

A Logical Framework for Anticipation of Network Incidents

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■ Introduction

- Network Evolution

- Attack Model Evolution

■ Anticipation game key features

- Dependency relations

- Player interaction

- Time

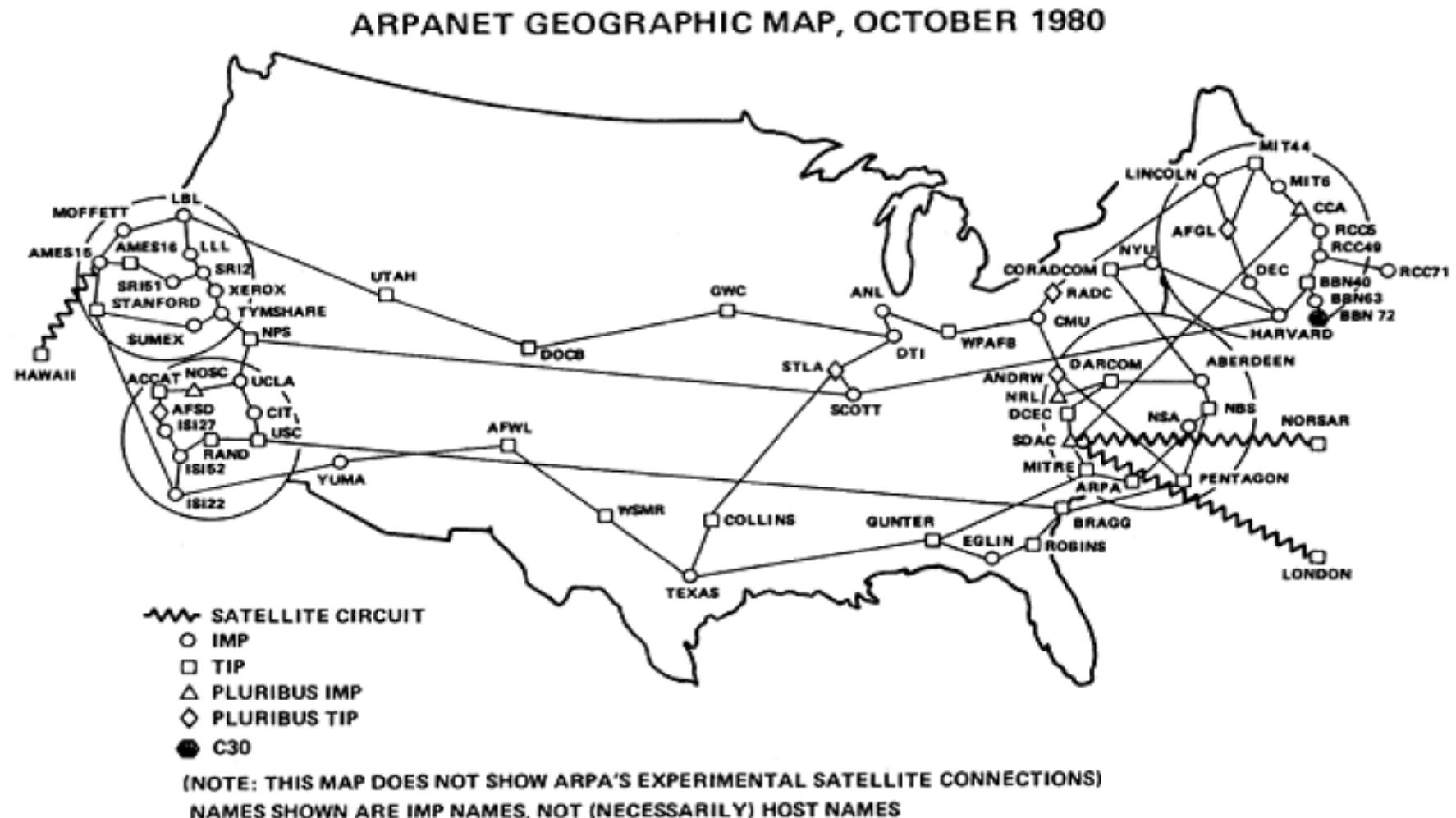
■ Model Logic

- Positional Logic

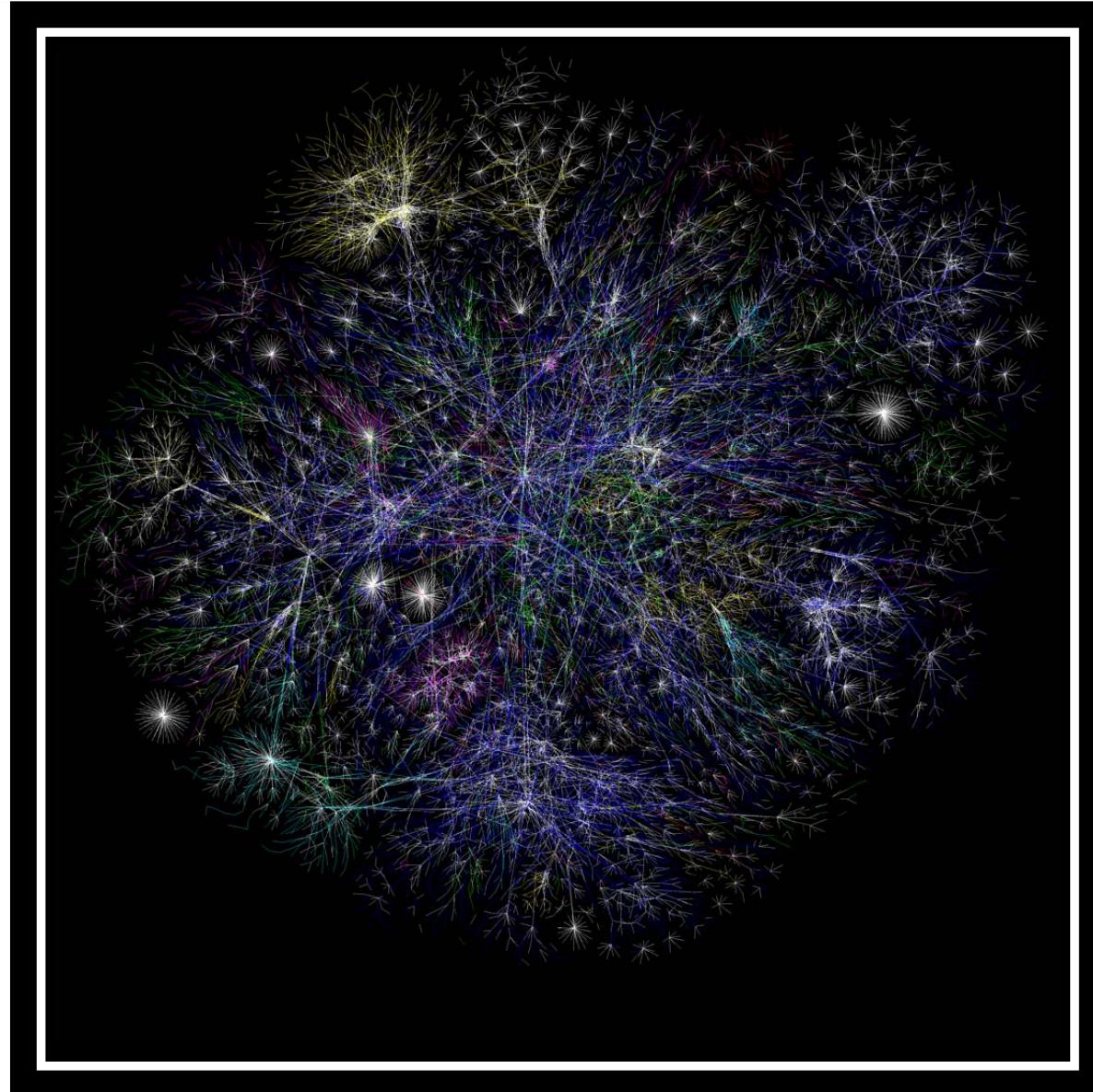
- Temporal Logic

■ Conclusion

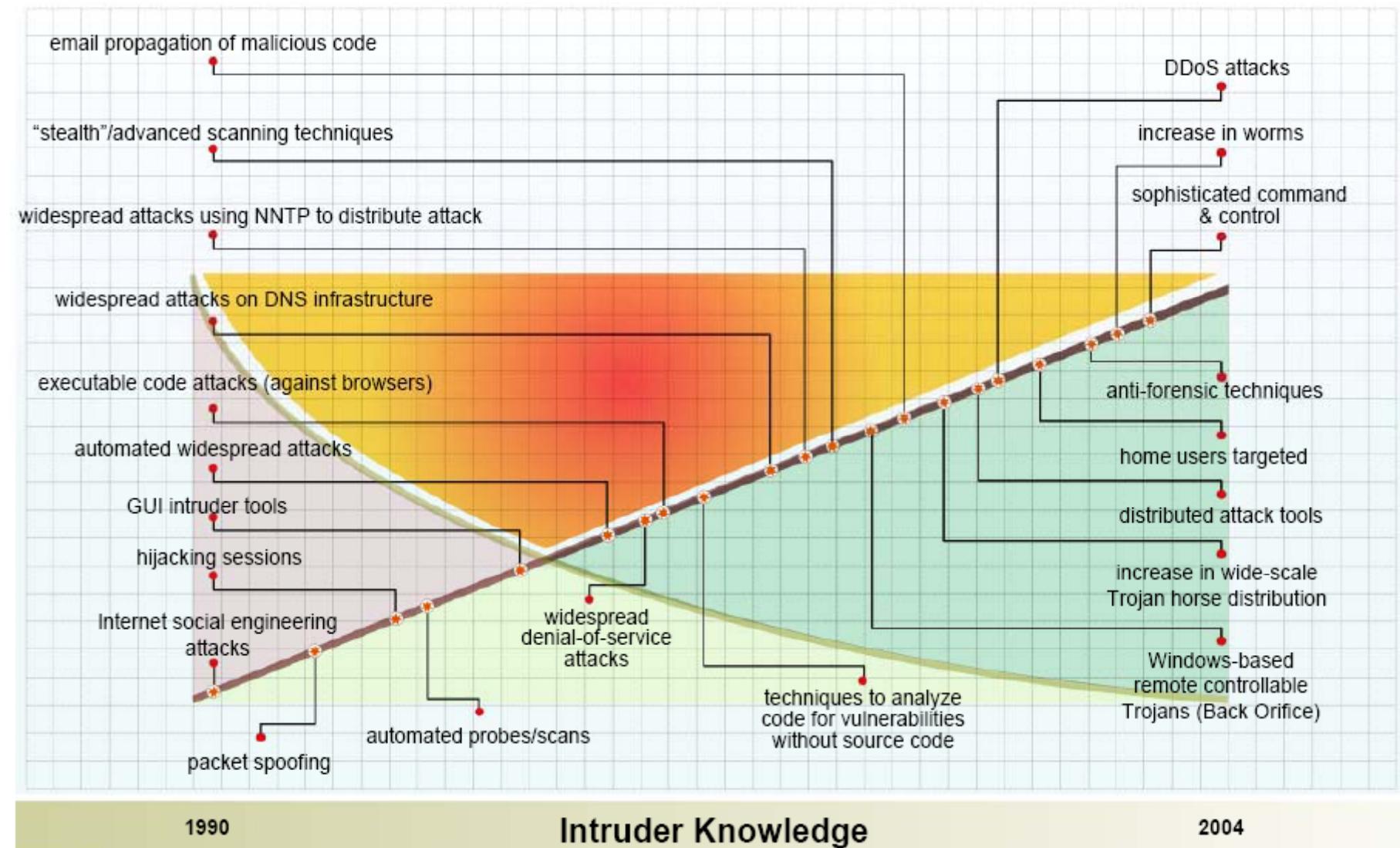
The Good Old Time



The Current Internet



Opte project

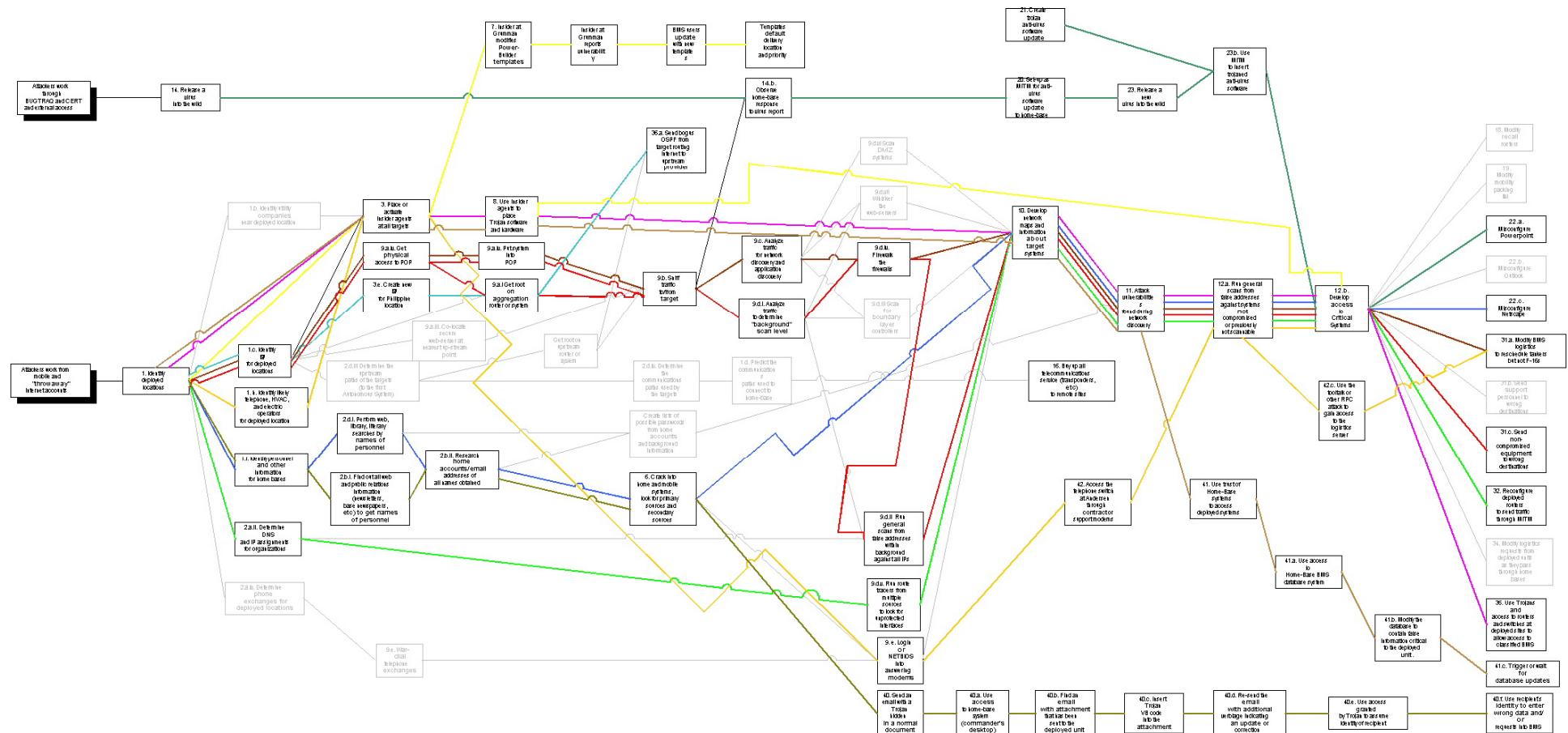


