

# Robotic Motion Planning: Bug Algorithms

(with some discussion on curve tracing and sensors)

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# What's Special About Bugs

- Many planning algorithms assume global knowledge
- Bug algorithms assume only *local* knowledge of the environment and a global goal
- Bug behaviors are simple:
  - 1) Follow a wall (right or left)
  - 2) Move in a straight line toward goal
- Bug 1 and Bug 2 assume essentially tactile sensing
- Tangent Bug deals with finite distance sensing

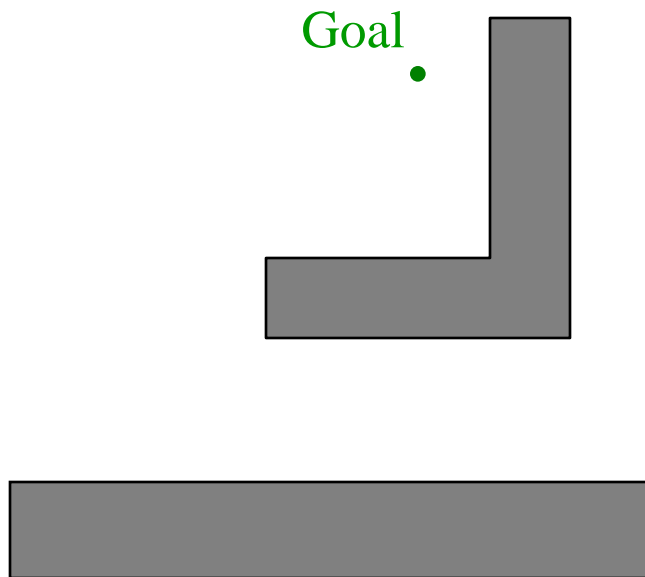
# A Few General Concepts

- Workspace  $W$ 
  - $\mathcal{R}(2)$  or  $\mathcal{R}(3)$  depending on the robot
  - could be infinite (open) or bounded (closed/compact)
- Obstacle  $WO_i$
- Free workspace  $W_{free} = W \setminus \cup_i WO_i$

# The *Bug* Algorithms

provable results...

## Insect-inspired



• Start

- **known direction to goal**

- **robot can measure distance  $d(x,y)$  between pts  $x$  and  $y$**

- **otherwise local sensing**

- walls/obstacles & encoders

- **reasonable world**

- 1) finitely many obstacles in any finite area

- 2) a line will intersect an obstacle finitely many times

- 3) Workspace is bounded

$$W \subset B_r(x), r < \infty$$

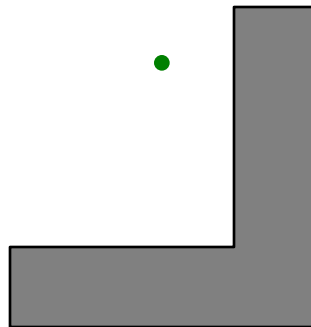
$$B_r(x) = \{ y \in \mathcal{R}(2) \mid d(x,y) < r \}$$

# Buginner Strategy

"Bug 0" algorithm

- known direction to goal
- otherwise local sensing

walls/obstacles & encoders



Some notation:

$q_{\text{start}}$  and  $q_{\text{goal}}$

"hit point"  $q_i^H$

"leave point"  $q_i^L$

A *path* is a sequence of hit/leave pairs bounded by  $q_{\text{start}}$  and  $q_{\text{goal}}$

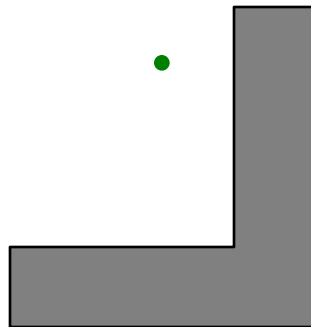
# Buginner Strategy

"Bug 0" algorithm

- known direction to goal

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walls/obstacles & encoders



1) head toward goal

2) follow obstacles until you can head toward the goal again

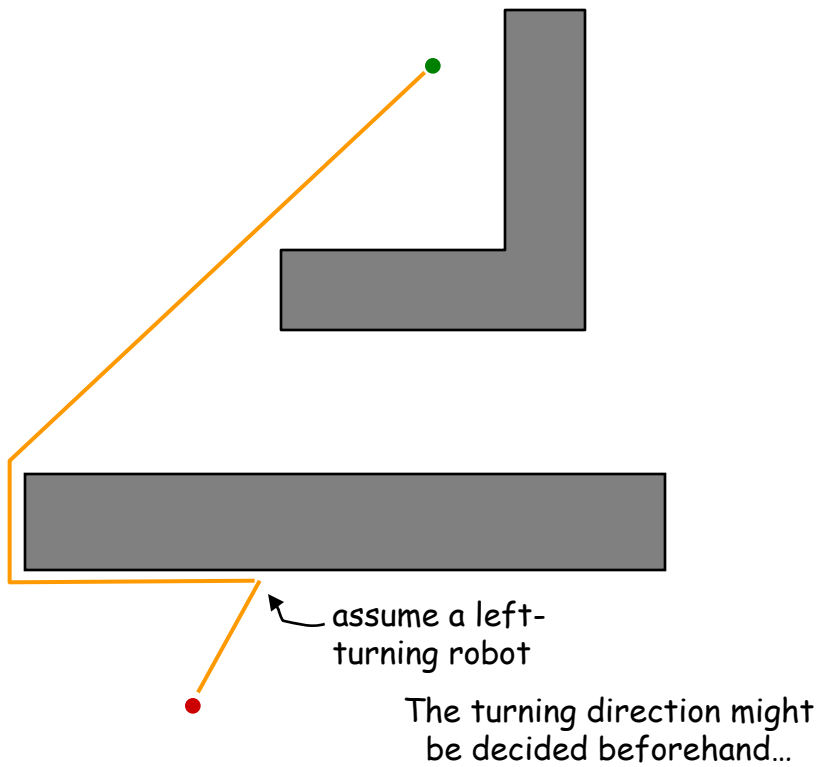
3) continue



# Buginner Strategy

## "Bug 0" algorithm

- 1) head toward goal
- 2) follow obstacles until you can head toward the goal again
- 3) continue



# Bug Zapper

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What map will foil Bug 0 ?

