

# Nonlinear systems of equations

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We saw earlier how to solve large systems of linear equations: collect them into a single matrix equation, and use an algorithm like Gaussian elimination to construct and solve a factorization.

We also saw how to make a linear approximation to a nonlinear function: if  $f \in \mathbb{R}^n \rightarrow \mathbb{R}^n$ , then we can get a first-order Taylor approximation by calculating the differential,

$$df(x) = f'(x)dx$$

These two tools can work together: suppose that we want to solve a nonlinear system of equations

$$f(x) = 0$$

If we start from a guess  $x_1$  at a solution, we can construct a first-order Taylor expansion

$$df = f'(x_1)dx$$

Holding  $x_1$  fixed, this is a *linear* equation for  $df$  in terms of  $dx$ . So we can ask to find  $dx$  that makes  $f(x_1) + df = 0$  — that is, we can solve a linear approximation to the original nonlinear equations. As before, Gaussian elimination or other factorizations can solve this linear system quickly and reliably.

*If  $f'(x_1)$  is singular, there are two possible cases. The first is that there might be multiple solutions. In this case we need to pick one; a good choice is the least-norm solution, which we can find using the SVD. The second case is that there might be no solutions. In this case a good choice is the least-squares solution, the one that minimizes  $\|f(x_1) + df\|^2$ . We can again find this solution using the SVD.*

With the solution  $dx$  in hand, we can construct a new guess

$$x_2 = x_1 + dx$$

We can then make a new Taylor expansion around  $x_2$ , leading to a new linear approximation to our system of equations,

$$f(x_2) + df = 0 \quad df = f'(x_2)dx$$

Repeating the process lets us construct  $x_3$ ,  $x_4$ , and so forth. Hopefully each successive  $x_t$  comes closer to satisfying  $f(x_t) = 0$ .

This process is called *Newton's method*, and it often converges rapidly to a solution of the nonlinear system  $f(x) = 0$ . In fact, the fixed points of Newton's method are strongly related to the solutions of our system: if  $f'(x)$  is nonsingular then a fixed point must satisfy both equations, and  $df = 0$  is equivalent to  $dx = 0$ . However, Newton's method isn't always stable: even if there are good fixed points near our initial guess, our sequence of guesses might diverge.

*If  $f'(x)$  is singular at a fixed point, then we might be in either of the two cases described above: we might satisfy the two equations but have multiple possible solutions for the second, or we might not be able to satisfy both equations and have to settle for the least-squares solution. In the first case, since we're at a fixed point we have to have  $dx = 0$ ; that means  $df = 0$  and  $f(x) = 0$ . In the second case we have  $f(x) + df \neq 0$ , so that we are at a fixed point that is not a solution.*

If Newton's method diverges, sometimes we can rescue it by *damping*, i.e., decreasing our step size: that is, we set  $x_{t+1} = x_t + \alpha_t dx$  for some  $\alpha_t \in (0, 1)$ . But tuning the step size (and other methods beyond damped Newton) are beyond the scope of this set of notes.

## Example

Let  $f(x) = e^x - 1$ , so that  $df = e^x dx$ . The solution to  $f(x) = 0$  is  $x = 0$ , but let's see if we can find this by Newton's method, starting from somewhere else.

$x$	$f$	$df$	Equation	$dx$
1	$e - 1$	$e$	$e dx = 1 - e$	$\frac{1-e}{e}$
-0.632	-0.468	0.532	$0.532 dx = 0.468$	0.880
0.248	0.281	1.281	$1.281 dx = -0.281$	-0.219

Quite rapidly we have reached  $x = 0.029$ , very close to the true solution.

## Unconstrained optimization

Solving optimization problems is strongly related to solving systems of equations. In an unconstrained optimization problem

$$\min_{\theta} L(\theta) \quad L \in \mathbb{R}^n \rightarrow \mathbb{R}$$

we can try to find the solution by looking for a critical point: a place where, locally, changes to  $\theta$  do not change  $L(\theta)$ .

