# Sparsity, Randomness and Compressed Sensing

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## **Sparsity**

## Why Sparsity

- Natural data and signals exhibit structure
- Sparsity often captures that structure
- Very general signal model
- Computationally tractable
- Wide range of applications in signal acquisition, processing, and transmission

Signal Representations

## Signal example: Images

- 2-D function f
- Idealized view

```
f \in \text{some function} \\ \text{space defined} \\ \text{over } [0,1] \times [0,1]
```



## Signal example: Images

- 2-D function f
- Idealized view

$$f \in \text{some function} \\ \text{space defined} \\ \text{over } [0,1] \times [0,1]$$



In practice

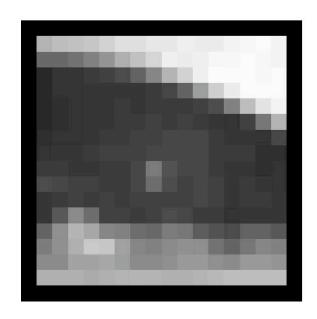
$$f \in \mathbb{R}^{N \times N}$$

ie: an  $N \times N$  matrix

## Signal example: Images

- 2-D function f
- Idealized view

$$f \in \text{some function}$$
 space defined over  $[0,1] \times [0,1]$ 

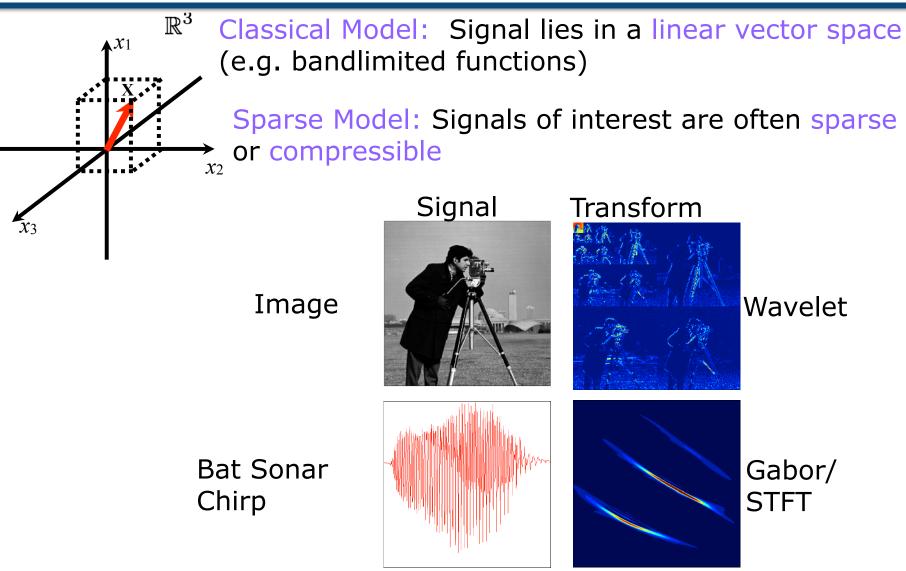


In practice

$$f \in \mathbb{R}^{N \times N}$$

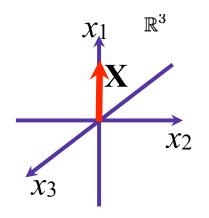
ie: an  $N \times N$  matrix (pixel average)

## Signal Models

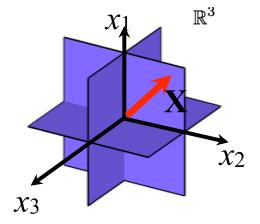


i.e., very few large coefficients, many close to zero.

## Sparse Signal Models

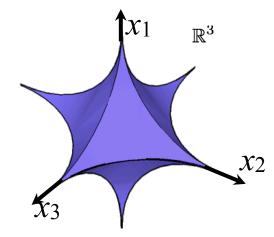


1-sparse



2-sparse

Sparse signals have few non-zero coefficients



Compressible signals have few significant coefficients.

The coefficients decay as a power law.

Compressible ( $\ell_p$  ball, p < 1)

**Sparse Approximation** 

## Computational Harmonic Analysis

Representation

$$f = \sum_{k} a_k b_k$$
coefficients basis, frame

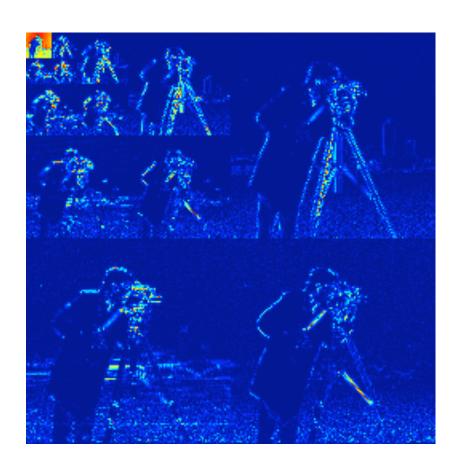
• Analysis: study f through structure of  $\{a_k\}$   $\{b_k\}$  should extract features of interest

• Approximation:  $\widehat{f}_N$  uses just a few terms N exploit sparsity of  $\{a_k\}$ 

## Wavelet Transform Sparsity



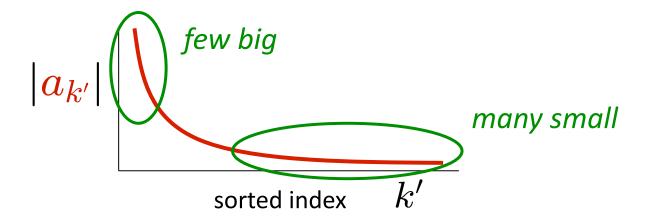
$$f = \sum_{k} a_k b_k$$



• Many 
$$a_k \approx 0$$
 (blue)

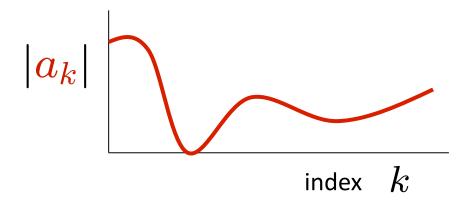
## Sparseness ⇒ Approximation

$$f = \sum_{k} a_k b_k$$



## Linear Approximation

$$f = \sum_{k} a_k b_k$$



## Linear Approximation

$$f = \sum_{k} a_k b_k$$

• *N*-term approximation: use "first"  $a_k$ 

$$\widetilde{f}_N := \sum_{k=1}^N a_k \, \mathbf{b}_k$$

$$|a_k|$$

$$N \quad \text{index } k$$

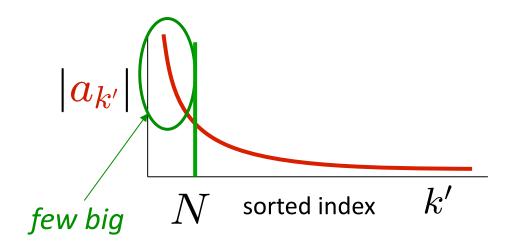
## Nonlinear Approximation

$$f = \sum_{k} a_k b_k$$

• N-term approximation: use largest  $a_k$  independently

$$\widehat{f}_N := \sum_{k'=1}^N a_{k'} \mathbf{b}_{k'}$$

Greedy / thresholding



## **Error Approximation Rates**

$$f = \sum_{k} a_{k} b_{k}$$

$$\widehat{f}_{N} = \sum_{k'=1}^{N} a_{k'} b_{k'}$$

$$\|f-\widehat{f}_N\|_2^2 < CN^{-\alpha}$$
 as  $N \to \infty$ 

- Optimize asymptotic *error decay rate*  $\, lpha \,$
- Nonlinear approximation works better than linear

