

15-213

Machine-Level Programming II

Control Flow

Feb. 3, 2000

Topics

- **Condition Codes**
 - Setting
 - Testing
- **Control Flow**
 - If-then-else
 - Varieties of Loops
 - Switch Statements

Condition Codes

Single Bit Registers

CF	Carry Flag
ZF	Zero Flag
SF	Sign Flag
OF	Overflow Flag

Implicit Setting By Arithmetic Operations

`addl Src, Dest`

C analog: $t = a + b$

- CF set if carry out from most significant bit
 - Used to detect unsigned overflow
- ZF set if $t == 0$
- SF set if $t < 0$
- OF set if two's complement overflow
 $(a > 0 \ \&\& \ b > 0 \ \&\& \ t < 0) \ || \ (a < 0 \ \&\& \ b < 0 \ \&\& \ t > 0)$

Not Set by `leal` instruction

Setting Condition Codes (cont.)

Explicit Setting by Compare Instruction

`cmp1 Src2,Src1`

- `cmp1 b,a` like computing `a-b` without setting destination
- CF set if carry out from most significant bit
 - Used for unsigned comparisons
- ZF set if `a == b`
- SF set if `(a-b) < 0`
- OF set if two's complement overflow
 - `(a>0 && b<0 && (a-b)<0) || (a<0 && b>0 && (a-b)>0)`

Explicit Setting by Test instruction

`test1 Src2,Src1`

- Sets condition codes based on value of `Src1` & `Src2`
 - Useful to have one of the operands be a mask
- `test1 b,a` like computing `a&b` without setting destination
- ZF set when `a&b == 0`
- SF set when `a&b < 0`

Reading Condition Codes

SetX Instructions

- Set single byte based on combinations of condition codes

SetX	Condition	Description
sete	ZF	Equal / Zero
setne	~ZF	Not Equal / Not Zero
sets	SF	Negative
setns	~SF	Nonnegative
setg	~(SF^OF) & ~ZF	Greater (Signed)
setge	~(SF^OF)	Greater or Equal (Signed)
setl	(SF^OF)	Less (Signed)
setle	(SF^OF) ZF	Less or Equal (Signed)
seta	~CF & ~ZF	Above (unsigned)
setb	CF	Below (unsigned)

Reading Condition Codes (Cont.)

SetX Instructions

- Set single byte based on combinations of condition codes
- One of 8 addressable byte registers
 - Embedded within first 4 integer registers
 - Does not alter remaining 3 bytes
 - Typically use `andl 0xFF,%eax` to finish job

```
int gt (int x, int y)
{
    return x > y;
}
```

Body

```
movl 12(%ebp),%eax # eax = y
cmpl %eax,8(%ebp)  # Compare x : eax ←
setg %al           # al = x > y
andl $255,%eax    # Zero rest of %eax
```

Note
inverted
ordering!

%eax	%ah	%al
%edx	%dh	%dl
%ecx	%ch	%cl
%ebx	%bh	%bl
%esi		
%edi		
%esp		
%ebp		

Jumping

jX Instructions

- Jump to different part of code depending on condition codes

jX	Condition	Description
jmp	1	Unconditional
je	ZF	Equal / Zero
jne	\sim ZF	Not Equal / Not Zero
js	SF	Negative
jns	\sim SF	Nonnegative
jg	$\sim (SF \wedge OF) \& \sim ZF$	Greater (Signed)
jge	$\sim (SF \wedge OF)$	Greater or Equal (Signed)
jl	$(SF \wedge OF)$	Less (Signed)
jle	$(SF \wedge OF) \mid ZF$	Less or Equal (Signed)
ja	$\sim CF \& \sim ZF$	Above (unsigned)
jb	CF	Below (unsigned)

Conditional Branch Example

```
int max(int x, int y)
{
    if (x > y)
        return x;
    else
        return y;
}
```

_max:

```
    pushl %ebp
    movl  %esp,%ebp
```

} Set Up

```
    movl  8(%ebp),%edx
    movl  12(%ebp),%eax
    cmpl  %eax,%edx
    jle  L9
    movl  %edx,%eax
```

} Body

L9:

```
    movl  %ebp,%esp
    popl  %ebp
    ret
```

} Finish

Conditional Branch Example (Cont.)

```
int goto_max(int x, int y)
{
    int rval = y;
    int ok = (x <= y);
    if (ok)
        goto done;
    rval = x;
done:
    return rval;
}
```

