

15-122: Principles of Imperative Computation

Recitation 5

Josh Zimmerman

Basic linear search: recap

(Note: I assume the `is_in` and `is_sorted` functions exist as defined in class.)

```
1 int lin_search(int x, int[] A, int n)
2 //@requires 0 <= n && n <= \length(A);
3 //@requires is_sorted(A, 0, n);
4 /*@ensures (-1 == \result && !is_in(x, A, 0, n))
5     || ((0 <= \result && \result < n) && A[\result] == x); @*/
6 {
7     for (int i = 0; i < n; i++)
8         //@loop_invariant 0 <= i && i <= n;
9         //@loop_invariant !is_in(x, A, 0, i);
10    {
11        if (A[i] == x) return i; // We found what we were looking for!
12        else if (x < A[i]) return -1; // Can't possibly be to the right
13        //@assert A[i] < x;
14    }
15    return -1;
16 }
```

Now, let's look at this code and see if we can prove that it works. Work on your own or with other people to follow the four-step process to proving that linear search works. (Remember: Show that the loop invariants hold initially, that they are preserved, that the loop invariants and the negation of the loop condition imply the postcondition, and that the loop terminates.)

I claim we can search a sorted array faster than this. We'll discuss why in lecture tomorrow, but for now try to think about how you could improve on this search method.

Linear search for integer square root

Recall linear search from lecture. We can apply the same concept to find the *integer* (since `c0` doesn't have floats!) square root of a given number. The integer square root of n is defined to be the greatest non-negative integer, m , such that $m^2 \leq n$.

```
1 int isqrt (int n)
2 //@requires n >= 0;
3 //@ensures \result * \result <= n;
4 //@ensures n < (\result+1) * (\result+1) || (\result+1) * (\result+1) < 0;
5 {
6     int i = 0;
7     int k = 0;
8     while (0 <= k && k <= n)
9         //@loop_invariant i * i == k;
10        //@loop_invariant i == 0 || (i > 0 && (i-1)*(i-1) <= n);
11    {
12        // Note: (i + 1)*(i + 1) == i * i + 2*i + 1 and k == i * i
13        k = k + 2*i + 1;
14        i = i + 1;
15    }
16    // This subtraction is necessary since we know k > n now
17    // and i * i == k. i is barely too large to be the square root of n
18    return i - 1;
19 }
```

Note that this function is very similar to the linear search function we discussed. It's essentially equivalent to searching through a sorted array containing all non-negative ints less than n , looking for the square root of n . There is a similar improved algorithm that we'll discuss on Friday.

Contract Checking

A water main break in GHC has, confusingly, broken the C0 compiler's `-d` option! C0 contracts are now being treated as comments, and the only way to generate assertion failures is with the `assert()` statements.

Insert `assert()` statements into the code below so that, when the code runs, all operations (C0 statements, conditional checks, and assertions) are performed at runtime in the *exact same sequence* that would have occurred if we compiled with `-d`. Not all of the blanks need to be filled in.

```
1 int mult(int x, int y)
2 //@requires x >= 0 && y >= 0;
3 //@ensures \result == x*y;
4 {
5   /* 1 */                               /* 1 */_____
6   int k = x; int n = y;
7   int res = 0;
8
9   /* 2 */                               /* 2 */_____
10  while (n != 0)
11  //@loop_invariant x * y == k * n + res;
12  {
13    /* 3 */                               /* 3 */_____
14    if ((k & 1) == 1) res = res + n;
15    k = k >> 1;
16    n = n << 1;
17    /* 4 */                               /* 4 */_____
18  }
19  /* 5 */                               /* 5 */_____
20
21  /* 6 */                               /* 6 */_____
22  return res;
23  /* 7 */                               /* 7 */_____
24 }
25
26 int main() {
27   int a;
28
29   /* 8 */                               /* 8 */_____
30   a = mult(3,4);
31
32   /* 9 */                               /* 9 */_____
33   return a;
34 }
```