

UNIT 7B

Data Representation: Compression

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Fixed-Width Encoding

- In a fixed-width encoding scheme, each character is given a binary code with the same number of bits.
 - Example:
Standard ASCII is a fixed width encoding scheme, where each character is encoded with 7 bits.
This gives us $2^7 = 128$ different codes for characters.

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Fixed-Width Encoding

- Given a character set with n characters, what is the minimum number of bits needed for a fixed-width encoding of these characters?
 - Since a fixed width of k bits gives us n unique codes to use for characters, where $n = 2^k$.
 - So given n characters, the number of bits needed is given by $k = \lceil \log_2 n \rceil$. (We use the ceiling function since $\log_2 n$ may not be an integer.)
 - Example: To encode just the alphabet A-Z using a fixed-width encoding, we would need $\lceil \log_2 26 \rceil = 5$ bits:
e.g. A => 00000, B => 00001, C => 00010, ..., Z => 11001.

Using Fixed-Width Encoding

- If we have a fixed-width encoding scheme using n bits for a character set and we want to transmit or store a file with m characters, we would need mn bits to store the entire file.
- Can we do better?
 - If we assign fewer bits to more frequent characters, and more bits to less frequent characters, then the overall length of the message might be shorter.

Huffman Coding

- We can use an encoding scheme named after David A. Huffman to compress our text without losing any information.
- Based on the idea that some characters occur more frequently than others.
- Huffman codes are not fixed-width.



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Huffman Coding: the process

1. Assign character codes
 - a. Obtain character frequencies
 - b. Use frequencies to build a *Huffman tree*
 - c. Use tree to assign variable-length codes to characters (store them in a table)
2. Use table to encode (compress) ASCII source file to variable-length codes
3. Use tree to decode (decompress) to ASCII

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The Hawaiian Alphabet



- The Hawaiian alphabet consists of 13 characters.
 - ' is the okina which sometimes occurs between vowels (e.g. **KAMA'AINA**)
- The table to the right shows each character along with its relative frequency in Hawaiian words.

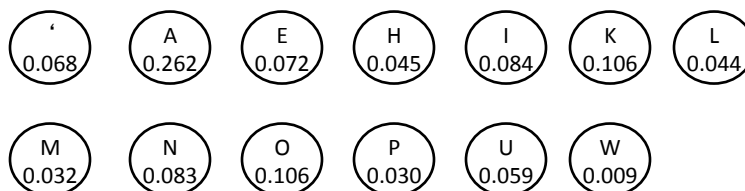
| | |
|---|-------|
| ' | 0.068 |
| A | 0.262 |
| E | 0.072 |
| H | 0.045 |
| I | 0.084 |
| K | 0.106 |
| L | 0.044 |
| M | 0.032 |
| N | 0.083 |
| O | 0.106 |
| P | 0.030 |
| U | 0.059 |
| W | 0.009 |

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The Huffman Tree

- We use a tree structure to develop the unique binary code for each letter.
- Start with each letter/frequency as its own node:

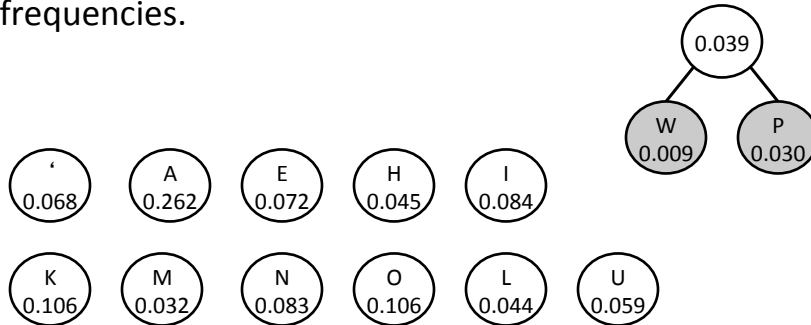


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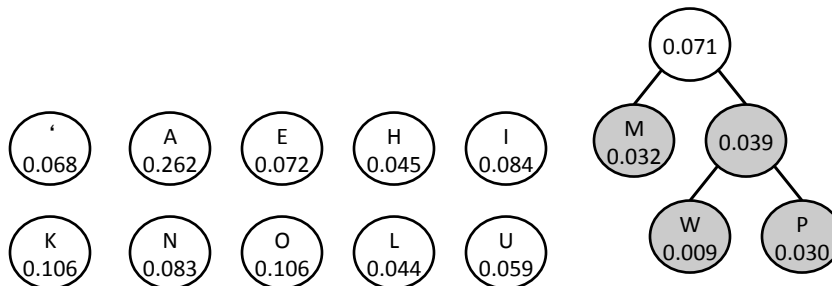
The Huffman Tree

- Combine lowest two frequency nodes into a tree with a new parent with the sum of their frequencies.



