System-Level I/O

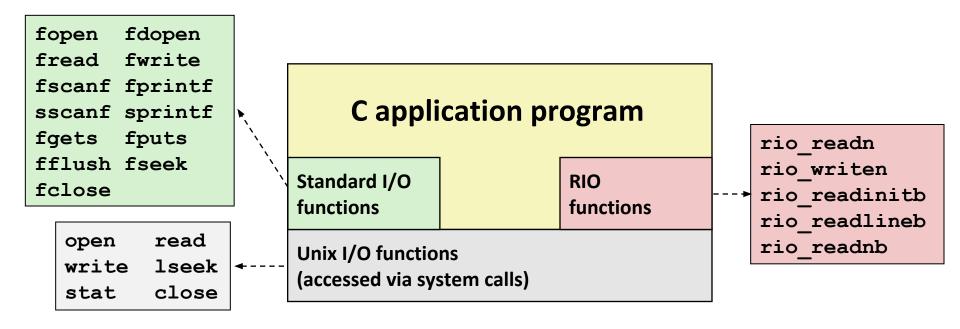
15-213: Introduction to Computer Systems 16th Lecture, June 28, 2018

Instructors:

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Today: Unix I/O and C Standard I/O

- Two sets: system-level and C level
- Robust I/O (RIO): 15-213 special wrappers good coding practice: handles error checking, signals, and "short counts"



Today

- Unix I/O
- RIO (robust I/O) package
- Standard I/O
- Which I/O when
- Metadata, sharing, and redirection

Unix I/O Overview

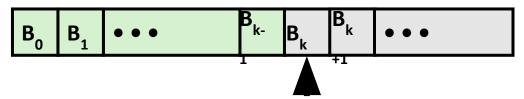
■ A Linux *file* is a sequence of *m* bytes:

$$B_0, B_1, \dots, B_k, \dots, B_{m-1}$$

- Cool fact: All I/O devices are represented as files:
 - /dev/sda2 (/usr disk partition)
 - /dev/tty2 (terminal)
- Even the kernel is represented as a file:
 - boot/vmlinuz-3.13.0-55-generic (kernel image)
 - /proc (kernel data structures)

Unix I/O Overview

- Elegant mapping of files to devices allows kernel to export simple interface called *Unix I/O*:
 - Opening and closing files
 - open() and close()
 - Reading and writing a file
 - read() and write()
 - Changing the current file position (seek)
 - indicates next offset into file to read or write
 - lseek()



Current file position = k

File Types

- Each file has a type indicating its role in the system
 - Regular file: Contains arbitrary data
 - Directory: Index for a related group of files
 - Socket: For communicating with a process on another machine

Other file types beyond our scope

- Named pipes (FIFOs)
- Symbolic links
- Character and block devices

Regular Files

- A regular file contains arbitrary data
- Applications often distinguish between text files and binary files
 - Text files are regular files with only ASCII or Unicode characters
 - Binary files are everything else
 - e.g., object files, JPEG images
 - Kernel doesn't know the difference!
- Text file is sequence of text lines
 - Text line is sequence of chars terminated by newline char ('\n')
 - Newline is 0xa, same as ASCII line feed character (LF)
- End of line (EOL) indicators in other systems
 - Linux and Mac OS: '\n' (0xa)
 - line feed (LF)
 - Windows and Internet protocols: '\r\n' (0xd 0xa)
 - Carriage return (CR) followed by line feed (LF)

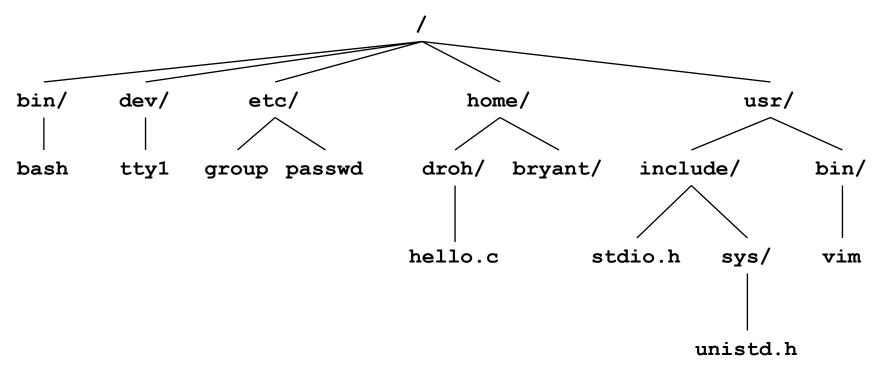


Directories

- Directory consists of an array of *links*
 - Each link maps a filename to a file
- Each directory contains at least two entries
 - . (dot) is a link to itself
 - . . (dot dot) is a link to the parent directory in the directory hierarchy (next slide)
- Commands for manipulating directories
 - mkdir: create empty directory
 - 1s: view directory contents
 - rmdir: delete empty directory

