15-213 Recitation Shell Lab

Your TAs Friday, November 8th

Reminders

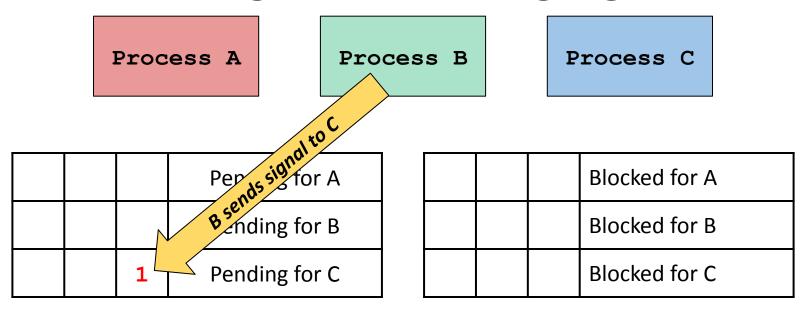
- **tshlab** released, due on *November 19th*.
- Code Reviews:
 - Watch your inbox for a malloc Final code review email!

Agenda

- Signals
- File I/O

Signals

Recall: Sending and Receiving Signals



- Pending signals represented by a single bit, one for each kind of signal.
- Kernel computes pnb (pending and not blocked) to determine which signals can be delivered.
- If we don't block them, signals can interrupt our program at any time!

Example: "Counting" with Signals

```
volatile int counter = 0;
void handler(int sig) { counter++; }
int main(void) {
    signal(SIGCHLD, handler);
    for (int i = 0; i < 10; i++) {
         if (fork() == 0) { exit(0); }
    }
    while (counter < 10) { // Do nothing }</pre>
    printf("Terminated :-)\n");
```

- What happens when we run this program? Will it terminate?
 - It might not, since signals can coalesce.

Example: "Counting" with Signals

```
volatile int counter = 0;
                                                                     Problem:
void handler(int sig) { counter++; }
                                                                   Signals Coalesce
int main(void) {
    signal(SIGCHLD, handler);
    for (int i = 0; i < 10; i++) {
         if (fork() == 0) { exit(0); }
                                                                     Problem:
                                                                  Tight loop used to
    while (counter < 10) { // Do nothing }
                                                                   wait for signals
    printf("Terminated :-)\n");
```

- What are some problems with the above code?
- How can we fix them?

Problem: Signals Coalesce

- # Times Signal Handler Called != # Times Signal was Sent
- What can we do instead?

```
void handler(int sig) { counter++; }

void handler(int sig) {
   pid_t pid;

while ((pid = waitpid(-1, NULL, WNOHANG)) > 0) {
      counter++;
   }
}
```

counter++ is not atomic. Why aren't there race conditions?

Problem: Efficiently waiting for Signals

Tight loop is inefficient, and is forbidden in Shell Lab.

```
while (counter < 10) { // Do nothing }

Problem:
Tight Loop
```

Use sigsuspend instead!

```
int sigsuspend(const sigset_t *mask);
```

- Temporarily installs mask, then pauses until a signal is received.
- Atomic version of:

```
sigprocmask(SIG_SETMASK, &mask, &prev);
pause();
sigprocmask(SIG_SETMASK, &prev, NULL);
```

Recall: Blocking Signals

Textbook: p764

Allows us to control when our program receives signals.

```
int <u>sigprocmask</u>(int how, sigset_t *mask, sigset_t *prev_mask);
```

Recommended approach:

```
sigset_t mask, prev;
sigemptyset(&mask);
sigaddset(&mask, SIGINT);
sigprocmask(SIG_BLOCK, &mask, &prev);
// ...
sigprocmask(SIG_SETMASK, &prev, NULL);
```

- Don't use SIG_UNBLOCK
 - Don't want to unblock a signal if it was already blocked.
 - Allows procedure to be nested multiple times.