## Compression is Approximation

Lossy compression of an image creates an approximation

$$f = \sum_k a_k \, \mathbf{b}_k$$
 $\uparrow$ 
 $\uparrow$ 
coefficients basis, frame

 $quantize \mid to \, R \, total \, bits$ 
 $\widehat{f}_R = \sum_k a_k^q \, \mathbf{b}_k$ 

## Sparse approximation ≠ Compression

 Sparse approximation chooses coefficients but does not quantize or worry about their locations

$$f = \sum_{k} a_{k} b_{k}$$

$$f_{N} = \sum_{k'=1}^{N} a_{k'} b_{k'}$$

## Location, Location

• Nonlinear approximation selects N largest  $a_k$  to minimize error (easy – threshold)

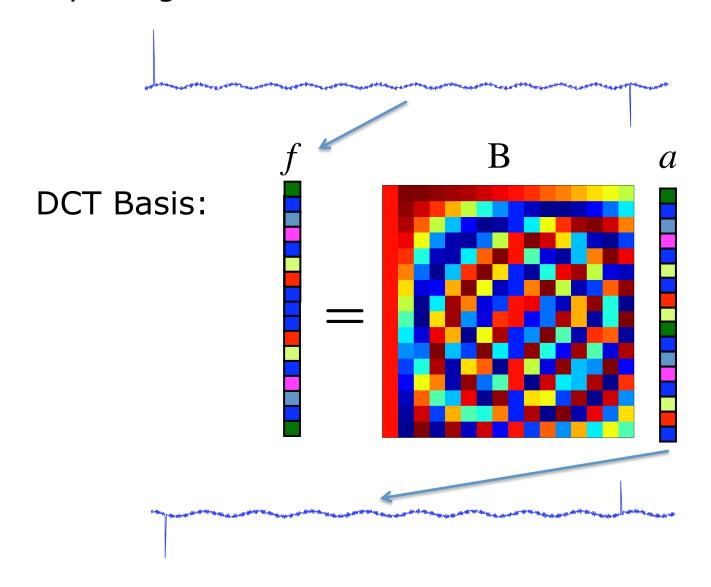
 Compression algorithm must encode both a set of ak and their locations (harder)



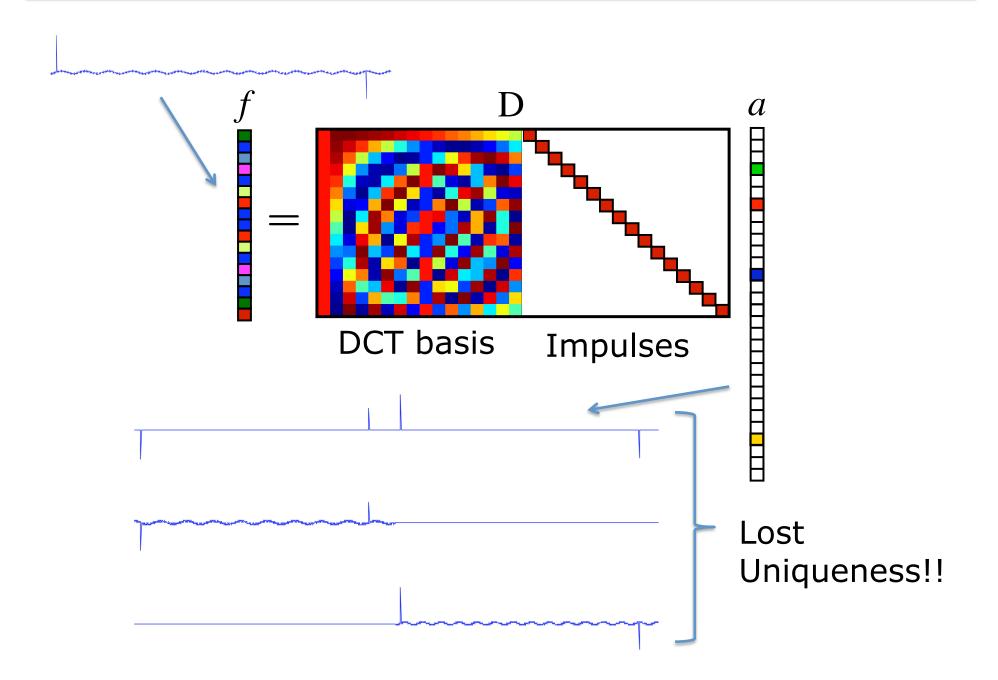
**Exposing Sparsity** 

## Spikes and Sinusoids example

Example Signal Model: Sinusoidal with a few spikes.

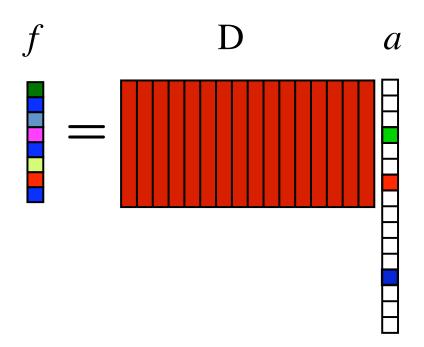


## Spikes and Sinusoids Dictionary



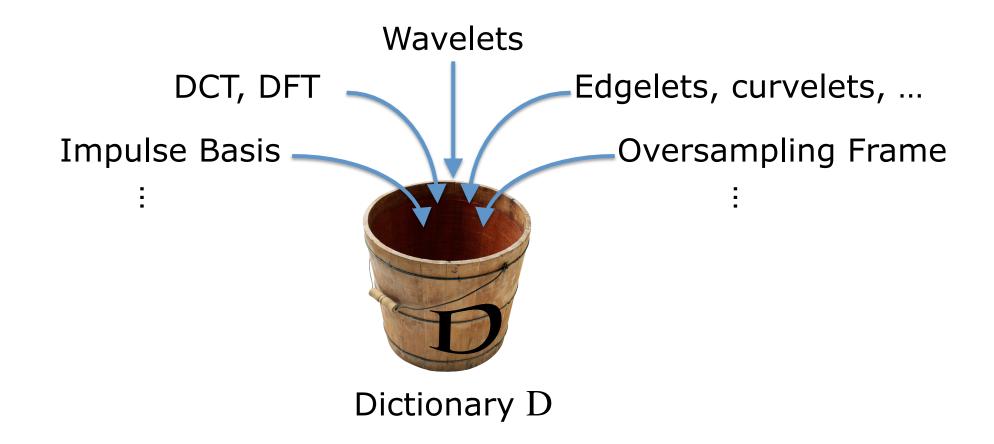
## Overcomplete Dictionaries

**Strategy: Improve** sparse approximation by constructing a large **dictionary.** 



How do we **design** a dictionary?