- C allows “goto” as means of transferring control
 - Closer to machine-level programming style
- Generally considered bad coding style

```
    movl 8(%ebp),%edx    # edx = x
    movl 12(%ebp),%eax   # eax = y
    cmpl %eax,%edx      # x : y
    jle L9              # if <= goto L9
    movl %edx,%eax      # eax = x } Skipped when x > y
L9:                    # Done:
```


“Do-While” Loop Example

C Code

```
int fact_do
(int x)
{
    int result = 1;
    do {
        result *= x;
        x = x-1;
    } while (x > 1);
    return result;
}
```

Goto Version

```
int fact_goto(int x)
{
    int result = 1;
loop:
    result *= x;
    x = x-1;
    if (x > 1)
        goto loop;
    return result;
}
```

- Use backward branch to continue looping
- Only take branch when “while” condition holds

“Do-While” Loop Compilation

Goto Version

```
int fact_goto
(int x)
{
    int result = 1;
loop:
    result *= x;
    x = x-1;
    if (x > 1)
        goto loop;
    return result;
}
```

Registers

```
%edx    x
%eax    result
```

Assembly

```
_fact_goto:
    pushl %ebp                # Setup
    movl %esp,%ebp           # Setup
    movl $1,%eax              # eax = 1
    movl 8(%ebp),%edx         # edx = x

L11:
    imull %edx,%eax           # result *= x
    decl %edx                  # x--
    cmpl $1,%edx              # Compare x : 1
    jg L11                     # if > goto loop

    movl %ebp,%esp           # Finish
    popl %ebp                 # Finish
    ret                        # Finish
```

General “Do-While” Translation

C Code

```
do  
  Body  
while (Test);
```

Goto Version

```
loop:  
  Body  
  if (Test)  
    goto loop
```

- *Body* can be any C statement
 - Typically compound statement:

```
{  
  Statement1;  
  Statement2;  
  ...  
  Statementn;  
}
```

- *Test* is expression returning integer
 - = 0 interpreted as false ≠0 interpreted as true

“While” Loop Example #1

C Code

```
int fact_while
(int x)
{
    int result = 1;
    while (x > 1) {
        result *= x;
        x = x-1;
    };
    return result;
}
```

First Goto Version

```
int fact_while_goto
(int x)
{
    int result = 1;
loop:
    if (!(x > 1))
        goto done;
    result *= x;
    x = x-1;
    goto loop;
done:
    return result;
}
```

- Is this code equivalent to the do-while version?
- Must jump out of loop if test fails

Actual “While” Loop Translation

C Code

```
int fact_while(int x)
{
    int result = 1;
    while (x > 1) {
        result *= x;
        x = x-1;
    };
    return result;
}
```

- Uses same inner loop as do-while version
- Guards loop entry with extra test

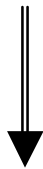
Second Goto Version

```
int fact_while_goto2
(int x)
{
    int result = 1;
    if (!(x > 1))
        goto done;
loop:
    result *= x;
    x = x-1;
    if (x > 1)
        goto loop;
done:
    return result;
}
```

General “While” Translation

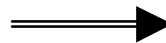
C Code

```
while (Test)  
  Body
```



Do-While Version

```
if (!Test)  
  goto done;  
do  
  Body  
  while(Test);  
done:
```



Goto Version

```
if (!Test)  
  goto done;  
loop:  
  Body  
  if (Test)  
    goto loop;  
done:
```

“While” Loop Example #2

```

/* Compute x raised to nonnegative power p */
int ipwr_while(int x, unsigned p)
{
    int result = 1;
    while (p) {
        if (p & 0x1)
            result *= x;
        x = x*x;
        p = p>>1;
    }
    return result;
}

```

Algorithm

- Exploit property that $p = p_0 + 2p_1 + 4p_2 + \dots + 2^{n-1}p_{n-1}$
- Gives: $x^p = z_0 \cdot z_1^2 \cdot (z_2^2)^2 \cdot \dots \cdot \underbrace{(\dots((z_{n-1}^2)^2)\dots)^2}_{n \text{ times}}$
 $z_i = 1$ when $p_i = 0$
 $z_i = x$ when $p_i = 1$
- Complexity $O(\log p)$