*Critical points can be minima or maxima, and either type of optimum can be local or global. In addition, critical points can be neither minima nor maxima: they can be places where the function flattens out temporarily, or places where it looks like a saddle, curving upward in some directions and downward in others. For now, we won't be concerned with checking which is which.*

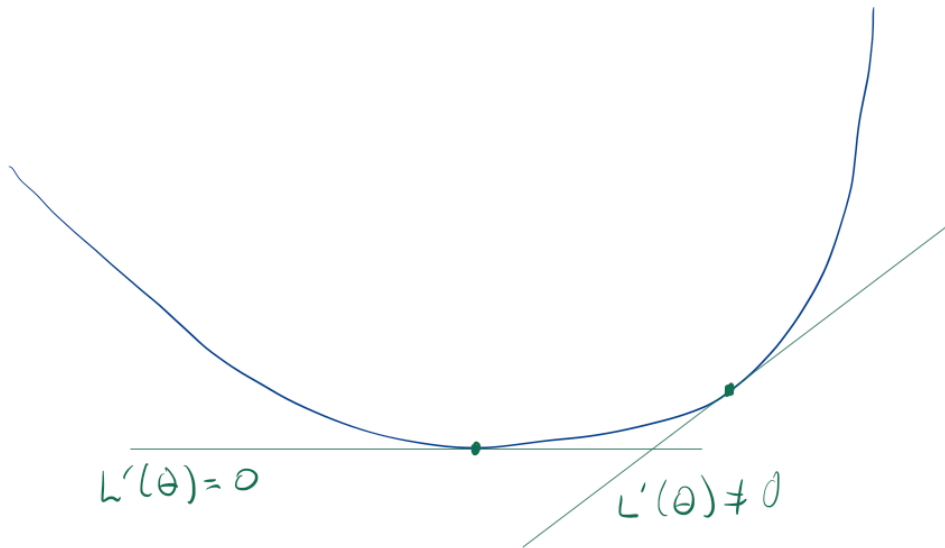
To find a critical point, we can look at the first order Taylor expansion of  $L$ :

$$dL = L'(\theta) d\theta$$

At a critical point, all possible changes  $d\theta$  should leave  $dL = 0$ . That means we must have

$$L'(\theta) = 0$$

Since  $L'(\theta) \in \mathbb{R}^{1 \times n}$ , this is a system of  $n$  equations. These equations are the *first-order optimality conditions* for  $L(\theta)$ . Geometrically, they mean that the Taylor expansion is flat: a constant function of  $d\theta$ .



Of course, the system of equations  $L'(\theta) = 0$  could be nonlinear. So, we can apply Newton's method — that is, we can set a first-order Taylor approximation of  $L'$  to zero and solve for  $d\theta$ :

$$L'(\theta) + dL' = 0 \quad dL' = L''(\theta)d\theta$$

We can find  $L'$  and  $L''$  by differentiating  $L$  twice.

This is such a common application of Newton's method that it shares the same name. If

necessary, we can distinguish by calling the two algorithms "Newton's method for solving a system of equations" and "Newton's method for optimizing a function".

