

# Privacy Enhancing Credential



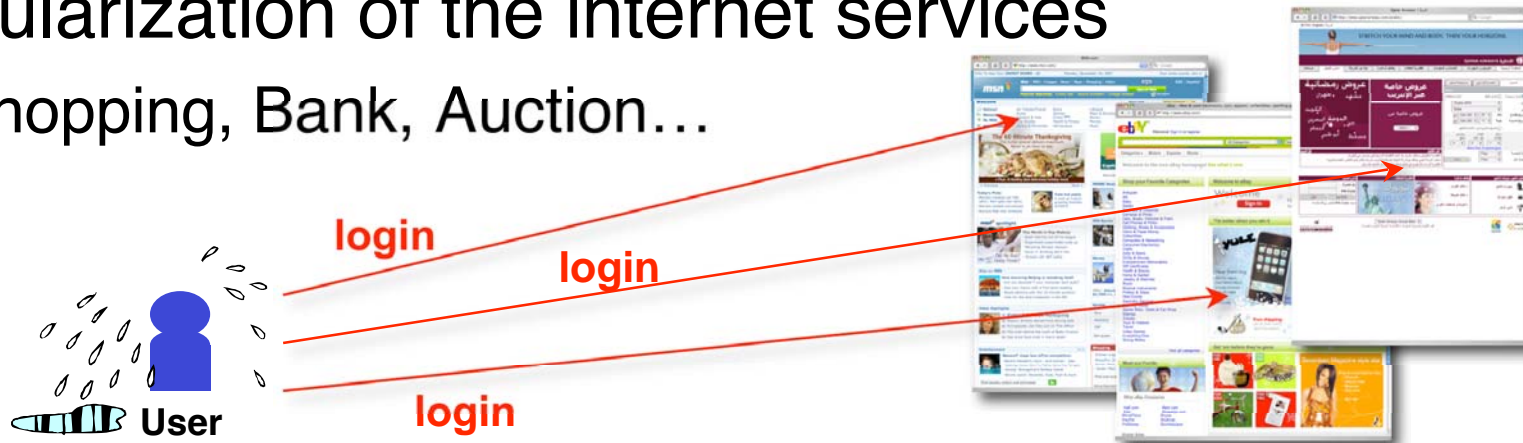
National Institute of  
Information and Communications Technology  
(Japan)

○ Junji Nakazato  
Lihua Wang  
Akihiro Yamamura

# Motivation

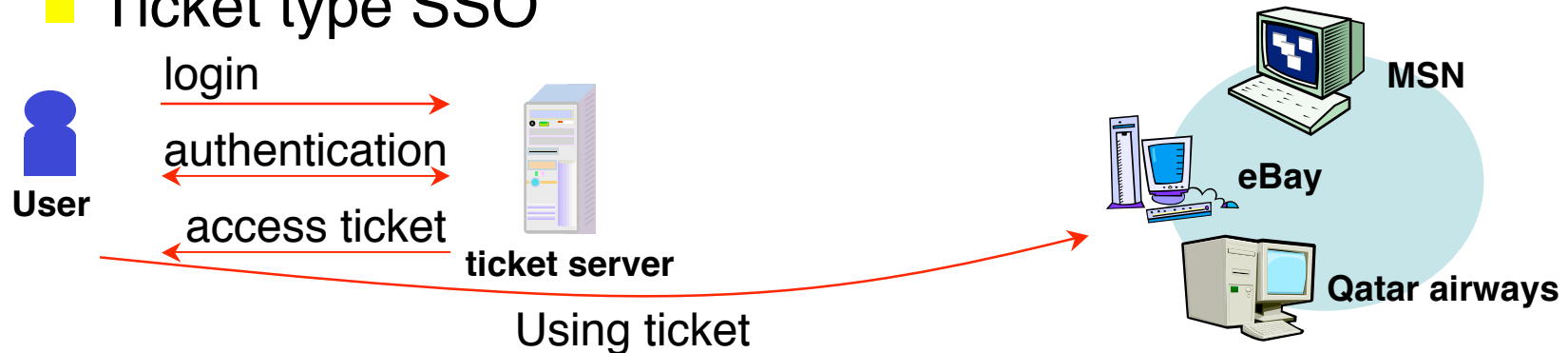
◆ Popularization of the internet services

- Shopping, Bank, Auction...



