

15-744: Computer Networking

L-8 Routers



Forwarding and Routers



- Forwarding
- IP lookup
- High-speed router architecture
- Readings
 - [McK97] A Fast Switched Backplane for a Gigabit Switched Router
 - [KCY03] Scaling Internet Routers Using Optics
 - Know RIP/OSPF
- Optional
 - [D+97] Small Forwarding Tables for Fast Routing Lookups
 - [BV01] Scalable Packet Classification

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Outline



- **IP router design**
- IP route lookup
- Variable prefix match algorithms
- Alternative methods for packet forwarding

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What Does a Router Look Like?



- Currently:
 - Network controller
 - Line cards
 - Switched backplane
- In the past?
 - Workstation
 - Multiprocessor workstation
 - Line cards + shared bus

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Line Cards



- Network interface cards
- Provides parallel processing of packets
- Fast path per-packet processing
 - Forwarding lookup (hardware/ASIC vs. software)

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Network Processor



- Runs routing protocol and downloads forwarding table to line cards
 - Some line cards maintain two forwarding tables to allow easy switchover
- Performs “slow” path processing
 - Handles ICMP error messages
 - Handles IP option processing

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Switch Design Issues



- Have N inputs and M outputs
 - Multiple packets for same output – output contention
 - Switch contention – switch cannot support arbitrary set of transfers
 - Crossbar
 - Bus
 - High clock/transfer rate needed for bus
 - Banyan net
 - Complex scheduling needed to avoid switch contention
- Solution – buffer packets where needed

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Switch Buffering



- Input buffering
 - Which inputs are processed each slot – schedule?
 - Head of line packets destined for busy output blocks other packets
- Output buffering
 - Output may receive multiple packets per slot
 - Need speedup proportional to # inputs
- Internal buffering
 - Head of line blocking
 - Amount of buffering needed

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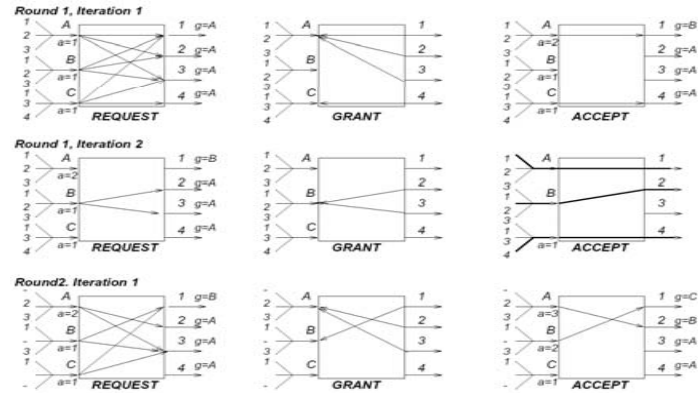
Line Card Interconnect



- Virtual output buffering
 - Maintain per output buffer at input
 - Solves head of line blocking problem
 - Each of MxN input buffer places bid for output
- Crossbar connect
- Challenge: map of bids to schedule for crossbar

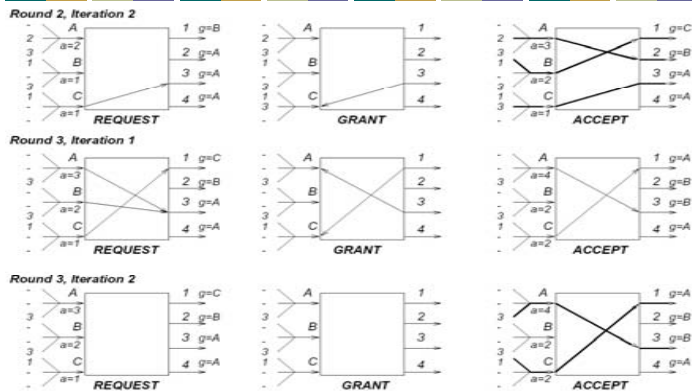
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ISLIP



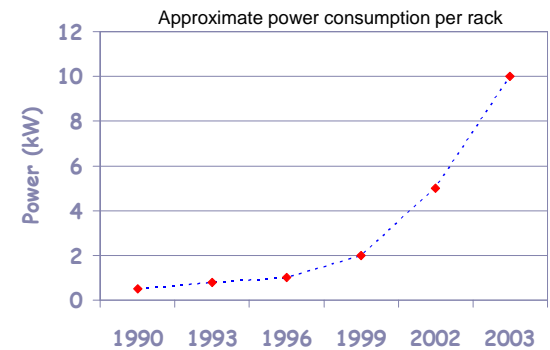
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ISLIP (cont.)



