## Lecture 3 Activity Solution

## Model 0: Review of Addition / Positive

1. 10110
2. 5 bits are required.
3. The number of bits in the result is one more than the number of bits of the operands.
4. You could truncate overflow bits, resulting in 0110.

## Model 0: Review of Negative Integers

1. The leftmost bit in a non-negative number in two's complement is 0 .

| Bits | Most Positive | Most Negative |
| :--- | :--- | :--- |
| 1 | 0 | -1 |
| 2 | 1 | -2 |
| 3 | 3 | -4 |
| 4 | 7 | -8 |

3. $2^{N-1}-1$
4. $-\left(2^{N-1}\right)$
5. 10011111. If the two numbers are unsigned, the result is correct $\left(1111000_{2}=120_{10}, 0100111_{2}=\right.$ $\left.39_{10}, 10011111_{2}=159_{10}\right)$, but the result is not correct for signed numbers $\left(1111000_{2}=-8_{10}, 0100111_{2}=\right.$ $\left.39_{10}, 10011111_{2}=-97_{10}\right)$.
1. No, but the difference in expected results for signed integers comes from improper handling of overflow or sign extension.

## Model 1: Bit-Level Operations

1.     - 0x3501

- 0xC3C3
- 0xFFFF

|  | OP0 | OP1 | AND | OR | XOR |
| :--- | :--- | :--- | :--- | :--- | :--- |
| 0 | 0 | 0 | 0 | 0 |  |
| 2. | 1 | 0 | 0 | 1 | 1 |
| 0 | 1 | 0 | 1 | 1 |  |
| 1 | 1 | 1 | 1 | 0 |  |


| Dec | Bin | X \& 0x1 |
| :--- | :--- | :--- |
| -2 | 1110 | 0000 |
| -1 | 1111 | 0001 |
| 0 | 0000 | 0000 |
| 1 | 0001 | 0001 |
| 2 | 0010 | 0000 |

4. The decimal numbers -1 and 1 , which both are odd and therefore have a 1 in the rightmost (least-significant) bit.
5. for each bit in X :
```
    if that bit is set in FLAG but not set in X:
        return false
return true
```

6. The OR (|) operation is setting the relevant bits in the file access mode to create a flag with the bits set for all of O_WRONLY, O_CREAT, and O_TRUNC.
7. 

| x | y | $\sim(\mathrm{x} \& \mathrm{y})$ | $(\sim \mathrm{x}) \mid(\sim \mathrm{y})$ | equal? |
| :---: | :---: | :---: | :---: | :---: |
| 0 xF | 0 x 1 | 1110 | 1110 | Y |
| 0 x 5 | 0 x 7 | 1010 | 1010 | Y |
| 0 x 3 | 0 xC | 1111 | 1111 | Y |

## Model 2: Logical Operations

1. 1 value is false and 15 values are true.
2. $0 \mathrm{x} 3 \& \& 0 \mathrm{xC}=0001,0 \mathrm{x} 3 \& 0 \mathrm{xC}=0000$, so $0 \mathrm{x} 3 \& \& 0 \mathrm{xC}==0 \mathrm{x} 3 \& 0 \mathrm{xC}$ is false.

| X | $!\mathrm{X}$ | $!!\mathrm{X}$ | $!!\mathrm{X}==\mathrm{X}$ |
| :---: | :---: | :---: | :---: |
| -1 | 0 | 1 | 0 |
| 0 | 1 | 0 | 1 |
| 1 | 0 | 1 | 1 |
| 2 | 0 | 1 | 0 |

4. Yes, the results differ-every $\sim \sim X=X$. Note that $\sim \sim X$ is a no-op (gives $X$ back for all $X$ ) while !! X is not.

## Model 2: Multiplication and Division

1. 

| Value | $\ll$ | Result |
| :---: | :---: | :---: |
| 0x30 | 1 | $0 \times 60$ |
| 0x5A | 4 | $0 \times 5 \mathrm{~A} 0$ |
| 0x11D | 31 | $0 \times 80000000$ |

2. $\mathrm{X}=6_{10}=0110_{2}$
3. Two acceptable answers: $\mathrm{x} \ll 2+\mathrm{x} \ll 1$, or $(\mathrm{x}+\mathrm{x}+\mathrm{x}) \ll 1$.
4. The largest 3 -bit unsigned integer is $111_{2}=7_{10}$, its value squared is 49 , which requires 6 bits.
5. $001_{2}=1_{10}$, if truncating excess bits.
6. | Value | $\gg$ | Result |
| :---: | :---: | :---: |
| $0 \times 30$ | 1 | $0 \times 18$ |
| $0 \times 5 \mathrm{~A}$ | 4 | $0 \times 5$ |
| $0 \times 11$ | 3 | 0 x 2 |
7. | Value | $\gg$ | Result |
| :---: | :---: | :---: |
| 48 | 1 | 24 |
| 90 | 4 | 5 |
| 17 | 3 | 2 |

A single right shift is equivalent to dividing by 2 , so right shifting by $N$ is equivalent to dividing by $2^{N}$.
8. $0 \mathrm{xA} \gg 1=0 \mathrm{x} 5$
9. We expect that $-2 \gg 1=-1$.
10. $-2_{10}=1110_{2}$ in two's complement. After right shifting by 1 , we get $0111_{2}=7_{10}$.
11. We could replicate the most significant (leftmost) bit so all bits shifted "in" would be copies of the leftmost bit instead of zeroes.
12. while ( x ! $=0$ )
\{
saveNextBit(x \& 0x1);
$\mathrm{x}=\mathrm{x} \gg 1$;
\}

