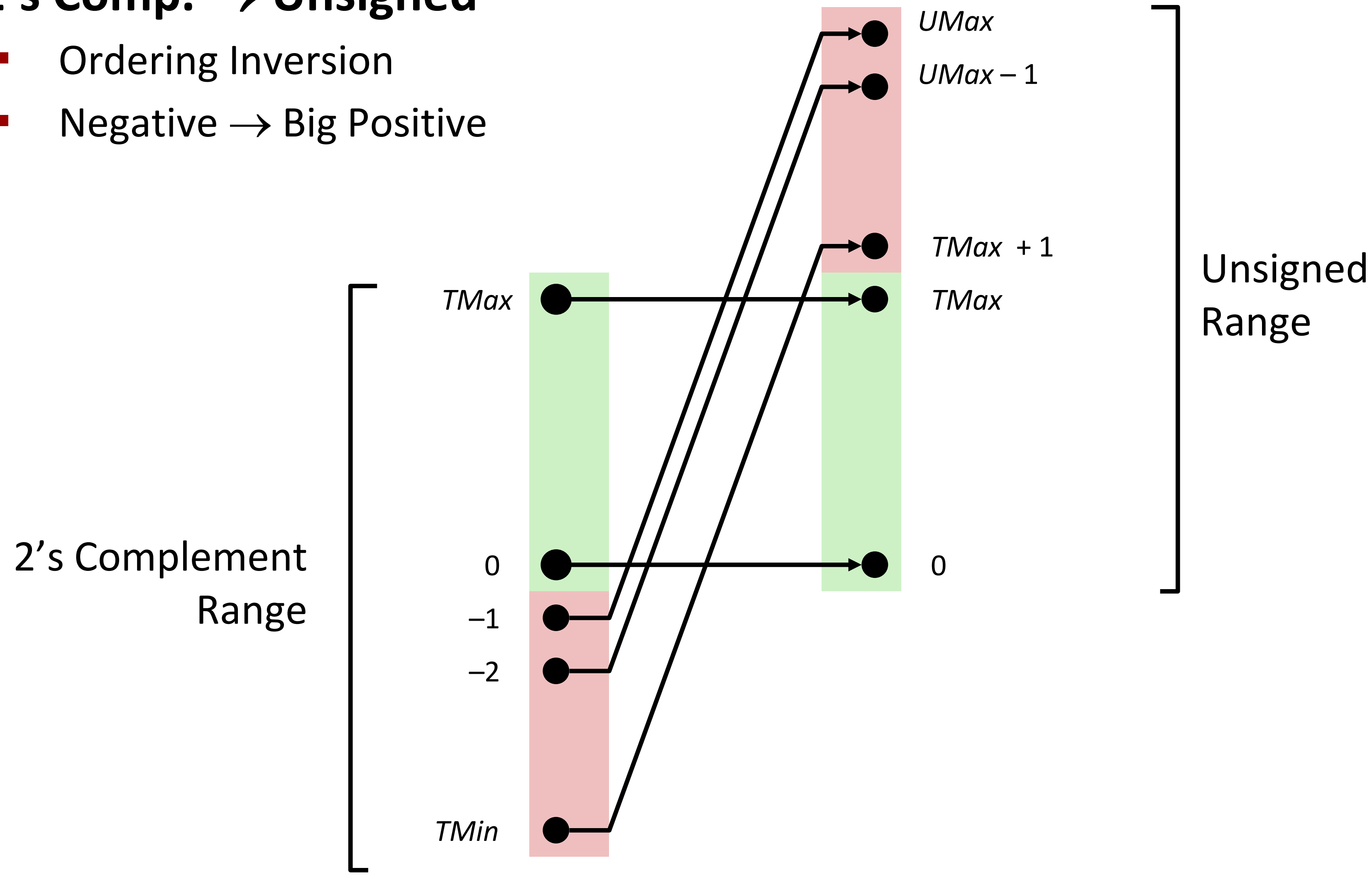
Large positive weight
becomes
Large negative weight

Mapping Signed ↔ Unsigned

Bits	Signed		Unsigned
0000	0		0
0001	1		1
0010	2		2
0011	3		3
0100	4	=	4
0101	5		5
0110	6		6
0111	7		7
1000	-8		8
1001	-7		9
1010	-6		10
1011	-5	± 16	11
1100	-4		12
1101	-3		13
1110	-2		14
1111	-1		15

Conversion Visualized

- **2's Comp. → Unsigned**
  - Ordering Inversion
  - Negative → Big Positive



Signed vs. Unsigned in C

Constants

- By default are considered to be signed integers
- Unsigned if have “U” as suffix

```
0U, 4294967259U
```

Casting

- Explicit casting between signed & unsigned same as U2T and T2U

```
int tx, ty;
unsigned ux, uy;
tx = (int) ux;
uy = (unsigned) ty;
```

- Implicit casting also occurs via assignments and procedure calls

```
tx = ux;                  int fun(unsigned u);
uy = ty;                  uy = fun(tx);
```

Casting Surprises

Expression Evaluation

- If there is a mix of unsigned and signed in single expression, ***signed values implicitly cast to unsigned***
- Including comparison operations **<, >, ==, <=, >=**
- Examples:

Constant 1	Constant 2	Relation	Evaluation
`0`	`0U`	`==`	Unsigned
`-1`	`0`	`<`	Signed
`-1`	`0U`	`>`	Unsigned
`INT_MAX`	`INT_MIN`	`>`	Signed
`(unsigned)INT_MAX`	`INT_MIN`	`<`	Unsigned
`-1`	`-2`	`>`	Signed
`(unsigned)-1`	`-2`	`>`	Unsigned
`INT_MAX`	`((unsigned)INT_MAX) + 1`	`<`	Unsigned
`INT_MAX`	`(int)(((unsigned)INT_MAX) + 1)`	`>`	Signed

Summary
Casting Signed ↔ Unsigned: Basic Rules

- Bit pattern is maintained
- But reinterpreted
- Can have unexpected effects: adding or subtracting 2^w

- **Expression containing signed and unsigned int**
  - `int` is cast to `unsigned`!!