◆ Using single sign on (SSO) technique

- Ticket type SSO

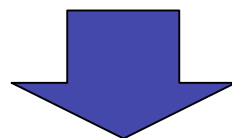


# Motivation

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## ◆ Disadvantages of SSO

- The same ticket is used for multiple services
- The user's privacy is obtained by the collusion with services
  - Use frequency of service
  - An order of using service
- The user can transfer ticket to another user

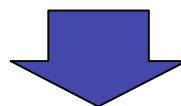


**Needs to think about privacy!!**

# Our goal

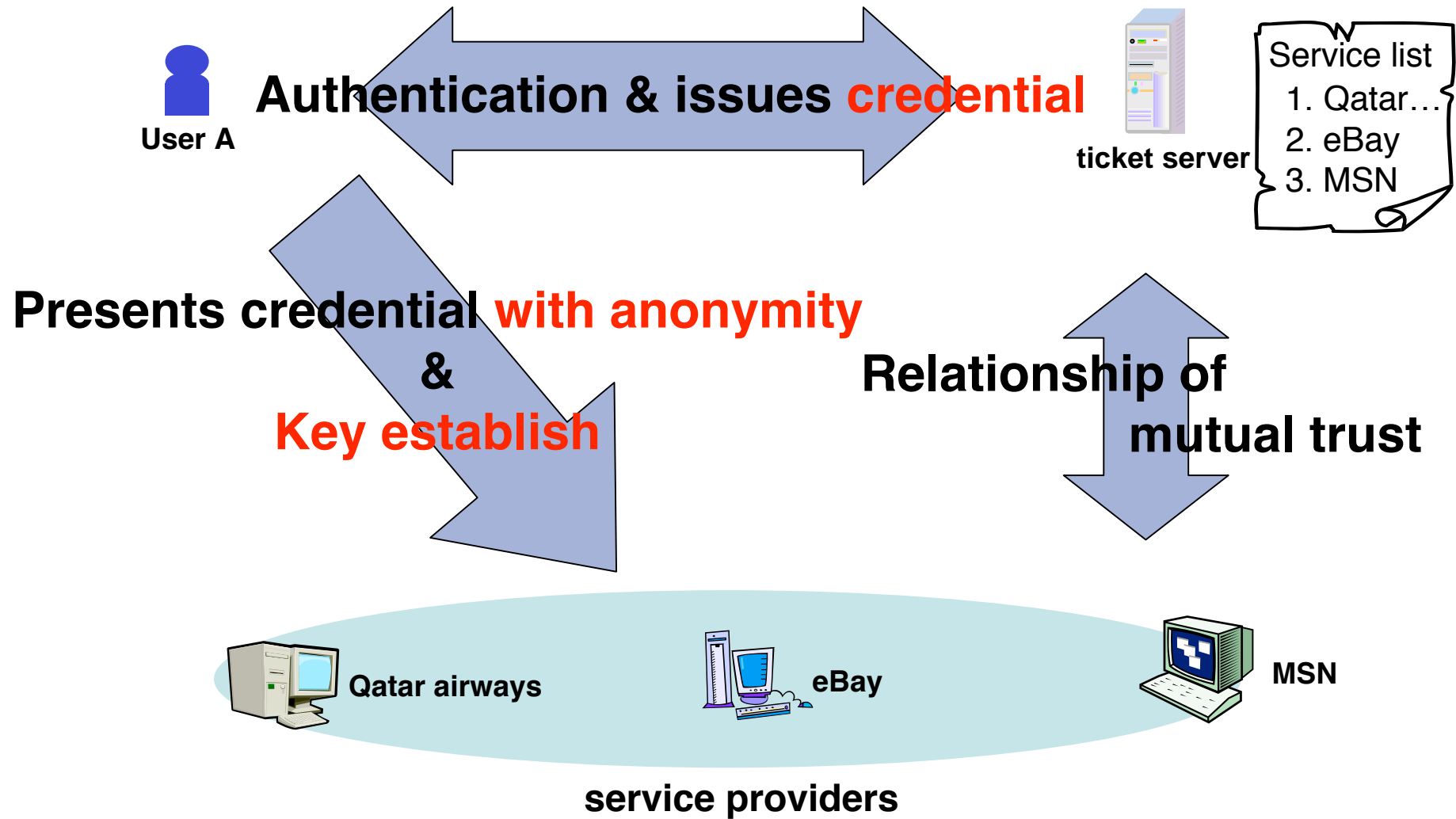
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- ◆ Multiple logins are not needed
  - The user only presents credential of whether he has the right to access service
- ◆ User's privacy is concealed
  - The user can access services with anonymity
- ◆ Credential cannot be transferred to anyone
  - No one can transfer correct credential to others
- ◆ Authenticated key exchange
  - To provide secure channel between user and service provider