Directory Hierarchy

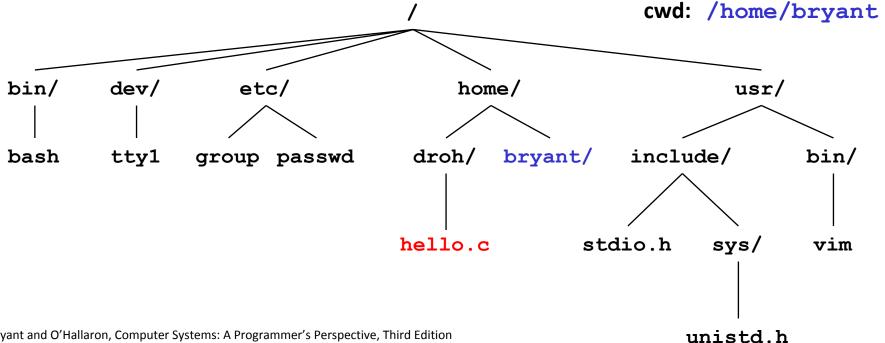
 All files are organized as a hierarchy anchored by root directory named / (slash)



- Kernel maintains current working directory (cwd) for each process
 - Modified using the cd command

Pathnames

- Locations of files in the hierarchy denoted by *pathnames*
 - Absolute pathname starts with '/' and denotes path from root
 - home/droh/hello.c
 - Relative pathname denotes path from current working directory
 - ../home/droh/hello.c



Opening Files

 Opening a file informs the kernel that you are getting ready to access that file

```
int fd; /* file descriptor */
if ((fd = open("/etc/hosts", O_RDONLY)) < 0) {
   perror("open");
   exit(1);
}</pre>
```

- Returns a small identifying integer *file descriptor*
 - fd == -1 indicates that an error occurred
- Each process created by a Linux shell begins life with three open files associated with a terminal:
 - 0: standard input (stdin)
 - 1: standard output (stdout)
 - 2: standard error (stderr)

Closing Files

Closing a file informs the kernel that you are finished accessing that file

```
int fd;  /* file descriptor */
int retval; /* return value */

if ((retval = close(fd)) < 0) {
   perror("close");
   exit(1);
}</pre>
```

- Closing an already closed file is a recipe for disaster in threaded programs (more on this later)
- Moral: Always check return codes, even for seemingly benign functions such as close()

Reading Files

 Reading a file copies bytes from the current file position to memory, and then updates file position

- Returns number of bytes read from file fd into buf
 - Return type ssize_t is signed integer
 - nbytes < 0 indicates that an error occurred
 - Short counts (nbytes < sizeof (buf)) are possible and are not errors!</p>

Writing Files

 Writing a file copies bytes from memory to the current file position, and then updates current file position

- Returns number of bytes written from buf to file fd
 - nbytes < 0 indicates that an error occurred</p>
 - As with reads, short counts are possible and are not errors!

Simple Unix I/O example

Copying stdin to stdout, one byte at a time

```
#include "csapp.h"
int main(void)
{
    char c;
    while(! (STDIN_FILENO, &c, 1) != 0)
        write STDOUT_FILENO, &c, 1);
    exit(0);
}
```

Always check return codes from system calls!

Simple Unix I/O example

Copying stdin to stdout, one byte at a time

```
#include "csapp.h"
int main(void)
{
    char c;

    while(Read(STDIN_FILENO, &c, 1) != 0)
        Write(STDOUT_FILENO, &c, 1);
    exit(0);
}
```

On Short Counts

- Short counts can occur in these situations:
 - Encountering (end-of-file) EOF on reads
 - Reading text lines from a terminal
 - Reading and writing network sockets
- Short counts never occur in these situations:
 - Reading from disk files (except for EOF)
 - Writing to disk files
- Best practice is to always allow for short counts.

Today

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- RIO (robust I/O) package
- Standard I/O
- Which I/O when
- Metadata, sharing, and redirection

The RIO Package (15-213/CS:APP Package)

RIO is a set of wrappers that provide efficient and robust I/O in apps, such as network programs that are subject to short counts

- RIO provides two different kinds of functions
 - Unbuffered input and output of binary data
 - rio_readn and rio_writen
 - Buffered input of text lines and binary data
 - rio_readlineb and rio_readnb
 - Buffered RIO routines are thread-safe and can be interleaved arbitrarily on the same descriptor
- Download from http://csapp.cs.cmu.edu/3e/code.html
 - → src/csapp.c and include/csapp.h