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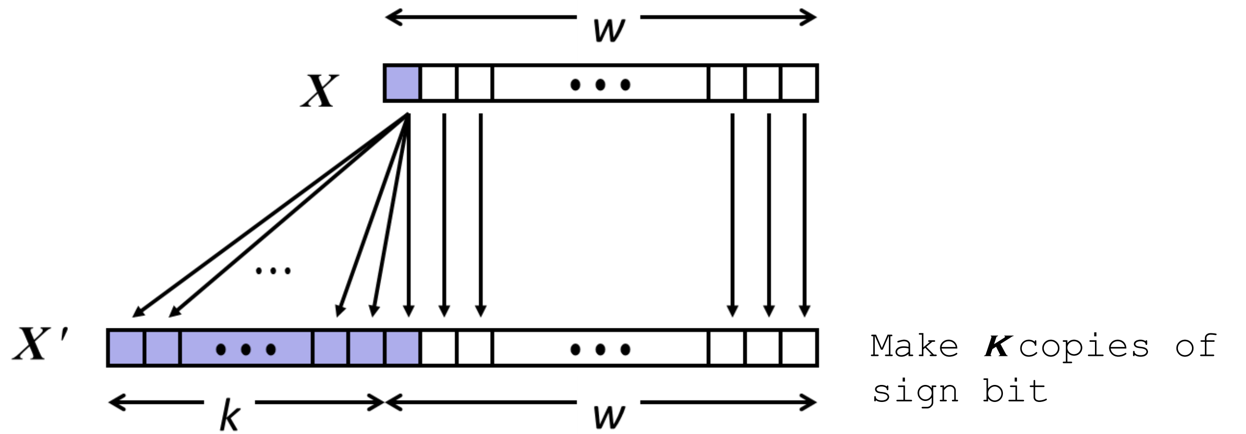
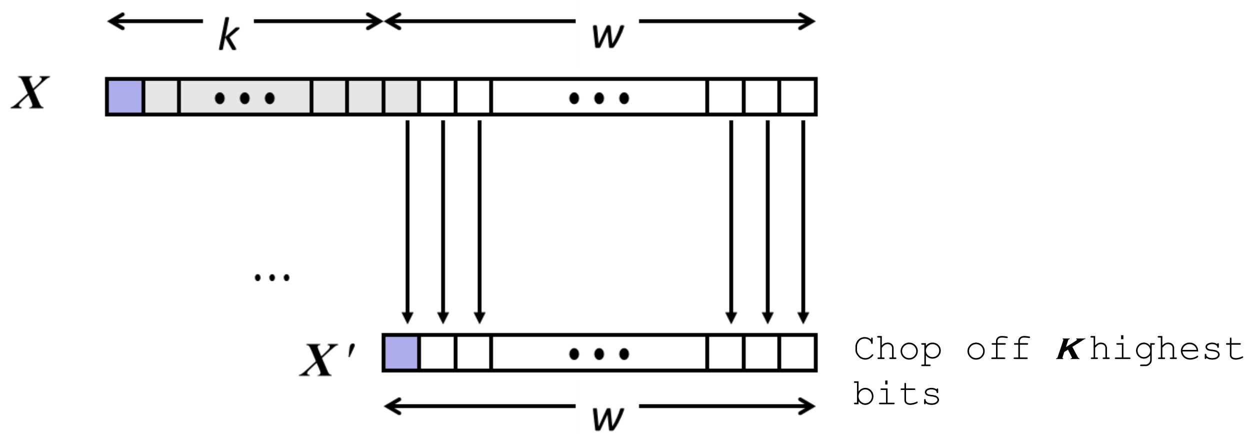
Sign Extension and Truncation

- **Sign Extension**



- **Truncation**



Sign Extension: Simple Example

Positive number

	-16	8	4	2	1
10 =	0	1	0	1	0

	-32	16	8	4	2	1
10 =	0	0	1	0	1	0

Negative number

	-16	8	4	2	1
-10 =	1	0	1	1	0

	-32	16	8	4	2	1
-10 =	1	1	0	1	1	0

Truncation: Simple Example

No sign change

	-16	8	4	2	1
2 =	0	0	0	1	0

	-8	4	2	1
2 =	0	0	1	0

`2 mod 16 = 2`

	-16	8	4	2	1
-6 =	1	1	0	1	0

	-8	4	2	1
-6 =	1	0	1	0

`-6 mod 16 = 26U mod 16 = 10U = -6`

Sign change

	-16	8	4	2	1
10 =	0	1	0	1	0

	-8	4	2	1
-6 =	1	0	1	0

`10 mod 16 = 10U mod 16 = 10U = -6`

	-16	8	4	2	1
-10 =	1	0	1	1	0

	-8	4	2	1
6 =	0	1	1	0

`-10 mod 16 = 22U mod 16 = 6U = 6`

Shifting

- **Left Shift: `x << y`**
  - Shift bit-vector `x` left `y` positions
  - Throw away extra bits on left
  - Fill with 0's on right
  - Equivalent to multiplying by 2^y
- **Right Shift: `x >> y`**
  - Shift bit-vector `x` right `y` positions
  - Throw away extra bits on right
  - Two kinds:
    - "Logical": Fill with 0's on left
    - "Arithmetic": Replicate most significant bit on left
  - *Almost* equivalent to dividing by 2^y
- **Undefined Behavior (in C)**
  - Shift amount < 0 or ≥ word size

Argument `x`	01100010
<< 3	00010000
Logical >> 2	00011000
Arithmetic >> 2	00011000

Argument `x`	10100010
<< 3	00010000
Logical >> 2	00101000
Arithmetic >> 2	11101000

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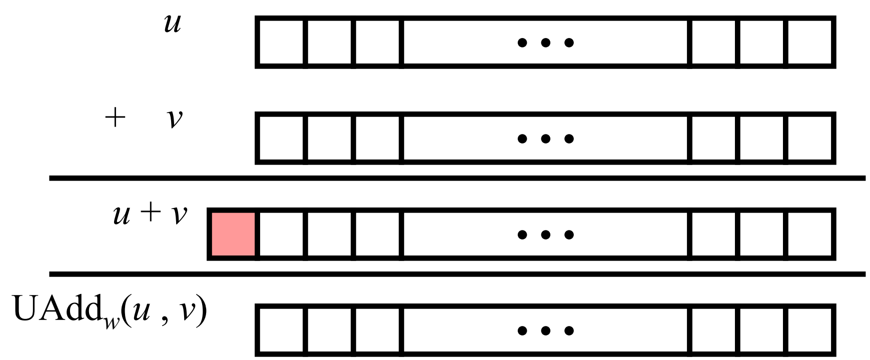
Unsigned Addition