**Try to apply “credential system”**

# Our goal



# Credential system

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## ◆ Previous work

- “Designated Group Credentials”
- Ching Yu Ng, Willy Susilo, Yi Mu
- ASIACCS 2006
  
- Using pairing technique
- Authority can designate the verifiers
  - Ticket issuer can designate the service providers
- The user **authentication is necessary** for the outside
- The authority **can trace user’s movement**
- The user **can transfer correct credential** to others

# Comparison of requirements

	Group credential [4]	Our proposal
Unforgeability	Yes	Yes
Designated	Yes	Yes
Non-transferability	No	Yes
Anonymity	No	Yes
Unlinkability	No	Yes

# Technique

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## ◆ Based on pairing

- Bilinear: Given any  $Q, R$  in  $G_1$  and  $a, b$  in  $Z_q$ , we have  $e(aQ, bR) = e(Q, R)^{ab}$
- Non-degenerate:  $e(P, P) \neq 1$
- Computable: There is an efficient algorithm to compute  $e(Q, R)$  for any  $Q, R$  in  $G_1$

## ◆ Non-transferability

- private key of user is included in the credential

## ◆ Unlinkability

- Randomize credential when he uses it

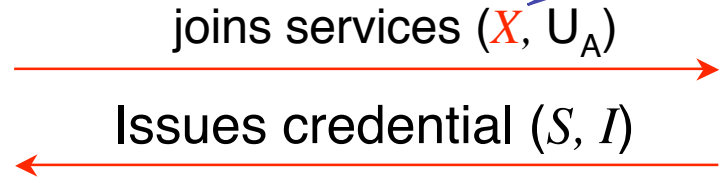


# Our proposed scheme

Including identity  
(private key)



User A



ticket server

- Service list
1. Qatar...
  2. eBay
  3. MSN

Computes credentials for each service and **randomize** them.

$(S, I) \rightarrow$

- $(\sigma_1, M_1, A_1, A_2, B_1, B_2, seed_1)$
- $(\sigma_2, M_2, A_1, A_2, B_1, B_2, seed_2)$

Computes credential ( $S, I$ ) **form**  $X$  to use services 1 and 3 for user A (designated).



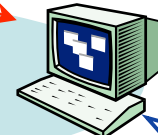
Authenticated key sharing using  $M(\text{credential})$



Qatar airways



eBay



MSN

service providers

Check the validity of credential

Check the validity of credential

# Conclusion

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- ◆ Propose privacy enhancing credential
  - We preserved unlinkability (anonymity).
  - We satisfied non-transferability.
  - We achieved authenticated key exchange.
  
- ◆ We can provide time restriction function
  - It can be achieved by a few modification.
  - Change generator  $F$  to  $h(t)$
  - $h()$  : hash function ( $h(*) \rightarrow G_1$ )

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# Thanks

*NiCT*

独立行政法人

情報通信研究機構

National Institute of Information and Communications Technology