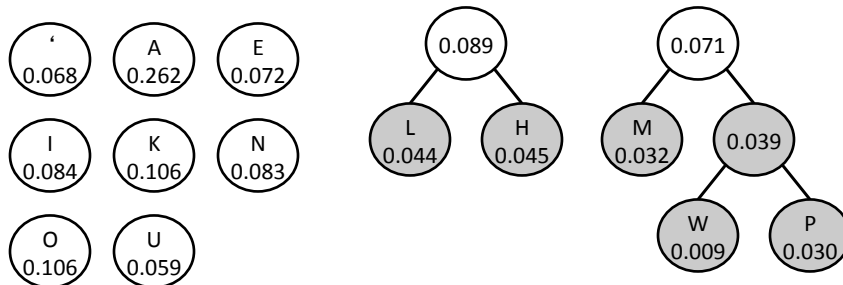
The Huffman Tree

- Combine lowest two frequency nodes (including the new node we just created) into a tree with a new parent with the sum of their frequencies.



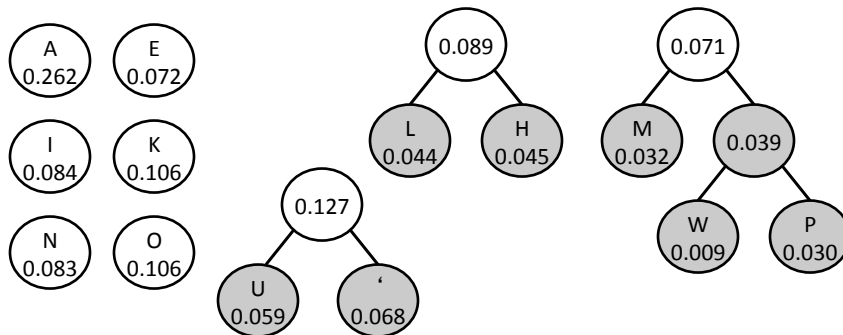
The Huffman Tree

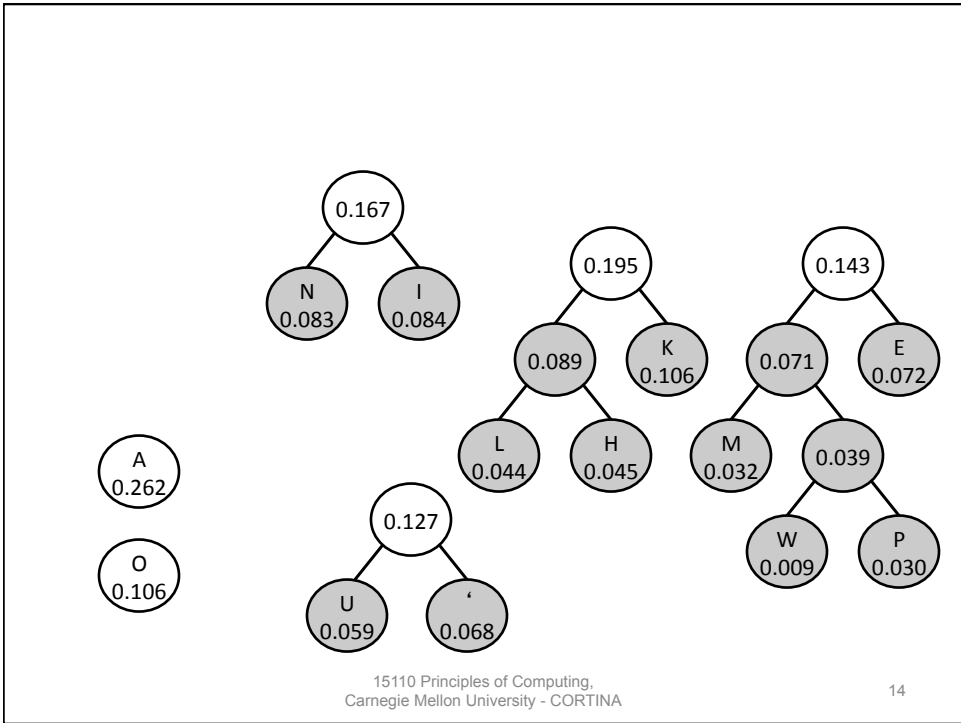
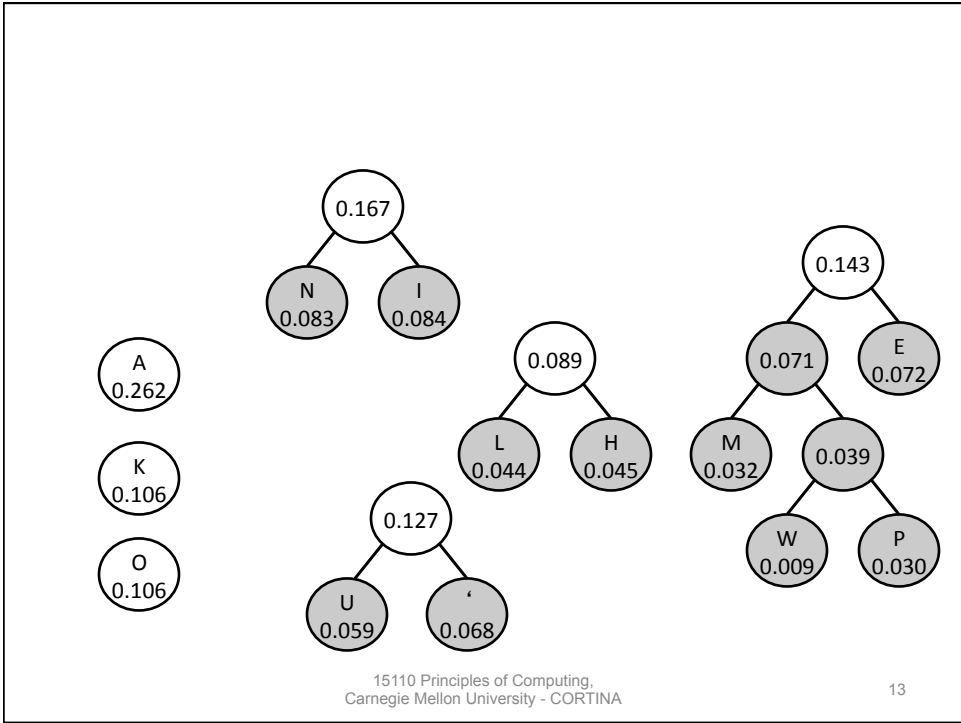
- Combine lowest two frequency nodes (including the new node we just created) into a tree with a new parent with the sum of their frequencies.