Operands: *w* bits

True Sum: *w*+1 bits

Discard Carry: *w* bits

- **Standard Addition Function**
  - Ignores carry output
- **Implements Modular Arithmetic**

$$s \quad = \quad \mathrm{UAdd}_w(u\,,v) \quad = \quad u + v \;\bmod 2^w$$

```
unsigned char      1110 1001      E9      233
                 +  1101 0101   + D5    + 213
                 ____________   ____    _____

                 ____________   ____    _____
```

Hex	Decimal	Binary
0	0	0000
1	1	0001
2	2	0010
3	3	0011
4	4	0100
5	5	0101
6	6	0110
7	7	0111
8	8	1000
9	9	1001
A	10	1010
B	11	1011
C	12	1100
D	13	1101
E	14	1110
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Unsigned Addition

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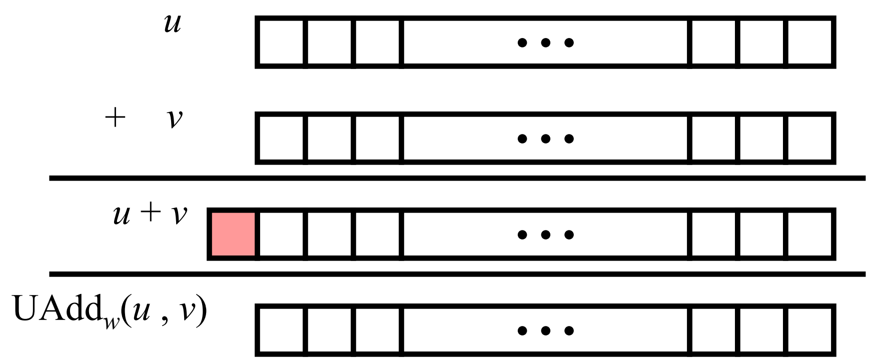

- **Standard Addition Function**
  - Ignores carry output
- **Implements Modular Arithmetic**

$$s \quad = \quad \text{UAdd}_w(u\ ,\ v) \quad = \quad u + v \mod 2^w$$

`unsigned char`	1110 1001	E9	233
	+ 1101 0101	+ D5	+ 213
	1 1011 1110	1BE	446
	1011 1110	BE	190

Hex	Decimal	Binary
0	0	0000
1	1	0001
2	2	0010
3	3	0011
4	4	0100
5	5	0101
6	6	0110
7	7	0111
8	8	1000
9	9	1001
A	10	1010
B	11	1011
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D	13	1101
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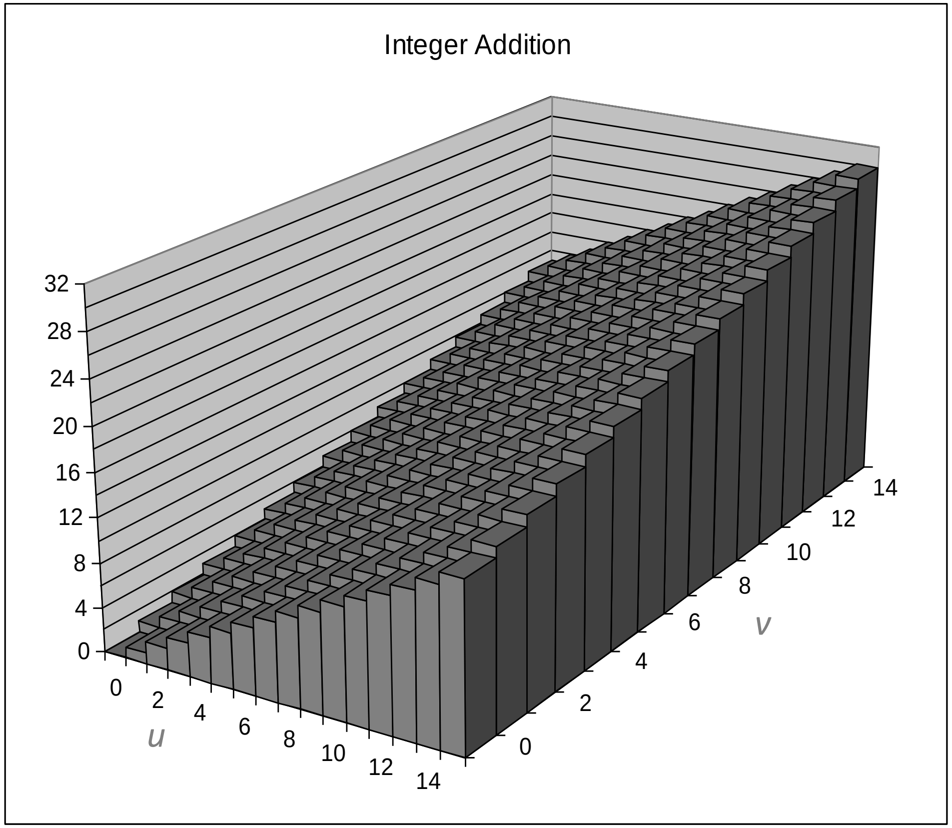
F	15	1111

Visualizing (Mathematical) Integer Addition

Integer Addition

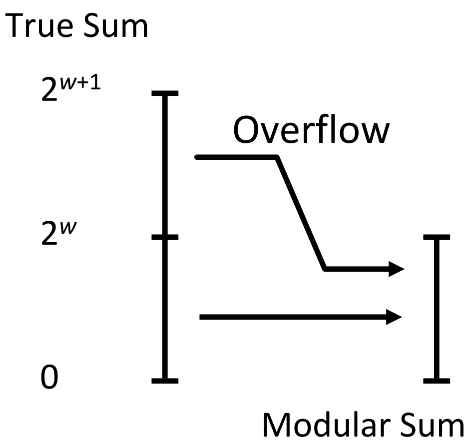
- 4-bit integers *u*, *v*
- Compute true sum $\text{Add}_4(u, v)$
- Values increase linearly with *u* and *v*
- Forms planar surface