Unbuffered RIO Input and Output

- Same interface as Unix read and write
- Especially useful for transferring data on network sockets

```
#include "csapp.h"
ssize_t rio_readn(int fd, void *usrbuf, size_t n);
ssize_t rio_writen(int fd, void *usrbuf, size_t n);
Return: num. bytes transferred if OK, 0 on EOF (rio_readn only), -1 on error
```

- rio_readn returns short count only if it encounters EOF
 - Only use it when you know how many bytes to read
- rio_writen never returns a short count
- Calls to rio_readn and rio_writen can be interleaved arbitrarily on the same descriptor

Implementation of rio readn

```
/*
* rio readn - Robustly read n bytes (unbuffered)
*/
ssize t rio readn(int fd, void *usrbuf, size t n)
   size t nleft = n;
   ssize t nread;
   char *bufp = usrbuf;
   while (nleft > 0) {
   if ((nread = read(fd, bufp, nleft)) < 0) {</pre>
       if (errno == EINTR) /* Interrupted by sig handler return */
       nread = 0;  /* and call read() again */
       else
       return -1; /* errno set by read() */
   else if (nread == 0)
                         /* EOF */
       break;
   nleft -= nread;
   bufp += nread;
   return (n - nleft); /* Return >= 0 */
```

Buffered RIO Input Functions

 Efficiently read text lines and binary data from a file partially cached in an internal memory buffer

```
#include "csapp.h"

void rio_readinitb(rio_t *rp, int fd);

ssize_t rio_readlineb(rio_t *rp, void *usrbuf, size_t maxlen);
ssize_t rio_readnb(rio_t *rp, void *usrbuf, size_t n);

Return: num. bytes read if OK, 0 on EOF, -1 on error
```

- rio_readlineb reads a text line of up to maxlen bytes from file fd and stores the line in usrbuf
 - Especially useful for reading text lines from network sockets
- Stopping conditions
 - maxlen bytes read
 - EOF encountered
 - Newline ('\n') encountered

Buffered RIO Input Functions (cont)

```
#include "csapp.h"

void rio_readinitb(rio_t *rp, int fd);

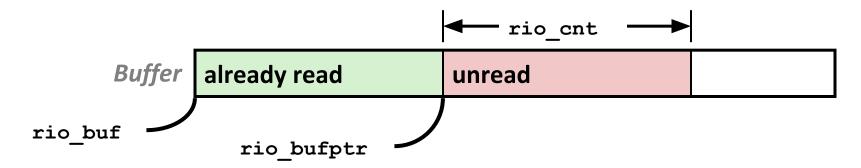
ssize_t rio_readlineb(rio_t *rp, void *usrbuf, size_t maxlen);
ssize_t rio_readnb(rio_t *rp, void *usrbuf, size_t n);

Return: num. bytes read if OK, 0 on EOF, -1 on error
```

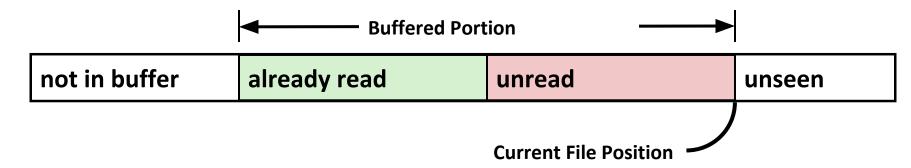
- rio readnb reads up to n bytes from file fd
- Stopping conditions
 - maxlen bytes read
 - EOF encountered
- Calls to rio_readlineb and rio_readnb can be interleaved arbitrarily on the same descriptor
 - Warning: Don't interleave with calls to rio_readn

Buffered I/O: Implementation

- For reading from file
- File has associated buffer to hold bytes that have been read from file but not yet read by user code

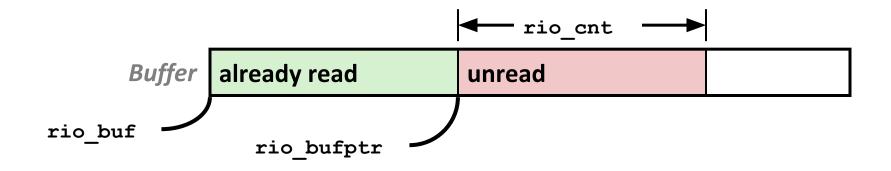


Layered on Unix file:



Buffered I/O: Declaration

All information contained in struct



RIO Example

 Copying the lines of a text file from standard input to standard output

```
#include "csapp.h"

int main(int argc, char **argv)
{
    int n;
    rio_t rio;
    char buf[MAXLINE];

    Rio_readinitb(&rio, STDIN_FILENO);
    while((n = Rio_readlineb(&rio, buf, MAXLINE)) != 0)
    Rio_writen(STDOUT_FILENO, buf, n);
    exit(0);
}
```

