**Approximations for Inference in Undirected Graphical Models** 

Pradeep Ravikumar

#### **Graphical Models, The History**

Rennaissance Period



I want to model the world, and I like graphs...

The world was not ready, and so Da Vinci hid clues in architecture and paintings.

#### Mid to Late Twentieth Century



Pioneering work of Conspiracy Theorists

The System, it is all **connected**...

Late Twentieth Century: people realize that existing scientific literature (Statistics, The Da Vinci Code) offers a marriage between probability theory and graph theory – which can be used to model the world.



Common Misconception: Called graphical models after Grafici Modeles, a sculptor protege of Da Vinci.

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 $X_1 \perp X_4 \,|\, X_2, \, X_3$ 



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- Global Markov Property
- Graph Encodes Conditional Independencies

#### Hammersley and Clifford Theorem

 $\mathcal C$  set of cliques in graph G.

Positive p over X is Markovian with respect to G iff p factorizes according to cliques in G,

$$p(X) = \frac{1}{Z} \prod_{C \in \mathcal{C}} \psi_C(X_C)$$

#### **Exponential Family**

$$p(X) = \frac{1}{Z} \prod_{C \in \mathcal{C}} \psi_C(X_C)$$
$$= \exp(\sum_{C \in \mathcal{C}} \log \psi_C(X_C) - \log Z)$$

Exponential family:  $p(X;\theta) = \exp\left(\sum_{\alpha \in \mathcal{C}} \theta_{\alpha} \phi_{\alpha}(X) - \Psi(\theta)\right)$ 

- $\{\phi_{\alpha}: \ \alpha \in \mathcal{C}\}$  features  $\{\theta_{\alpha}: \ \alpha \in \mathcal{C}\}$  parameters
- { $\Psi(\theta)$  : log partition function }

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Overcomplete potentials:

$$\begin{split} \mathcal{I}_{j}(x_{s}) &= \begin{cases} 1 & x_{s} = j \\ 0 & \text{otherwise} \end{cases} \\ \mathcal{I}_{j,k}(x_{s}, x_{t}) &= \begin{cases} 1 & x_{s} = j \text{ and } x_{t} = k \\ 0 & \text{otherwise} . \end{cases} \end{split}$$

$$p(x|\theta) = \exp\left(\sum_{s,j} \theta_{s,j} \mathcal{I}_j(x_s) + \sum_{s,t;j,k} \theta_{s,j;t,k} \mathcal{I}_{j,k}(x_s, x_t) - \Psi(\theta)\right)$$

#### Inference



# Inference

For undirected model  $p(x;\theta) = \exp\left(\sum_{\alpha \in I} \theta_{\alpha} \phi_{\alpha}(x) - \Psi(\theta)\right)$  key inference problems are:

- ▷ compute log partition function (normalization constant)  $\Psi(\theta)$
- ▷ marginals  $p(x_A) = \sum_{x_v, v \notin A} p(x)$
- ▷ most probable configurations  $x^* = \arg \max_x p(x | x_L)$

These problems are intractable in full generality.

Approximate inference techniques have focused on sampling and variational methods.

$$Z = \log \sum_{x} \prod_{\alpha \in I} \psi_{\alpha}(x_{\alpha})$$



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Exponential in tree-width.

### **Preconditioner Approximations**

$$Z = \log \sum_{x} \prod_{\alpha \in I} \psi_{\alpha}(x_{\alpha})$$

Recent scientific computing developments allow us to propose a preconditioner based approximation.



Matrix form:

$$\Theta := \begin{pmatrix} \theta_{11} & \dots & \theta_{1n} \\ \vdots & \vdots & \vdots \\ \theta_{n1} & \dots & \theta_{nn} \end{pmatrix} \quad \Phi(x) := \begin{pmatrix} \phi_{11}(x_1, x_1) & \dots & \phi_{1n}(x_1, x_n) \\ \vdots & \vdots & \vdots \\ \phi_{n1}(x_n, x_1) & \dots & \phi_{nn}(x_n, x_n) \end{pmatrix}$$

Notation:  $tr(AB) = \langle \langle A, B \rangle \rangle = \sum_{ij} A_{ij} B_{ij}$ .