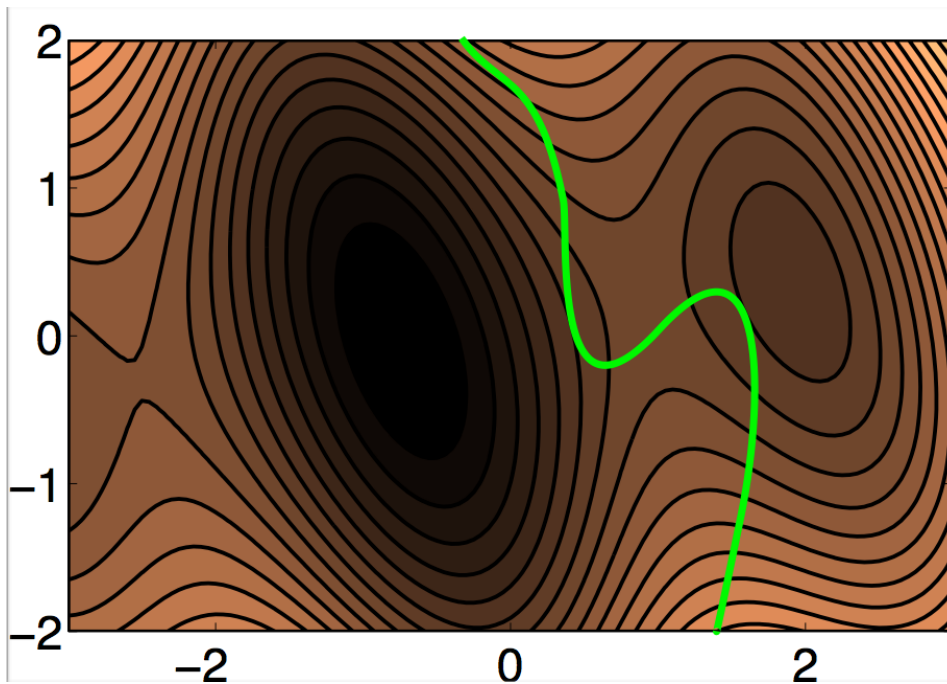
## Constrained optimization

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In a constrained problem

$$\min_{\theta} L(\theta) \quad \text{s.t.} \quad g(\theta) = 0 \quad L \in \mathbb{R}^n \rightarrow \mathbb{R}, g \in \mathbb{R}^n \rightarrow \mathbb{R}$$

we don't need  $L'(\theta)$  to be zero: it's OK if there's a direction of decrease in  $L(\theta)$  as long the constraint prevents us from moving in this direction.



To encode this condition, we need to be a bit clever. First note that the solutions to the following problem

$$\min_{\theta} [L(\theta) + \alpha g(\theta)] \quad \text{s.t.} \quad g(\theta) = 0$$

are the same as the solutions to our original problem, no matter what the value of  $\alpha$  is, since  $\alpha g(\theta) = 0$  for any feasible  $\theta$ .

Then note that, by choosing  $\alpha$  appropriately, we can rule out any direction of decrease in  $L$  that doesn't satisfy the constraint: if  $L$  would decrease on the side of the constraint where  $g(\theta) > 0$ , then we choose  $\alpha$  to be very positive, so that any motion in this direction would cause  $L(\theta) + \alpha g(\theta)$  to increase instead of decreasing. Similarly, if  $L$  would decrease on the other side of the constraint, where  $g(\theta) < 0$ , we choose  $\alpha$  to be very

negative.

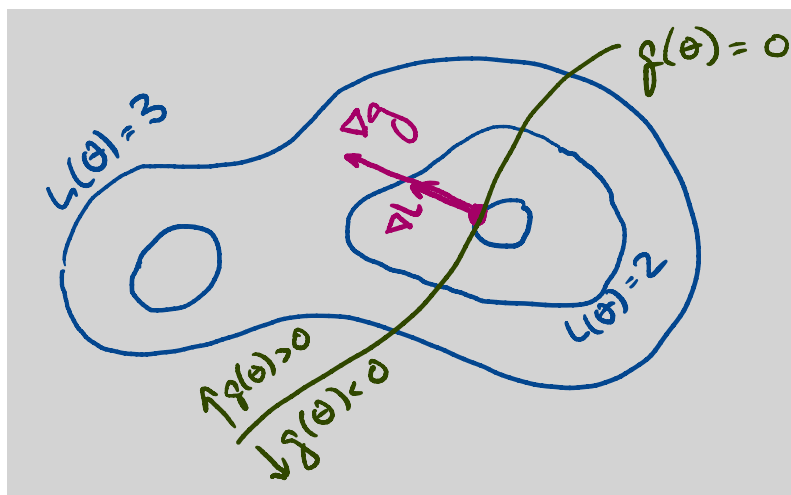
Given this new objective, we can use a Taylor expansion the same way as before. We ask for a critical point: a  $\theta$  where, to first order, the objective doesn't change as we change  $d\theta$ . That is,

$$0 = d(L(\theta) + \alpha g(\theta)) = L'(\theta)d\theta + \alpha g'(\theta)d\theta$$

which implies

$$L'(\theta) + \alpha g'(\theta) = 0$$

Geometrically, this equation tells us that at a critical point we can only change  $L(\theta)$  by changing  $\theta$  in a direction *orthogonal* to the constraint (parallel to  $g'(\theta)$ ): sliding in any direction along the constraint doesn't change  $L$ , at least to first order.