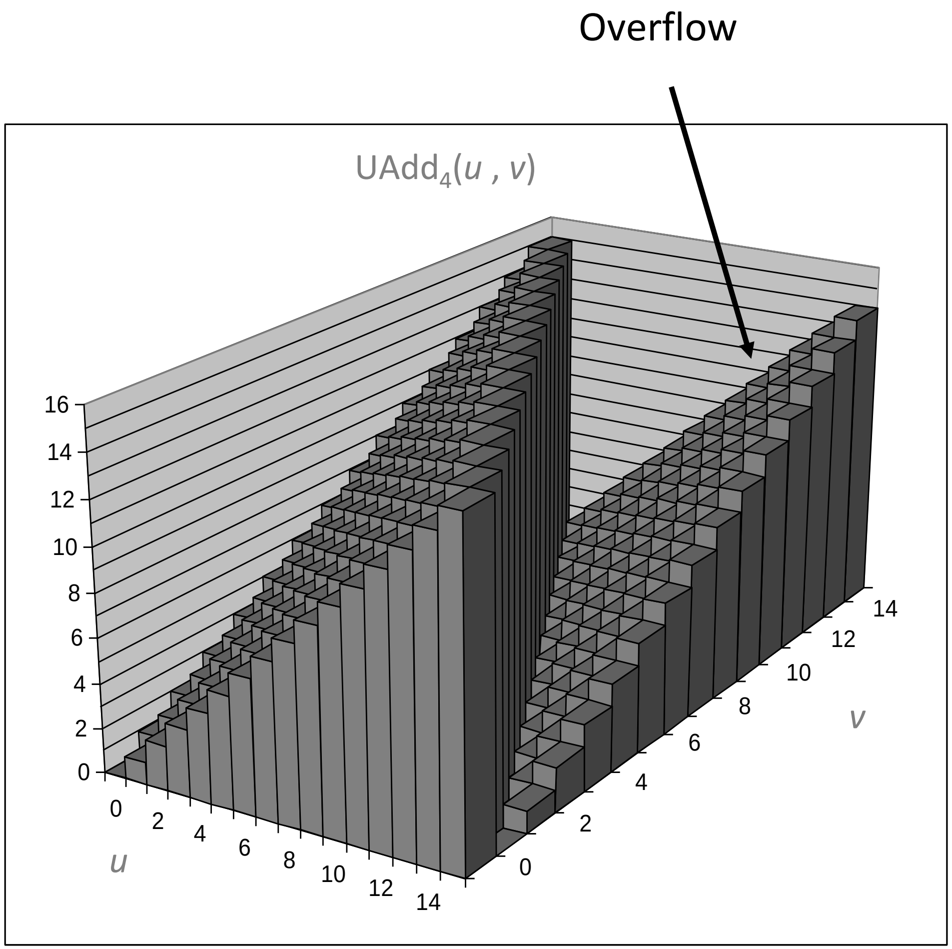
$\text{Add}_4(u, v)$



Visualizing Unsigned Addition

- **Wraps Around**
  - If true sum ≥ 2^w
  - At most once



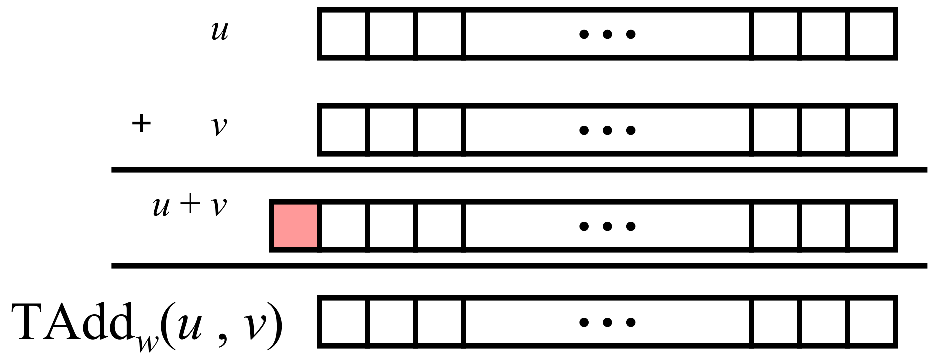


Two’s Complement Addition

Operands: *w* bits

True Sum: *w*+1 bits

Discard Carry: *w* bits



TAdd and UAdd have Identical Bit-Level Behavior

- Signed vs. unsigned addition in C:

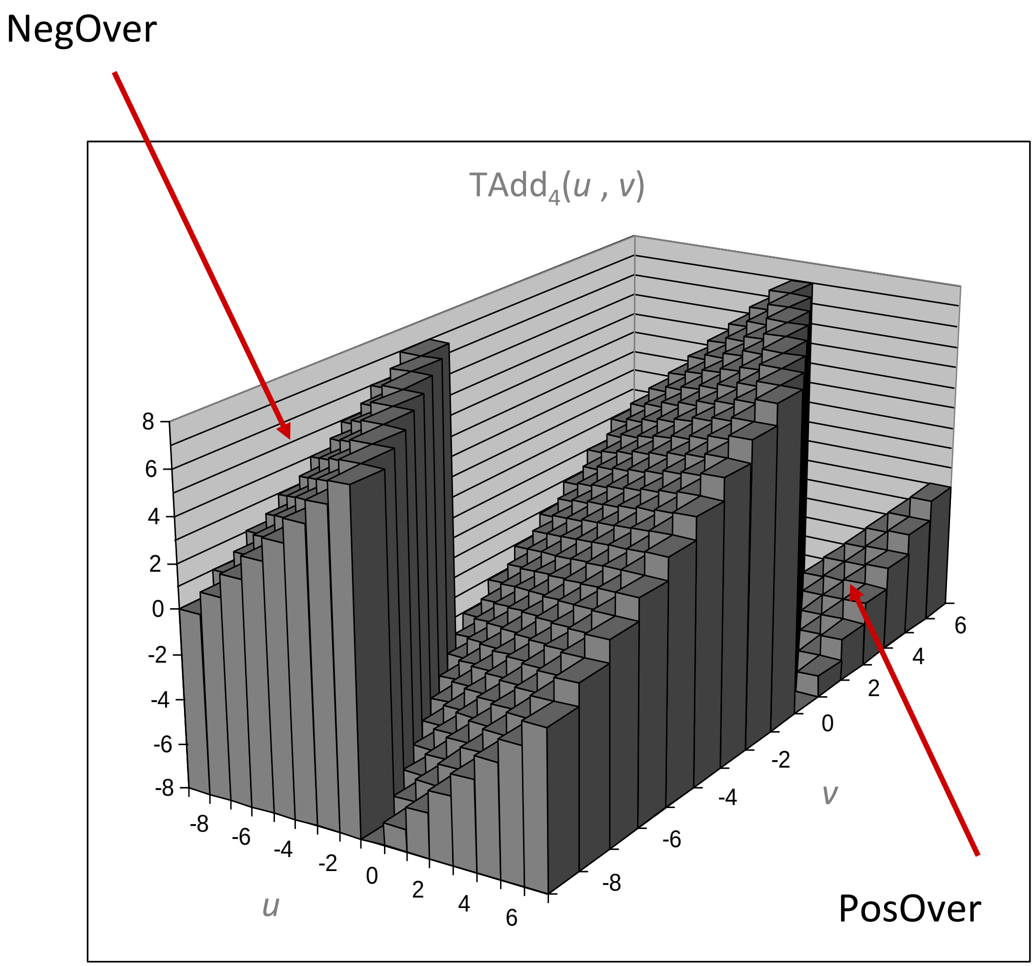
```
int s, t, u, v;
s = (int) ((unsigned) u + (unsigned) v);
t = u + v
```