Energy  $\sim \sum_{st} \theta_{st} \phi_{st}(x_s, x_t) = \langle\!\langle \Theta, \Phi(x) \rangle\!\rangle$ 

$$p(x;\theta) = \frac{\exp\left\langle\!\left\langle\Theta, \Phi(x)\right\rangle\!\right\rangle}{Z(\theta)} \qquad \mathcal{Z}(\theta) = \sum_{x \in \mathcal{X}} \exp\left\langle\!\left\langle\Theta, \Phi(x)\right\rangle\!\right\rangle$$



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A x = b,  $A n \times n$ , m non-zero entries

Direct methods  $\sim O(n^3)$ 

 $B^{-1}Ax = B^{-1}b$ 

 $B \sim A, B^{-1}A \sim I$ 

Simpler?

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Computation required for  $\epsilon$ -approximate solution:

$$T(A) = \sqrt{\kappa(A, B)} (m + T(B)) \log\left(\frac{1}{\epsilon}\right)$$

 $\kappa(A, B)$  is "quality of approximation" and T(B) is time to solve By = c.

**Objective**: Find "sparse" matrix B, e.g., a spanning tree, with small condition number  $\kappa(A, B)$ 

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### **Recent Developments in Preconditioners**

Nearly optimal linear system solvers:  $O\left(m \log^{O(1)} n\right)$ 

- ▷ Vaidya (1990)
- ▷ Gremban and Miller (1996)
- ▶ Boman and Hendrickson (2002)
- ▷ Spielman and Teng (2003)
- ▶ Elkin et al. (2004)

Matrix A  $\sim$  Laplacian of a graph

Methods require diagonally dominant matrices:  $A_{ii} \ge \sum_{j \neq i} |A_{ij}|$ 

#### **Basic Approach**

Intuition:  $\langle\!\langle \Theta, \Phi(x) \rangle\!\rangle \sim x^\top A x$ .

If  $B \sim A$ ,  $\kappa(A, B)$  is small:

$$\kappa(A,B) = \max_{x} \frac{x^{\top}Ax}{x^{\top}Bx} / \min_{x} \frac{x^{\top}Ax}{x^{\top}Bx} = \lambda_{max}(A,B) / \lambda_{min}(A,B)$$

B chosen to minimize condition number, rather than KL
 Scale B appropriately

- If *B* is a tree, sums can be computed efficiently
- B can be used to approximate event probabilities

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- If B is a tree, sums can be computed efficiently
- $\bullet~B$  can be used to approximate event probabilities

Bounds on energy:  $\langle \langle C, \Phi(x) \rangle \rangle \leq \langle \langle \Theta, \Phi(x) \rangle \rangle \leq \langle \langle B, \Phi(x) \rangle \rangle$ Imply bounds on partition function and probabilities:

$$\frac{\exp\left\langle\!\left\langle C, \Phi(x)\right\rangle\!\right\rangle}{Z(B)} \le p(x; \Theta) \le \frac{\exp\left\langle\!\left\langle B, \Phi(x)\right\rangle\!\right\rangle}{Z(C)}$$

Want B to be as similar to  $\Theta$  as possible. For upper bound, this leads to optimization problem

$$B^{\star} = \underset{B}{\operatorname{argmax}} \min_{x} \quad \frac{\langle\!\langle \Theta, \Phi(x) \rangle\!\rangle}{\langle\!\langle B, \Phi(x) \rangle\!\rangle}$$
  
such that 
$$\frac{\langle\!\langle \Theta, \Phi(x) \rangle\!\rangle}{\langle\!\langle B, \Phi(x) \rangle\!\rangle} \leq 1$$

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#### Reductions

Have to solve constrained minimax and maximin problems!