"Bug 0" algorithm

- 1) head toward goal
- 2) follow obstacles until you can head toward the goal again
- 3) continue

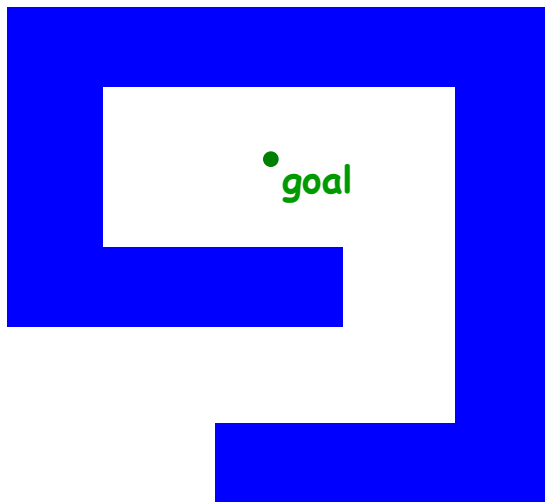


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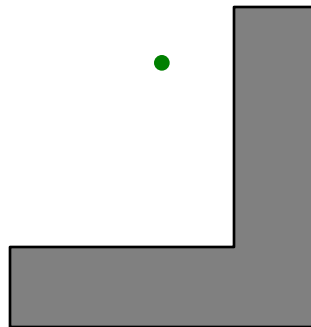


• start

# A better bug?

But add some memory!

- known direction to goal
  - otherwise local sensing
- walls/obstacles & **encoders**



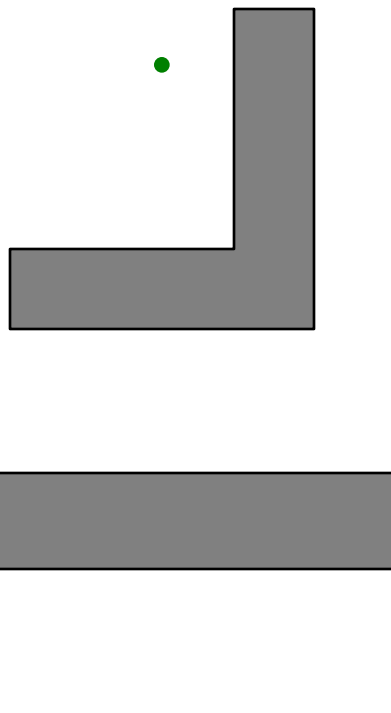


# Bug 1

But some computing power!

- known direction to goal
- otherwise local sensing

walls/obstacles & **encoders**



## "Bug 1" algorithm

- 1) head toward goal
- 2) if an obstacle is encountered, circumnavigate it *and* remember how close you get to the goal
- 3) return to that closest point (by wall-following) and continue

Vladimir Lumelsky & Alexander Stepanov: Algorithmica 1987

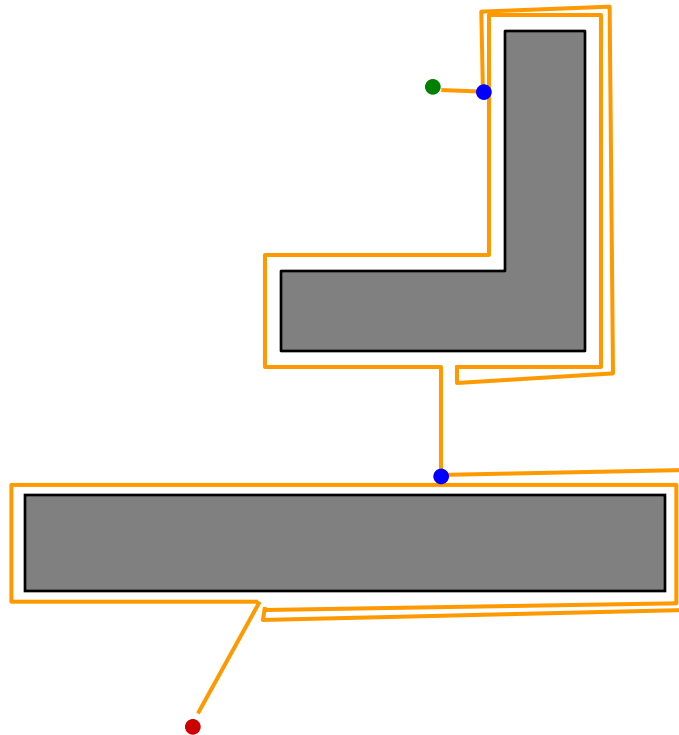


# Bug 1

But some computing power!

- known direction to goal
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## "Bug 1" algorithm

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16-735, Howie Choset with slides from G.D. Hager and Z. Dodds

# BUG 1 More formally

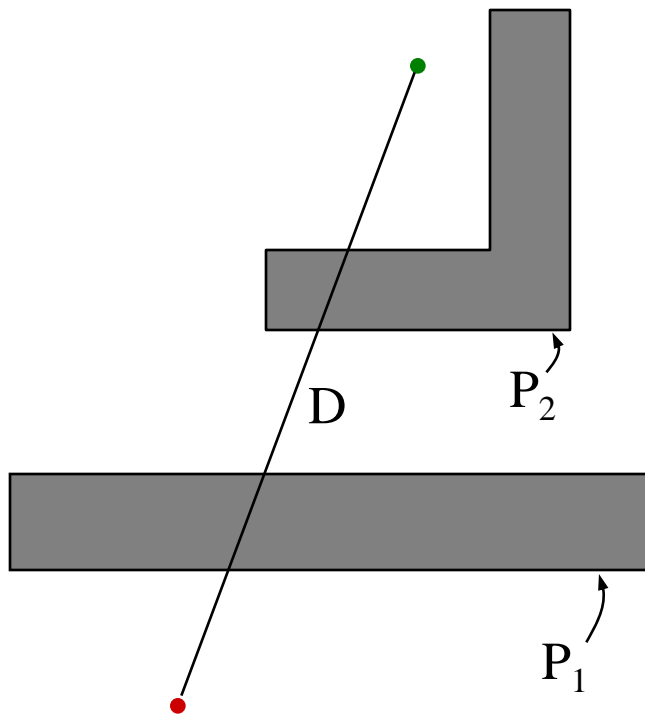
- Let  $q_0^L = q_{\text{start}}$ ;  $i = 1$
- repeat
  - repeat
    - from  $q_{i-1}^L$  move toward  $q_{\text{goal}}$
  - until goal is reached or obstacle encountered at  $q_i^H$
  - if goal is reached, exit
  - repeat
    - follow boundary recording pt  $q_i^L$  with shortest distance to goal
  - until  $q_{\text{goal}}$  is reached or  $q_i^H$  is re-encountered
  - if goal is reached, exit
  - Go to  $q_i^L$
  - if move toward  $q_{\text{goal}}$  moves into obstacle
    - exit with failure
  - else
    - $i=i+1$
    - continue