## Dictionary Design

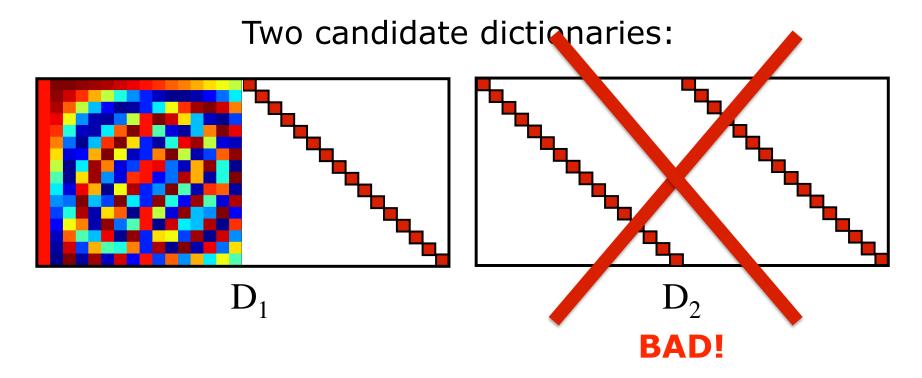


Can we just throw in the bucket everything we know?

## Dictionary Design Considerations

- Dictionary Size:
  - Computation and storage increases with size
- Fast Transforms:
  - FFT, DCT, FWT, etc. dramatically decrease computation and storage
- Coherence:
  - Similarity in elements makes solution harder

## **Dictionary Coherence**



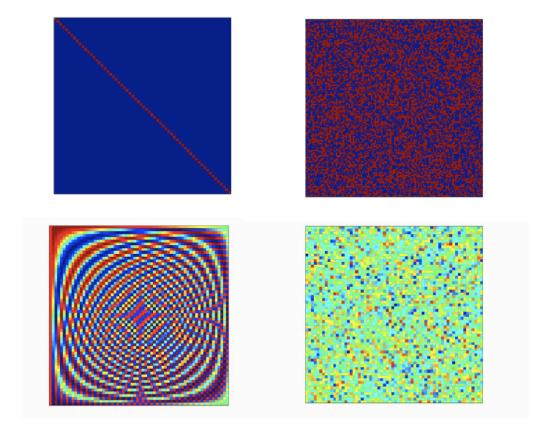
Intuition:  $D_2$  has too many similar elements. It is very coherent.

**Coherence** (similarity) between elements:  $|\langle d_1, d_2 \rangle|$ 

Dictionary coherence:  $\mu = \max_{i,j} |\langle d_i, d_j \rangle|$ 

## **Incoherent Bases**

- "Mix" well the signal components
  - Impulses and Fourier Basis
  - Anything and Random Gaussian
  - Anything and Random 0-1 basis



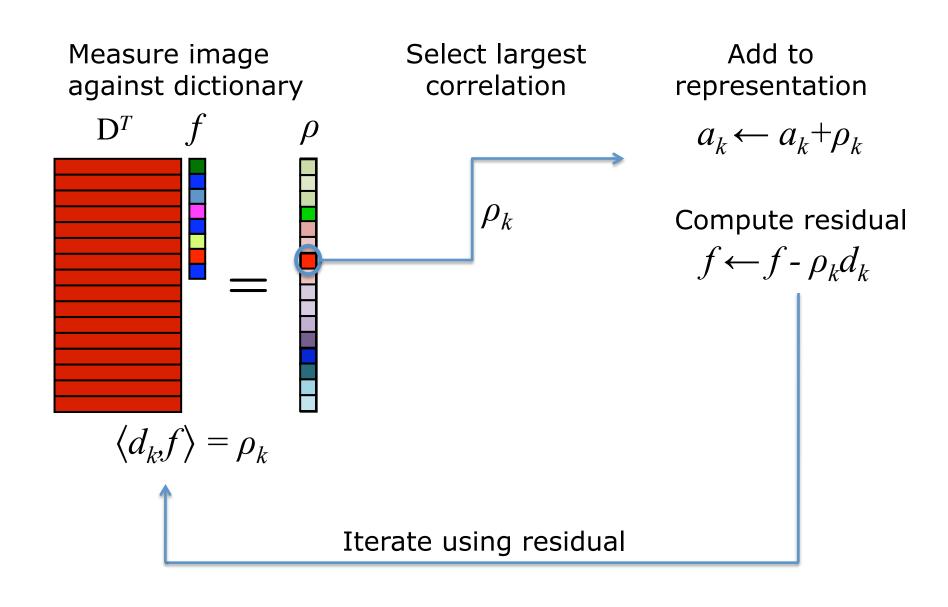
## Computing Sparse Representations

## Thresholding

Zero out Compute set of coefficients small ones  $a=D^{\dagger}f$ 

Computationally efficient Good for small and very incoherent dictionaries

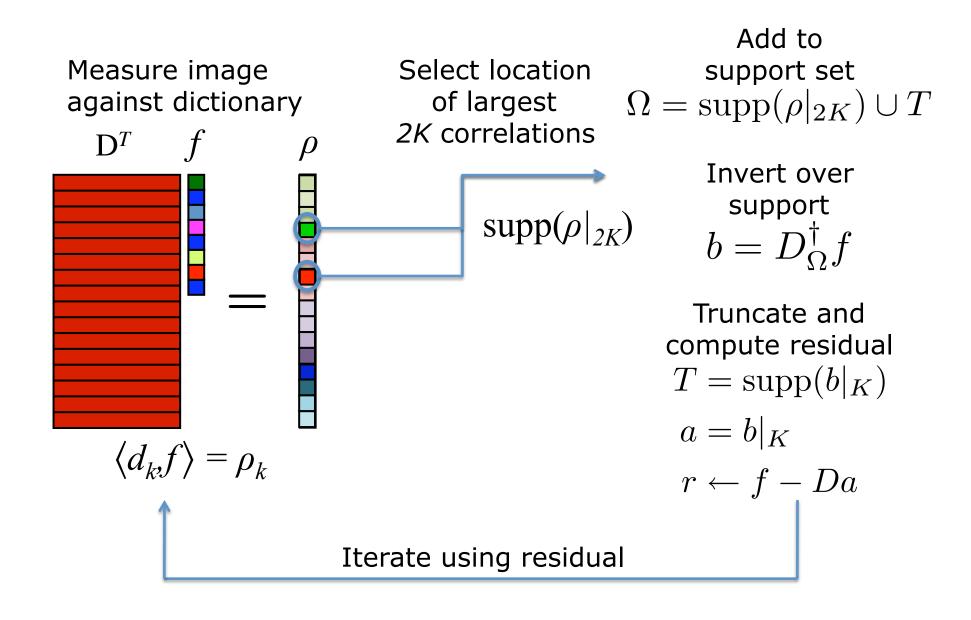
## Matching Pursuit



## Greedy Pursuits Family

- Several Variations of MP:
   OMP, StOMP, ROMP, CoSaMP, Tree MP, ...
   (You can create an AndrewMP if you work on it...)
- Some have provable guarantees
- Some improve dictionary search
- Some improve coefficient selection