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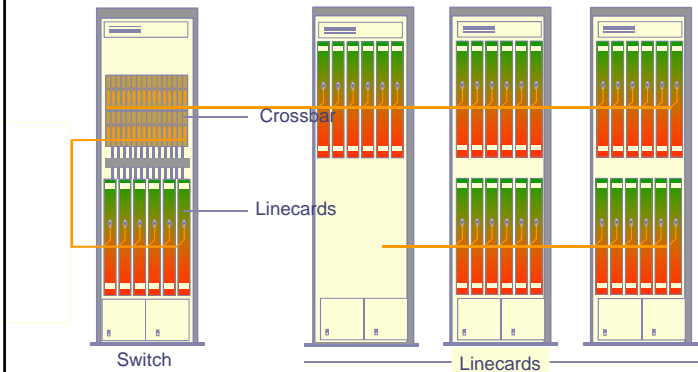
What Limits Router Capacity?



Power density is the limiting factor today

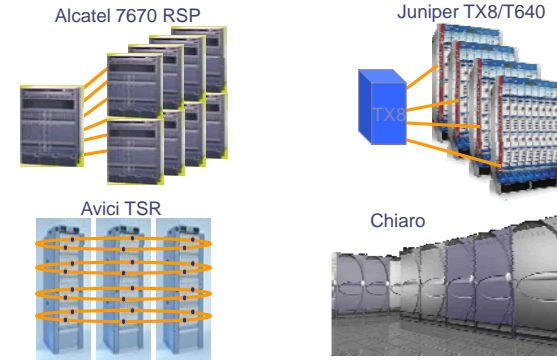
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Multi-rack Routers Reduce Power Density



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Examples of Multi-rack Routers



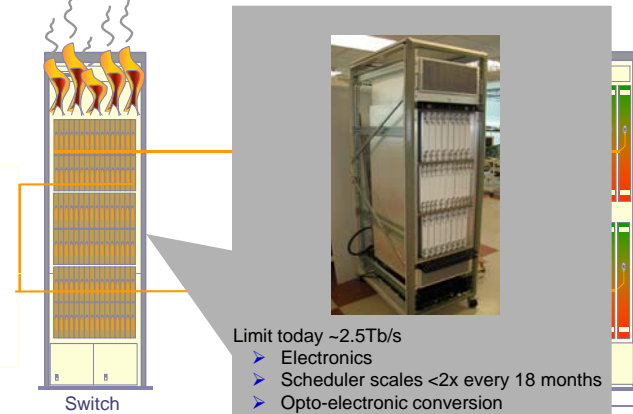
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Limits to Scaling

- Overall power is dominated by linecards
 - Sheer number
 - Optical WAN components
 - Per packet processing and buffering.
- But power *density* is dominated by switch fabric

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Multi-rack Routers Reduce Power Density



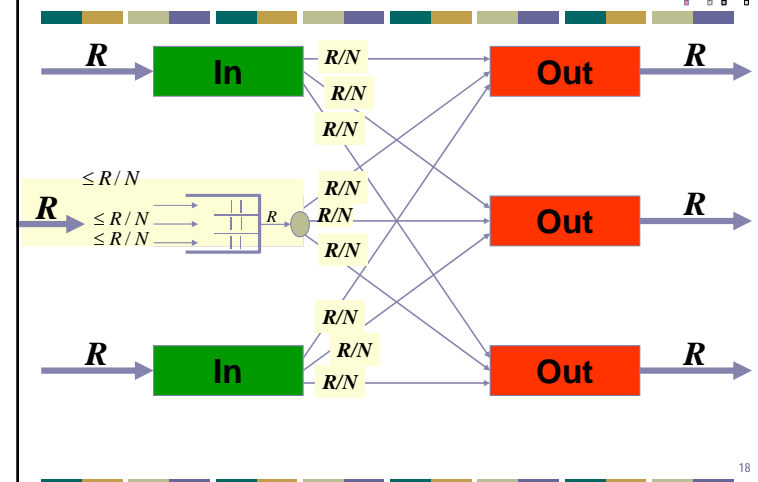
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Question