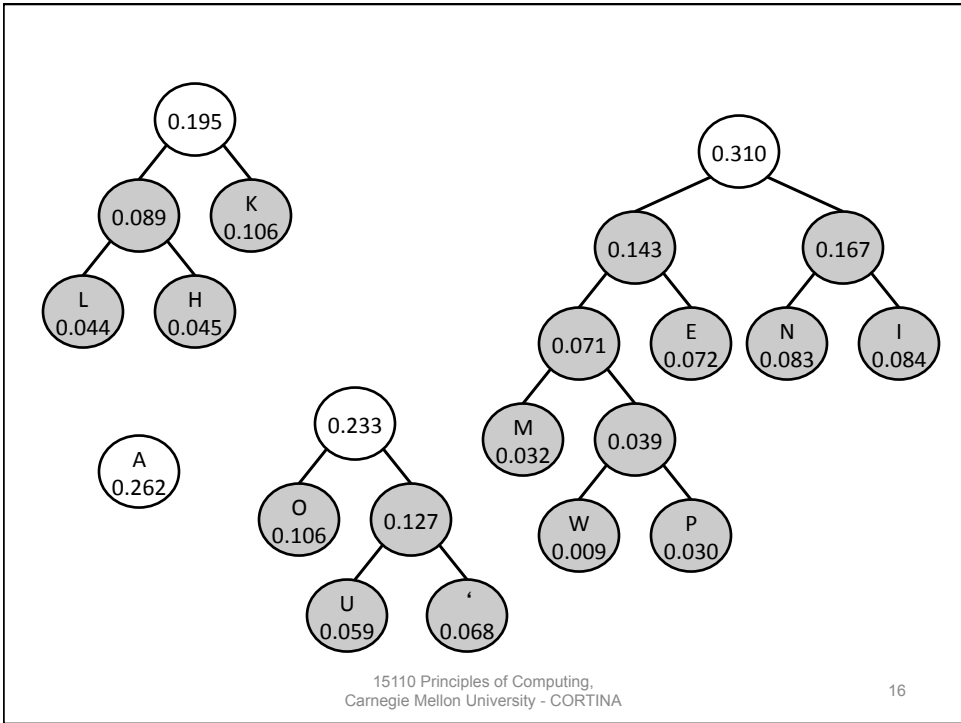
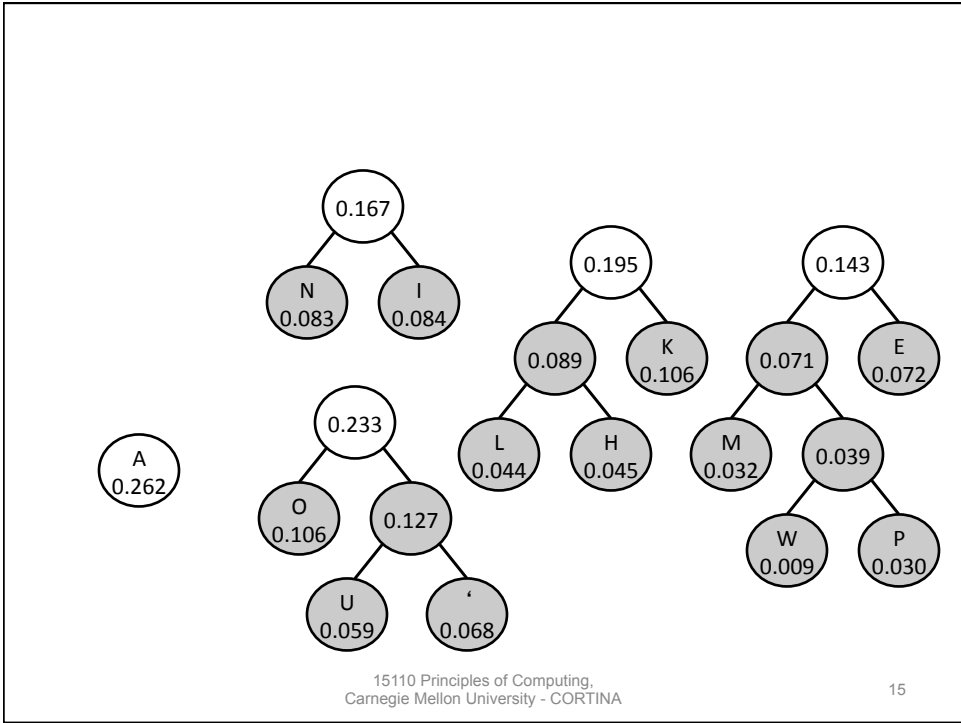