# “Quiz”

# Bug 1 analysis

## Bug 1: Path Bounds

What are upper/lower bounds on the path length that the robot takes?



$D$  = straight-line distance from start to goal

$P_i$  = perimeter of the  $i$ th obstacle

Lower bound:

What's the shortest distance it might travel?

Upper bound:

What's the longest distance it might travel?

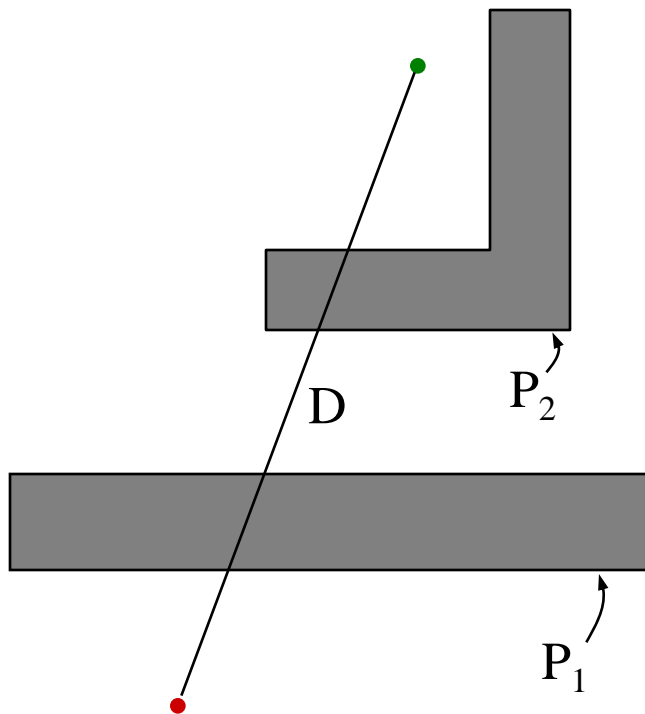
What is an environment where your upper bound is required?

# “Quiz”

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$D$

Upper bound:

What's the longest distance it might travel?

$$D + 1.5 \sum_i P_i$$

What is an environment where your upper bound is required?

# How Can We Show Completeness?

- An algorithm is *complete* if, in finite time, it finds a path if such a path exists or terminates with failure if it does not.
- Suppose BUG1 were incomplete
  - Therefore, there is a path from start to goal
    - By assumption, it is finite length, and intersects obstacles a finite number of times.
  - BUG1 does not find it
    - Either it terminates incorrectly, or, it spends an infinite amount of time
    - Suppose it never terminates
      - but each leave point is closer to the obstacle than corresponding hit point
      - Each hit point is closer than the last leave point
      - Thus, there are a finite number of hit/leave pairs; after exhausting them, the robot will proceed to the goal and terminate
    - Suppose it terminates (incorrectly)
    - Then, the closest point after a hit must be a leave where it would have to move into the obstacle
      - But, then line from robot to goal must intersect object even number of times (Jordan curve theorem)
      - But then there is another intersection point on the boundary closer to object. Since we assumed there is a path, we must have crossed this pt on boundary which contradicts the definition of a leave point.



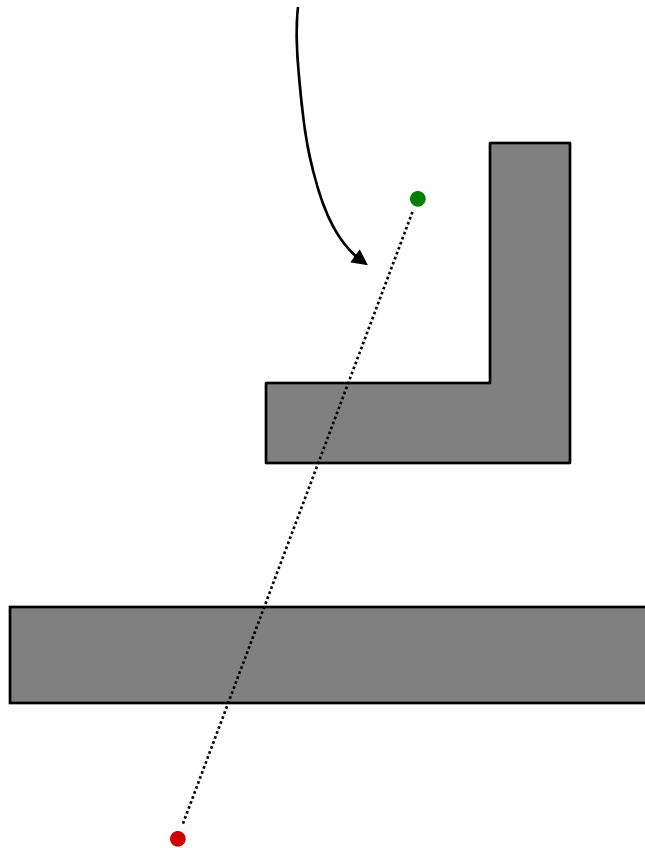
# Another step forward?

