Analysis of Algorithms: Solutions 4

Problem 1

For each of the following functions, give an asymptotically tight bound (Θ -notation).

(a)
$$(2n^6 + 6n^2)^3 = (\Theta(n^6) + o(n^6))^3 = \Theta((n^6)^3) = \Theta(n^{18})$$

(b)
$$(n+2)^2 \cdot (n+3)^3 \cdot (n+6)^6 = (\Theta(n))^2 \cdot (\Theta(n))^3 \cdot (\Theta(n))^6 = \Theta(n^2 \cdot n^3 \cdot n^6) = \Theta(n^{11})$$

(c)
$$3^{2n} + 2^{3n} = (3^2)^n + (2^3)^n = 9^n + 8^n = 9^n + o(9^n) = \Theta(9^n)$$

(d)
$$\sqrt{2n+2} \cdot \sqrt[3]{3n+3} \cdot \sqrt[6]{6n+6} = \Theta(n^{1/2} \cdot n^{1/3} \cdot n^{1/6}) = \Theta(n^{1/2+1/3+1/6}) = \Theta(n)$$

(e)
$$3^{6n} + n! + \sqrt{n^n} = (3^6)^n + n! + n^{n/2} = o(n!) + n! + o(n!) = \Theta(n!)$$

(f)
$$2^{\frac{\log_3 n}{\log_9 2}} = 2^{\frac{\lg n/\lg 3}{\lg 2/\lg 9}} = 2^{\frac{\lg n/\lg 3}{\lg 2/(2 \cdot \lg 3)}} = 2^{\frac{(2 \cdot \lg 3) \cdot \lg n}{\lg 3 \cdot \lg 2}} = 2^{2 \cdot \lg n} = \left(2^{\lg n}\right)^2 = n^2 = \Theta(n^2)$$

Problem 2

Give an example of functions f(n) and g(n) that satisfy all of the following conditions:

$$f(n) \neq \Theta(g(n))$$

$$f(n) \neq o(g(n))$$

$$f(n) \neq \omega(g(n))$$

Consider the following two functions:

$$f(n) = \begin{cases} n & \text{if } n \text{ is even;} \\ 1 & \text{if } n \text{ is odd.} \end{cases}$$

$$g(n) = \begin{cases} 1 & \text{if } n \text{ is even;} \\ n & \text{if } n \text{ is odd.} \end{cases}$$

For even n, f(n) grows asymptotically faster than g(n). On the other hand, for odd n, f(n) grows asymptotically slower. Therefore, g(n) is neither asymptotically lower bound nor asymptotically upper bound for f(n).