Example

$$\begin{aligned}
 3^{10} &= 3^2 * 3^8 \\
 &= 3^2 * ((3^2)^2)^2
 \end{aligned}$$

ipwr Computation

```
int ipwr(int x, unsigned p)
{
    int result = 1;
    while (p) {
        if (p & 0x1)
            result *= x;
        x = x*x;
        p = p>>1;
    }
    return result;
}
```

result	x	p
1	3	10
1	9	5
9	81	2
9	6561	1
531441	43046721	0

"While" ® "Do-While" ® "Goto"

```
int result = 1;
while (p) {
    if (p & 0x1)
        result *= x;
    x = x*x;
    p = p>>1;
}
```



```
int result = 1;
if (!p) goto done;
do {
    if (p & 0x1)
        result *= x;
    x = x*x;
    p = p>>1;
} while (p);
done:
```



```
int result = 1;
if (!p)
    goto done;
loop:
    if (!(p & 0x1))
        goto skip;
    result *= x;
skip:
    x = x*x;
    p = p>>1;
    if (p)
        goto loop;
done:
```

- Also converted conditional update into test and branch around update code

Example #2 Compilation

Goto Version

```
int result = 1;
if (!p)
    goto done;
loop:
    if (!(p & 0x1))
        goto skip;
    result *= x;
skip:
    x = x*x;
    p = p>>1;
    if (p)
        goto loop;
done:
```

Registers

```
%ecx    x
%edx    p
%eax    result
```

```
    pushl %ebp                # Setup
    movl %esp,%ebp           # Setup
    movl $1,%eax             # eax = 1
    movl 8(%ebp),%ecx        # ecx = x
    movl 12(%ebp),%edx       # edx = p
    testl %edx,%edx         # Test p
    je L36                   # If 0, goto done
L37:                          # Loop:
    testb $1,%dl            # Test p & 0x1
    je L38                   # If 0, goto skip
    imull %ecx,%eax         # result *= x
L38:                          # Skip:
    imull %ecx,%ecx         # x *= x
    shrl $1,%edx            # p >>= 1
    jne L37                 # if p goto Loop
L36:                          # Done:
    movl %ebp,%esp         # Finish
    popl %ebp              # Finish
    ret                    # Finish
```

"For" Loop Example

General Form

```
int result;
for (result = 1;
    p != 0;
    p = p>>1) {
    if (p & 0x1)
        result *= x;
    x = x*x;
}
```

```
for (Init; Test; Update )
    Body
```

Init

```
result = 1
```

Test

```
p != 0
```

Update

```
p = p >> 1
```

Body

```
{
    if (p & 0x1)
        result *= x;
    x = x*x;
}
```

"For"® "While"

For Version

```
for (Init; Test; Update
)
    Body
```

While Version

```
Init;
while (Test) {
    Body
    Update ;
}
```

Do-While Version

```
Init;
if (!Test)
    goto done;
do {
    Body
    Update ;
} while (Test)
done:
```

Goto Version

```
Init;
if (!Test)
    goto done;
loop:
    Body
    Update ;
    if (Test)
        goto loop;
done:
```

"For" Loop Compilation

Goto Version

```
Init;  
if (!Test)  
    goto done;  
loop:  
    Body  
    Update ;  
    if (Test)  
        goto loop;  
done:
```



```
result = 1;  
if (p == 0)  
    goto done;  
loop:  
    if (p & 0x1)  
        result *= x;  
    x = x*x;  
    p = p >> 1;  
    if (p != 0)  
        goto loop;  
done:
```

Init

```
result = 1
```

Test

```
p != 0
```

Update

```
p = p >> 1
```

Body

```
{  
    if (p & 0x1)  
        result *= x;  
    x = x*x;  
}
```

Switch Statements

Implementation Options

```
typedef enum
{ADD, MULT, MINUS, DIV, MOD, BAD}
  op_type;

char unparse_symbol(op_type op)
{
  switch (op) {
  case ADD :
    return '+';
  case MULT:
    return '*';
  case MINUS:
    return '-';
  case DIV:
    return '/';
  case MOD:
    return '%';
  case BAD:
    return '?';
  }
}
```

- **Series of conditionals**
 - Good if few cases
 - Slow if many
- **Jump Table**
 - Lookup branch target
 - Avoids conditionals
 - Possible when cases are small integer constants
- **GCC**
 - Picks one based on case structure
- **Bug in example code**
 - No default given

Jump Table Structure

Switch Form

```
switch(op) {  
  case 0:  
    Block 0  
  case 1:  
    Block 1  
    . . .  
  case n-1:  
    Block n-1  
}
```