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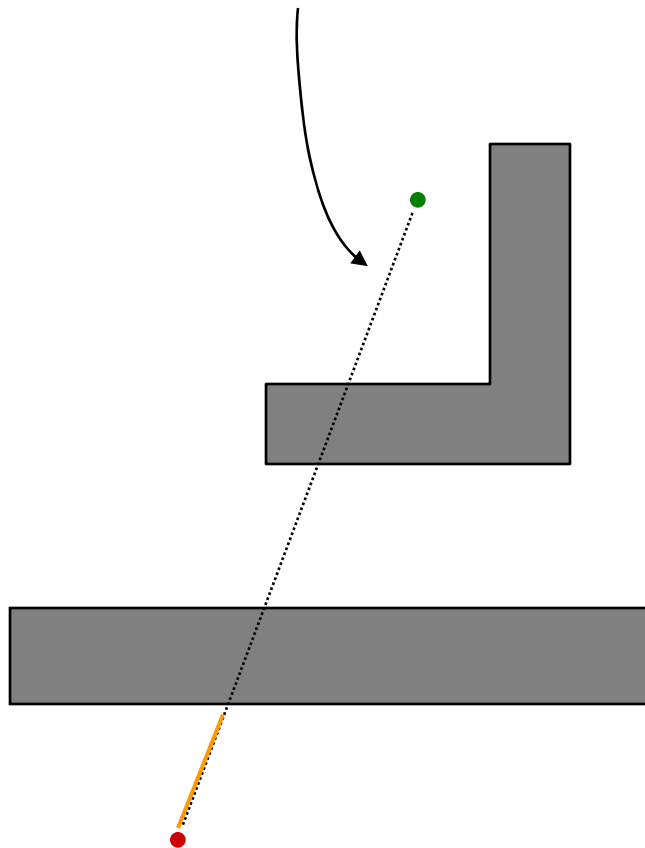
Call the line from the starting point to the goal the *m-line*

"Bug 2" Algorithm



# A better bug?

Call the line from the starting point to the goal the *m-line*

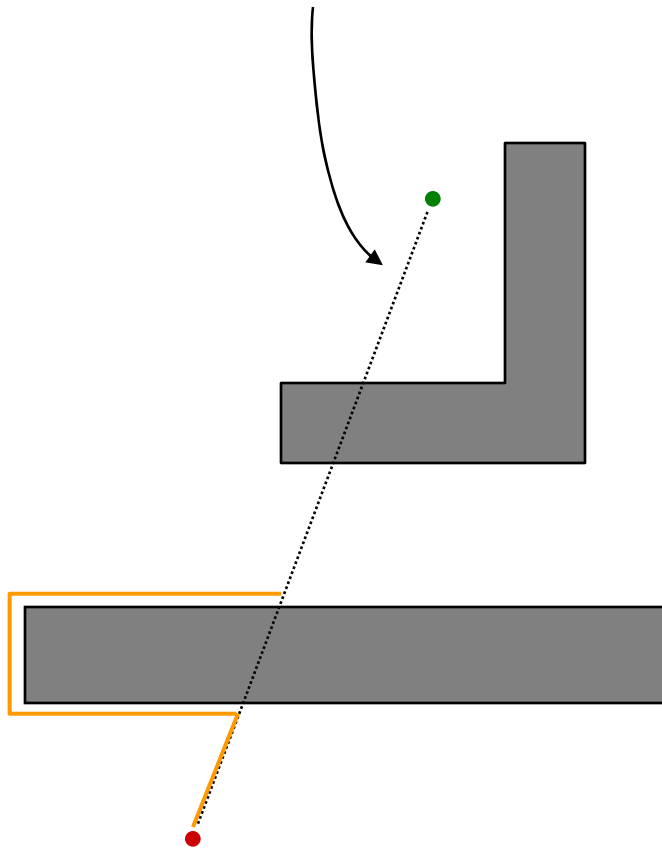


"Bug 2" Algorithm

1) head toward goal on the *m-line*

# A better bug?

Call the line from the starting point to the goal the *m-line*

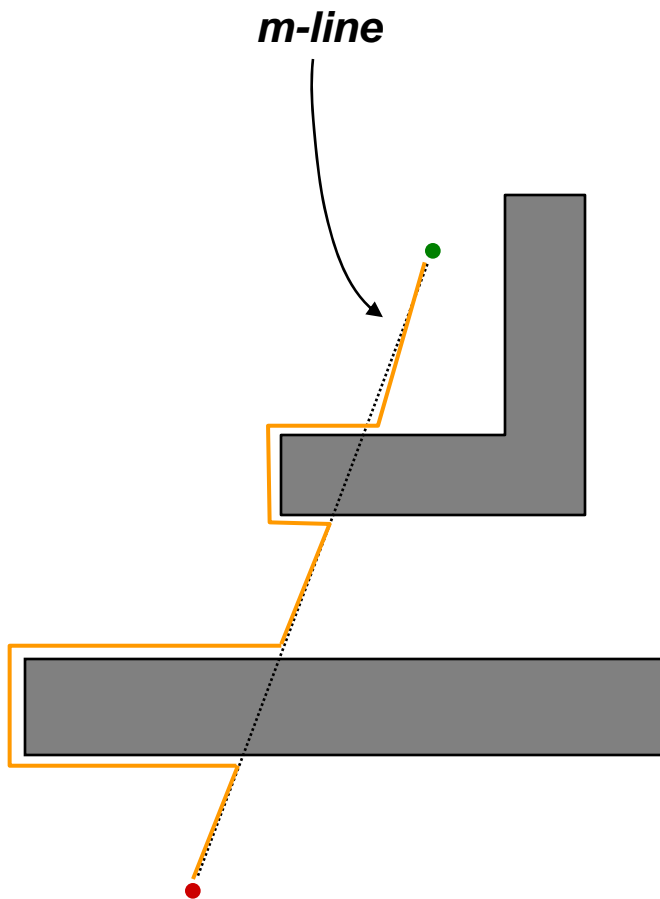


## "Bug 2" Algorithm

- 1) head toward goal on the *m-line*
- 2) if an obstacle is in the way, follow it until you encounter the *m-line* again.