Recall: Signal Handlers

- Parent process sends SIGINT to child process. What is the behavior of the child?
 - No handler specified: default action.
 - Can catch the signal with a signal handler.

```
void sigchld handler(int sig)
                                                                   Save and restore
    int olderrno = errno;
                                                                       errno
    sigset t mask all, prev all;
    pid t pid;
    sigfillset(&mask all);
                                                                   Temporarily block
    while ((pid = waitpid(-1, NULL, 0)) > 0) { /* Reap */}
         sigprocmask(SIG BLOCK, &mask all, &prev all);
                                                                   signals to protect
         deletejob(pid); /* Delete from job list */
                                                                     shared data
         sigprocmask(SIG SETMASK, &prev all, NULL);
                                                                       Call only
    if (pid != 0 && errno != ECHILD)
                                                                   async-signal-safe
         sio eprintf("waitpid error");
                                                                      functions
    errno = olderrno;
```

Example

```
int main(void) {
    char *tgt = "child";
    sigset t mask, old mask;
    sigemptyset(&mask);
    sigaddset(&mask, SIGINT);
    sigprocmask(SIG BLOCK, &mask, &old mask); // Block
    pid t pid = fork();
    if (pid == 0) {
        pid = getppid(); // Get parent pid
        tqt = "parent";
    kill(pid, SIGINT);
    sigprocmask(SIG SETMASK, &old mask, NULL); // Unblock
    printf("Sent SIGINT to %s:%d\n", tgt, pid);
    exit(0);
```

- How many different lines could be printed?
 - 0 or 1 line. Parent and child try to terminate each other.

File I/O

I/O Functions: open and close

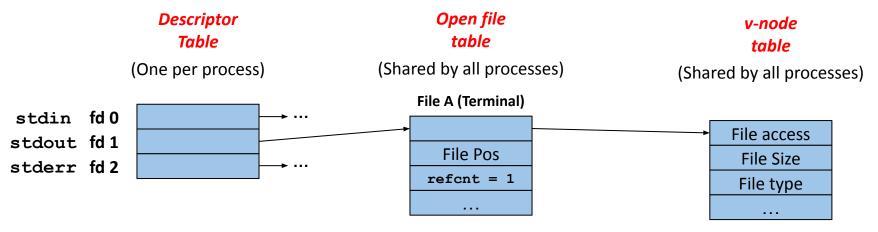
Textbook: p893

int open(const char *pathname, int flags, mode_t mode);

- pathname path to file
- flags:
 - File Creation: O CREAT, O TRUNC, etc.
 - Access modes: O RDONLY, O WRONLY, O RDWR
- mode
 - Specifies who else can read/write the new file.
 - Unless you have reasons not to, use **DEF_MODE** from textbook.

