

Daniel Leeds, 15-212 R02, September 5, 2007

### Conventions:

*Right associative:*  $\rightarrow, ::, @$

E.g.:  $'a \rightarrow 'b \rightarrow 'c = 'a \rightarrow ('b \rightarrow 'c)$

*Left associative:* orelse, andalso

E.g.:  $e_1 \text{ orelse } e_2 \text{ andalso } e_3 = (e_1 \text{ orelse } e_2) \text{ andalso } e_3$

*Binding strength: (in decreasing order)*

list  $> *$   $> \rightarrow$

$+,-,*,\text{div} > ::,@ > =,<$

E.g.: int  $\rightarrow$  int \* int = int  $\rightarrow$  (int \* int)

### Passing functions:

fun useful a [] = [] (\* What does **this** do? \*)

| useful a (b::L) = (a b)::(useful a L)

### Proofs:

fun rev [] = []

| rev (h::L) = (rev L) @ [h]

fun rev2 [] acc = acc

| rev2 (h::L) acc = rev2 L (h::acc)

Prove:  $(\text{rev } L) @ A = \text{rev2 } L A$

```

fun ins x [] = [x]
| ins x (y::L) = if (x<=y) then x::y::L else y::(ins x L)

fun sort [] = []
| sort (x::L) = ins x (sort L)

fun listEq [] [] = true
| listEq (x::xs) (y::ys) = (x=y) andalso (listEq xs ys)

LEMMA: ins x (ins y L) = ins y (ins x L)

By induction on length of list L.

Assume WLOG x < y

Base Case: L = []
ins x (ins y []) =
  ins x [y]
  ins x (y::[])
  if (x<=y) then x::y::[] else ...
  x::y::[]
  x::([y])
  x::(ins y [])
  if (y<=x) then ... else x::(ins y [])
  ins y (x::[])
  ins y ([x])
  ins y (ins x [])

Inductive Hyp: Assume true for some list L of length n
Inductive Case: L'=z::L of length n+1
ins x (ins y (z::L)) =
  ins x (if (y<=z) then y::z::L else z::(ins y L))

CASE 1: x < y <= z
ins x (y::z::L)
  if (x<=y) then x::y::z::L else ...
  x::y::z::L
  x::(y::z::L)
  x::(ins y z::L)
  if (y<=x) then ... else x::(ins y z::L)
  ins y (x::z::L)
  ins y (if (x<=z) then x::z::L else ...)
  ins y (ins x (z::L))

CASE 2: x <= z < y
ins x (z::(ins y L))
  if (x<=z) then x::z::(ins y L) else ...
  x::z::(ins y L)
  x::(z::(ins y L))
  x::(if (y<=z) then ... else z::(ins y L))
  x::(ins y z::L)
  if (y<=x) then ... else x::(ins y z::L)
  ins y (x::z::L)
  ins y (if (x<=z) then x::z::L else ...)
  ins y (ins x (z::L))

CASE 3: z < x < y
ins x (z::(ins y L))
  if (x<=z) then ... else z::(ins x (ins y L))
  z::(ins x (ins y L))
  z::(ins y (ins x L))
  if (y<=z) then ... else z::(ins y (ins x L))
  ins y (z::(ins x L))
  ins y (if (x<=z) then ... else z::(ins x L))
  ins y (ins x (z::L))

The case where x>y is now trivial:
ins y (ins x L) = ins x (ins y L)

Similarly x=y is easy:
ins x (ins y L) = ins x (ins x L) = ins y (ins x L)

```