# A better bug?

## "Bug 2" Algorithm

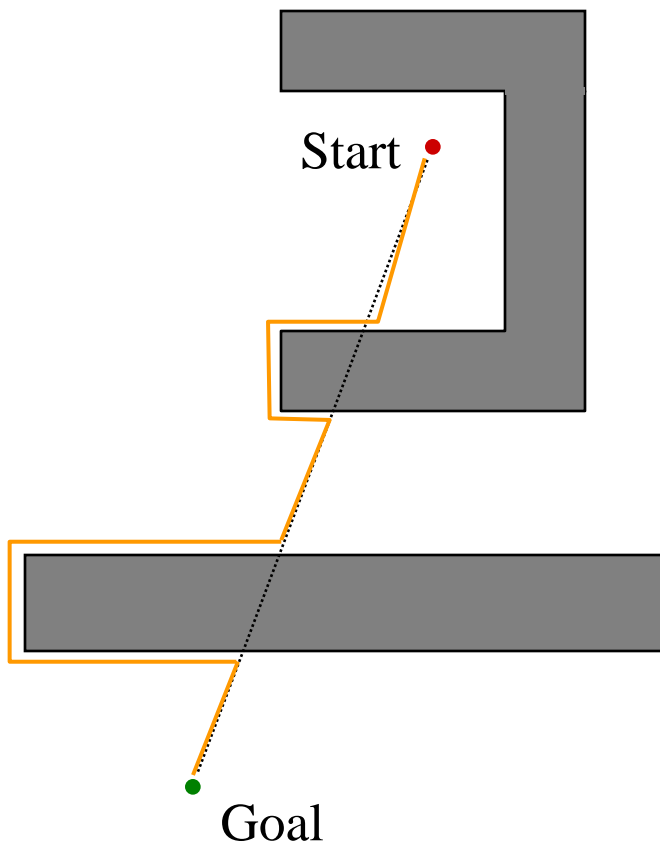


- 1) head toward goal on the *m-line*
- 2) if an obstacle is in the way, follow it until you encounter the *m-line* again.
- 3) Leave the obstacle and continue toward the goal



# A better bug?

## "Bug 2" Algorithm



- 1) head toward goal on the *m-line*
- 2) if an obstacle is in the way, follow it until you encounter the *m-line* again *closer to the goal*.
- 3) Leave the obstacle and continue toward the goal

Better or worse than Bug1?

# BUG 2 More formally

- Let  $q_0^L = q_{\text{start}}$ ;  $i = 1$
- repeat
  - repeat
    - from  $q_{i-1}^L$  move toward  $q_{\text{goal}}$  along the m-line
  - until goal is reached or obstacle encountered at  $q_i^H$
  - if goal is reached, exit
  - repeat
    - follow boundary
  - until  $q_{\text{goal}}$  is reached or  $q_i^H$  is re-encountered or m-line is re-encountered,  $x$  is not  $q_i^H$ ,  $d(x, q_{\text{goal}}) < d(q_i^H, q_{\text{goal}})$  and way to goal is unimpeded
  - if goal is reached, exit
  - if  $q_i^H$  is reached, return failure
  - else
    - $q_i^L = m$
    - $i=i+1$
    - continue

# head-to-head comparison

or thorax-to-thorax, perhaps

Draw worlds in which Bug 2 does better than Bug 1 (and vice versa).

**Bug 2** beats **Bug 1**

**Bug 1** beats **Bug 2**

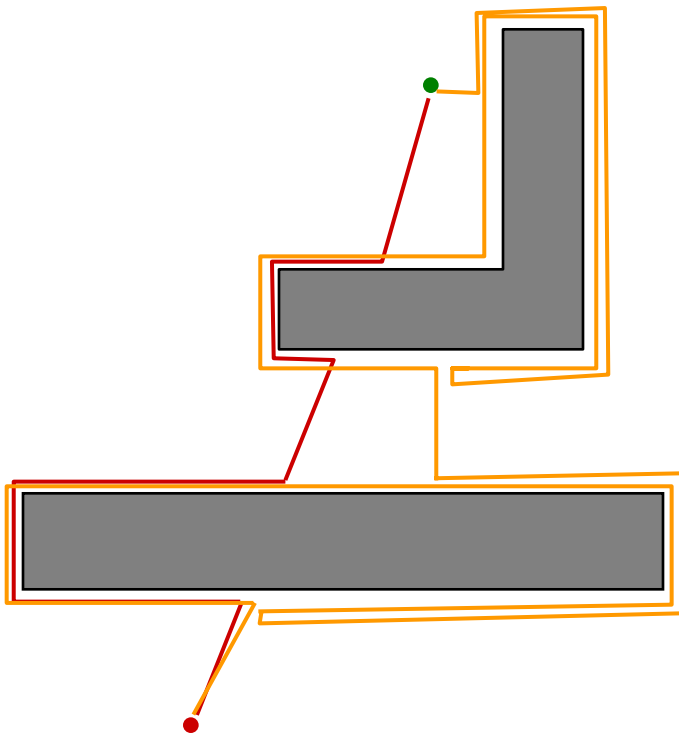


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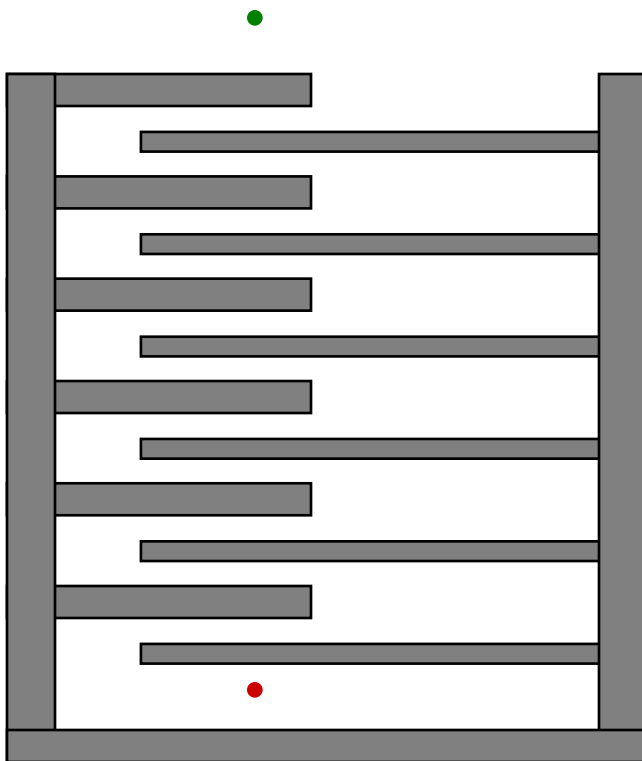
# BUG 1 vs. BUG 2

- BUG 1 is an *exhaustive search algorithm*
  - it looks at all choices before committing
- BUG 2 is a *greedy* algorithm
  - it takes the first thing that looks better
- In many cases, BUG 2 will outperform BUG 1, but
- BUG 1 has a more predictable performance overall

# “Quiz”

# Bug 2 analysis

## Bug 2: Path Bounds



What are upper/lower bounds on the path length that the robot takes?