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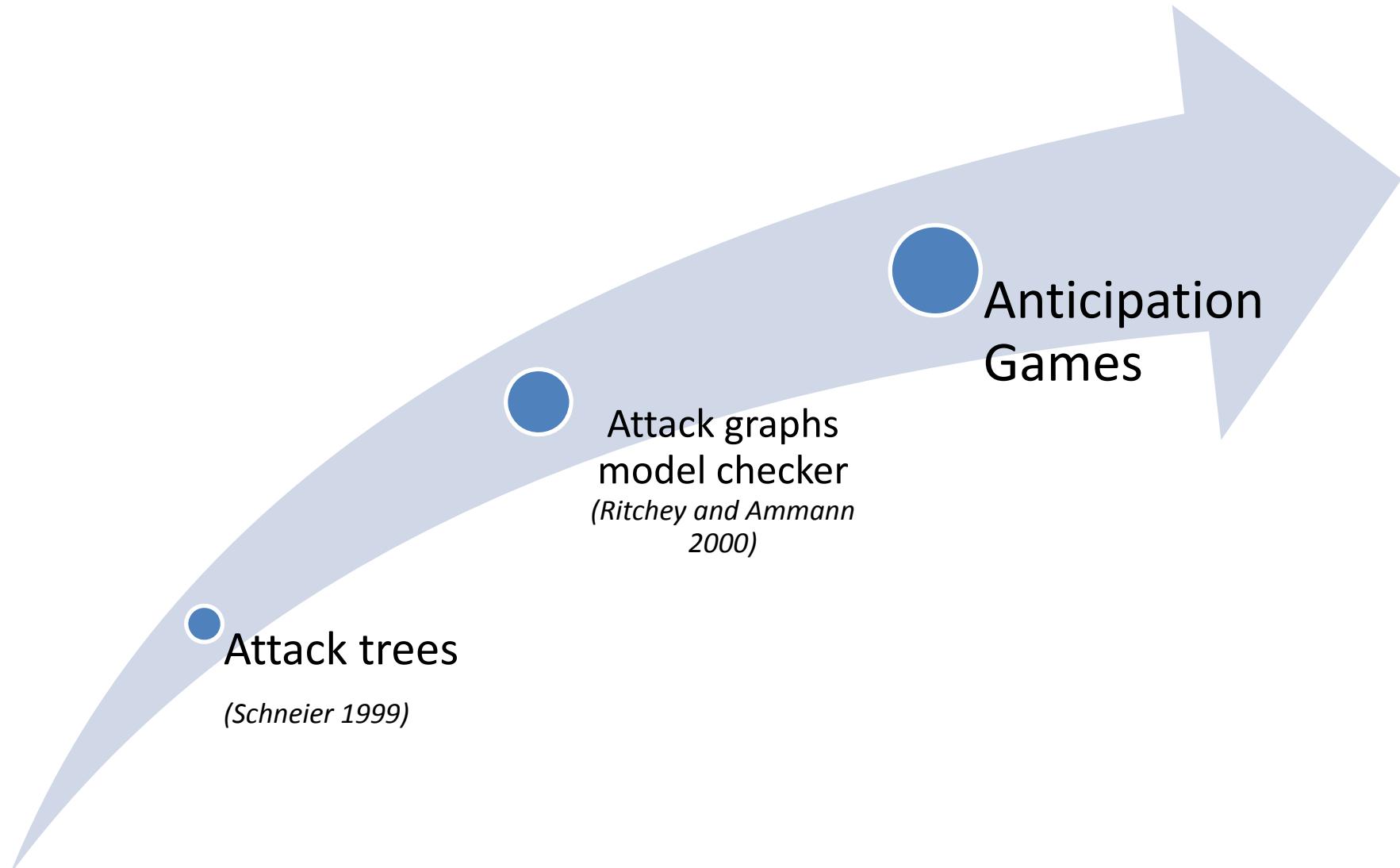
- Large network may suffer multiple vulnerabilities
- Patches and counter-measures need to be prioritized
- A minor vulnerability can turn into a major hole when used as a step-stone