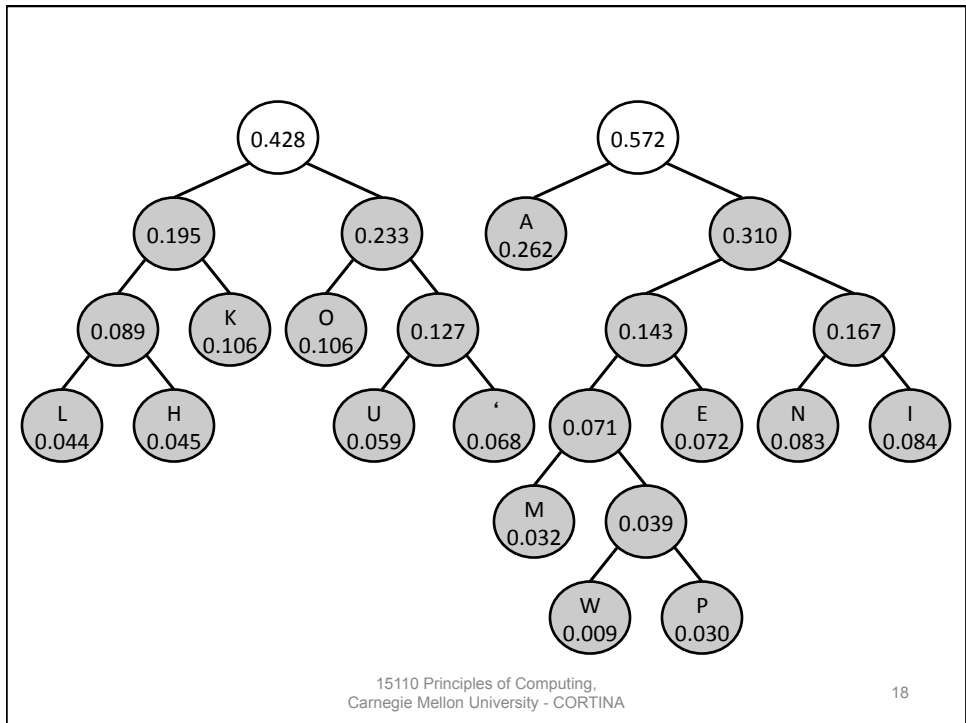
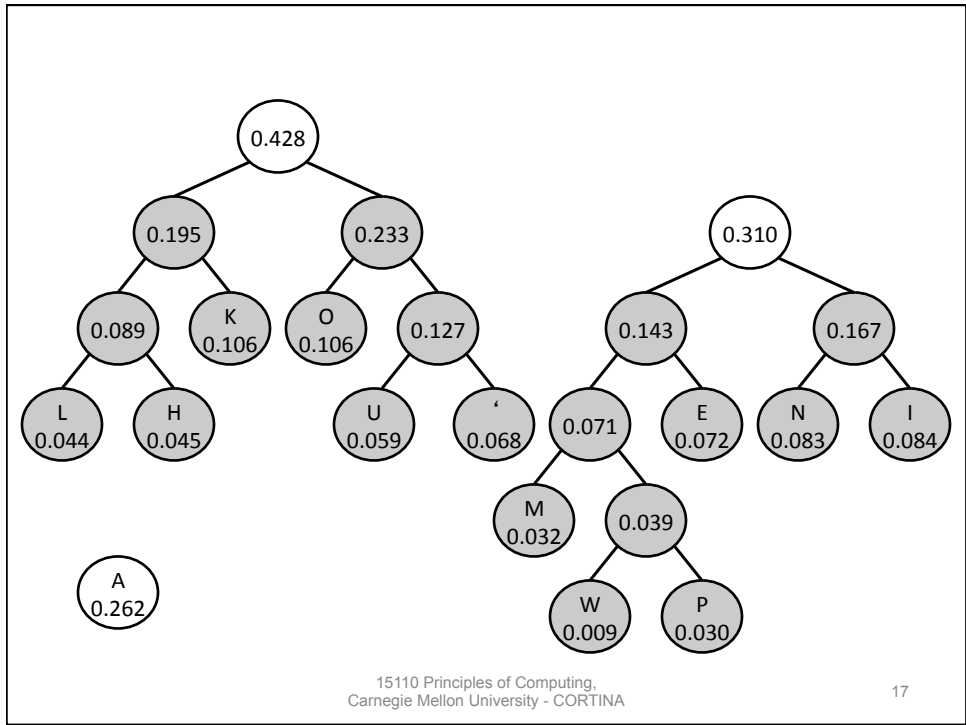
The Huffman Tree