For a graphical model with potential function matrix  $\Phi$  and a pair of matrices (A, B), we define

Generalized graphical model eigenvalues:

$$\lambda_{\max}^{\Phi}(A,B) = \max_{\substack{x: \langle\!\langle B, \Phi(x) \rangle\!\rangle \neq 0}} \frac{\langle\!\langle A, \Phi(x) \rangle\!\rangle}{\langle\!\langle B, \Phi(x) \rangle\!\rangle}$$
$$\lambda_{\min}^{\Phi}(A,B) = \min_{\substack{x: \langle\!\langle B, \Phi(x) \rangle\!\rangle \neq 0}} \frac{\langle\!\langle A, \Phi(x) \rangle\!\rangle}{\langle\!\langle B, \Phi(x) \rangle\!\rangle} = \frac{1}{\lambda_{\max}^{\Phi}(B,A)}$$

Graphical model condition number:

$$\kappa^{\Phi}(A,B) = \frac{\lambda^{\Phi}_{\max}(A,B)}{\lambda^{\Phi}_{\min}(A,B)}$$
#### **Condition Number Relaxation**

**Proposition 1.** Let  $C = \operatorname{argmin}_{B} \kappa^{\Phi}(\Theta, B)$ . Then the optimal upper bound matrix  $B^{U} = \lambda_{\max}(\Theta, C) C$  and the optimal lower bound matrix  $B^{L} = \lambda_{\min}(\Theta, C) C$ .

Optimizing the constrained minimax and maximin problems can be reduced to minimizing (graphical model) condition numbers!

#### **Condition Number Relaxation**

**Proposition 2.** For a potential function matrix  $\Phi(x) \succeq 0$  then  $\lambda_{\max}^{\Phi}(\Theta, B) \leq \lambda_{\max}(\Theta, B)$ . If  $B \succeq 0$  then  $\kappa^{\Phi}(\Theta, B) \leq \kappa(\Theta, B)$ .

Optimizing graphical model CN can be reduced to minimizing classical CN!

Recipe for Graphical Model Preconditioner Approximations:

- $B^* = \operatorname{arg\,min}_B \kappa(\Theta, B)$
- Upper bound matrix  $\lambda_{\max}(\Theta, B^*) \; B^*$
- Lower bound matrix  $\lambda_{\min}(\Theta, B^*) \ B^*$

#### **Reduction to Ising Form**

For the Ising potential function,  $\phi_{ij}(x_i, x_j) = x_i x_j$ 

$$\kappa^{\Phi}(A,B) = \kappa(A,B)$$

Can reduce any potential function  $\phi$  over discrete k-ary random variables  $X_i$ , to a binary Ising potential function:

$$\mathcal{E}(x) = \sum_{(i,j)\in E} \theta_{ij}\phi_{ij}(x_i, x_j) = \bar{x}^{\mathsf{T}}A(\theta, \phi)\bar{x}$$

where  $\bar{x}_{(i,l)} = \delta(x_i, l)$  and

$$A(\theta,\phi) = \left[A_{(i,l),(j,m)}(\theta,\phi)\right] = \left[\theta_{ij}\phi_{ij}(l,m)\right]$$

Solve for preconditioner bounds of the matrix  $A(\theta, \phi)$ 

### **Simple Preconditioners**

Vaidya's Spanning Tree Preconditioner

- Matrix corresponding to the maximum spanning tree of the given graph.
- Requires the parameter matrix to be Laplacian.
- Gremban-Miller Support Tree Preconditioner





# **Experiments:** $\log Z(\theta)$ Lower Bounds, 2D Grid



#### **Experiments**



# **Preconditioner Approximations**

- The framework yields upper and lower bounds on energy, and consequently on the log-partition function, general event probabilities, and the MAP energy.
- The procedure has a low time complexity: both the construction of a sparse preconditioner and inference using a sparse (e.g. tree-based) preconditioner matrix are typically linear.