## CoSaMP (Compressive Sampling MP)



## Optimization (Basis Pursuit)

### **Sparse approximation:**

Minimize non-zeros in representation s.t.: representation is close to signal

$$\min \|a\|_{\mathfrak{d}} \text{ s.t. } f \approx \mathrm{D}a$$

Number of non-zeros (sparsity measure)

Data Fidelity (approximation quality)

Combinatorial complexity. Very hard problem!

## Optimization (Basis Pursuit)

### **Sparse approximation:**

Minimize non-zeros in representation s.t.: representation is close to signal

min 
$$\|a\|_{\mathbf{X}}$$
 s.t.  $f \approx \mathrm{D}a$ 

Convex Relaxation

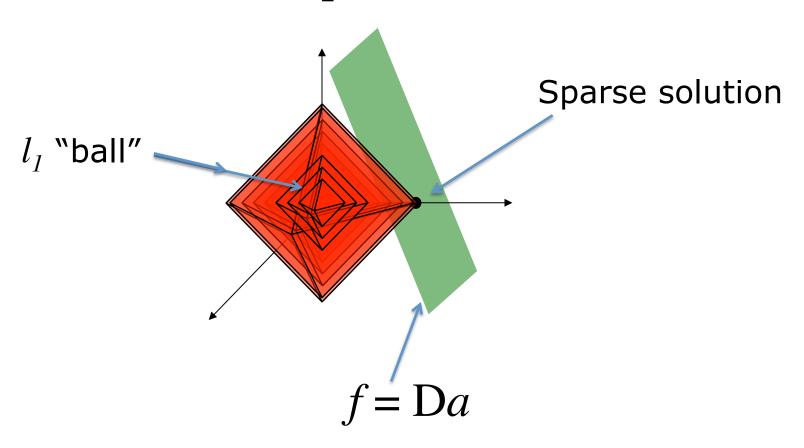
min  $\|a\|_1$  s.t.  $f \approx \mathrm{D}a$ 

Ploynomial complexity.

Solved using linear programming.

## Why $l_1$ relaxation works

min 
$$||a||_1$$
 s.t.  $f \approx Da$ 



#### **Basis Pursuits**

- Have provable guarantees
  - Finds sparsest solution for incoherent dictionaries
- Several variants in formulation:

BPDN, LASSO, Dantzig selector, ...

Variations on fidelity term and relaxation choice

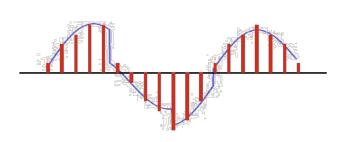
Several fast algorithms:

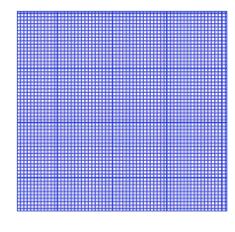
FPC, GPSR, SPGL, ...

Compressed Sensing:
Sensing, Sampling and
Data Processing

#### **Data Acquisition**

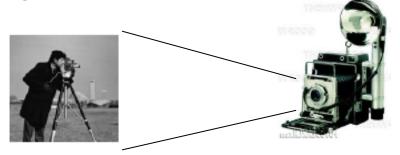
- Usual acquisition methods sample signals uniformly
  - Time: A/D with microphones, geophones, hydrophones.
  - Space: CCD cameras, sensor arrays.
- Foundation: Nyquist/Shannon sampling theory
  - Sample at twice the signal bandwidth.
  - Generally a projection to a complete basis that spans the signal space.





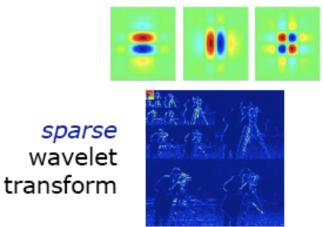
#### Data Processing and Transmission

- Data processing steps:
  - Sample Densely



Signal x, *N* coefficients

Transform to an informative domain (Fourier, Wavelet)



*K*<<*N* significant coefficients

Process/Compress/Transmit

Sets small coefficients to zero (sparsification)

### Sparsity Model

• Signals can usually be **compressed** in some basis

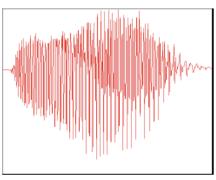
 $N \ {
m pixels}$ 

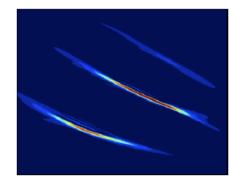




 $K \ll N$  large wavelet coefficients

N wideband signal samples

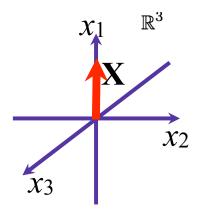




 $K \ll N$  large Gabor coefficients

Sparsity: good prior in picking from a lot of candidates

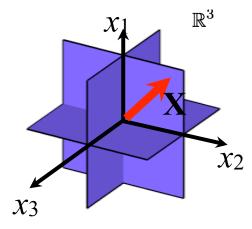
### Compressive Sensing Principles



If a signal is sparse, do not waste effort sampling the empty space.

1-sparse

Instead, use fewer samples and allow ambiguity.

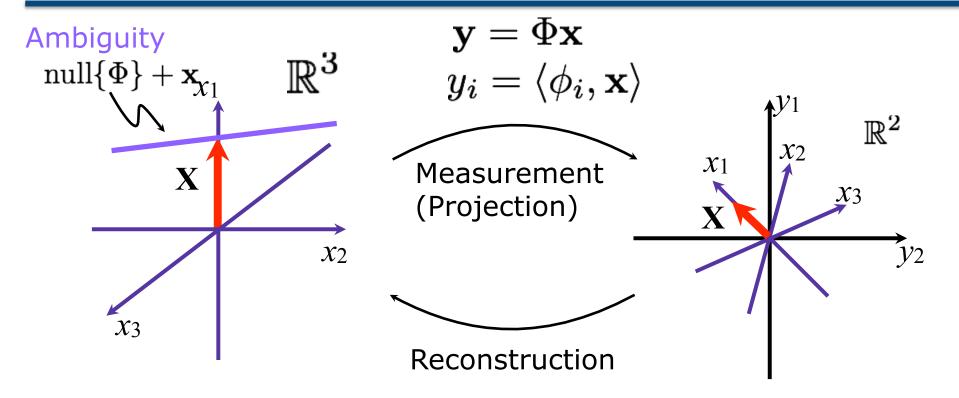


Use the sparsity model to reconstruct and uniquely resolve the ambiguity.