- Instead, can we use an **optical** fabric at 100Tb/s with 100% throughput?
- Conventional answer: **No**
 - Need to reconfigure switch too often
 - 100% throughput requires complex electronic scheduler.

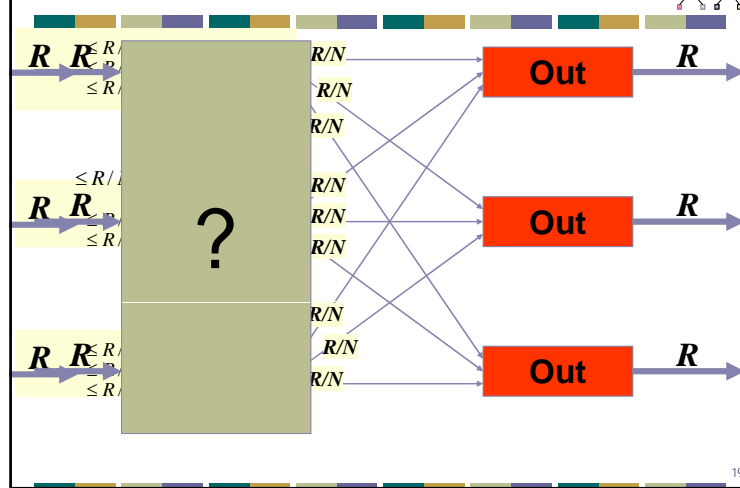
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If Traffic is Uniform...



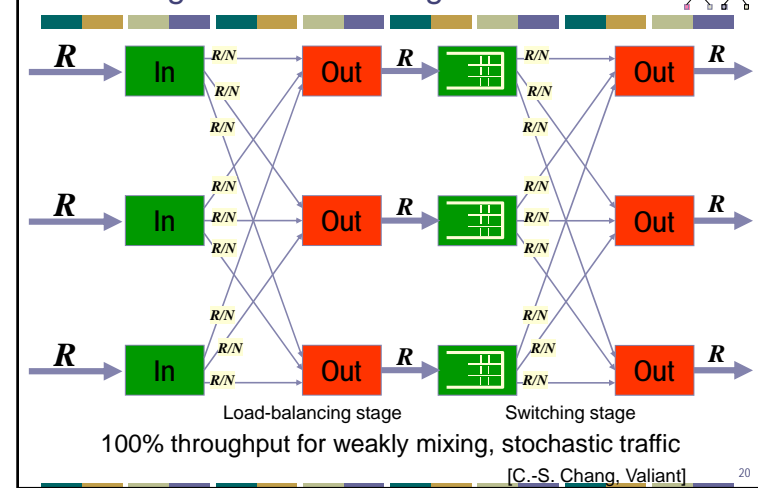
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Real Traffic is Not Uniform