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Standard I/O Functions

- The C standard library (libc.so) contains a collection of higher-level standard I/O functions
 - Documented in Appendix B of K&R
- Examples of standard I/O functions:
 - Opening and closing files (fopen and fclose)
 - Reading and writing bytes (fread and fwrite)
 - Reading and writing text lines (fgets and fputs)
 - Formatted reading and writing (fscanf and fprintf)

Standard I/O Streams

- Standard I/O models open files as streams
 - Abstraction for a file descriptor and a buffer in memory
- C programs begin life with three open streams (defined in stdio.h)
 - stdin (standard input)
 - stdout (standard output)
 - stderr (standard error)

```
#include <stdio.h>
extern FILE *stdin; /* standard input (descriptor 0) */
extern FILE *stdout; /* standard output (descriptor 1) */
extern FILE *stderr; /* standard error (descriptor 2) */
int main() {
   fprintf(stdout, "Hello, world\n");
}
```

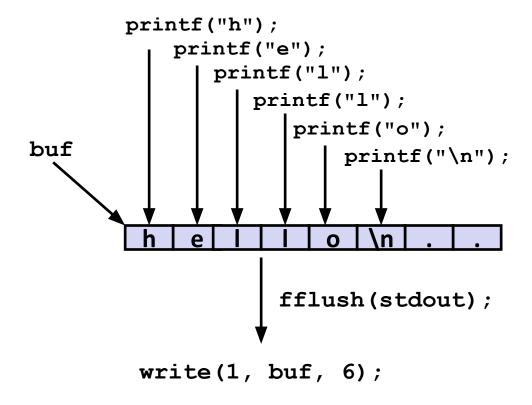
Buffered I/O: Motivation

- Applications often read/write one character at a time
 - getc, putc, ungetc
 - gets, fgets
 - Read line of text one character at a time, stopping at newline
- Implementing as Unix I/O calls expensive
 - read and write require Unix kernel calls
 - > 10,000 clock cycles
- Solution: Buffered read
 - Use Unix read to grab block of bytes
 - User input functions take one byte at a time from buffer
 - Refill buffer when empty



Buffering in Standard I/O

Standard I/O functions use buffered I/O



Buffer flushed to output fd on "\n", call to fflush or exit, or return from main.

Standard I/O Buffering in Action

You can see this buffering in action for yourself, using the always fascinating Linux strace program:

```
#include <stdio.h>
int main()
{
    printf("h");
    printf("e");
    printf("l");
    printf("l");
    printf("o");
    printf("\n");
    fflush(stdout);
    exit(0);
}
```

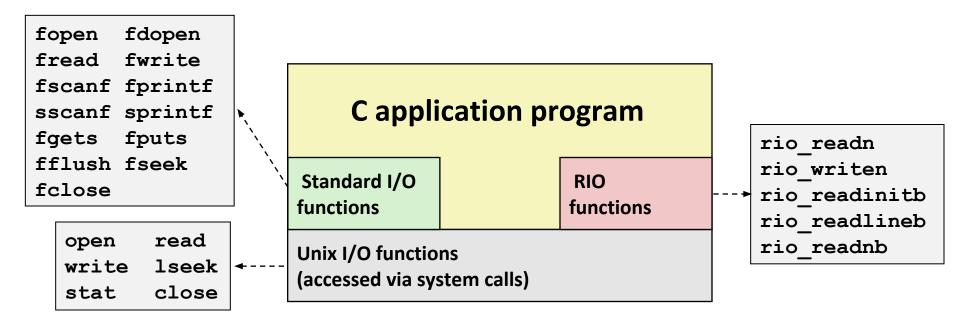
```
linux> strace ./hello
execve("./hello", ["hello"], [/* ... */]).
...
write(1, "hello\n", 6) = 6
...
exit_group(0) = ?
```

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Unix I/O vs. Standard I/O vs. RIO

Standard I/O and RIO are implemented using low-level Unix I/O



Which ones should you use in your programs?

Pros and Cons of Unix I/O

Pros

- Unix I/O is the most general and lowest overhead form of I/O
 - All other I/O packages are implemented using Unix I/O functions
- Unix I/O provides functions for accessing file metadata
- Unix I/O functions are async-signal-safe and can be used safely in signal handlers

Cons

- Dealing with short counts is tricky and error prone
- Efficient reading of text lines requires some form of buffering, also tricky and error prone
- Both of these issues are addressed by the standard I/O and RIO packages

Pros and Cons of Standard I/O

Pros:

- Buffering increases efficiency by decreasing the number of read and write system calls
- Short counts are handled automatically

Cons:

- Provides no function for accessing file metadata
- Standard I/O functions are not async-signal-safe, and not appropriate for signal handlers
- Standard I/O is not appropriate for input and output on network sockets
 - There are poorly documented restrictions on streams that interact badly with restrictions on sockets (CS:APP3e, Sec 10.11)