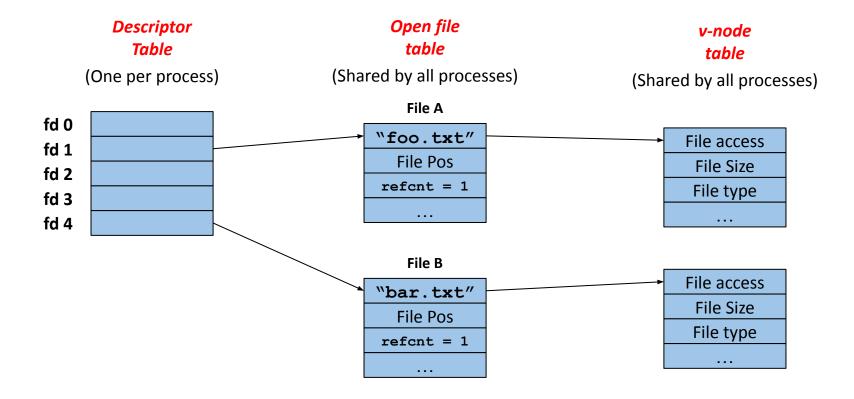
```
int close(int fd);
```

Std File Descriptors

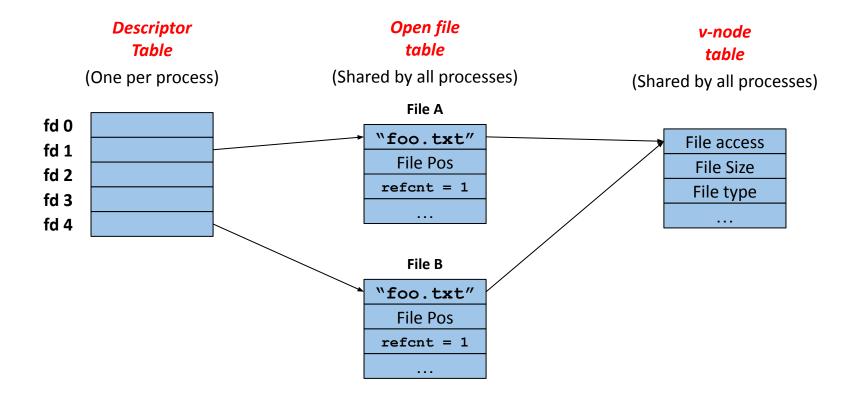


- stdin, stdout, and stderr are set up automatically.
- Closed by normal termination or exit().

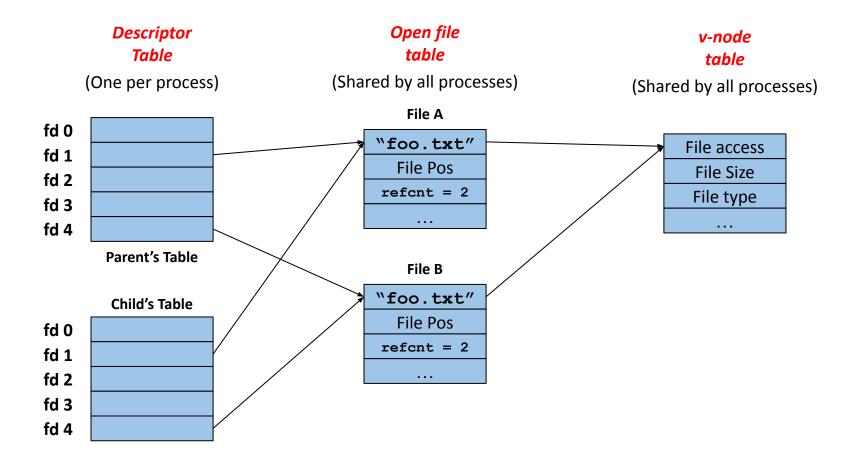
File Descriptors (File A != File B)



File Descriptors (File A == File B)



File Descriptors: What Happens After fork()?



Example: I/O and fork()

```
int main(int argc, char** argv)
{
    int i;
    for (i = 0; i < 4; i++)
    {
        int fd = open("foo", O_RDONLY);
        pid_t pid = fork();
        if (pid == 0)
        {
            int ofd = open("bar", O_RDONLY);
            execve(...);
        }
    }
}
// How many file descriptors are open in the parent?</pre>
```

- How many file descriptors are open in the parent process at the indicated point?
- How many does each child have open at the call to

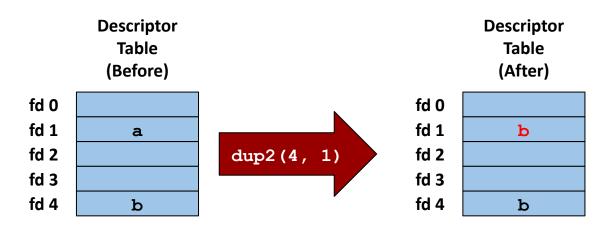
```
execve()?
```

I/O Redirection: dup2()

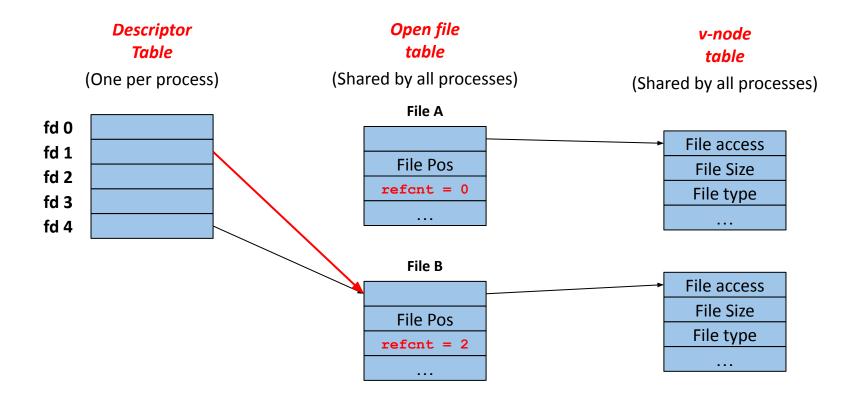
Textbook: p909

int dup2 (int oldfd, int newfd);