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Lower bound:

What's the shortest distance it might travel?

**D**

Upper bound:

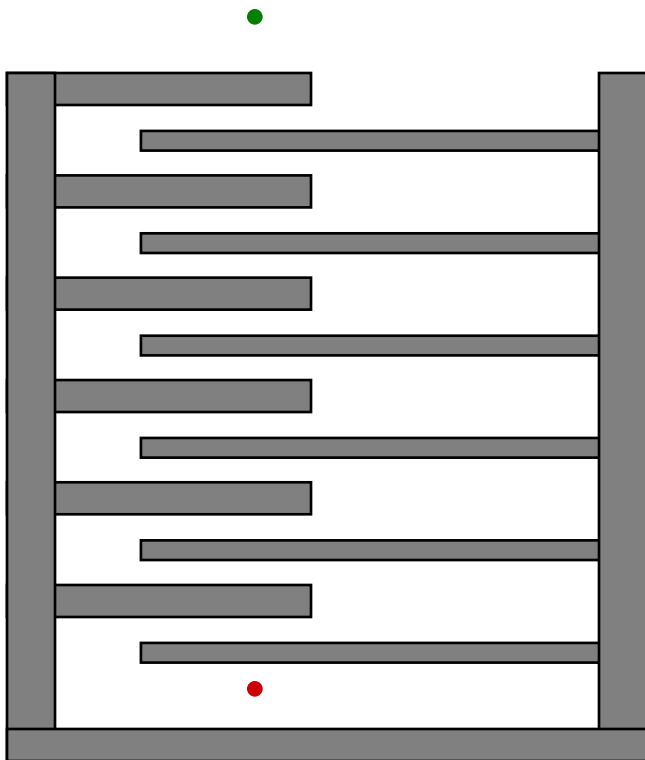
What's the longest distance it might travel?

What is an environment where your upper bound is required?

# “Quiz”

# Bug 2 analysis

## Bug 2: Path Bounds



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$D$

Upper bound:

What's the longest distance it might travel?

$$D + \sum_i \frac{n_i}{2} P_i$$

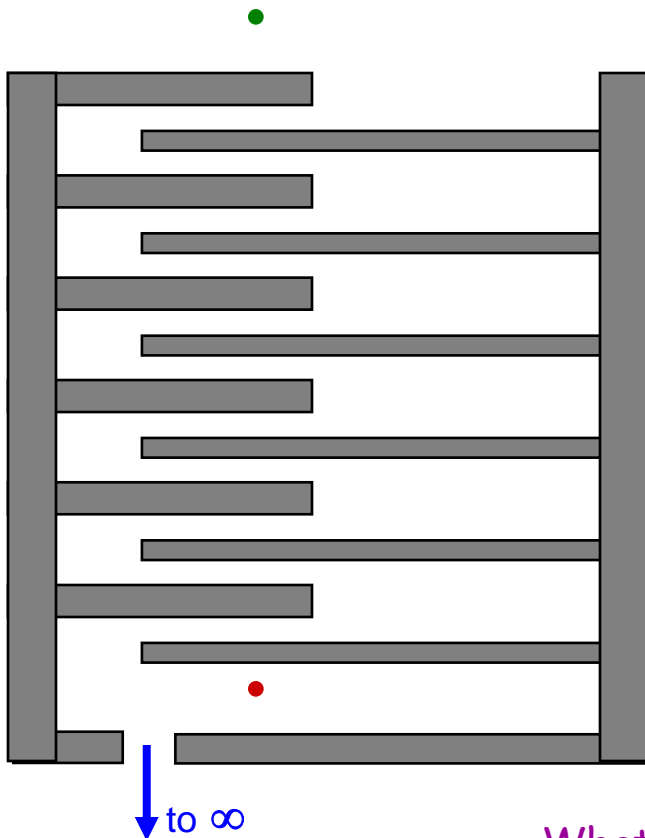
$n_i$  = # of s-line intersections of the  $i$ th obstacle

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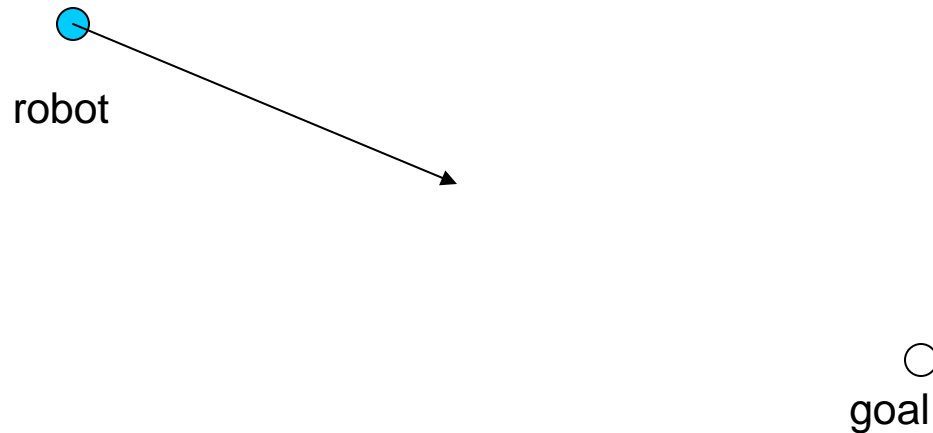
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# A More Realistic Bug