Choosing I/O Functions

- General rule: use the highest-level I/O functions you can
 - Many C programmers are able to do all of their work using the standard I/O functions
 - But, be sure to understand the functions you use!
- When to use standard I/O
 - When working with disk or terminal files
- When to use raw Unix I/O
 - Inside signal handlers, because Unix I/O is async-signal-safe
 - In rare cases when you need absolute highest performance
- When to use RIO
 - When you are reading and writing network sockets
 - Avoid using standard I/O on sockets

Aside: Working with Binary Files

- Functions you should never use on binary files
 - Text-oriented I/O: such as fgets, scanf, rio_readlineb
 - Interpret EOL characters.
 - Use functions like rio_readn or rio_readnb instead
 - String functions
 - strlen, strcpy, strcat
 - Interprets byte value 0 (end of string) as special

Today

- Unix I/O
- RIO (robust I/O) package
- Standard I/O
- Metadata, sharing, and redirection
- Closing remarks

File Metadata

- Metadata is data about data, in this case file data
- Per-file metadata maintained by kernel
 - accessed by users with the stat and fstat functions

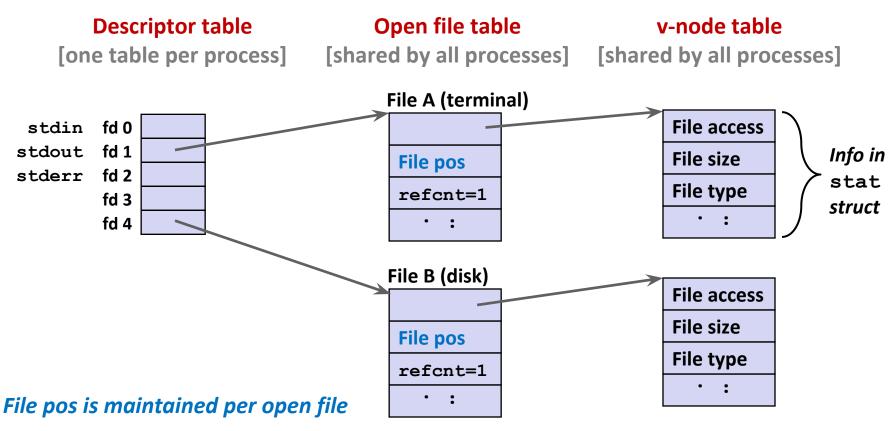
```
/* Metadata returned by the stat and fstat functions */
struct stat {
               st dev; /* Device */
   dev t
               st ino; /* inode */
   ino t
              st_mode; /* Protection and file type */
   mode t
   nlink_t st_nlink; /* Number of hard links */
              st uid; /* User ID of owner */
   uid t
              st gid; /* Group ID of owner */
   gid t
   dev t
              st_rdev; /* Device type (if inode device) */
   off t
              st size; /* Total size, in bytes */
   unsigned long st blksize; /* Blocksize for filesystem I/O */
   unsigned long st blocks; /* Number of blocks allocated */
   time_t st_atime; /* Time of last access */
   time t st mtime; /* Time of last modification */
               st ctime; /* Time of last change */
   time t
```

Example of Accessing File Metadata

```
linux> ./statcheck statcheck.c
int main (int argc, char **argv)
                                      type: regular, read: yes
                                      linux> chmod 000 statcheck.c
   struct stat stat:
                                       linux> ./statcheck statcheck.c
   char *type, *readok;
                                      type: regular, read: no
                                      linux> ./statcheck ...
   Stat(argv[1], &stat);
                                      type: directory, read: yes
    if (S ISREG(stat.st mode)) /* Determine file type */
   type = "regular";
   else if (S ISDIR(stat.st mode))
   type = "directory";
   else
       type = "other";
    if ((stat.st mode & S IRUSR)) /* Check read access */
   readok = "yes";
   else
       readok = "no";
   printf("type: %s, read: %s\n", type, readok);
   exit(0);
                                                     statcheck.c
```

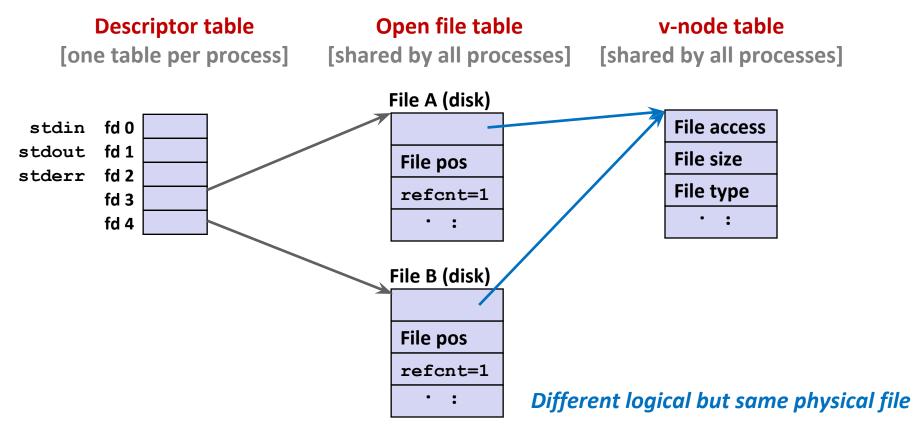
How the Unix Kernel Represents Open Files