2-sparse

Measuring Sparse Signals

#### Compressive Measurements



 $\Phi$  has rank  $M \ll N$ 

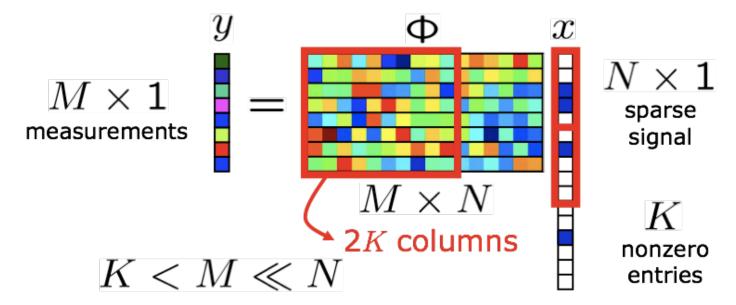
*K* = Signal sparsity

N = Signal dimensionality M = Number of measurements (dimensionality of y)

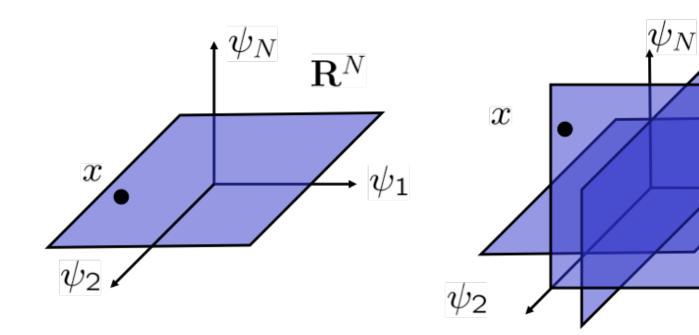
$$N \gg M \gtrsim K$$

#### One Simple Question

- When is it possible to distinguish K-sparse signals?
  - require  $\Phi x_1 \neq \Phi x_2$  for all K-sparse  $x_1 \neq x_2$
- - otherwise there exist K-sparse  $x_1, x_2$  s.t.  $\Phi(x_1-x_2)=0$
- Sufficient: Gaussian Φ with 2K rows



# Geometry of Sparse Signal Sets



Linear

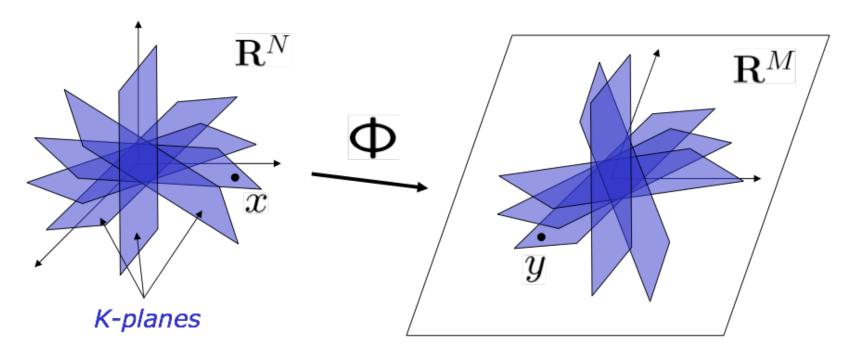
K-plane

Sparse, Nonlinear

 $\mathbf{R}^N$ 

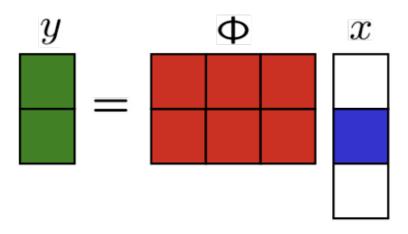
Union of K-planes

# Geometry: Embedding in R<sup>M</sup>



- Φ(K-plane) = K-plane in general
- M ≥ 2K measurements
  - necessary for injectivity
  - sufficient for injectivity when  $\Phi$  Gaussian
  - but not enough for efficient, robust recovery
- See also FROI [Vetterli et al., Lu and Do]