- Will give `s == t`

1110 1001	E9	-23
+ 1101 0101	+ D5	+ -43
1 1011 1110	1BE	-66
1011 1110	BE	-66

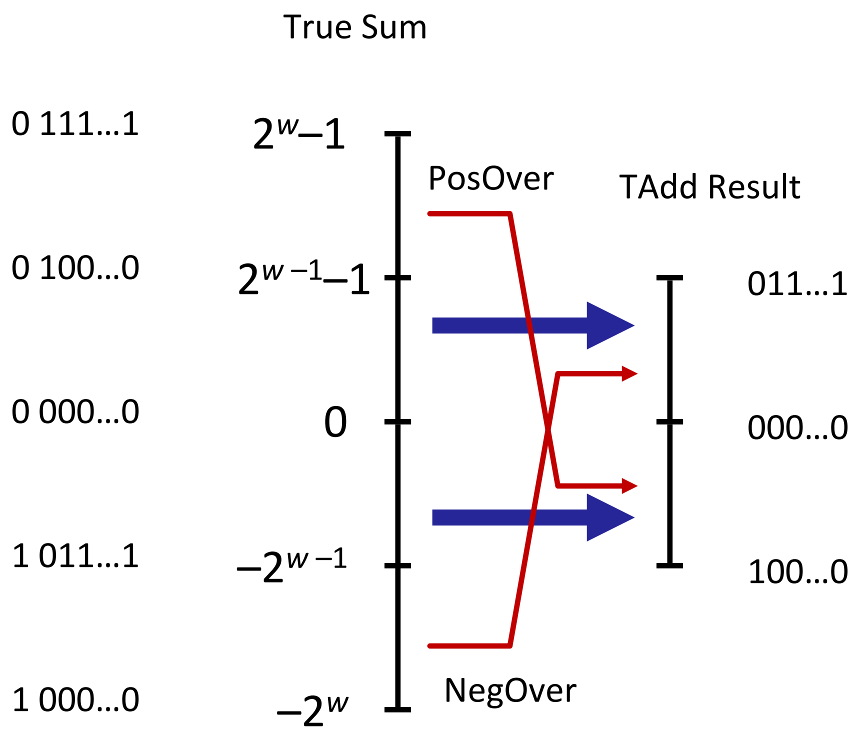
Visualizing 2's Complement Addition

- **Values**
  - 4-bit two's comp.
  - Range from -8 to +7
- **Wraps Around**
  - If sum $\geq 2^{w-1}$
    - Becomes negative
    - At most once
  - If sum $< -2^{w-1}$
    - Becomes positive
    - At most once



TAdd Overflow

- **Functionality**
  - True sum requires w+1 bits
  - Drop off MSB
  - Treat remaining bits as 2's comp. integer



Multiplication

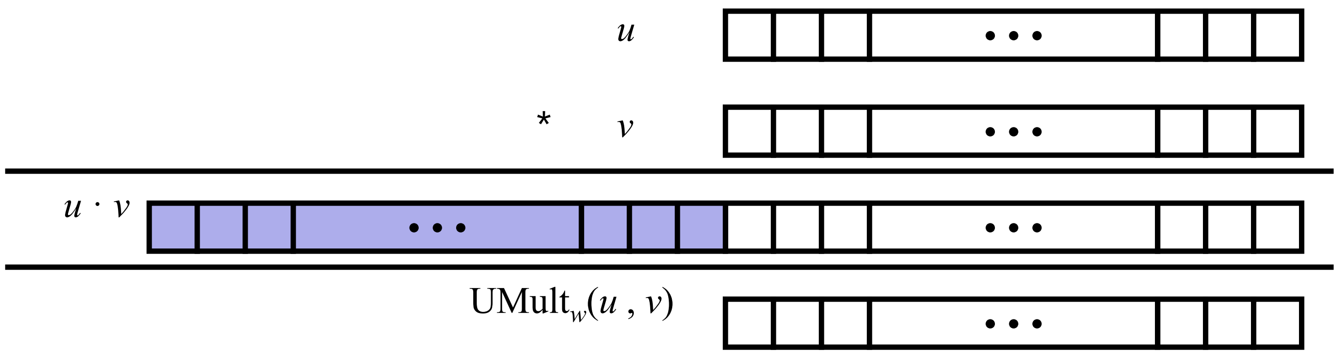
- **Goal: Computing Product of *w*-bit numbers *x*, *y***
  - Either signed or unsigned
- **But, exact results can be bigger than *w* bits**
  - Unsigned: up to 2*w* bits
    - Result range: $0 \le x * y \le (2^w - 1)^2 = 2^{2w} - 2^{w+1} + 1$
  - Two's complement min (negative): Up to 2*w*-1 bits
    - Result range: $x * y \ge (-2^{w-1})*(2^{w-1}-1) = -2^{2w-2} + 2^{w-1}$
  - Two's complement max (positive): Up to 2*w* bits, but only for $(TMin_w)^2$
    - Result range: $x * y \le (-2^{w-1})^2 = 2^{2w-2}$
- **So, maintaining exact results…**
  - would need to keep expanding word size with each product computed
  - is done in software, if needed
    - e.g., by "arbitrary precision" arithmetic packages