Two descriptors referencing two distinct open files.
 Descriptor 1 (stdout) points to terminal, and descriptor 4 points to open disk file



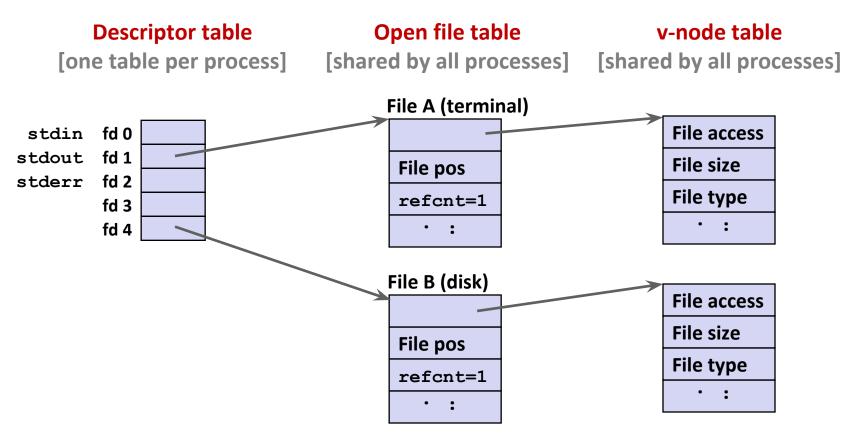
File Sharing

- Two distinct descriptors sharing the same disk file through two distinct open file table entries
 - E.g., Calling open twice with the same filename argument



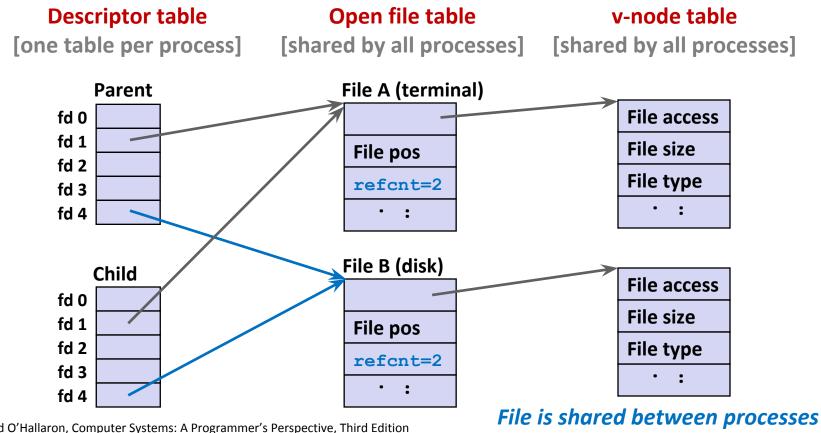
How Processes Share Files: fork

- A child process inherits its parent's open files
 - Note: situation unchanged by exec functions (use fcntl to change)
- Before fork call:



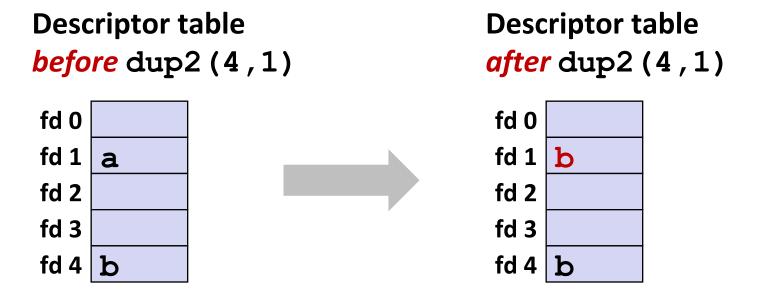
How Processes Share Files: fork

- A child process inherits its parent's open files
- After fork:
 - Child's table same as parent's, and +1 to each refent