- As presented: global beacons plus contact-based wall following
- The reality: we typically use some sort of range sensing device that lets us look ahead (but has finite resolution and is noisy).
- Let us assume we have a range sensor
- distance fn:  $\rho(x, \theta) = \min_{\lambda \geq 0} d(x, x + \lambda[c_\theta, s_\theta])$   
s.t.  $x + \lambda[c_\theta, s_\theta] \in \cup_i WO_i$
- Note we write  $\rho: \mathcal{R}(2) \times S(1) \rightarrow \mathcal{R}$ 
  - what is  $S(1)$  ?
- Saturated distance:  $\rho_R(x, \theta) = \rho(x, \theta)$  if  $\rho(x, \theta) < R$ , else  $\infty$

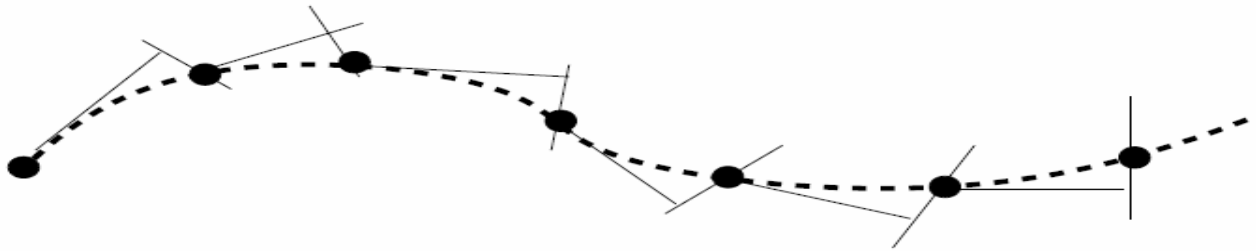
# Move to Goal

- Distance  $d(a,b) = ((a_x - b_x)^2 + (a_y - b_y)^2)^{1/2}$
- Gradient descent of  $d(a,b)$ , i.e., decrease distance to the goal





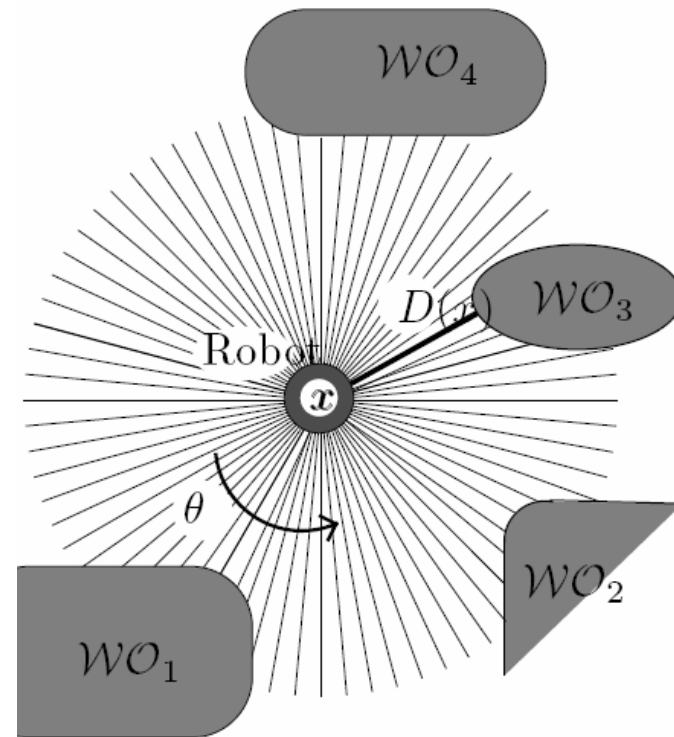
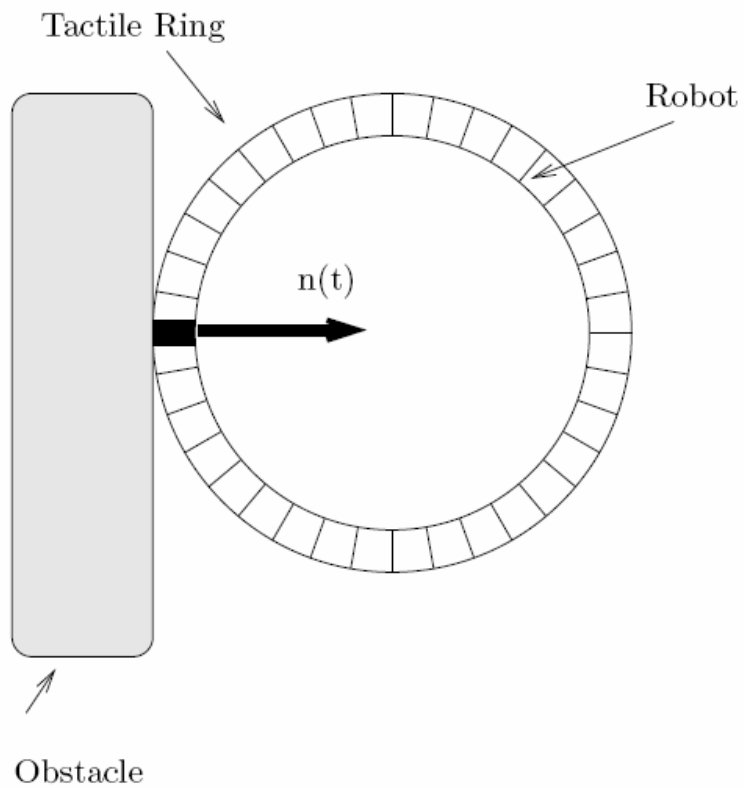
# Circumnavigating Obstacles: Curve Tracing