Attack graph allows to reason
about attack sequences





Sandia Red Team “White Board” attack graph from DARPA CC20008 Information battle space preparation experiment



Attack graph

- Model checker-based (Ritchey et. al S&P'00, Sheyner et. al S&P'02)
- Graph-based (Ammann et. al CCS'02, Ritchey et. al ACSAC'02, Noel et. al ACSAC'03, Wang et. al ESORICS'05, Wang et. al DBSEC'06)

Timed Game

- ATL (Alur et al. 97)
- The Element of Surprise in Timed Games (De Alfaro et al. CONCUR 2003)
- TATL (Henzinger et al 2006 Formats)

Dependency

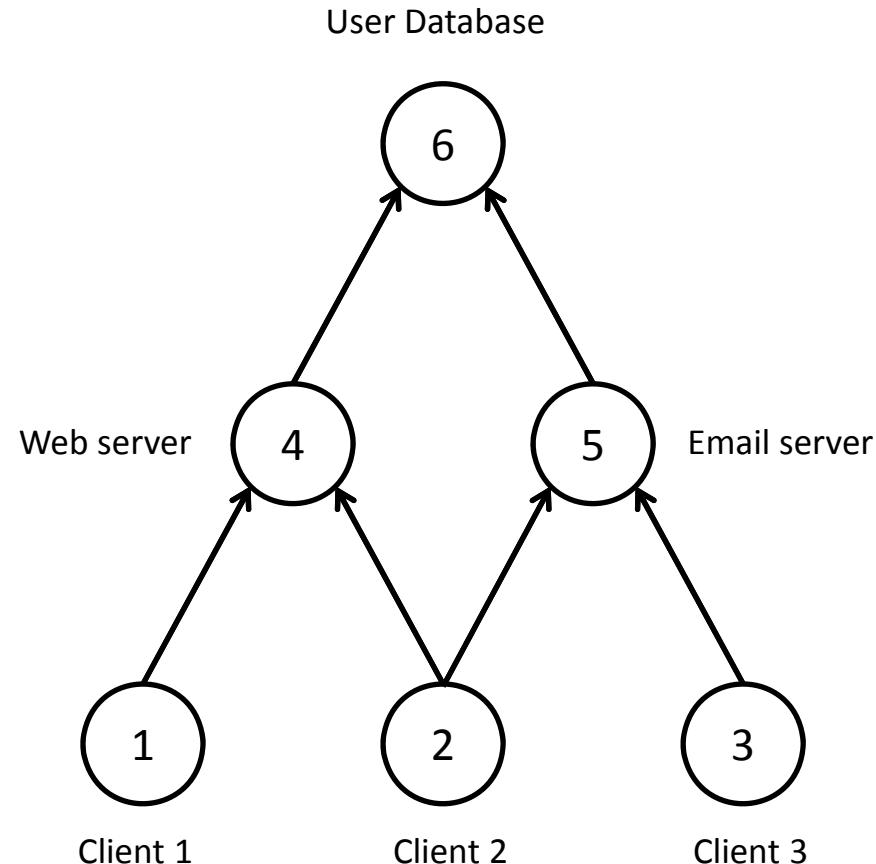
- Collateral effects
- Trust relations

Interaction

- Administrator
- Intruder

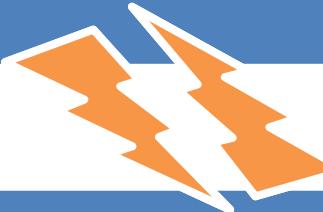
Time

- Action take time

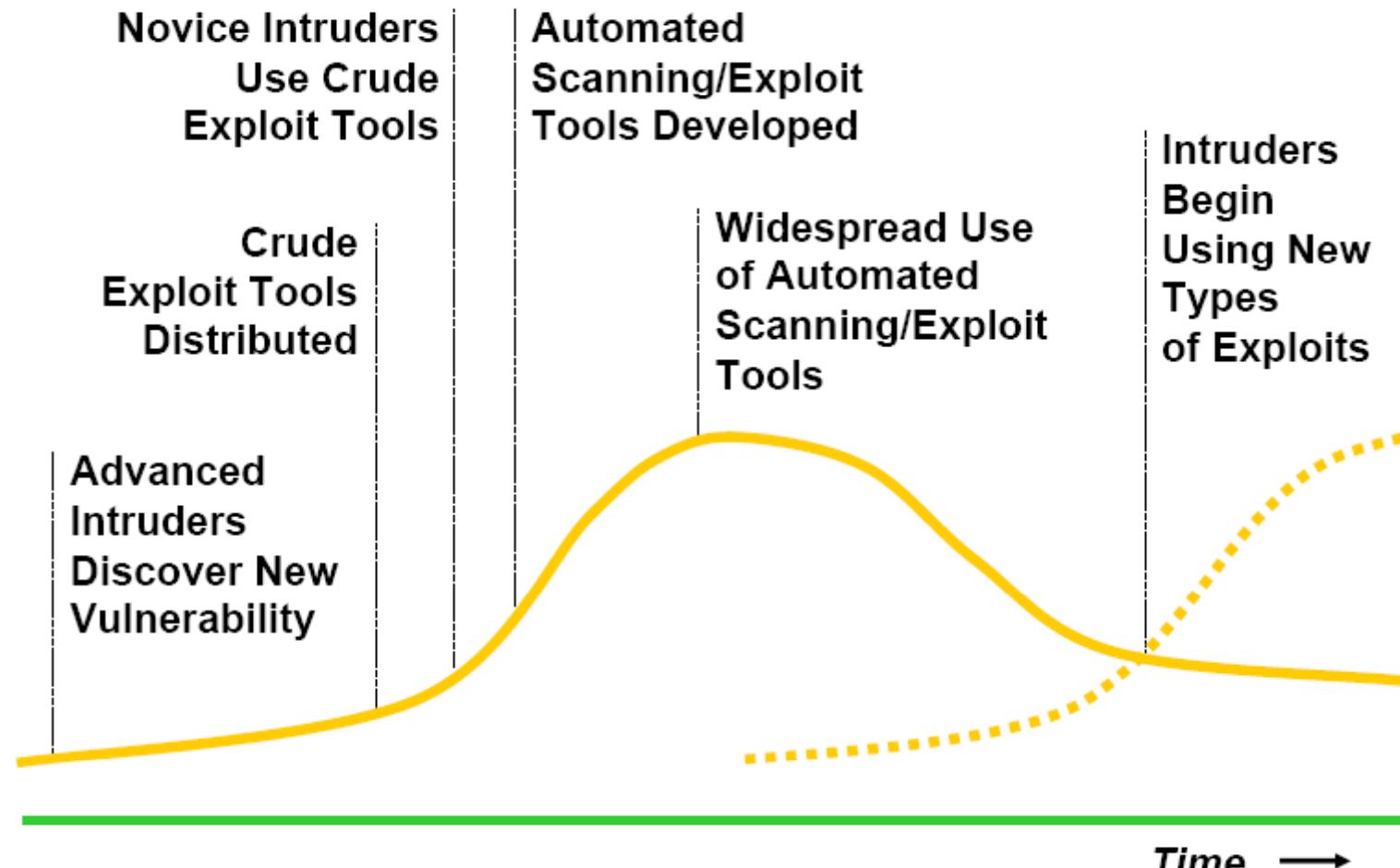




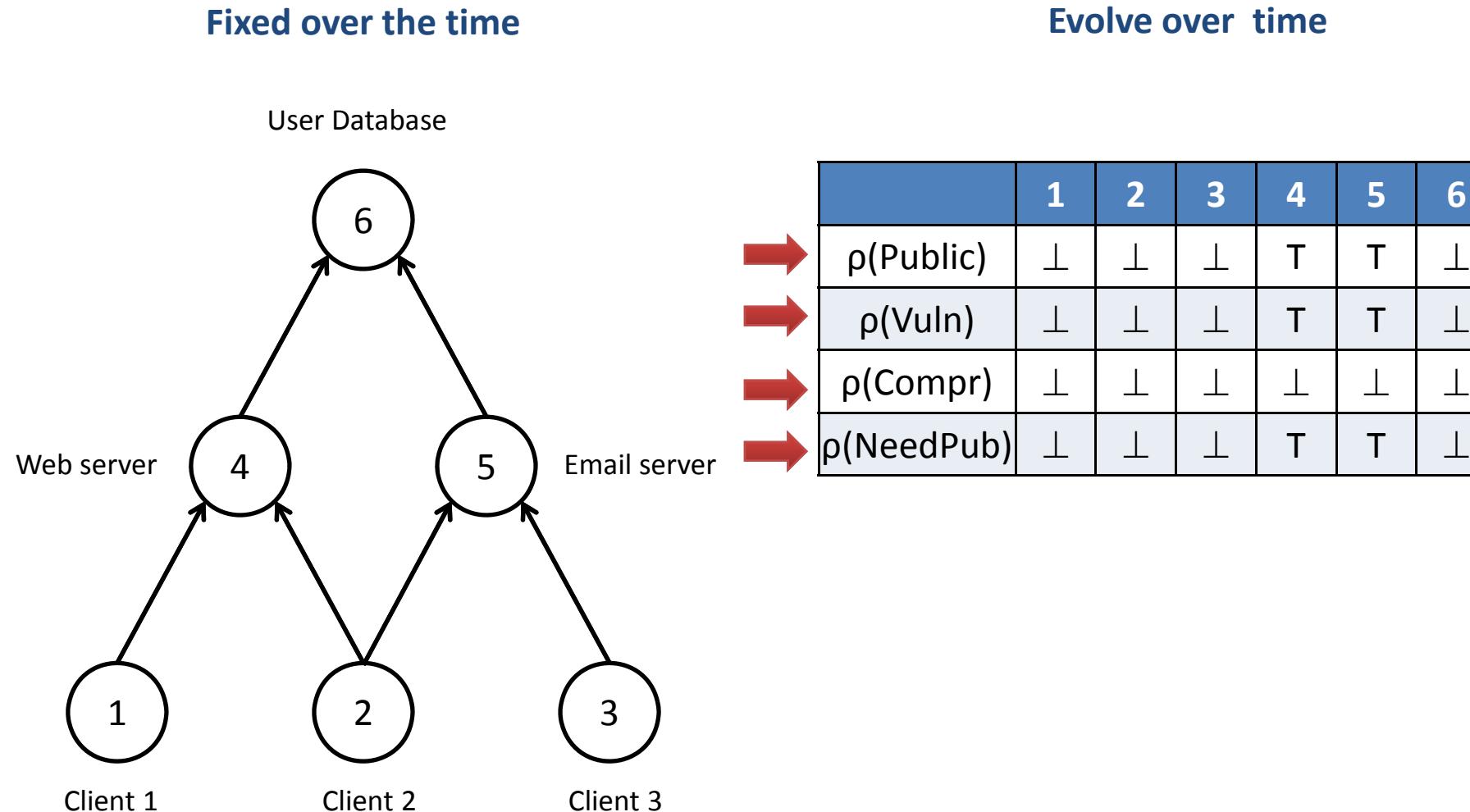
Exploit vulnerabilities
Abuse trust relations