PROB = 0.01



PROB = 0.2



Most Probable Configuration?

$$p(x|\theta) \propto \exp\left(\sum_{s,j} \theta_{s;j} \mathcal{I}_j(x_s) + \sum_{s,t;j,k} \theta_{s,j;t,k} \mathcal{I}_{j,k}(x_s, x_t)\right).$$

$$x^* = \underset{x}{\operatorname{arg\,max}} \sum_{s,j} \theta_{s;j} \mathcal{I}_j(x_s) + \sum_{s,t;j,k} \theta_{s,j;t,k} \mathcal{I}_{j,k}(x_s, x_t).$$

#### **Integer Linear Program**

 $I_j(x_s) \sim \mu_1(s;j)$  $I_{j,k}(x_s, x_t) \sim \mu_2(s,j;t,k)$ 

max

such that

$$\sum_{s;j} \theta_{s;j} \mu_1(s;j) + \sum_{s,t;j,k} \theta_{s,j;t,k} \mu_2(s,j;t,k)$$
$$\sum_k \mu_2(s,j;t,k) = \mu_1(s;j)$$
$$\sum_k \mu_1(s;j) = 1$$
$$\mu_1(s;j) \in \{0,1\}$$
$$\mu_2(s,j;t,k) \in \{0,1\}.$$

## **Linear Relaxation**

$$\begin{split} \max & \sum_{s;j} \theta_{s;j} \, \mu_1(s;j) + \sum_{s,t;j,k} \theta_{s,j;t,k} \, \mu_2(s,j;t,k) \\ \text{such that} & \sum_k \mu_2(s,j;t,k) = \mu_1(s;j) \\ & \sum_j \mu_1(s;j) = 1 \\ & 0 \leq \mu_1(s;j) \leq 1 \\ & 0 \leq \mu_2(s,j;t,k) \leq 1. \end{split}$$

#### **Linear Relaxation**

 $LP \sim Chekuri, Khanna, Naor, Zosin$ 

Tree-reweighted Belief Propagation, Dual of LP  $\sim$  Wainwright, Jaakkola, Willsky

# **Quadratic Relaxation**

$$I_{j,k}(x_s, x_t) = I_j(x_s)I_k(x_t)$$

$$\mu_2(s, j; t, k) \sim \mu_1(x_s; j) \mu_1(x_t; k)$$

$$O(|E|K^2) \rightarrow O(nK)$$
 variables!

## **Quadratic Integer Program**

max

$$\begin{array}{ll} \max & \sum_{s;j} \theta_{s;j} \mu(s;j) + \sum_{s,t;j,k} \theta_{s,j;t,k} \, \mu(s;j) \, \mu(t;k) \\ \text{subject to} & \sum_{j} \mu(s;j) = 1 \\ & \mu(s;j) \in \{0,1\} \end{array}$$

# **Quadratic Relaxation**

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Theorem: QP Relaxation is equivalent to the MAP problem.

- Relaxation is tight!
- MAP is in P when  $\{\Theta_{s,j;t,k}\}$  is negative semi-definite

### **Quadratic Relaxation**

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Want  $-\Theta\sim$  Positive semi-definite

Want  $-\Theta\sim$  Diagonally Dominant

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 $-\Theta_{s,j;s,j} \ge \sum_{t,k} |\Theta_{s,j;t,k}|$ 

Solution: Substract from diagonal?

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### **Diagonally Dominant**



# **Diagonally Dominant**

$$\Theta' = \Theta - \mathsf{diag}(d(s;j))$$

$$\theta'_{s;j} = \theta_{s;j} + d(s;j)$$

$$\begin{split} \max_{\mu} & \sum_{s;j} \theta_{s;j}' \, \mu(s;j) + \sum_{s,t;j,k} \theta_{s,j;t,k}' \, \mu(s;j) \, \mu(t;k) \\ \text{such that} & \sum_{j} \mu(s;j) = 1 \\ \mu(s;j) \in \{0,1\} \end{split}$$

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- Convex QP (with simple box constraints); polynomial time
- Not tight!
- Additive Guarantee:  $E(y^{CVX}) \ge E^* \frac{1}{4} \sum_{s,i} d(s;i)$

#### **Iterative Procedure**

$$\begin{array}{ll} \max & \sum_{s;j} \theta_{s;j} \mu(s;j) + \sum_{s,t;j,k} \theta_{s,j;t,k} \, \mu(s;j) \, \mu(t;k) \\ \text{subject to} & \sum_{j} \mu(s;j) = 1 \\ & 0 \leq \mu(s;j) \leq 1 \end{array}$$