Interestingly, that means that we didn't have to choose  $\alpha$  a priori: any  $\theta$  and  $\alpha$  that satisfy

$$g(\theta) = 0 \quad L'(\theta) + \alpha g'(\theta) = 0$$

will represent a critical point. So, as before, we've turned our optimization problem into a possibly-nonlinear system of equations. We can solve this system with Newton's method or any other appropriate tool.

The new variable  $\alpha$  is called a *Lagrange multiplier* or a *dual variable*. We can interpret  $-L'(\theta)$  as a force that wants to push our current point  $\theta$  downhill, toward a minimum of  $L$ . We can then think of  $-\alpha g'(\theta)$  as a force that pushes back, keeping  $\theta$  from violating the constraint. At a solution, the two forces balance exactly.

By introducing the dual variable, we've transformed our optimization problem into a

system of simultaneous equations, where the objective and the constraints are treated the same way. This transformation was what let us apply Newton's method.

Exercise: solve the following problem by introducing a Lagrange multiplier.

$$\min_{x,y} \frac{1}{2}(x^2 + y^2) \quad \text{s.t.} \quad x + 2y = 1$$

## Multiple constraints

Suppose we have more than one constraint:

$$\min_{\theta} L(\theta) \quad \text{s.t.} \quad g(\theta) = 0 \quad L \in \mathbb{R}^n \rightarrow \mathbb{R}, g \in \mathbb{R}^n \rightarrow \mathbb{R}^d$$

where the output of  $g(\theta)$  is in  $\mathbb{R}^d$  instead of  $\mathbb{R}$ . The solution in this case is almost identical: we can still solve

$$g(\theta) = 0 \quad L'(\theta) + \alpha g'(\theta) = 0$$

But now, instead of  $\alpha \in \mathbb{R}$ , we need  $\alpha \in \mathbb{R}^{1 \times d}$ , so that  $\alpha g'(\theta)$  is the same shape as  $L'(\theta)$ : a  $1 \times n$  matrix.

Each coordinate  $\alpha_i$  is still called a Lagrange multiplier. The geometric interpretation is only slightly different from before: we think of each  $\alpha_i$  as controlling a separate force, in a direction that's normal to the corresponding constraint  $g_i(\theta) = 0$ . At a critical point, all of the forces  $\alpha_i g'_i(\theta)$  combine to cancel out  $L'(\theta)$ .

Exercise: solve the following problem by introducing two Lagrange multipliers.

$$\min_{x,y,z} \frac{1}{2}(x^2 + y^2 + z^2) \quad \text{s.t.} \quad x + y = 1, y + z = 1$$

Solution: the loss derivative is  $(x, y, z)$ . The constraint derivatives are  $(1, 1, 0)$  and  $(0, 1, 1)$ . So,  $(x, y, z) + \alpha(1, 1, 0) + \beta(0, 1, 1) = 0$ . So, the first-order optimality conditions are:

$$\begin{pmatrix} 1 & 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 1 & 1 \\ 0 & 0 & 1 & 0 & 1 \\ 1 & 1 & 0 & 0 & 0 \\ 0 & 1 & 1 & 0 & 0 \end{pmatrix} \begin{pmatrix} x \\ y \\ z \\ \alpha \\ \beta \end{pmatrix} = \begin{pmatrix} 0 \\ 0 \\ 0 \\ 1 \\ 1 \end{pmatrix}$$

The first three rows implement  $(x, y, z) + \alpha(1, 1, 0) + \beta(0, 1, 1) = 0$ . The last two rows implement  $x + y = 1, y + z = 1$ . We can solve this system of equations by hand or with

assistance from a computer. Either way, we get

$$\begin{pmatrix} x \\ y \\ z \\ \alpha \\ \beta \end{pmatrix} = \begin{pmatrix} \frac{1}{3} \\ \frac{2}{3} \\ \frac{1}{3} \\ -\frac{1}{3} \\ -\frac{1}{3} \end{pmatrix}$$