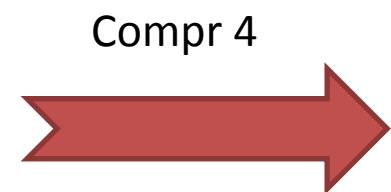
Patch
Firewall
Restore



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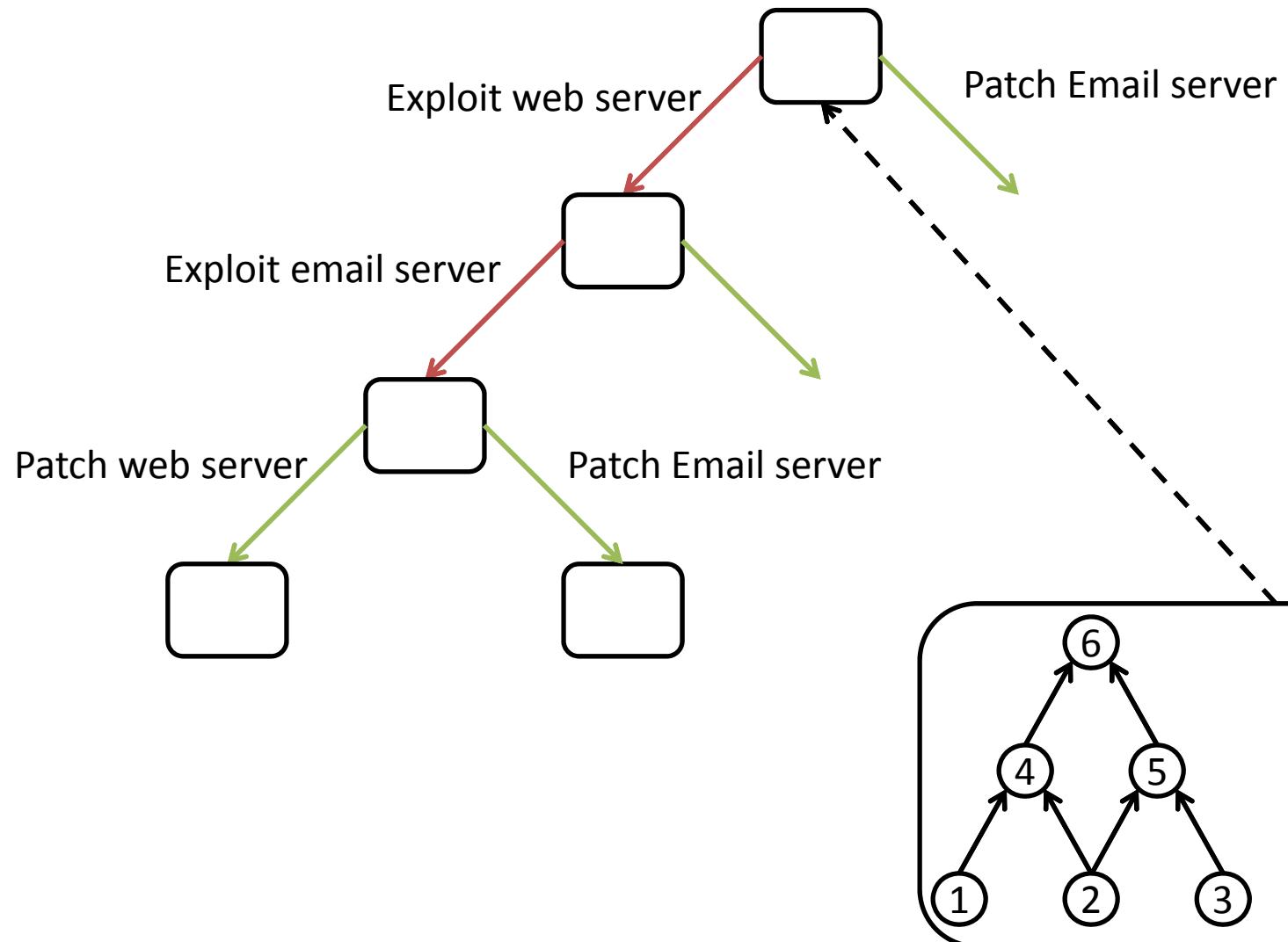


| | 1 | 2 | 3 | 4 | 5 | 6 |
|---------------------------------|----------|----------|----------|----------|----------|----------|
| $\rho(\text{Public})$ | \perp | \perp | \perp | T | T | \perp |
| $\rho(\text{Vuln})$ | \perp | \perp | \perp | T | T | \perp |
| $\rho(\text{Compr})$ | \perp | \perp | \perp | \perp | \perp | \perp |
| $\rho(\text{NeedPu} \text{ b})$ | \perp | \perp | \perp | T | T | \perp |