Unsigned Multiplication in C

Operands: *w* bits

True Product: 2**w* bits

Discard *w* bits: *w* bits



Standard Multiplication Function

- Ignores high order *w* bits

Implements Modular Arithmetic

$\text{UMult}_w(u, v) \quad = \quad u \cdot v \bmod 2^w$

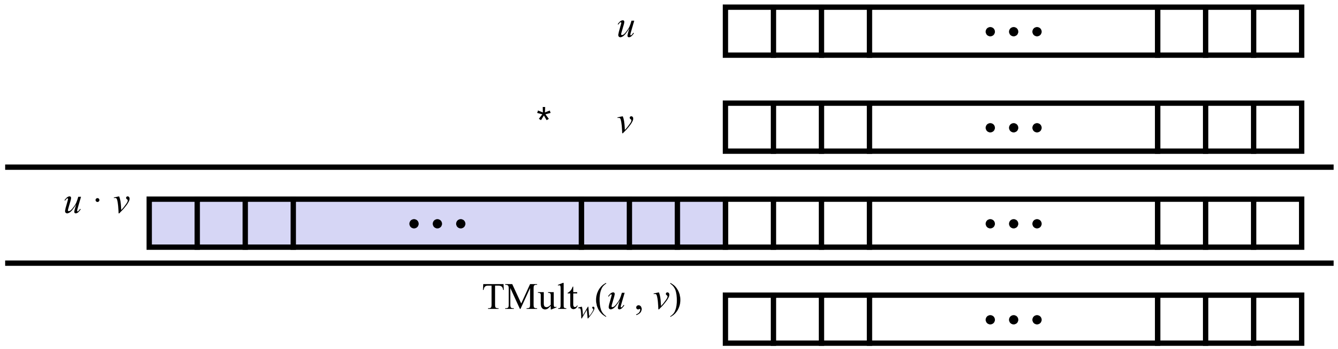
```
                1110 1001        E9        233
*               1101 0101     *  D5     *  213
  1100 0001 1101 1101          C1DD      49629
                1101 1101        DD        221
```

Signed Multiplication in C

Operands: *w* bits

True Product: 2**w* bits

Discard *w* bits: *w* bits



Standard Multiplication Function

- Ignores high order *w* bits
- Some of which are different for signed vs. unsigned multiplication
- Lower bits are the same

	1110 1001	E9	-23
*	1101 0101	* D5	* -43
	0000 0011 1101 1101	03DD	989
	1101 1101	DD	-35

Power-of-2 Multiply with Shift

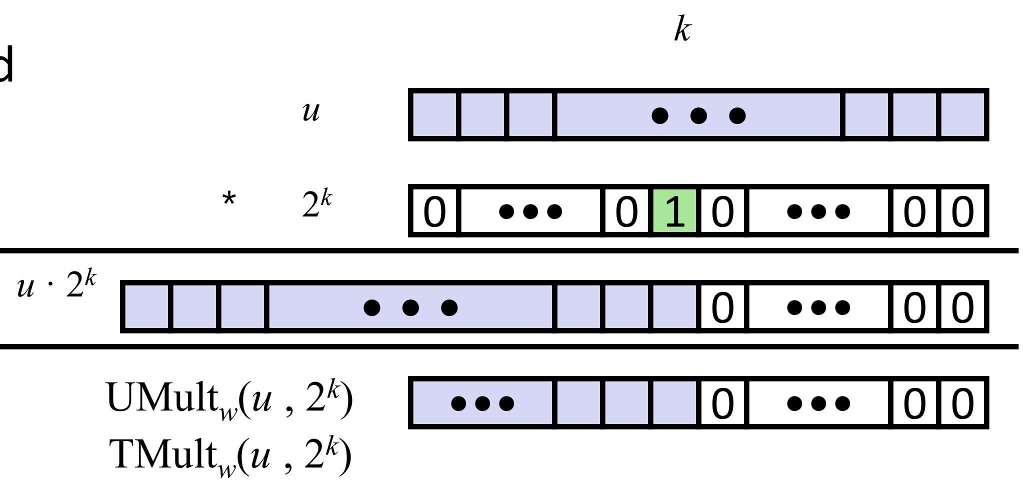
Operation

- `u << k` gives `u * 2`^{k}
- Both signed and unsigned

Operands: *w* bits

True Product: *w*+*k* bits

Discard *k* bits: *w* bits



Examples

- `u << 3     ==     u * 8`
- `(u << 5) - (u << 3)     ==     u * 24`
- Most machines shift and add faster than multiply
  - Compiler generates this code automatically

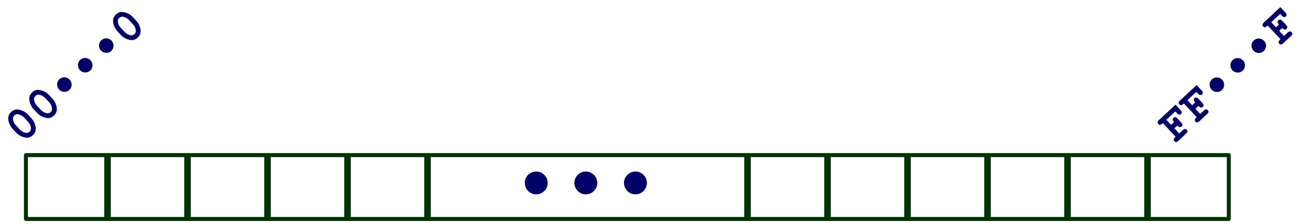
Activity: https://www.cs.cmu.edu/afs/cs/academic/class/15213-m22/www/activities/213_lecture3.pdf

Do “model 3”, then stop.

Today: Bits, Bytes, and Integers

- Representing information as bits
- Bit-level manipulations
- Integers
  - Representation: unsigned and signed; negation
  - Conversion, casting
  - Extension, truncation, shifting
  - Addition, multiplication
- **Representations in memory, pointers, strings**

Byte-Oriented Memory Organization



- **Programs refer to data by address**
  - Imagine all of RAM as an enormous array of bytes
  - An address is an index into that array
    - A pointer variable stores an address
- **System provides a private *address space* to each "process"**
  - A process is an instance of a program, being executed
  - An address space is one of those enormous arrays of bytes
  - Each program can see only its own code and data within its enormous array
  - We'll come back to this later ("virtual memory" classes)