- Copies descriptor table entry oldfd to descriptor table entry newfd, overwriting existing entry for newfd.
- If newfd is open, it is closed.



File Descriptors after dup2 (4,1)



Initial situation shown on "File A!= File B" slide.

Example: Redirecting I/O

```
int main(int argc, char** argv)
   int i, fd;
   fd = open("foo", O WRONLY);
   dup2(fd, STDOUT FILENO);
   // Point A
   close(fd);
   // Point B
```

- How many open file table entries are there at Point A?
- How many open file table entries are there at Point B?

Example: Redirecting I/O + fork()

```
int main(int argc, char** argv) {
    int i;
    for (i = 0; i < 4; i++)
        int fd = open("foo", O RDONLY);
        pid t pid = fork();
        if (pid == 0)
             int ofd = open("bar", O WRONLY);
             dup2(fd, STDIN FILENO);
             dup2(ofd, STDOUT FILENO);
            execve(...);
    // How many file descriptors are open in the parent?
```

How many file descriptors are open in the parent?

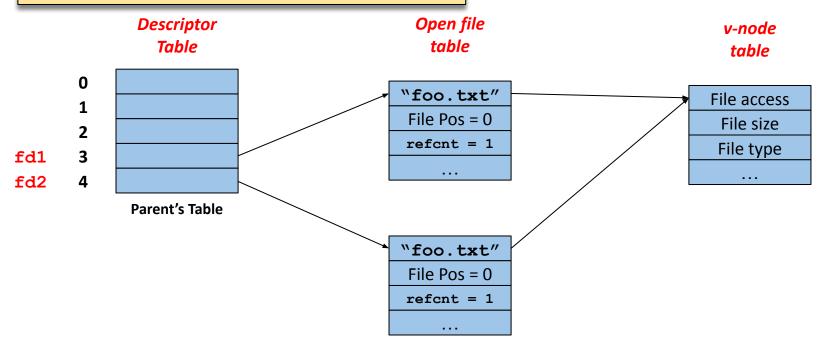
```
int main(int argc, char *argv[]) {
    int fd1 = open("foo.txt", O RDONLY);
    int fd2 = open("foo.txt", O RDONLY);
    read and print one(fd1);
    read and print one(fd2);
    if(!fork()) {
         read and print one(fd2);
         read and print one(fd2);
         close (fd2);
         fd2 = dup(fd1);
         read and print one(fd2);
    } else {
         wait(NULL);
         read and print one(fd1);
         read and print one(fd2);
         printf("\n");
    close (fd1);
    close (fd2);
    return 0;
```

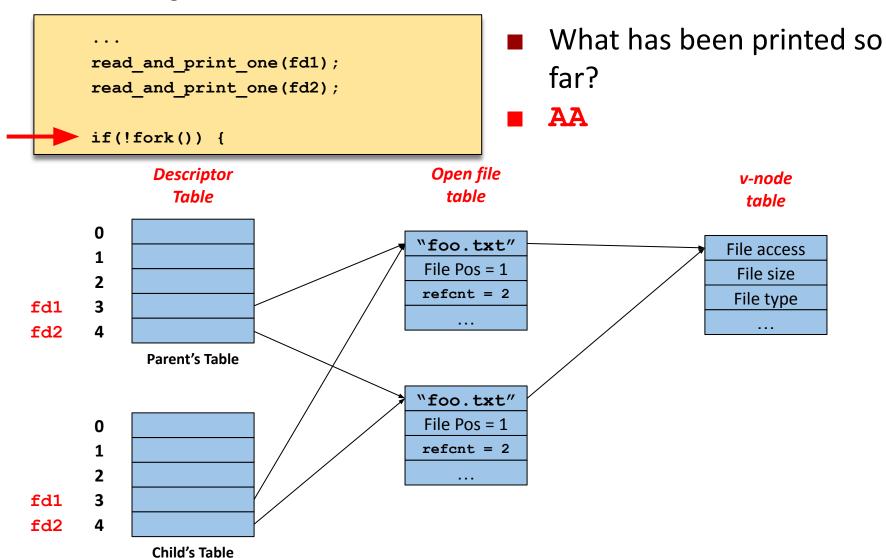
```
void read_and_print_one(int fd) {
    char c;
    read(fd, &c, 1);
    printf("%c", c);
    fflush(stdout);
}
```