## Illustrative Example

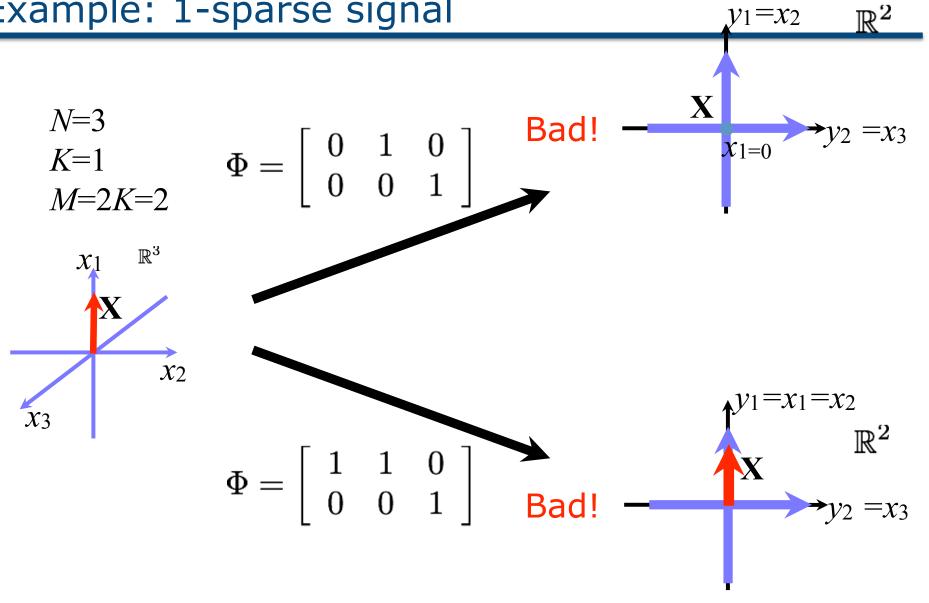


N = 3: signal length

K = 1: sparsity

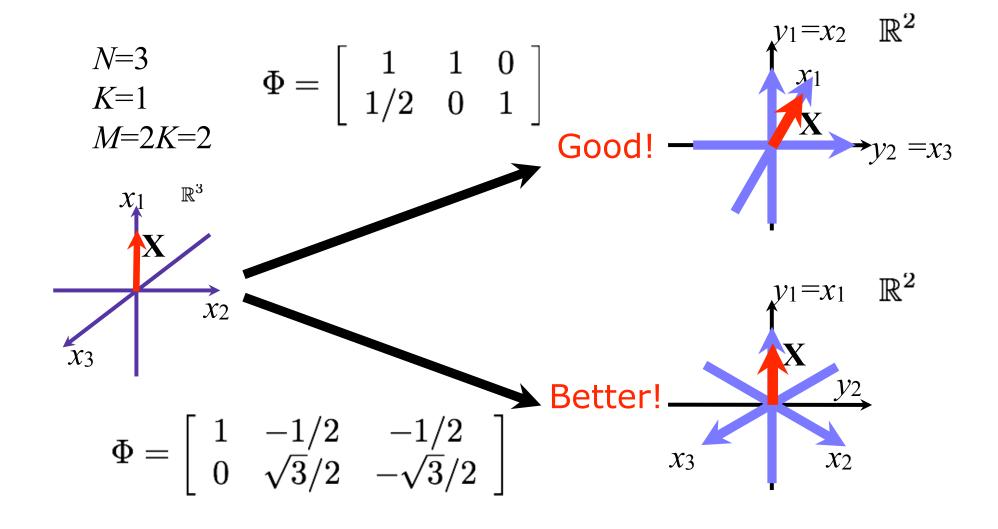
M = 2K = 2: measurements

# Example: 1-sparse signal



 $y_1 = x_2$ 

# Example: 1-sparse signal

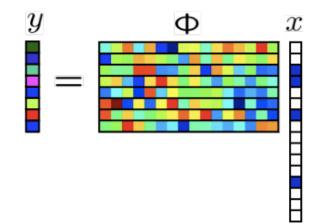


#### Restricted Isometry Property

[Candès, Romberg, Tao]

• Measurement matrix  $\Phi$  has **RIP of order** K if

$$(1 - \delta_K) \le \frac{\|\Phi x\|_2^2}{\|x\|_2^2} \le (1 + \delta_K)$$



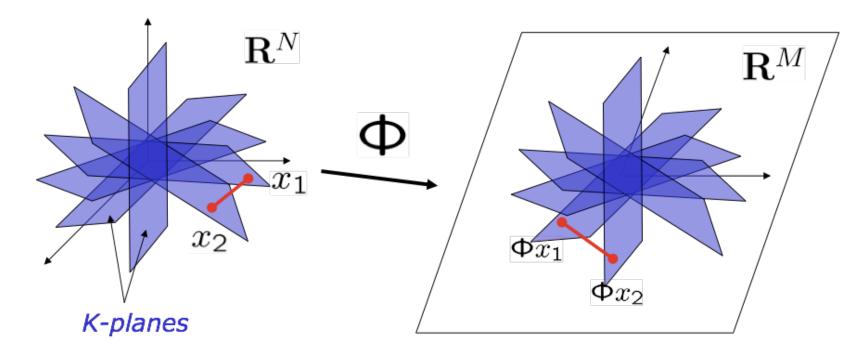
for all K-sparse signals x.

- Does *not* hold for K>M; may hold for smaller K.
- Implications: tractable, stable, robust recovery

#### RIP as a "Stable" Embedding

• RIP of order 2K implies: for all K-sparse  $x_1$  and  $x_2$ 

$$(1 - \delta_{2K}) \le \frac{\|\Phi x_1 - \Phi x_2\|_2^2}{\|x_1 - x_2\|_2^2} \le (1 + \delta_{2K})$$



(if  $\delta_{2K}$  < 1 have injectivity; smaller  $\delta_{2K}$  more stable)

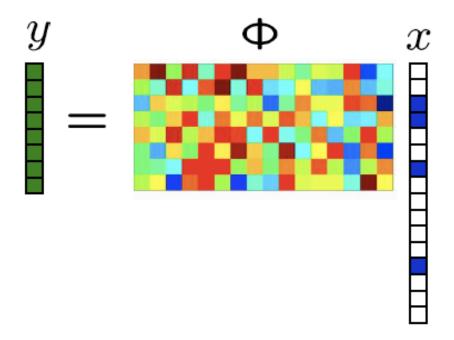
#### Verifying RIP

#### **How Many Measurements?**

- Want RIP of order 2K (say) to hold for MxN Φ
  - difficult to verify for a given Φ
  - requires checking eigenvalues of each submatrix
- Prove random Φ will work
  - iid Gaussian entries
  - iid Bernoulli entries (+/- 1)
  - iid subgaussian entries
  - random Fourier ensemble
  - random subset of incoherent dictionary
- In each case,  $M = O(K \log N)$  suffices
  - with very high probability, usually 1-O(e-CN)
  - slight variations on log term

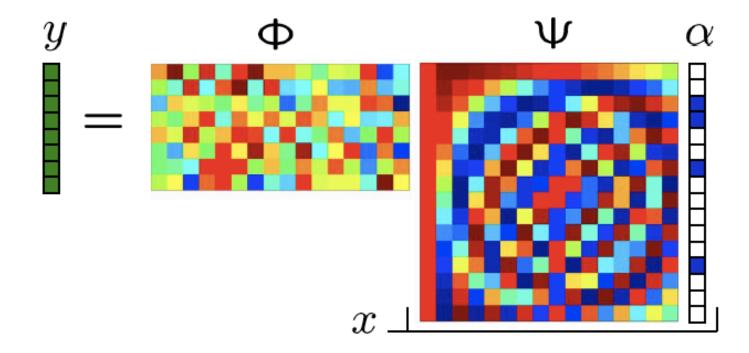
### Universality Property

- Gaussian white noise basis is incoherent with any fixed orthonormal basis (with high probability)
- Signal sparse in time domain:  $\Phi = I$



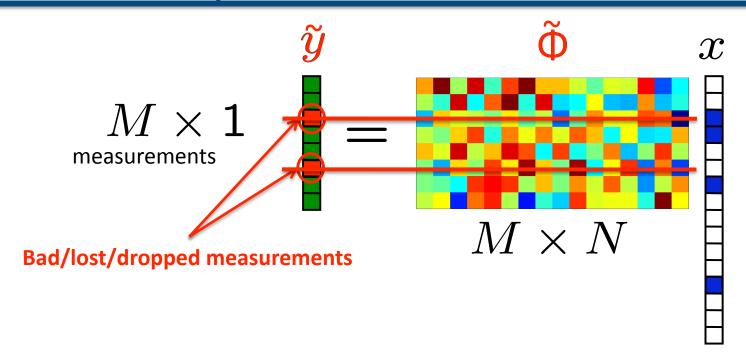
### Universality Property

- Gaussian white noise basis is incoherent with any fixed orthonormal basis (with high probability)
- Signal sparse in frequency domain:  $\Psi = idct$



ullet Product  $\Phi\Psi$  remains Gaussian white noise

# Democracy



- Measurements are democratic [Davenport, Laska, Boufounos, Baraniuk]
  - They are all equally important
  - We can loose some arbitrarily, (i.e. an adversary can choose which ones)
- ullet The  $ilde{\Phi}$  still satisfies RIP (as long as we don't drop too many)

#### Reconstruction

#### Requirements for Reconstruction

- Let  $x_1$ ,  $x_2$  be K-sparse signals (I.e.  $x_1$ - $x_2$  is 2K-sparse):
- Mapping  $y = \Phi x$  is **invertible** for K-sparse signals:

$$\Phi(x_1-x_2)\neq 0$$
 if  $x_1\neq x_2$ 

Mapping is robust for K-sparse signals:

$$||\Phi(x_1-x_2)||_2 \approx ||x_1-x_2||_2$$

- Restricted Isometry Property (RIP):
  - $\Phi$  preserves distance when projecting K-sparse signals
- Guarantees there exists a unique K-sparse signal explains the measurements, and is robust to noise.