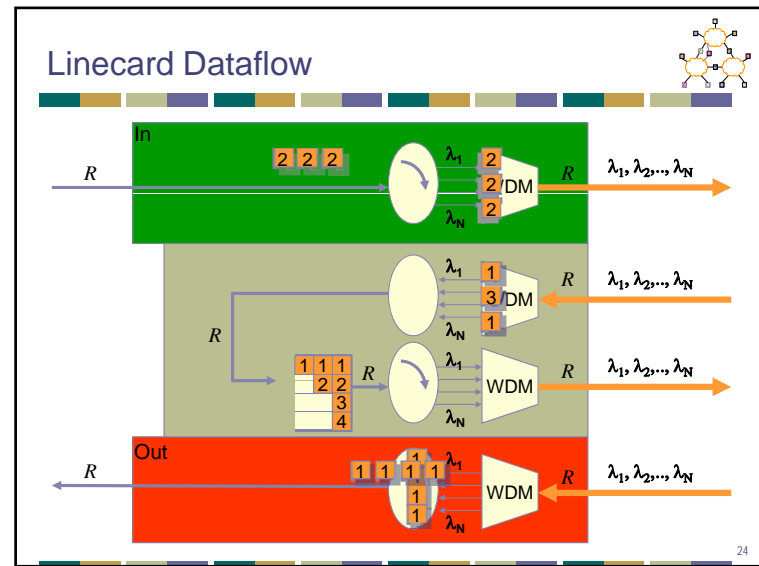
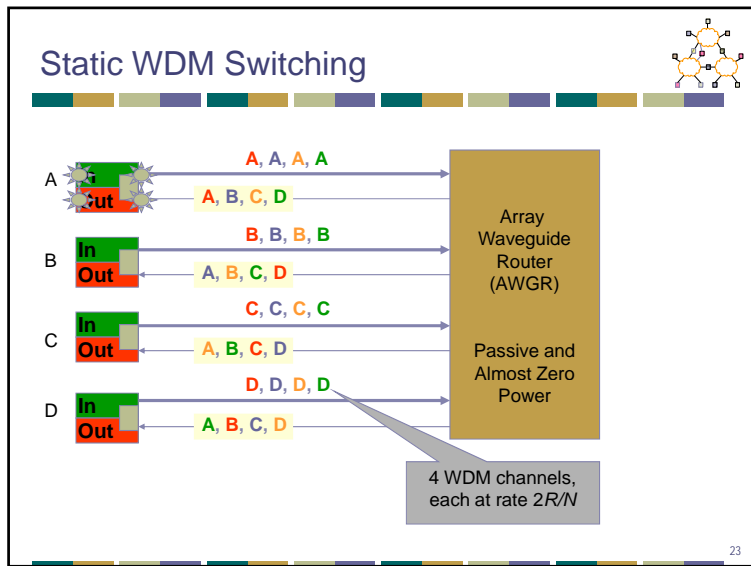
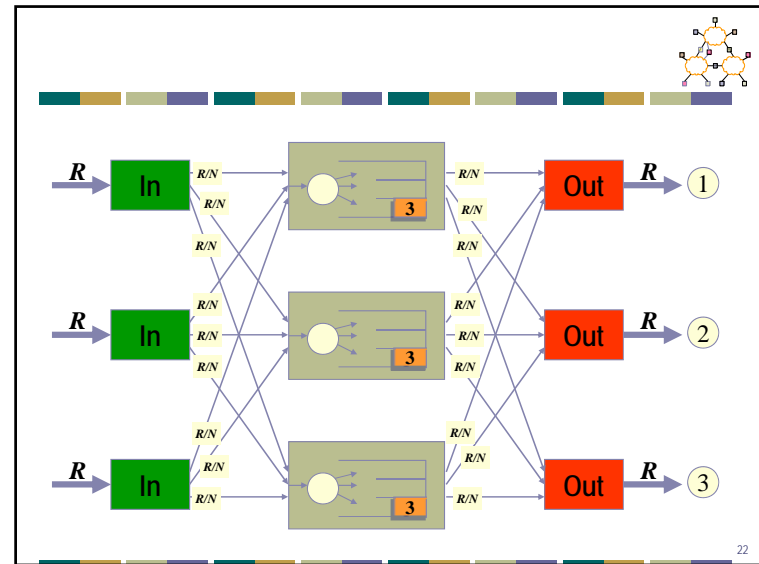
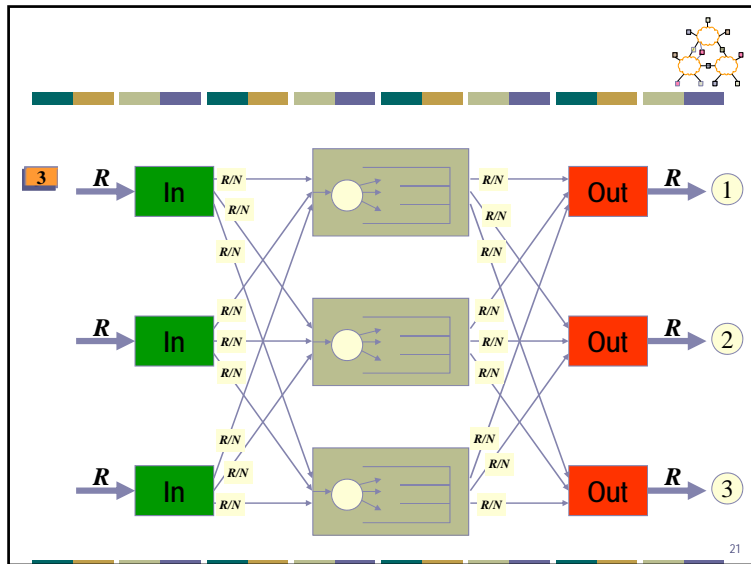