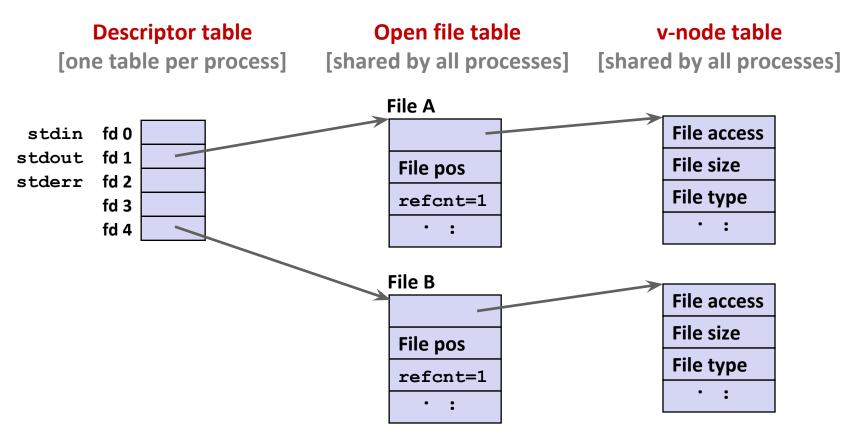
I/O Redirection

- Question: How does a shell implement I/O redirection?
 linux> ls > foo.txt
- Answer: By calling the dup2 (oldfd, newfd) function
 - Copies (per-process) descriptor table entry oldfd to entry newfd



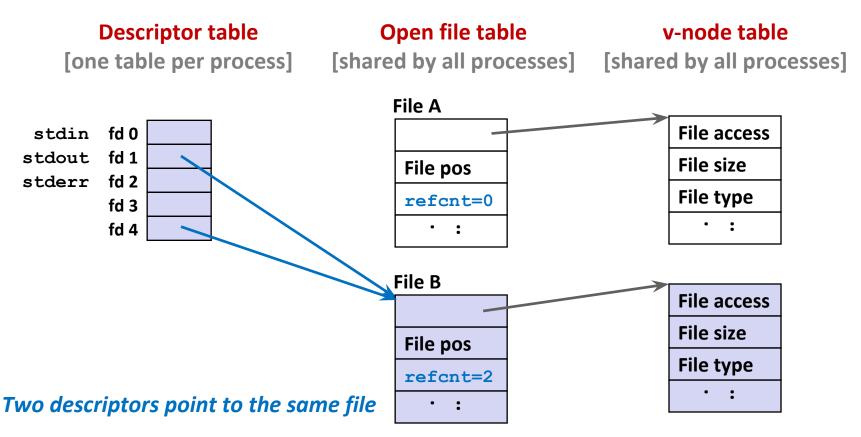
I/O Redirection Example

- Step #1: open file to which stdout should be redirected
 - Happens in child executing shell code, before exec



I/O Redirection Example (cont.)

- Step #2: call dup2 (4,1)
 - cause fd=1 (stdout) to refer to disk file pointed at by fd=4



Warm-Up: I/O and Redirection Example

```
#include "csapp.h"
int main(int argc, char *argv[])
    int fd1, fd2, fd3;
   char c1, c2, c3;
   char *fname = arqv[1];
    fd1 = Open(fname, O RDONLY, 0);
    fd2 = Open(fname, O RDONLY, 0);
    fd3 = Open(fname, O RDONLY, 0);
   Dup2 (fd2, fd3);
   Read(fd1, &c1, 1);
   Read(fd2, &c2, 1);
   Read(fd3, &c3, 1);
   printf("c1 = %c, c2 = %c, c3 = %c\n", c1, c2, c3);
   return 0;
                                              ffiles1.c
```

Warm-Up: I/O and Redirection Example

```
#include "csapp.h"
int main(int argc, char *argv[])
                                      c1 = a, c2 = a, c3 = b
   int fd1, fd2, fd3;
   char c1, c2, c3;
   char *fname = arqv[1];
   fd1 = Open(fname, O RDONLY, 0);
   fd2 = Open(fname, O RDONLY, 0);
   fd3 = Open(fname, O RDONLY, 0);
                                      dup2(oldfd, newfd)
   Dup2(fd2, fd3);
   Read(fd1, &c1, 1);
   Read(fd2, &c2, 1);
   Read(fd3, &c3, 1);
   printf("c1 = %c, c2 = %c, c3 = %c\n", c1, c2, c3);
   return 0;
                                             ffiles1.c
```

Master Class: Process Control and I/O

```
#include "csapp.h"
int main(int argc, char *argv[])
    int fd1;
    int s = getpid() & 0x1;
    char c1, c2;
    char *fname = argv[1];
    fd1 = Open(fname, O RDONLY, 0);
   Read(fd1, &c1, 1);
    if (fork()) { /* Parent */
        sleep(s);
        Read(fd1, &c2, 1);
        printf("Parent: c1 = %c, c2 = %c\n", c1, c2);
    } else { /* Child */
        sleep(1-s);
        Read(fd1, &c2, 1);
        printf("Child: c1 = %c, c2 = %c\n", c1, c2);
    return 0;
                                            ffiles2.c
```