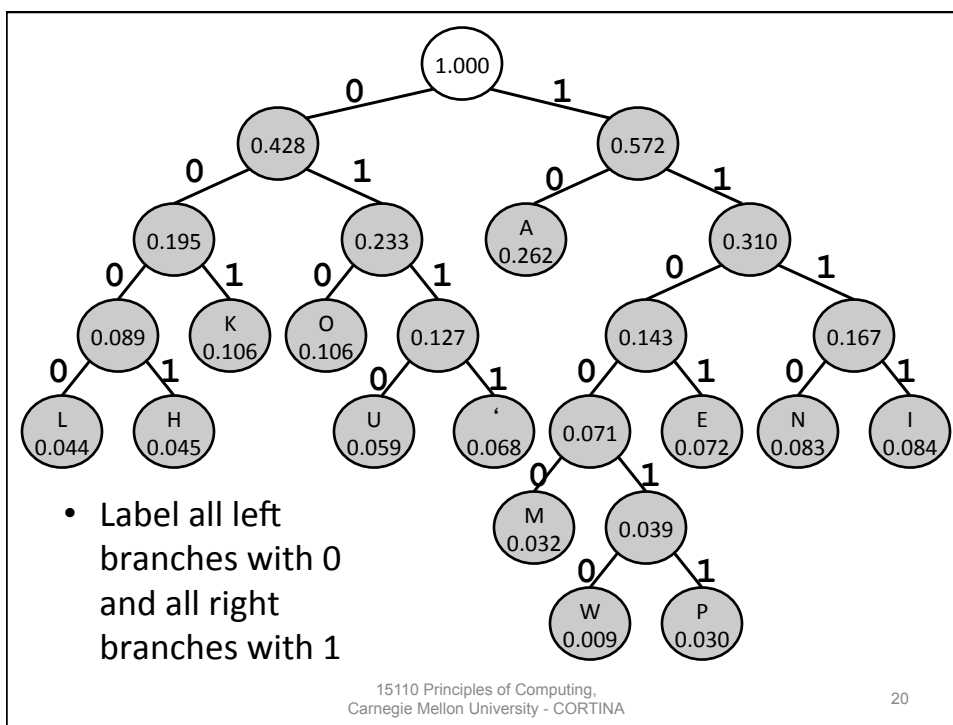
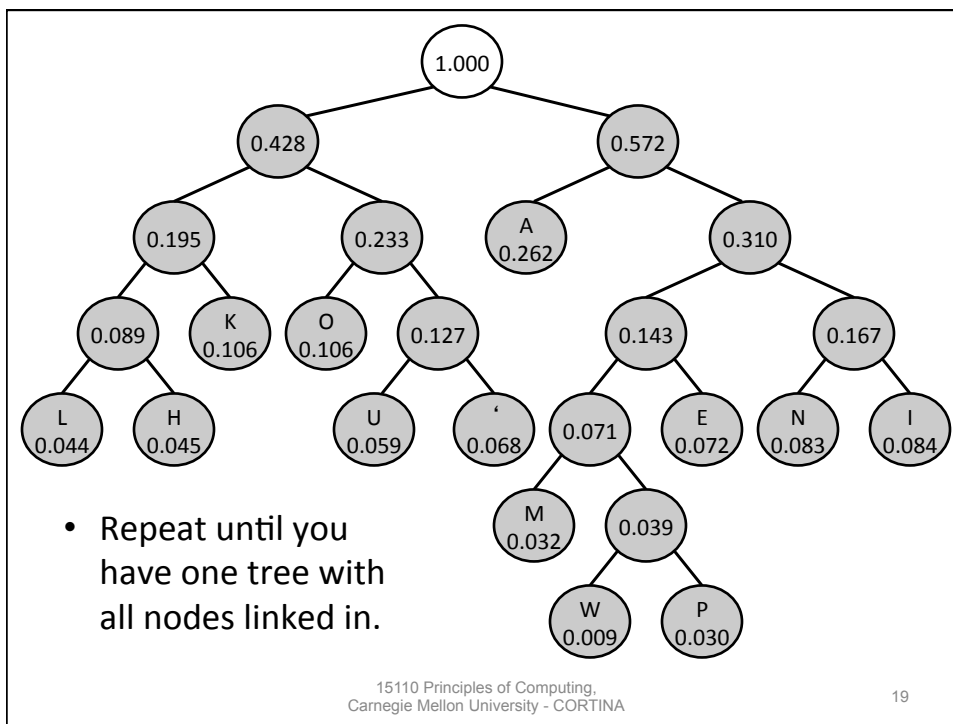
- Combine lowest two frequency nodes (including the new node we just created) into a tree with a new parent with the sum of their frequencies...