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Two-stage Load-Balancing Switch



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Original IP Route Lookup



- Address classes
 - A: 0 | 7 bit network | 24 bit host (16M each)
 - B: 10 | 14 bit network | 16 bit host (64K)
 - C: 110 | 21 bit network | 8 bit host (255)
- Address would specify prefix for forwarding table
 - Simple lookup

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Original IP Route Lookup – Example



- www.cmu.edu address 128.2.11.43
 - Class B address – class + network is 128.2
 - Lookup 128.2 in forwarding table
 - Prefix – part of address that really matters for routing
- Forwarding table contains
 - List of class+network entries
 - A few fixed prefix lengths (8/16/24)
- Large tables
 - 2 Million class C networks
- 32 bits does not give enough space encode network location information inside address – i.e., create a structured hierarchy

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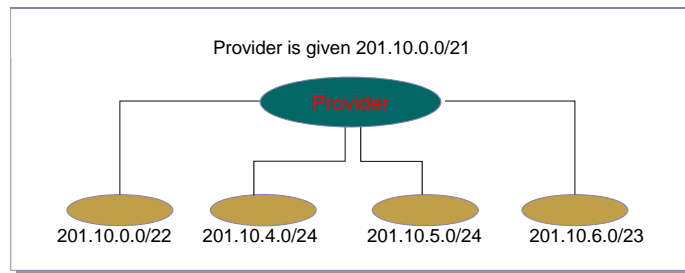
CIDR Revisited



- Supernet
 - Assign adjacent net addresses to same org
 - Classless routing (CIDR)
- How does this help routing table?
 - Combine routing table entries whenever all nodes with same prefix share same hop
 - Routing protocols carry prefix with destination network address
 - Longest prefix match for forwarding

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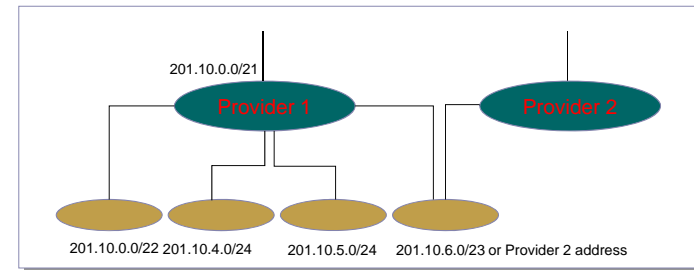
CIDR Illustration



