

Computational Foundations for ML

10-607

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Notes and reminders

- HW1 out
- Schedule is posted through Quiz I (on M 4/11)

$$(\lambda x \cdot x + 2)(17) \rightarrow 17 + 2$$

$$(\lambda x \cdot x(x + 2))(3) \rightarrow 3(3 + 2)$$

$$(\lambda x \cdot (\lambda x \cdot x + 2))(3x) \rightarrow \lambda x \cdot x + 2$$

$$(\lambda x \cdot (x, \lambda x \cdot x + 2))(3x) \rightarrow (3x, \lambda x \cdot x + 2)$$

$$(\lambda x \cdot (x, \lambda y \cdot y + 2))(3x) \rightarrow (3x, \lambda x \cdot \cancel{3x} + 2)$$

$$[\lambda x \cdot (\lambda x \cdot x + 2)(3x)](7)$$

$$\hookrightarrow (\lambda x \cdot x + 2)(3 \cdot 7) \rightarrow 3 \cdot 7 + 2$$

$p = \text{happy}(x)$ $p[x \rightarrow 17]$ $\text{happy}(17)$

$\text{first}(3, 7) \rightarrow 3$

$(\lambda x: \text{int. } \top) \mid$
 $(\lambda x: \text{char. } F) \mid (y)$

if $y = '9'$
we get F

if $y = 17$
we get \top

$\hookrightarrow (\lambda x: \text{char. } F) '9'$

$\hookrightarrow F$

pair = $\lambda x, y. \lambda f. f(x, y)$

first = $\lambda x, y. x$

second = $\lambda x, y. y$

(pair (3, 7)) (first)

↳ $(\lambda f. f(3, 7))$ (first)

↳ first (3, 7)

↳ 3

$E(\emptyset)$ \leftarrow $\text{prove: } E(x) \vee O(x) \xrightarrow{x \in \mathbb{N}}$

$O(S\emptyset)$

$E(\psi) \rightarrow E(SS\psi)$

$O(\psi) \rightarrow O(SS\psi)$

$E(\emptyset) \vee O(\emptyset)$

$\hookrightarrow \vee$ -intro

IH: $(E(x) \vee O(x)) \wedge$
 $(E(Sx) \vee O(Sx))$

$E(x) \vee O(x) \xrightarrow{?} E(Sx) \vee$
 $O(Sx)$
this doesn't work

base: $(E(\emptyset) \vee O(\emptyset)) \wedge (E(S\emptyset) \vee O(S\emptyset))$

ind: prove $(E(Sx) \vee O(Sx)) \wedge (E(SSx) \vee O(SSx))$

cases: $E(x) \rightarrow E(SSx) \rightarrow (E(SSx) \vee O(SSx))$

$$O(x) \rightarrow O(SSx) \rightarrow \underbrace{(E(SSx) \vee O(SSx))}$$

By assumption $O(Sx)$

\vee -intro $E(Sx) \vee O(Sx)$

$\exists x. \text{dog}(x)$

$\forall x. \text{dog}(x) \rightarrow \text{good}(x)$

$\exists x. (\text{dog}(x) \vee \forall y$
 $\text{friends}(x, y))$

$\exists x. (\text{dog}(x) \wedge \exists x. \text{cat}(x))$

$N(x)$

...

$\forall x. N(x)$

$\forall x:T. P(x)$

$\exists x:T. P(x)$

$\forall x. P(x)$

...

$P(\text{Fred})$

$P(\text{father}(y))$

$\hookrightarrow \forall y. P(y)$

$P(\text{Fred})$

$\exists x. P(x)$

$\exists x. P(x)$

$P(\text{foo})$

$a \vee b \vee \neg c$

$c \vee d$

$a \vee b \vee d$

$\forall z$

$\neg q(\text{Fred}, f(z))$

$q(x, f(\text{Bob}))$

$x = \text{Fred} \quad z = \text{Bob}$

$$a \wedge (b \vee c) \rightarrow (a \wedge b) \vee (a \wedge c)$$

a

$b \vee c$

$\therefore b$

$a \wedge b$

$(a \wedge b) \vee (a \wedge c)$

$b \rightarrow (a \wedge b) \vee (a \wedge c)$

$\therefore c$

$a \wedge c$

$(a \wedge b) \vee (a \wedge c)$

$c \rightarrow (\dots)$

from a question after class: how to show distributivity