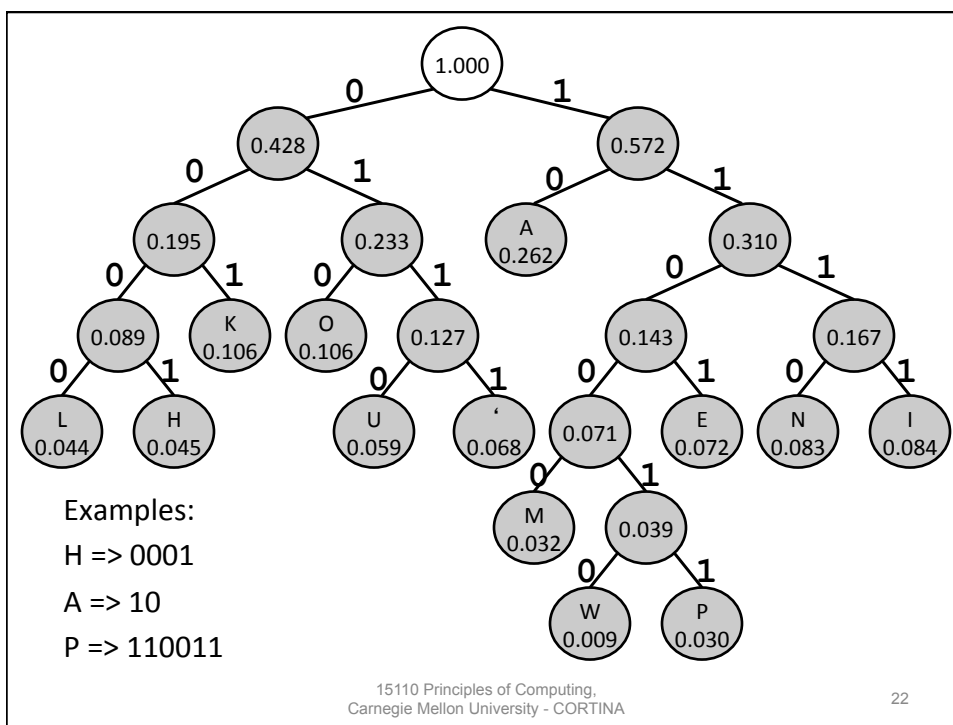
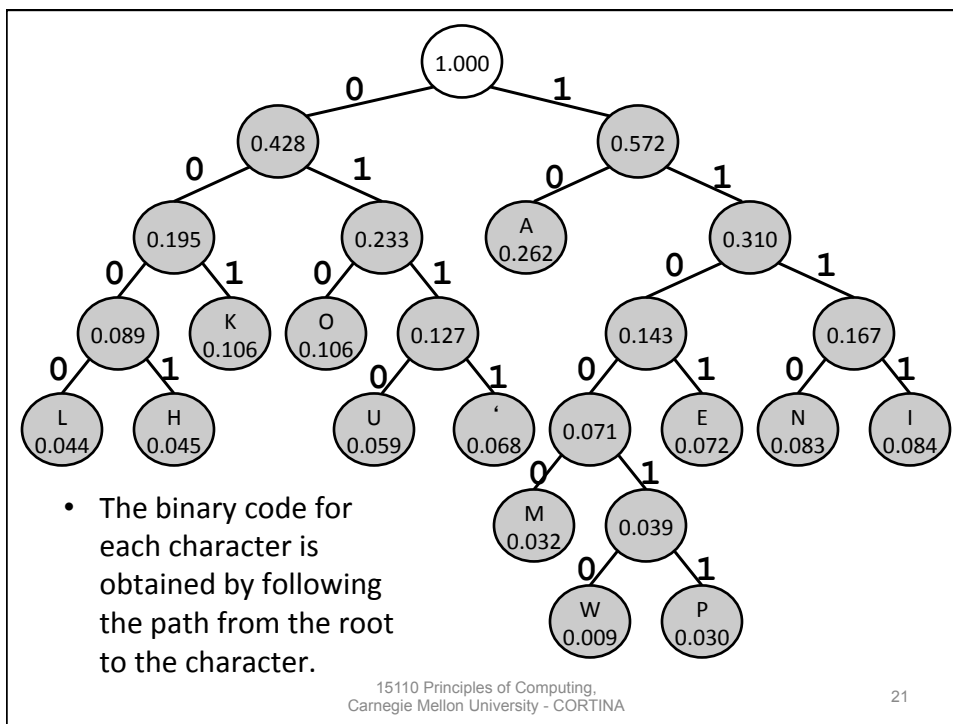