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CIDR Shortcomings

- Multi-homing
- Customer selecting a new provider



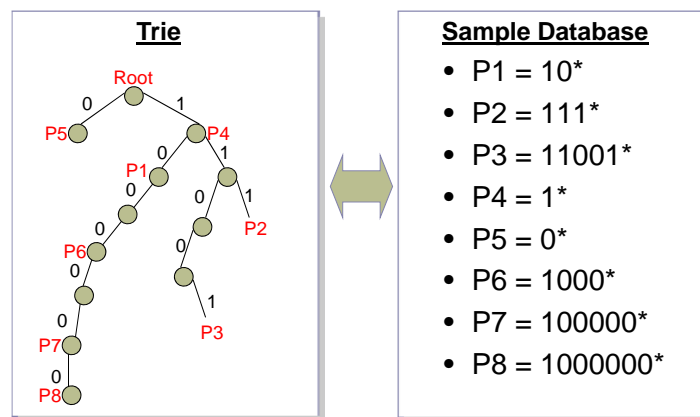
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Trie Using Sample Database



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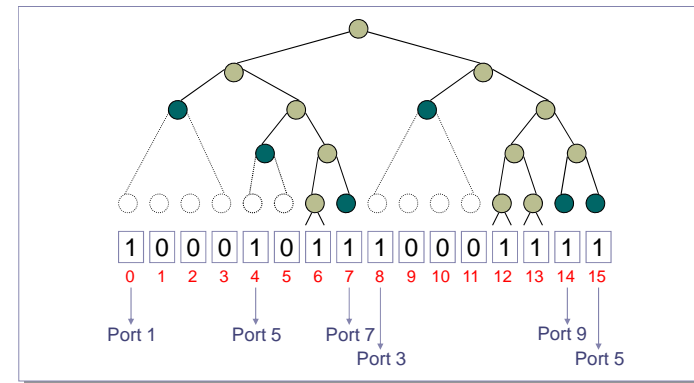
Speeding up Prefix Match (P+98)



- Cut prefix tree at 16 bit depth
 - 64K bit mask
 - Bit = 1 if tree continues below cut (root head)
 - Bit = 1 if leaf at depth 16 or less (genuine head)
 - Bit = 0 if part of range covered by leaf

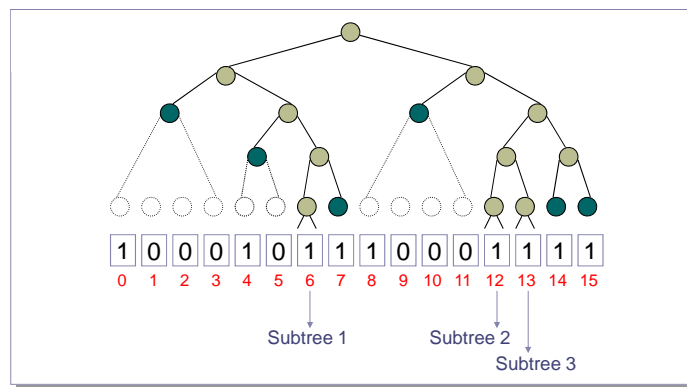
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Prefix Tree



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Prefix Tree



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Speeding up Prefix Match (P+98)



- Each 1 corresponds to either a route or a subtree
 - Keep array of routes/pointers to subtree
 - Need index into array – how to count # of 1s
 - Keep running count to 16bit word in base index + code word (6 bits)
 - Need to count 1s in last 16bit word
 - Clever tricks
- Subtrees are handled separately

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Speeding up Prefix Match (P+98)



- Scaling issues
 - How would it handle IPv6
- Update issues
- Other possibilities
 - Why were the cuts done at 16/24/32 bits?
 - Improve data structure by shuffling bits

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Speeding up Prefix Match - Alternatives



- Route caches
 - Temporal locality
 - Many packets to same destination
- Other algorithms
 - Waldvogel – Sigcomm 97
 - Binary search on prefixes
 - Works well for larger addresses
 - Bremner-Barr – Sigcomm 99
 - Clue = prefix length matched at previous hop
 - Why is this useful?
 - Lampson – Infocom 98
 - Binary search on ranges

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Speeding up Prefix Match - Alternatives



- Content addressable memory (CAM)
 - Hardware based route lookup
 - Input = tag, output = value associated with tag
 - Requires exact match with tag
 - Multiple cycles (1 per prefix searched) with single CAM
 - Multiple CAMs (1 per prefix) searched in parallel
- Ternary CAM
 - 0,1,don't care values in tag match
 - Priority (I.e. longest prefix) by order of entries in CAM