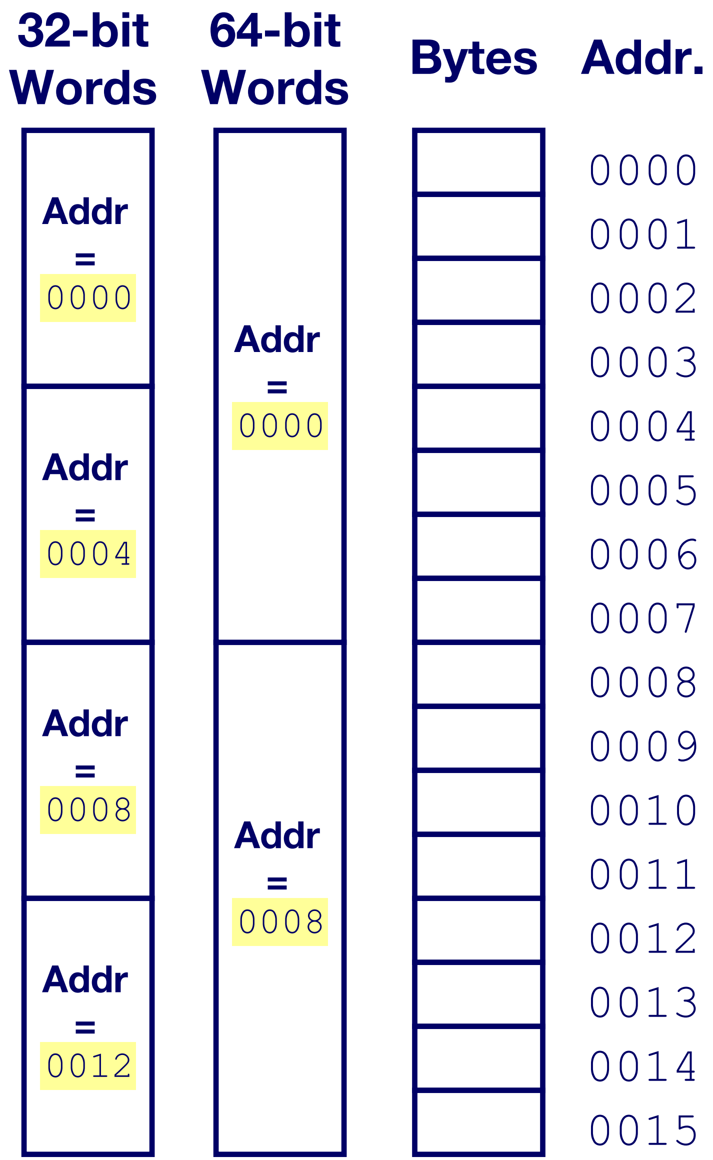
Machine Words

- **Any given computer has a "Word Size"**
  - Nominal size of integer-valued data
    - and of addresses
  - Until recently, most machines used 32 bits (4 bytes) as word size
    - Limits addresses to 4GB (2^{32} bytes)
  - Increasingly, machines have 64-bit word size
    - Potentially, could have 16 EB (exabytes) of addressable memory
    - That's 18.4×10^{18} bytes
  - Machines still support multiple data formats
    - Fractions or multiples of word size
    - Always integral number of bytes

Yes, both of these numbers are correct.
This discrepancy is known as the Great Storage Industry Marketing Lie. Ask me about it after class if you really want to know.

Addresses *Always* Specify Byte Locations

- Address of a word is address of the first byte in the word
- Addresses of successive words differ by 4 (32-bit) or 8 (64-bit)



Example Data Representations

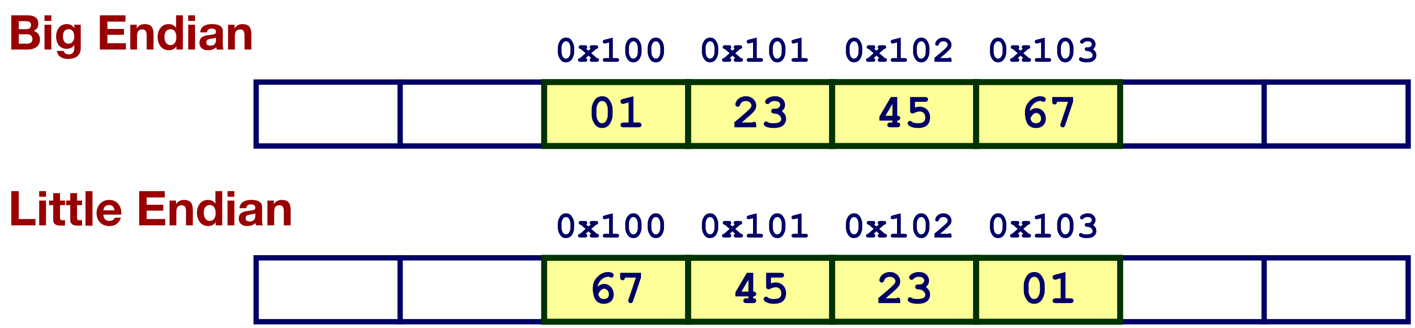
C Data Type	Typical 32-bit	Typical 64-bit	x86-64
`char`	1	1	1
`short`	2	2	2
`int`	4	4	4
`long`	4	8	8
`float`	4	4	4
`double`	8	8	8
pointer	4	8	8

Byte Ordering

- **So, how are the bytes within a multi-byte word ordered in memory?**
- **Conventions**
  - Big Endian: Sun, PPC Mac, *network packet headers*
    - Least significant byte has highest address
  - Little Endian: *x86*, ARM processors running Android, iOS, and Windows
    - Least significant byte has lowest address

Byte Ordering Example

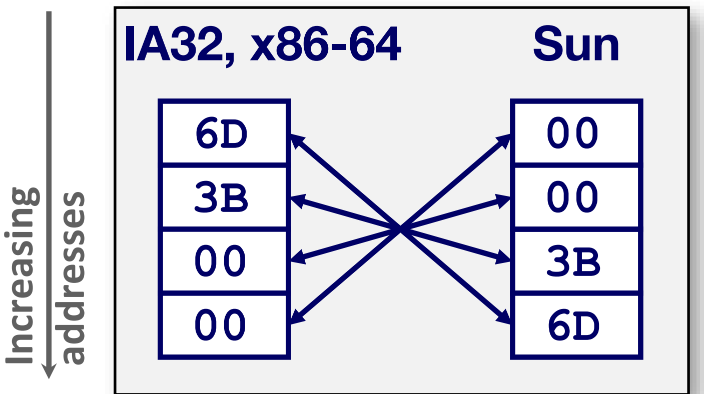
- **Example**
  - Variable x has 4-byte value of 0x01234567
  - Address given by &x is 0x100