Master Class: Process Control and I/O

```
#include "csapp.h"
                                       Child: c1 = a, c2 = b
int main(int argc, char *argv[])
                                      Parent: c1 = a, c2 = c
   int fd1;
   int s = getpid() & 0x1;
                                      Parent: c1 = a, c2 = b
   char c1, c2;
   char *fname = argv[1];
                                       Child: c1 = a, c2 = c
    fd1 = Open(fname, O RDONLY, 0);
   Read(fd1, &c1, 1);
                                      Bonus: Which way does it go?
   if (fork()) { /* Parent */
       sleep(s);
       Read(fd1, &c2, 1);
       printf("Parent: c1 = %c, c2 = %c\n", c1, c2);
    } else { /* Child */
       sleep(1-s);
       Read(fd1, &c2, 1);
       printf("Child: c1 = %c, c2 = %c\n", c1, c2);
    return 0;
                                          ffiles2.c
```

I/O Questions in Exams

Problem 10. (6 points):

Unix I/O.

A. Suppose that the disk file foobar.txt consists of the six ASCII characters "foobar". What is the output of the following program?

```
/* any necessary includes */
char buf[20] = \{0\}; /* init to all zeroes */
int main(int argc, char* argv[]) {
    int fd1 = open("foobar.txt", O_RDONLY);
    int fd2 = open("foobar.txt", O RDONLY);
    dup2(fd2, fd1);
    read(fd1, buf, 3);
    close (fd1);
    read(fd2, &buf[3], 3);
    close (fd2);
   printf("buf = %s\n", buf);
   return 0;
Output: buf = _____
```

```
*******

Problem 10

*******

A. Output: buf = foobar
```

Fall 2011 (model solution)

Extra Slides

For Further Information

The Unix bible:

- W. Richard Stevens & Stephen A. Rago, Advanced Programming in the Unix Environment, 2nd Edition, Addison Wesley, 2005
 - Updated from Stevens's 1993 classic text

The Linux bible:

- Michael Kerrisk, The Linux Programming Interface, No Starch Press, 2010
 - Encyclopedic and authoritative

Fun with File Descriptors (1)

```
#include "csapp.h"
int main(int argc, char *argv[])
    int fd1, fd2, fd3;
   char c1, c2, c3;
   char *fname = arqv[1];
    fd1 = Open(fname, O RDONLY, 0);
    fd2 = Open(fname, O RDONLY, 0);
    fd3 = Open(fname, O RDONLY, 0);
   Dup2 (fd2, fd3);
   Read(fd1, &c1, 1);
   Read(fd2, &c2, 1);
   Read(fd3, &c3, 1);
   printf("c1 = %c, c2 = %c, c3 = %c\n", c1, c2, c3);
   return 0;
                                              ffiles1.c
```

Fun with File Descriptors (2)

```
#include "csapp.h"
int main(int argc, char *argv[])
    int fd1;
    int s = getpid() & 0x1;
    char c1, c2;
    char *fname = argv[1];
    fd1 = Open(fname, O RDONLY, 0);
   Read(fd1, &c1, 1);
    if (fork()) { /* Parent */
        sleep(s);
        Read(fd1, &c2, 1);
        printf("Parent: c1 = %c, c2 = %c\n", c1, c2);
    } else { /* Child */
        sleep(1-s);
        Read(fd1, &c2, 1);
        printf("Child: c1 = %c, c2 = %c\n", c1, c2);
    return 0;
                                            ffiles2.c
```

Fun with File Descriptors (3)

```
#include "csapp.h"
int main(int argc, char *argv[])
    int fd1, fd2, fd3;
    char *fname = arqv[1];
    fd1 = Open(fname, O CREAT|O TRUNC|O RDWR, S IRUSR|S IWUSR);
   Write(fd1, "pqrs", 4);
    fd3 = Open(fname, O APPEND|O WRONLY, 0);
   Write(fd3, "jklmn", 5);
    fd2 = dup(fd1); /* Allocates descriptor */
   Write(fd2, "wxyz", 4);
   Write(fd3, "ef", 2);
    return 0;
                                                       ffiles3.c
```

What would be the contents of the resulting file?

Accessing Directories

- Only recommended operation on a directory: read its entries
 - dirent structure contains information about a directory entry
 - DIR structure contains information about directory while stepping through its entries

```
#include <sys/types.h>
#include <dirent.h>
 DIR *directory;
  struct dirent *de;
  if (!(directory = opendir(dir name)))
      error("Failed to open directory");
  while (0 != (de = readdir(directory))) {
      printf("Found file: %s\n", de->d name);
  closedir(directory);
```