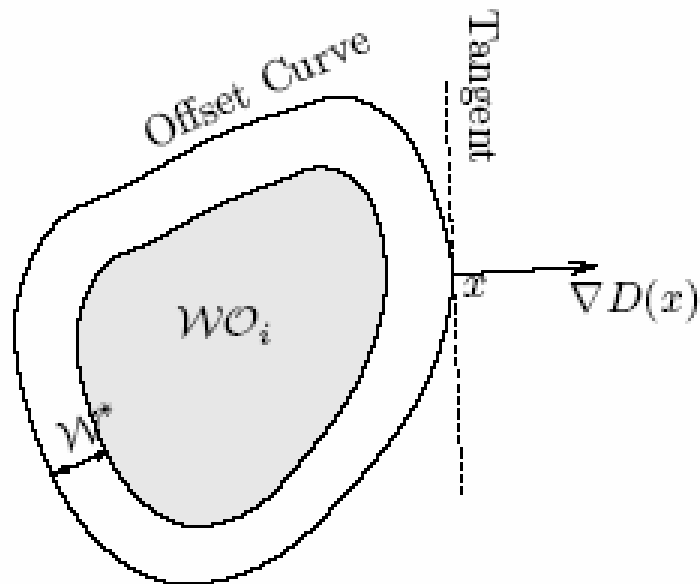
Predict: Tangent

Correct: Something else

# Normal (and hence Tangent) to Obstacle



# Circumnavigate Obstacles: Boundary Following



Safety distance  $W^*$

$$D(x) = \min d(x,c)$$

Normal is parallel to  $\nabla D(x)$

Increase/Decrease/Same

Tangent is orthogonal to both

$$\dot{c}(t) = v \quad v \text{ is in } (n(c(t)))^\perp$$

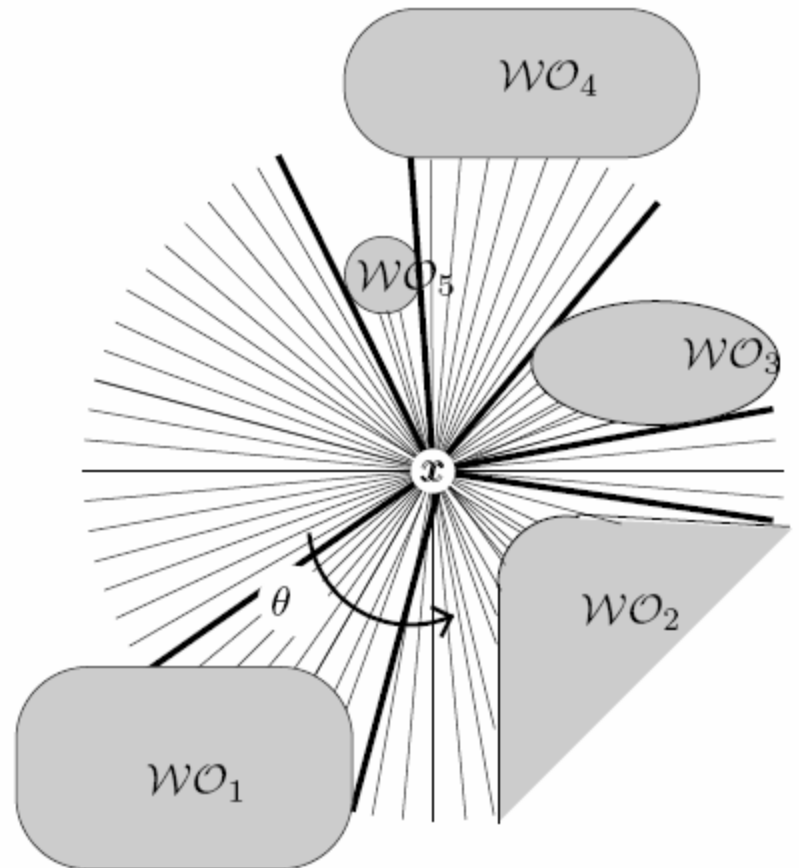
# Raw Distance Function

$$\rho(x, \theta) = \min_{\lambda \in [0, \infty]} d(x, x + \lambda[\cos \theta, \sin \theta]^T),$$

such that  $x + \lambda[\cos \theta, \sin \theta]^T \in \mathcal{C}$

Saturated raw distance function

$$\rho_R(x, \theta) = \begin{cases} \rho(x, \theta), & \text{if } \rho(x, \theta) < R \\ \infty, & \text{otherwise.} \end{cases}$$



# Implicit Function Theorem

$$G(x) = D(x) - W^*$$

Roots of  $G(x)$  trace the offset curve

$DG(x) = DD(x)$ , which is like a gradient in Euclidean spaces

Null of  $DG(x)$  is tangent, hence perp of  $DD(x)$  is too

**THEOREM D.1.1 (Implicit Function Theorem)** *Let  $f : \mathbb{R}^m \times \mathbb{R}^n \rightarrow \mathbb{R}^n$  be a smooth vector-valued function,  $f(x, y)$ . Assume that  $D_y f(x_0, y_0)$  is invertible for some  $x_0 \in \mathbb{R}^m$ ,  $y_0 \in \mathbb{R}^n$ . Then there exist neighborhoods  $X_0$  of  $x_0$  and  $Z_0$  of  $f(x_0, y_0)$  and a unique, smooth map  $g : X_0 \times Z_0 \rightarrow \mathbb{R}^n$  such that*

$$f(x, g(x, z)) = z$$

*for all  $x \in X_0$ ,  $z \in Z_0$ .*

# Correction

**THEOREM D.2.1 (Newton-Raphson Convergence Theorem)** *Let  $f : \mathbb{R}^n \rightarrow \mathbb{R}^n$  and  $f(y^*) = 0$ . For some  $\rho > 0$ , let  $f$  satisfy*

- *$Df(y^*)$  is nonsingular with bounded inverse, i.e.,  $\|(Df(y^*))^{-1}\| \leq \beta$*
- *$\|Df(x) - Df(y)\| \leq \gamma \|x - y\|$  for all  $x, y \in B_\rho(y^*)$ , where  $\gamma \leq \frac{2}{\rho\beta}$*

*Now consider the sequence  $\{y^h\}$  defined by*

$$y^{h+1} = y^h - (Df(y^h))^{-1} f(y^h),$$

*for any  $y^0 \in B_\rho(y^*)$ . Then  $y^h \in B_\rho(y^*)$  for all  $h > 0$ , and the sequence  $\{y^h\}$  quadratically converges onto  $y^*$ , i.e.,*

$$\|y^{h+1} - y^*\| \leq a \|y^h - y^*\|^2$$

$$\text{where } a = \frac{\beta\gamma}{2(1-\rho\beta\gamma)} < \frac{1}{\rho}.$$