Fixed Width vs. Huffman Coding

| | | | | |
|---|------|---|--------|----------------------|
| ' | 0000 | ' | 0111 | |
| A | 0001 | A | 10 | |
| E | 0010 | E | 1101 | <u>ALOHA</u> |
| H | 0011 | H | 0001 | |
| I | 0100 | I | 1111 | Fixed Width: |
| K | 0101 | K | 001 | 00010110100100110001 |
| L | 0110 | L | 0000 | 20 bits |
| M | 0111 | M | 11000 | |
| N | 1000 | N | 1110 | Huffman Code: |
| O | 1001 | O | 010 | 100000010000110 |
| P | 1010 | P | 110011 | 15 bits |
| U | 1011 | U | 0110 | |
| W | 1100 | W | 110010 | |

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Decoding

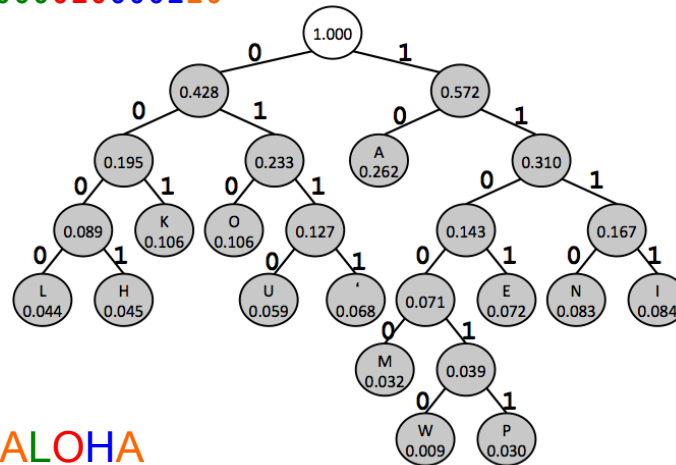
- In a fixed-width code, the boundaries between letters are fixed in advance:
0001 0110 1001 0011 0001
- With a variable-length code, the boundaries are determined by the letters themselves.
 - No letter's code can be a prefix of another letter.
 - Example: since A is "10", no other letter's code can begin with "10". All the remaining codes begin with "00", "01", or "11".

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100000010000110

Decoding



ALOHA

- To find the character use the bits to determine path from root

Programming the Huffman Tree

- Let's write Python code to produce a Huffman tree for a given alphabet.
- At each step we need to find the two nodes with the lowest frequency scores.
- This will be easy if nodes are kept in a list that is sorted by score value.
- Solution: use a **priority queue**.

Priority Queues

NOTE: For this unit, we use PythonLabs (on the linux server) and we need to include the following line in the code:

```
from PythonLabs.BitLab import PriorityQueue, Node,  
assign_codes, encode, decode
```

Priority Queue: a data structure

- A priority queue (PQ) is like a list that is automatically kept sorted.

```
>>> pq = PriorityQueue()
```

```
>>> pq
```

```
[]
```

- PQ methods: **insert** and **pop**

Priority Queue: insert

- To add an element into the priority queue in its correct position, we use the `insert` method:
- ```
>>> pq.insert("peach")
>>> pq.insert("apple")
>>> pq.insert("banana")
>>> pq
[apple, banana, peach]
```

## Priority Queue: pop

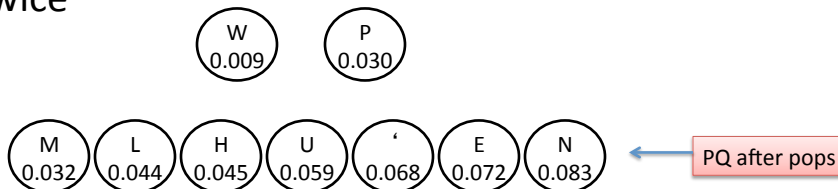
- To get the first (highest priority) element of the queue, use the `pop` method, which removes it as well:
- ```
>>> fruit1 = pq.pop()
>>> fruit1
'apple'
>>> pq
[banana, peach]
>>> fruit2 = pq.pop()
>>> fruit2
banana
>>> pq
[peach]
```

Using a PQ to build the tree

- Make a PQ of Nodes. Frequency = priority



To get the two lowest frequency nodes, pop twice



Making Tree Nodes

- Store the character and frequency data into a nested list:

```
table = [ ["'", 0.068], ["A", 0.262],
          ["E", 0.072], ["H", 0.045], ["I", 0.084],
          ["K", 0.106], ["L", 0.044], ["M", 0.032],
          ["N", 0.083], ["O", 0.106], ["P", 0.030],
          ["U", 0.059], ["W", 0.009] ]
```

- Making one of the tree nodes:

```
char = table[2][0]      # "E"
freq = table[2][1]     # 0.072
node = Node.new(char, freq)
```

`["E", 0.072]`



Building a PQ of Single Nodes

```
def make_pq(table):  
    pq = PriorityQueue()  
    for item in table:  
        char = item[0]  
        freq = item[1]  
        node = Node(char, freq)  
        pq.insert(node)  
    return pq
```

Remember: each item in the table is a 2-element list with a character and a frequency.