Co-ordinate Ascent: Optimize  $\mu(s;.)$  for node s, fixing values of other nodes,

$$\mu(s;.) = \max_{\mu(s;.)} \sum_{j} \theta_{s;j} \mu(s;j) + \sum_{t;j,k} \theta_{s,j;t,k} \, \mu(s;j) \, \mu(t;k)$$

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Solution of fixed point equation,

$$j^*(s) = \underset{j}{\operatorname{arg\,max}} \ \theta_{s;j} + \sum_{t;j,k} \theta_{s,j;t,k} \mu(t;k)$$

and setting  $\mu(s, j) = \mathcal{I}_{j^*(s)}(j)$ .

This is the Iterative Conditional Modes algorithm! (Besag 86)

Alternatively, use conjugate gradient for the convex approximation.

Consider Wainwright et al's polytope formulation of MAP,

$$\mu^* = \max_x \left< \theta, \phi \right> = \sup_{\mu \in \mathcal{M}} \left< \theta, \mu \right>$$



If  $M_I \subset \mathcal{M}$  is any subset of the marginal polytope that includes all of the vertices,

$$\mu^* = \max_x \left< \theta, \phi \right> = \sup_{\mu \in M_I} \left< \theta, \mu \right>$$



For the given graph G and a subgraph H, let

$$\mathcal{E}(H) = \{ \theta' \,|\, \theta'_{st} = \theta_{st} \,\mathbf{1}_{(s,t)\in H} \}$$

$$\mathcal{M}(G; H) = \{ \mu \, | \, \mu = E_{\theta}[\phi(x)] \text{ for some } \theta \in \mathcal{E}(H) \} .$$
  
$$\mathcal{M}(G; H) \subseteq \mathcal{M}(G)$$



Mean Field parameters,

 $\mathcal{M}(G; H_0) = \{ \mu(s; j), \mu(s, j; t, k) \mid 0 \le \mu(s; j) \le 1, \mu(s, j; t, k) = \mu(s; j) \mu(t; k) \}$ 

Mean Field Relaxation,

$$\sup_{\mu \in \mathcal{M}(G;H_0)} \langle \theta, \mu \rangle$$
  
= 
$$\sup_{\mu \in \mathcal{M}(G;H_0)} \sum_{s;j} \theta_{s;j} \mu(s;j) + \sum_{st;jk} \theta_{s,j;t,k} \mu(s,j;t,k)$$
  
= 
$$\sup_{\mu \in \mathcal{M}(G;H_0)} \sum_{s;j} \theta_{s;j} \mu(s;j) + \sum_{st;jk} \theta_{s,j;t,k} \mu(s;j) \mu(t;k)$$

#### **Experiments**

- $10 \times 10$  grid graphs; n = 100.
- Number of labels, k = 4.
- Potential Functions: Ising, Quadratic, Linear, Uniform.
- Methods Compared:
  - Iterative Conditional Modes
  - Chekuri LP
  - Tree-reweighted Max-Product
  - Convex Approximation to QP



Comparison on  $10 \times 10$  grid graphs using Ising potentials.


Comparison on  $10\times 10$  grid graphs using linear potentials.



Comparison on  $10 \times 10$  grid graphs using quadratic potentials.



Comparison on  $10 \times 10$  grid graphs using uniform potentials.

## **Experiments**



Figure 5: Comparison of ICM and TRW on larger graphs, using Ising potentials with mixed coupling. The right plot shows  $(e_{ICM} - e_{QP})/e_{ICM}$  and  $(e_{TRW} - e_{QP})/e_{TRW}$ .

## **QP** Relaxation

- The QP has O(nk) variables in contrast to  $O(|E|k^2)$  variables for linear relaxation.
- The QP more accurately represents objective function: Relaxation is tight!
- There exists a convex approximation with an additive guarantee.
- The QP can be extended to variational inner polytope relaxations.

## **Structure Selection**