- Suppose the contents of foo.txt are "ABCDEFG".
- What is the output of this program?

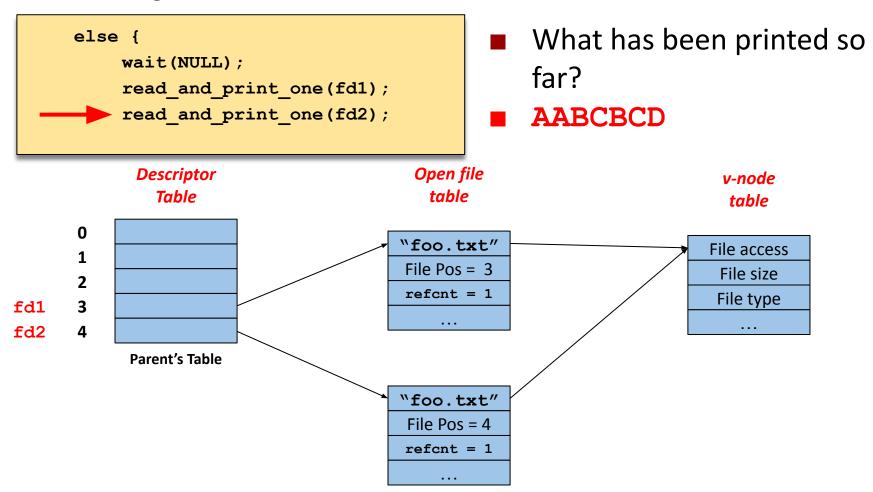
```
int main(int argc, char *argv[]) {
   int fd1 = open("foo.txt", O_RDONLY);
   int fd2 = open("foo.txt", O_RDONLY);
   ...
```

What does the diagram look like at this point?





```
if(!fork()) {
                                                           What has been printed so
            read and print one(fd2);
                                                            far?
            read and print one(fd2);
            close(fd2);
                                                            AABCB
            fd2 = dup(fd1);
            read and print one(fd2);
              Descriptor
                                              Open file
                                                                                  v-node
                                                table
                Table
                                                                                   table
       0
                                             "foo.txt"
                                                                                 File access
       1
                                              File Pos = 2
                                                                                  File size
       2
                                             refcnt = 3
                                                                                  File type
fd1
fd2
       4
             Parent's Table
                                             "foo.txt"
                                              File Pos = 3
       0
                                             refcnt = 1
       1
fd1
fd2
       4
              Child's Table
```



Wrapping Up

- **tshlab** is due on *November 19th*.
- Code Reviews:
 - Watch your inbox for a malloc Final code review email!
- Getting Started:
 - Lecture Slides
 - Textbook (Chapter 8)
 - man pages
 - Develop *Incrementally*: see roadmap in write-up.

The End

Shell Lab Reference Sheet

Page numbers refer to CS:APP3e

Loading Programs

<u>execve</u> () [p750]

Waiting and Reaping

- waitpid() [p743]
- sigsuspend() [p781]
- Exit status macros(WIFEXITED, etc.) [p745]

Blocking/Unblocking Signals

■ sigprocmask() [p764]

Sending Signals

■ <u>kill</u>() [p761]

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- open(), close() [p893]
- dup2 () [p909]