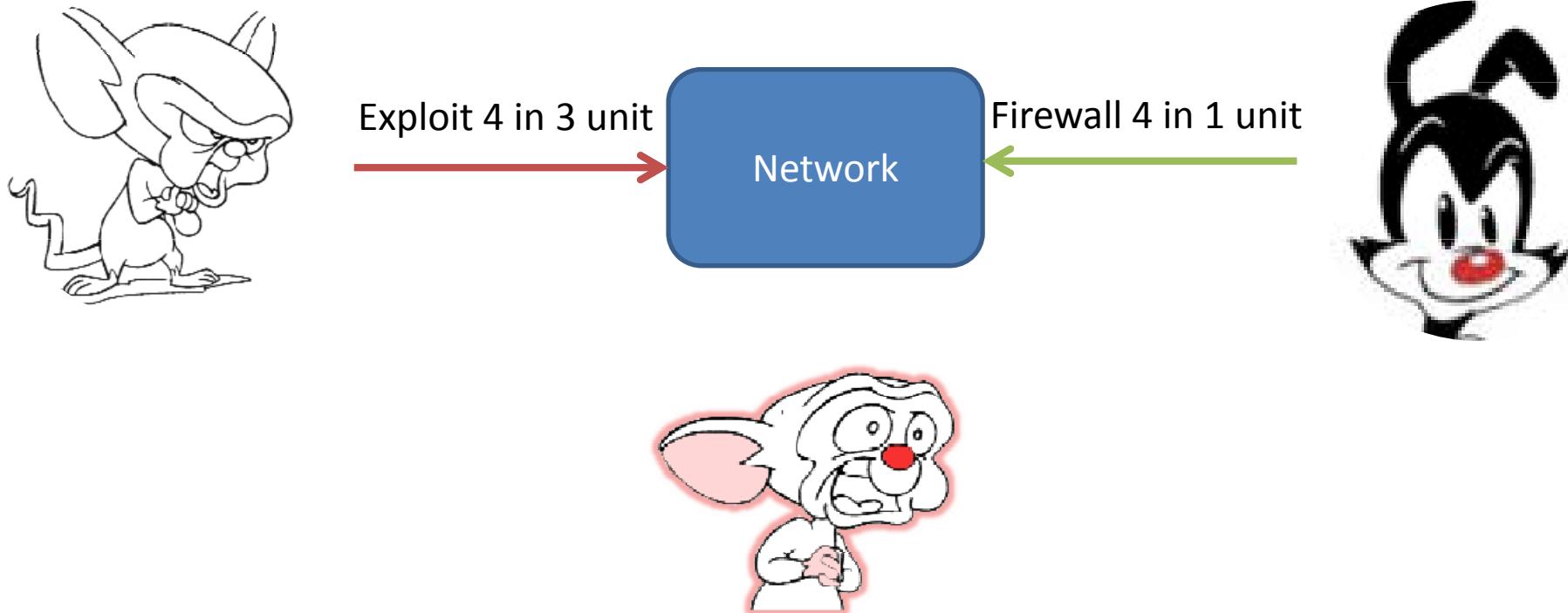
| | 1 | 2 | 3 | 4 | 5 | 6 |
|---------------------------------|----------|----------|----------|----------|----------|----------|
| $\rho(\text{Public})$ | \perp | \perp | \perp | T | T | \perp |
| $\rho(\text{Vuln})$ | \perp | \perp | \perp | T | T | \perp |
| $\rho(\text{Compr})$ | \perp | \perp | \perp | T | \perp | \perp |
| $\rho(\text{NeedPu} \text{ b})$ | \perp | \perp | \perp | T | T | \perp |

A Incomplete Game Example



- Each action requires a different amount of time
 - Patching a service: Download, extract, apply, restart
 - Exploit a service
 - Firewalling a service
- In anticipation games as in TATL the fastest action win
- Player can be taken by surprise

The element of surprise



- Anticipation games allows to model
 - Denial of service
 - Buffer overflow execution
 - Permission abuse
 - Cross-scripting
 - Information leak
 -

| | | | |
|-----|-------|-----------------------|---------------------------------------|
| F | $::=$ | A | atomic propositions, in \mathcal{A} |
| | | \top | true |
| | | $\neg F$ | negation |
| | | $F \wedge F$ | conjunction |
| | | $\Diamond F$ | |
| | | $\Diamond_{\equiv} F$ | |

$\vdash \Diamond Compr$

A successor node is compromised



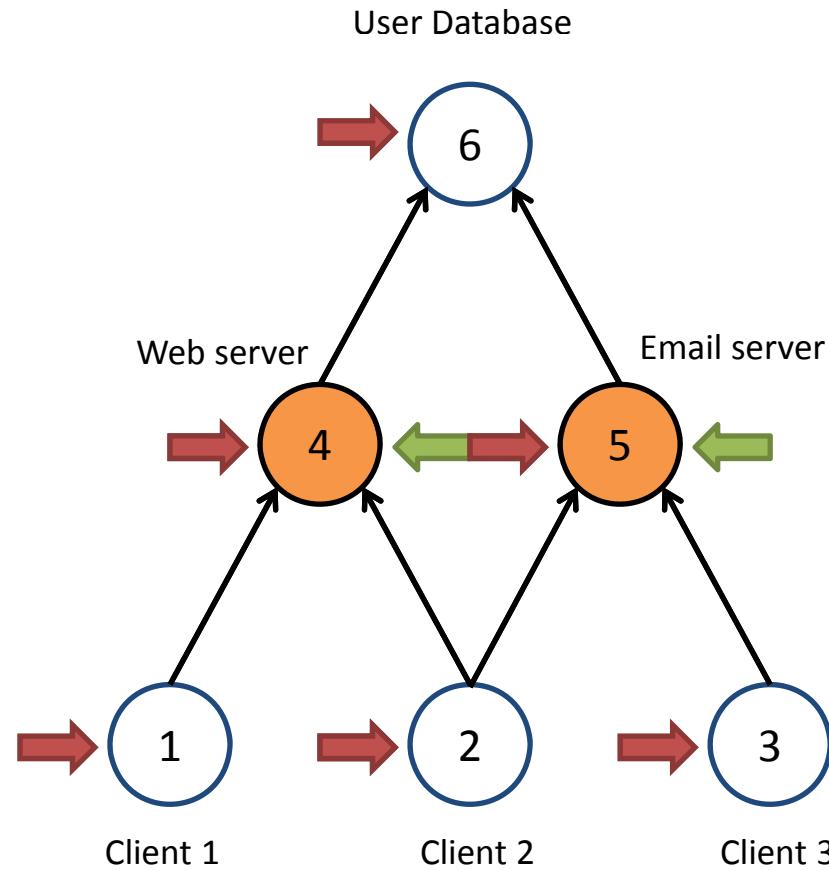
$\vdash \Diamond_{\equiv} Public$

At least, one of the nodes the belongs to the equivalence is public



| | | | |
|---|------------------------------|------------------------------|-------------------------------|
| Pre $Vuln \wedge Public \wedge \neg Compr$ | $(2,I,Compromise\ 0day)$ | $\xrightarrow{\hspace{1cm}}$ | $Compr$ |
| Pre $Vuln \wedge Public \wedge \neg Compr$ | $(7,I,Compromise\ public)$ | $\xrightarrow{\hspace{1cm}}$ | $Compr$ |
| Pre $\neg Compr \wedge \Diamond Compr$ | $(4,I,Compromise\ backward)$ | $\xrightarrow{\hspace{1cm}}$ | $Compr$ |
| Pre $Compr \wedge \Diamond \neg Compr$ | $(4,I,Compromise\ forward)$ | $\xrightarrow{\hspace{1cm}}$ | $\Diamond Compr$ |
| Pre $Public \wedge Vuln$ | $(1,A,Firewall)$ | $\xrightarrow{\hspace{1cm}}$ | $\neg Public$ |
| Pre $Public \wedge \neg Vuln \wedge NeedPub$ | $(1,A,UnFirewall)$ | $\xrightarrow{\hspace{1cm}}$ | $Public$ |
| Pre $Vuln \wedge \neg Compr$ | $(3,A,Patch)$ | $\xrightarrow{\hspace{1cm}}$ | $\neg Vuln \wedge \neg Compr$ |