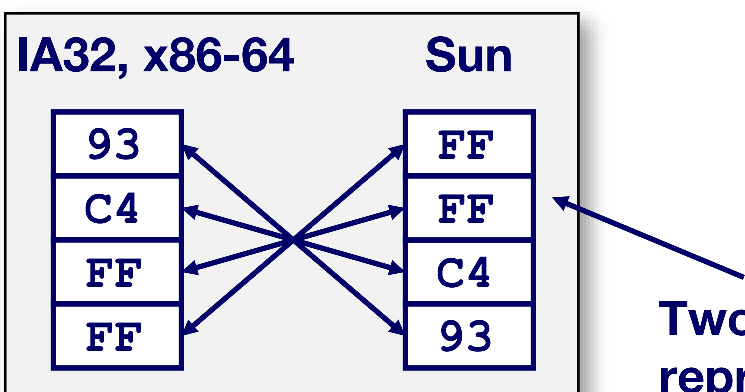
Representing Integers

Decimal:	15213			
Binary:	0011	1011	0110	1101
Hex:	3	B	6	D

`int A = 15213;`

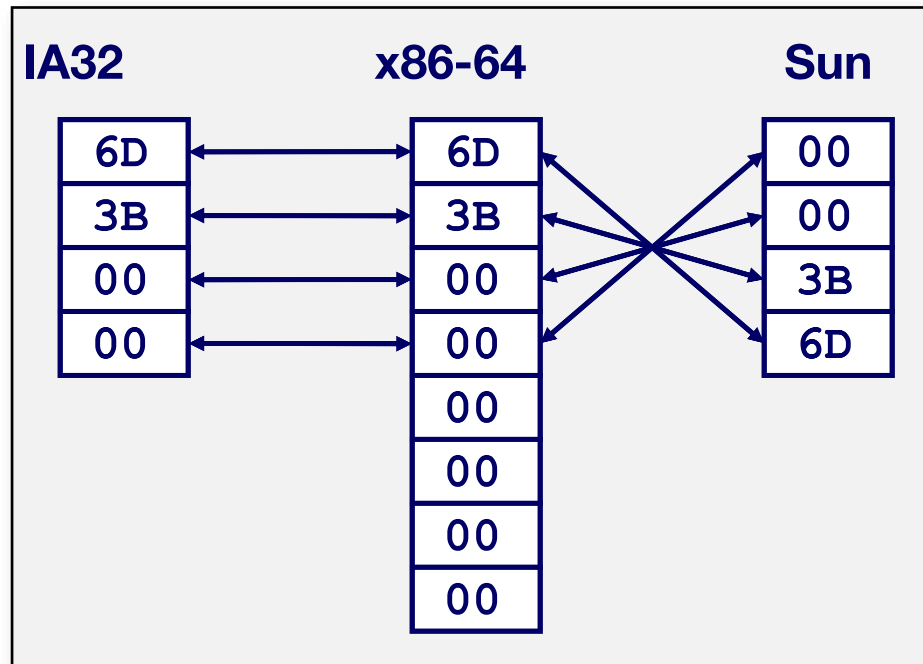


`int B = -15213;`



Two’s complement representation

`long int C = 15213;`



Examining Data Representations

- **Code to Print Byte Representation of Data**
  - Casting pointer to unsigned char * allows treatment as a byte array

```
typedef unsigned char *pointer;

void show_bytes(pointer start, size_t len){
  size_t i;
  for (i = 0; i < len; i++)
    printf(”%p\t0x%.2x\n",start+i, start[i]);
  printf("\n");
}
```

Printf directives:

%p: Print pointer

%x: Print Hexadecimal

show_bytes Execution Example

```
int a = 15213;
printf("int a = 15213;\n");
show_bytes((pointer) &a, sizeof(int));
```

Result (Linux x86-64):

```
int a = 15213;
0x7fffb7f71dbc    6d
0x7fffb7f71dbd    3b
0x7fffb7f71dbe    00
0x7fffb7f71dbf    00
```

Representing Pointers

```
int B = -15213;
int *P = &B;
```

Sun	IA32	x86-64
EF	AC	3C
FF	28	1B
FB	F5	FE
2C	FF	82
		FD
		7F
		00
		00

- Different compilers & machines assign different locations to objects
- Even get different results each time run program

Representing Strings

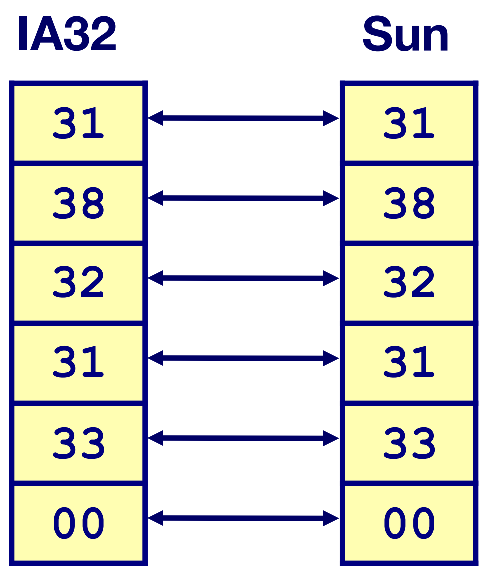
```
char S[6] = "18213";
```

Strings in C

- Represented by array of characters
- Each character encoded in ASCII format
  - Standard 7-bit encoding of character set
  - Character "0" has code 0x30
    - Digit *i* has code 0x30+*i*
- String should be null-terminated
  - Final character = 0

Compatibility

- Byte ordering not an issue



A note about x86 machine code

- **x86 machine code is a sequence of *bytes***
  - Grouped into variable-length instructions, which look like strings…
  - But they contain embedded little-endian numbers…
- **Example Fragment**

Address	Instruction Code	Assembly Rendition
`8048365:`	`5b`	`pop %ebx`
`8048366:`	`81 c3 ab 12 00 00`	`add $0x12ab,%ebx`
`804836c:`	`83 bb 28 00 00 00 00`	`cmpl $0x0,0x28(%ebx)`

- **Deciphering Numbers**
  - Value: `0x12ab`
  - Pad to 32 bits: `0x000012ab`
  - Split into bytes: `00 00 12 ab`
  - Reverse: `ab 12 00 00`