#### Reconstruction Ambiguity

Solution should be consistent with measurements

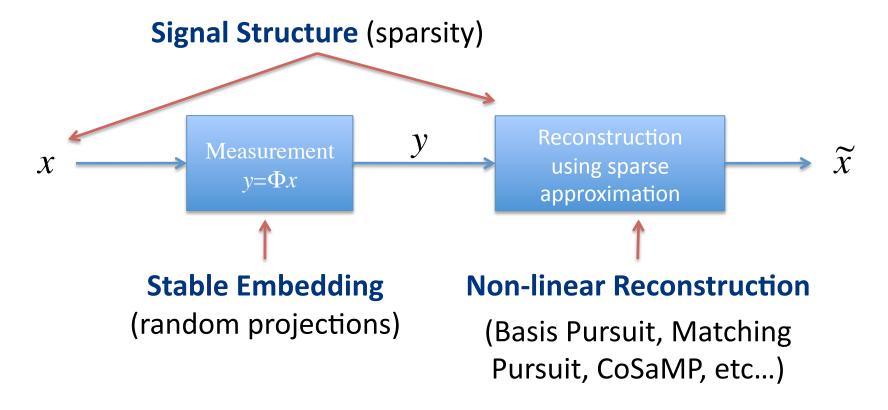
$$\hat{\mathbf{x}}$$
 s.t.  $\mathbf{y} = \Phi \hat{\mathbf{x}}$  or  $\mathbf{y} \approx \Phi \hat{\mathbf{x}}$ 

- Projections imply that an infinite number of solutions are consistent!
- Classical approach: use the pseudoinverse (minimize  $l_2$  norm)
- Compressive sensing approach: pick the sparsest.
- RIP guarantee: sparsest solution unique and reconstructs the signal.

Becomes a sparse approximation problem!

Putting everything together

### Compressed Sensing Coming Together

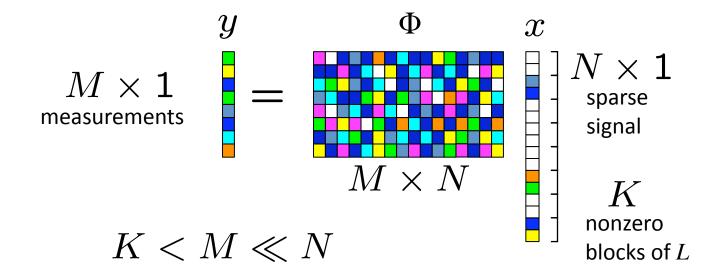


- Signal model: Provides prior information; allows undersampling
- Randomness: Provides robustness/stability; makes proofs easier
- Non-linear reconstruction: Incorporates information through computation

Beyond: Extensions,
Connections, Generalizations

# **Sparsity Models**

### **Block Sparsity**

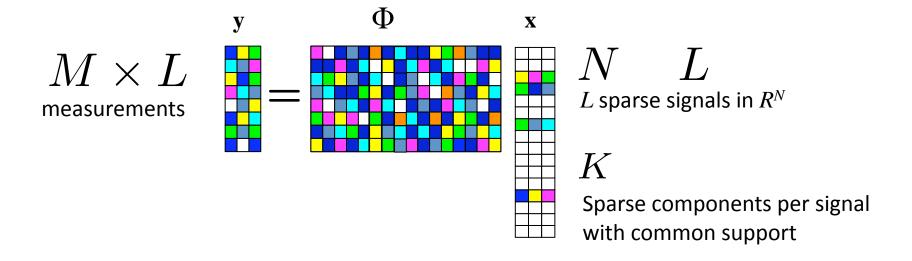


Mixed  $l_1/l_2$  norm—sum of  $l_2$  norms:  $\sum_i \|\mathbf{x}_{B_i}\|_2$ 

Basis pursuit becomes:  $\min_{\mathbf{x}} \sum_{i} \|\mathbf{x}_{B_i}\|_2$  s.t.  $y \approx \Phi x$ 

Blocks are not allowed to overlap

### Joint Sparsity



Mixed 
$$l_1/l_2$$
 norm—sum of  $l_2$  norms:  $\sum_i \|\mathbf{x}_{(i,\cdot)}\|_2$ 

Basis pursuit becomes:  $\min_{\mathbf{x}} \sum_{i} \|\mathbf{x}_{(i,\cdot)}\|_2 \text{ s.t. } \mathbf{y} \approx \Phi \mathbf{x}$ 

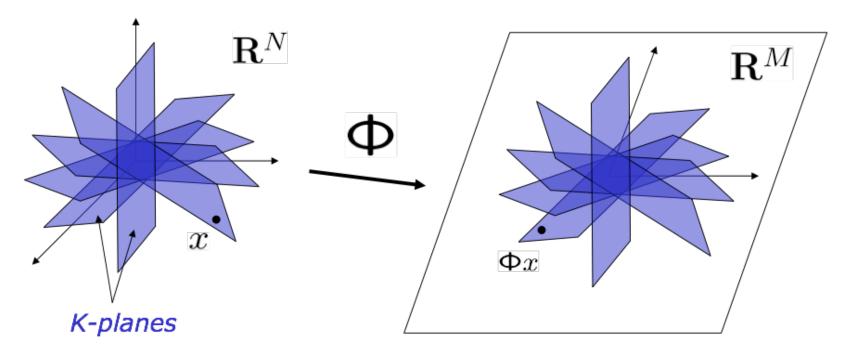
# Randomized Embeddings

#### Stable Embeddings

#### Recall: RIP

RIP of order K requires: for all K-sparse x,

$$(1 - \delta_K) \le \frac{\|\Phi x\|_2^2}{\|x\|_2^2} \le (1 + \delta_K)$$

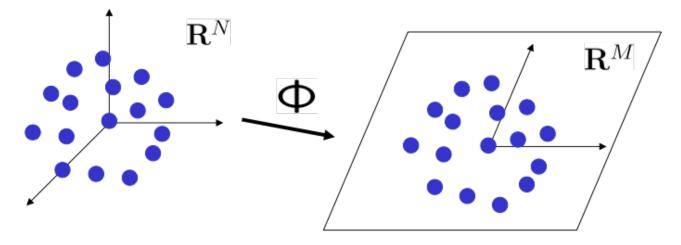


#### Johnson-Lindenstrauss Lemma

[see also Dasgupta, Gupta; Frankl, Maehara; Achlioptas; Indyk, Motwani]