A Play example



| Player | Action | Rule | Target | Succ |
|----------|---------|--------------------------------|--------|------|
| Admin | Execute | PathForward | 5 | |
| Intruder | Execute | Compromise Edoward Backward | 5 | 5 |

| | |
|---|---------------------------------------|
| $\varphi ::= A$ | atomic propositions, in \mathcal{A} |
| | |
| $\neg\varphi$ | |
| | |
| $\varphi \wedge \varphi$ | |
| | |
| $\Diamond\varphi$ | |
| | |
| $\Diamond_{\equiv}\varphi$ | |
| | |
| $x + d_1 \leq y + d_2$ | clock constraints |
| | |
| $x \cdot \varphi$ | freeze |
| | |
| $\langle\!\langle \mathfrak{P} \rangle\!\rangle \blacksquare \varphi$ | invariant |
| | |
| $\langle\!\langle \mathfrak{P} \rangle\!\rangle \varphi_1 \cup \varphi_2$ | eventually |

We abbreviate $\langle\!\langle \mathfrak{P} \rangle\!\rangle \text{TRUE} \cup \varphi$ as $\langle\!\langle \mathfrak{P} \rangle\!\rangle \blacklozenge \varphi$.

$\vdash \langle\!\langle A \rangle\!\rangle \varphi$

The player A have a strategy to satisfy
the property φ

$\vdash \blacksquare Compr$

In every future the node will be
compromised

$$\langle\!\langle A \rangle\!\rangle \blacksquare \lozenge \equiv \neg \text{Compr}$$

TATL Formula model checking in
Anticipation game is decidable and
EXPTIME-complete

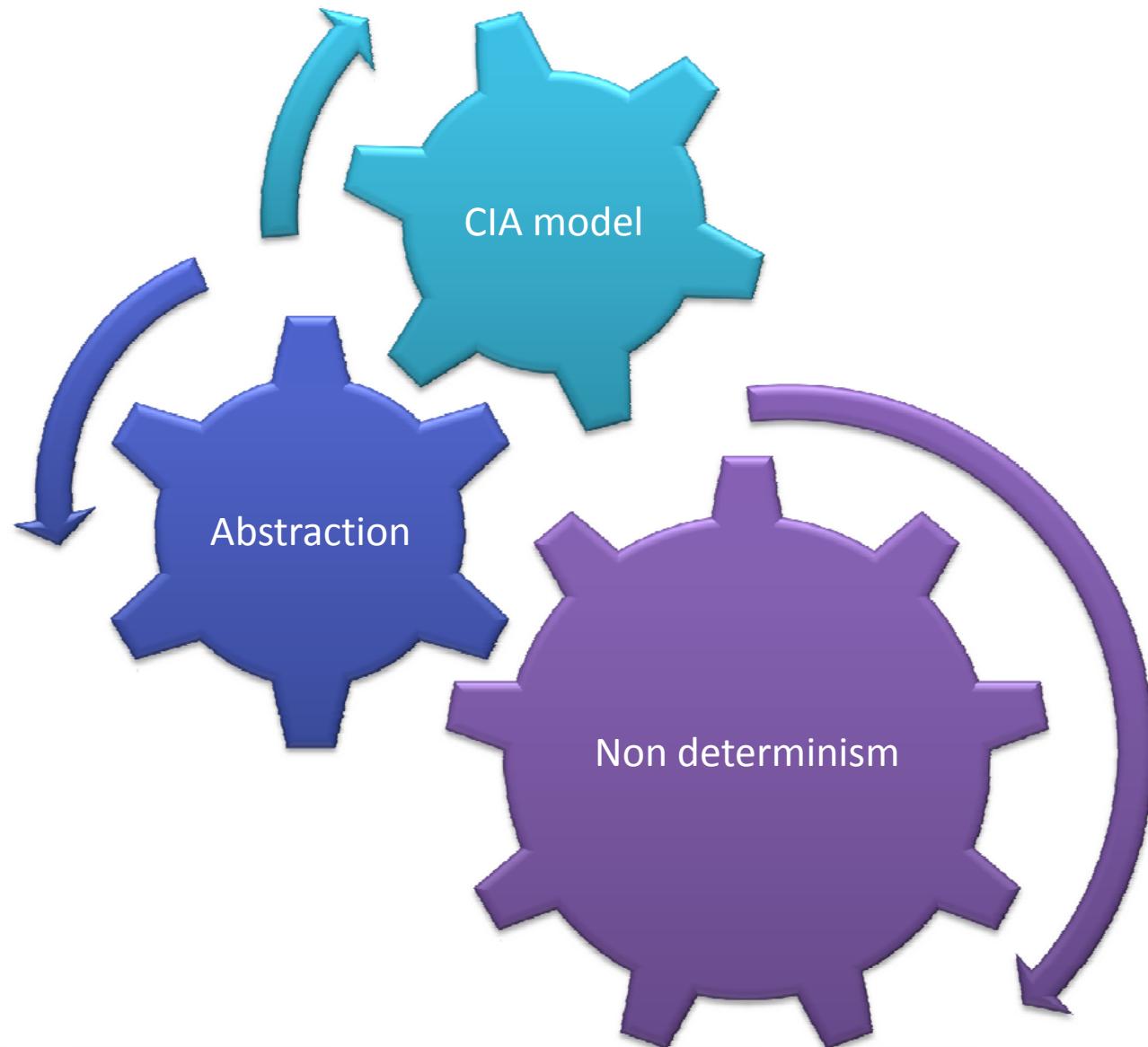
One More Thing !

- Model and Strategies are fully implemented in C
- The talk example cannot be analyzed by hand
 - 4011 plays
 - 40825 states



Analyzer Demo





- Game and Time provide a richer model for intrusion analysis
- Many directions to explore



During this work no network service was injured or tortured.