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Outline



- IP router design
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- Variable prefix match algorithms
- **Alternative methods for packet forwarding**

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Techniques for Forwarding Packets



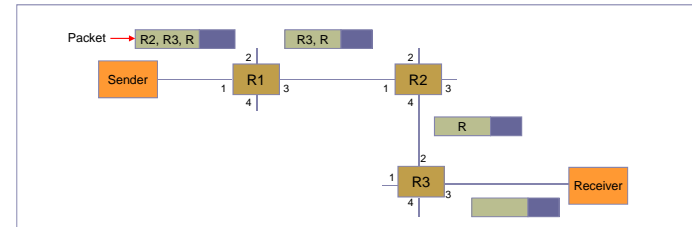
- Source routing
 - Packet carries path
- Table of virtual circuits
 - Connection routed through network to setup state
 - Packets forwarded using connection state
- Table of global addresses (IP)
 - Routers keep next hop for destination
 - Packets carry destination address

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Source Routing



- List entire path in packet
 - Driving directions (north 3 hops, east, etc..)
- Router processing
 - Examine first step in directions
 - Strip first step from packet
 - Forward to step just stripped off



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Source Routing



- Advantages
 - Switches can be very simple and fast
- Disadvantages
 - Variable (unbounded) header size
 - Sources must know or discover topology (e.g., failures)
- Typical use
 - Ad-hoc networks (DSR)
 - Machine room networks (Myrinet)

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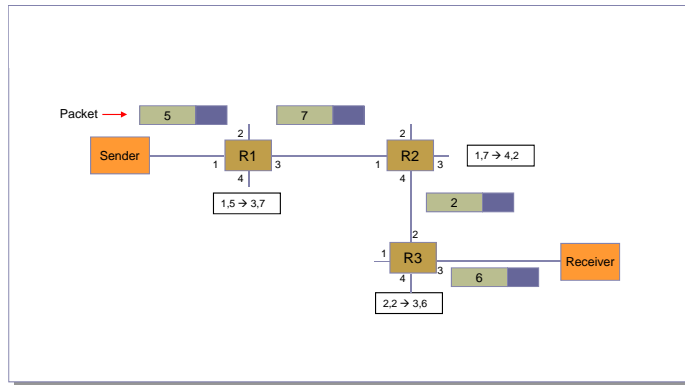
Virtual Circuits/Tag Switching



- Connection setup phase
 - Use other means to route setup request
 - Each router allocates flow ID on local link
 - Creates mapping of inbound flow ID/port to outbound flow ID/port
- Each packet carries connection ID
 - Sent from source with 1st hop connection ID
- Router processing
 - Lookup flow ID – simple table lookup
 - Replace flow ID with outgoing flow ID
 - Forward to output port

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Virtual Circuits Examples



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Virtual Circuits



- Advantages
 - More efficient lookup (simple table lookup)
 - More flexible (different path for each flow)
 - Can reserve bandwidth at connection setup
 - Easier for hardware implementations
- Disadvantages
 - Still need to route connection setup request
 - More complex failure recovery – must recreate connection state
- Typical uses
 - ATM – combined with fix sized cells
 - MPLS – tag switching for IP networks

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IP Datagrams on Virtual Circuits



- Challenge – when to setup connections
 - At bootup time – permanent virtual circuits (PVC)
 - Large number of circuits
 - For every packet transmission
 - Connection setup is expensive
 - For every connection
 - What is a connection?
 - How to route connectionless traffic?

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IP Datagrams on Virtual Circuits



- Traffic pattern
 - Few long lived flows
 - Flow – set of data packets from source to destination
 - Large percentage of packet traffic
 - Improving forwarding performance by using virtual circuits for these flows
- Other traffic uses normal IP forwarding

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Summary: Addressing/Classification



- Router architecture carefully optimized for IP forwarding
- Key challenges:
 - Speed of forwarding lookup/classification
 - Power consumption
- Some good examples of common case optimization
 - Routing with a clue
 - Classification with few matching rules
 - Not checksumming packets