Consider a point set  $Q \subset R^N$  and random\*  $M \times N \Phi$  with  $M = O(\log(\#Q) \epsilon^{-2})$ . With high prob., for all  $x_1, x_2 \in Q$ ,

$$(1-\epsilon) \le \frac{\|\Phi x_1 - \Phi x_2\|_2^2}{\|x_1 - x_2\|_2^2} \le (1+\epsilon).$$



Proof via *concentration inequality*: For any  $x \in R^N$ 

$$\mathbf{P}(\|\Phi x\|_{2}^{2}-\|x\|_{2}^{2})\geq \epsilon\|x\|_{2}^{2})\leq 2e^{-\frac{M}{2}(\epsilon^{2}/2-\epsilon^{3}/3)}.$$

#### Favorable JL Distributions

Gaussian

$$\phi_{i,j} \sim \mathcal{N}\!\left(\mathsf{0}, rac{\mathsf{1}}{M}
ight)$$

Bernoulli/Rademacher [Achlioptas]

$$\phi_{i,j} := \begin{cases} +\frac{1}{\sqrt{M}} & \text{with probability} & \frac{1}{2}, \\ -\frac{1}{\sqrt{M}} & \text{with probability} & \frac{1}{2} \end{cases}$$

"Database-friendly" [Achlioptas]

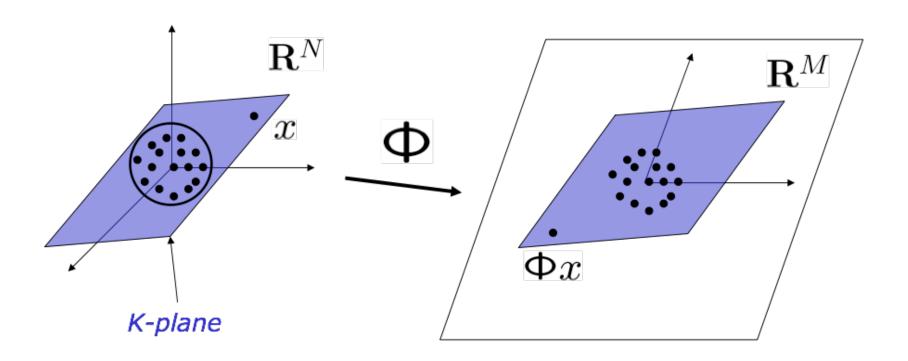
$$\phi_{i,j} := \begin{cases} +\sqrt{\frac{3}{M}} & \text{with probability} \quad \frac{1}{6}, \\ 0 & \text{with probability} \quad \frac{2}{3}, \\ -\sqrt{\frac{3}{M}} & \text{with probability} \quad \frac{1}{6} \end{cases}$$

Random Orthoprojection to R<sup>M</sup> [Gupta, Dasgupta]

### Connecting JL to RIP

#### Consider effect of random JL $\Phi$ on each K-plane

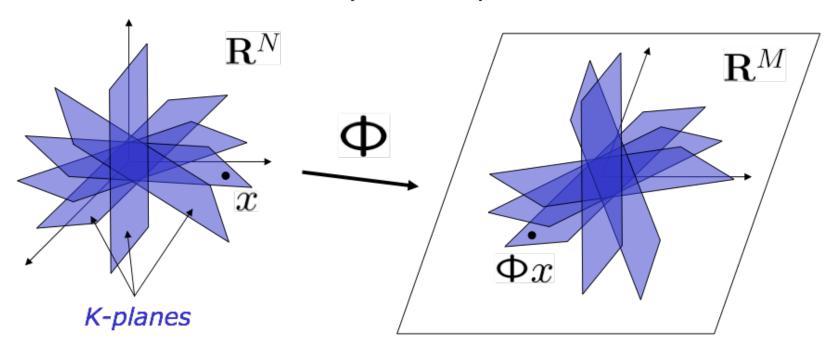
- construct covering of points Q on unit sphere
- JL: isometry for each point with high probability
- union bound → isometry for all q ∈ Q
- extend to isometry for all x in K-plane



#### Connecting JL to RIP

#### Consider effect of random JL $\Phi$ on each K-plane

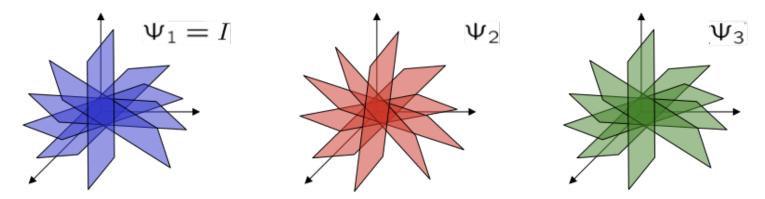
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- extend to isometry for all x in K-plane
- union bound → isometry for all K-planes



Theorem: Supposing Φ is drawn from a JL-favorable distribution,\* then with probability at least 1-e-C\*M,
 Φ meets the RIP with

$$K \le C \cdot \frac{M}{\log(N/M) + 1}.$$

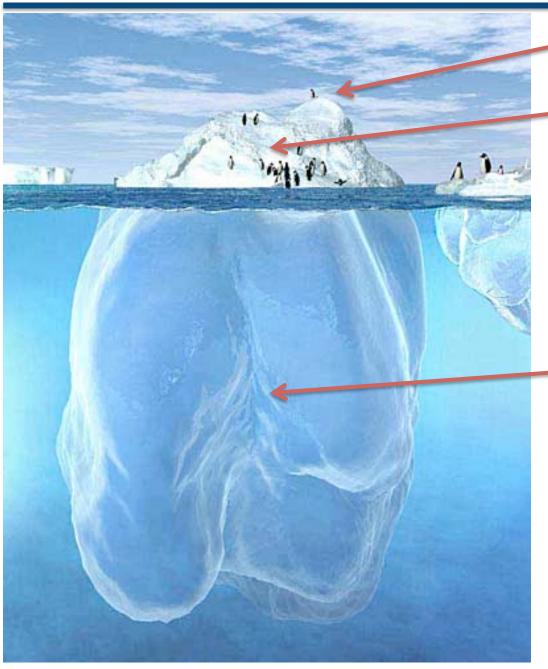
- \* Gaussian/Bernoulli/database-friendly/orthoprojector
- Bonus: *universality* (repeat argument for any  $\Psi$ )



See also Mendelson et al. concerning subgaussian ensembles

More?

# The tip of the iceberg



Today's lecture

Compressive Sensing Repository dsp.rice.edu/cs

Blog on CS nuit-blanche.blogspot.com/

Yet to be discovered... Start working on it ☺