Building our Priority Queue

```
>>> pq = make_pq(table)  
pq
```

```
[ ( W: 0.009 ), ( P: 0.030 ),  
  ( M: 0.032 ), ( L: 0.044 ),  
  ( H: 0.045 ), ( U: 0.059 ),  
  ( ' : 0.068 ), ( E: 0.072 ),  
  ( N: 0.083 ), ( I: 0.084 ),  
  ( K: 0.106 ), ( O: 0.106 ),  
  ( A: 0.262 ) ]
```

One tree node

Priority queue showing the 13 nodes in sorted order based on frequency.

Building a Huffman Tree

```
def build_tree(pq):  
    while len(pq) > 1:  
        node1 = pq.pop()  
        node2 = pq.pop()  
        pq.insert(Node(node1, node2))  
    return pq[0]
```

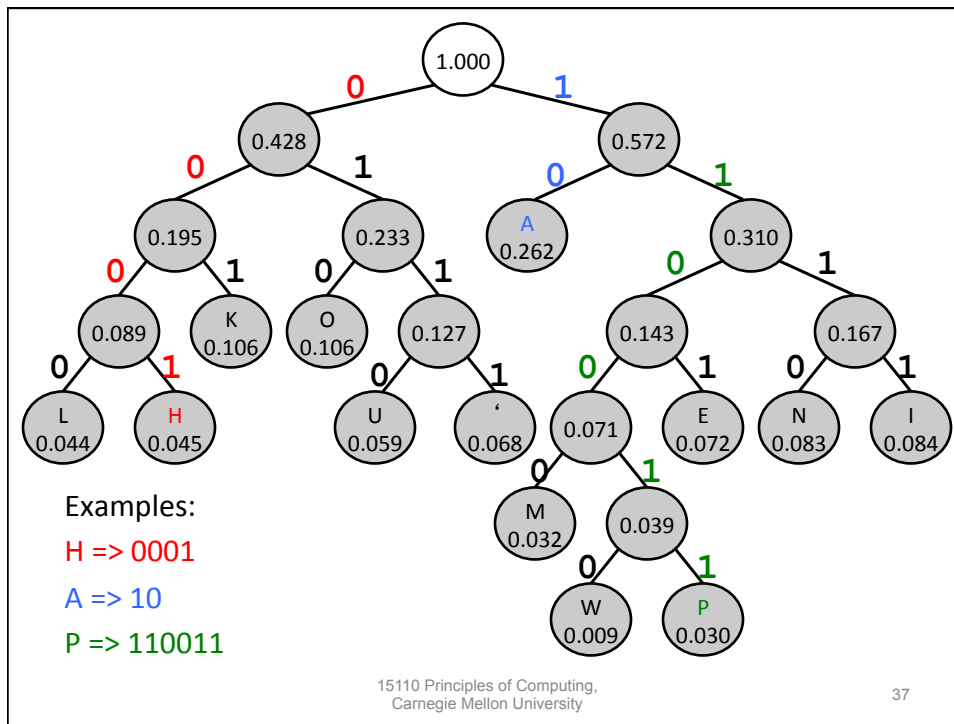
Creates a new node
with node1 as its left child
and node2 as its right child

(Unlike book version we already created the pq)

Building our Huffman Tree

```
tree = build_tree(pq)  
=> ( 1.000  
    ( 0.428  
      ( 0.195  
        ( 0.089 ( L: 0.044 ) ( H: 0.045  
          ( K: 0.106 ) )  
        ( 0.233  
          ( O: 0.106 )  
          ( 0.127 ( U: 0.059 ) ( ': 0.068 ) ) ) )  
      ( 0.572  
        ( A: 0.262 )  
        ( 0.310  
          ( 0.143  
            ( 0.071 ( M: 0.032 )  
              ( 0.039 ( W: 0.009 ) ( P: 0.030 ) ) )  
            ( E: 0.072 ) )  
          ( 0.167 ( N: 0.083 ) ( I: 0.084 ) ) ) ) ) ) )
```

This is just our Huffman
tree expressed using
recursively nested
parenthetical components:
(root (left)
 (right))



Assigning Codes, Encoding & Decoding

```
>>> ht = assign_codes(tree)
```

from BitLab
takes a Huffman tree and returns a hash table that maps each letter to its binary code

```
>>> ht["W"]  
110010
```

```
>>> ht["A"]  
10
```

Note the [] syntax.
This returns the code associated with the character from the hash table.

```
>>> msg = encode("ALOHA", tree)  
100000010000110
```

```
>>> decode(msg, tree)  
"ALOHA"
```

from BitLab
encode and decode functions