Jump Table

jtab:

Targ0
Targ1
Targ2
•
•
•
Targn-1

Jump Targets

Targ0:

Code Block 0

Targ1:

Code Block 1

Targ2:

Code Block 2

•

•

•

Targn-1:

Code Block n-1

Approx. Translation

```
target = JTab[op];  
goto *target;
```

Switch Statement Example

Branching Possibilities

```
typedef enum
  {ADD, MULT, MINUS, DIV, MOD,
  BAD}
  op_type;

char unparse_symbol(op_type op)
{
  switch (op) {
    • • •
  }
}
```

Enumerated Values

ADD	0
MULT	1
MINUS	2
DIV	3
MOD	4
BAD	5

Setup:

<code>pushl %ebp</code>	<code># Setup</code>
<code>movl %esp,%ebp</code>	<code># Setup</code>
<code>movl 8(%ebp),%eax</code>	<code># eax = op</code>
<code>cmpl \$5,%eax</code>	<code># Compare op : 5</code>
<code>ja .L64</code>	<code># If > goto done</code>
<code>jmp *.L72(,%eax,4)</code>	<code># goto Table[op]</code>

Assembly Setup Explanation

Symbolic Labels

- Labels of form `.LXX` translated into addresses by assembler

Table Structure

- Each target requires 4 bytes
- Base address at `.L72`

Jumping

```
jmp .L64
```

- Jump target is denoted by label `.L64`

```
jmp *.L72(,%eax,4)
```

- Start of jump table denoted by label `.L72`
- Register `%eax` holds `op`
- Must scale by factor of 4 to get offset into table
- Fetch target from effective Address `.L72 + op*4`

Jump Table

Targets & Completion

Table Contents

```
.L72:  
    .long .L66 #Op = 0  
    .long .L67 #Op = 1  
    .long .L68 #Op = 2  
    .long .L69 #Op = 3  
    .long .L70 #Op = 4  
    .long .L71 #Op = 5
```

Enumerated Values

ADD	0
MULT	1
MINUS	2
DIV	3
MOD	4
BAD	5

```
.L66:  
    movl $43,%eax # '+'  
    jmp .L64  
.L67:  
    movl $42,%eax # '*'  
    jmp .L64  
.L68:  
    movl $45,%eax # '-'  
    jmp .L64  
.L69:  
    movl $47,%eax # '/'  
    jmp .L64  
.L70:  
    movl $37,%eax # '%'  
    jmp .L64  
.L71:  
    movl $63,%eax # '?'  
    # Fall Through to .L64
```

Switch Statement Completion

```
.L64:                # Done:
    movl %ebp,%esp   # Finish
    popl %ebp       # Finish
    ret             # Finish
```

Puzzle

- What value returned when `op` is invalid?

Answer

- Register `%eax` set to `op` at beginning of procedure
- This becomes the returned value

Advantage of Jump Table

- Can do k -way branch in $O(1)$ operations

Object Code

Setup

- Label `.L64` becomes address `0x80487b5`
- Label `.L72` becomes address `0x8048770`

```
804875d: 89 e5          movl    %esp,%ebp
804875f: 8b 45 08       movl    0x8(%ebp),%eax
8048762: 83 f8 05       cmpl   $0x5,%eax
8048765: 77 4e         ja     80487b5
<unparse_symbol+0x59>
8048767: ff 24 85 70 87 jmp     *0x8048770(,%eax,4)
```

Object Code (cont.)

Jump Table

- Disassembler tries to interpret byte sequence as instructions
- Very strange results!

```
804876c: 04 08
804876e: 89 f6          movl    %esi,%esi
8048770: 88 87 04 08 90  movb    %al,0x87900804(%edi)
8048775: 87
8048776: 04 08          addb    $0x8,%al
8048778: 98            cwtl
8048779: 87 04 08       xchgl   %eax,(%eax,%ecx,1)
804877c: a0 87 04 08 a8  movb    0xa8080487,%al
8048781: 87 04 08       xchgl   %eax,(%eax,%ecx,1)
8048784: b0 87          movb    $0x87,%al
8048786: 04 08          addb    $0x8,%al
```

Decoding Jump Table

Known

- Starts at 8048770
- 4 bytes / entry
- Little Endian byte ordering

804876c: 04 08
804876e: 89 f6
8048770: 88 87 04 08 90
8048775: 87
8048776: 04 08
8048778: 98
8048779: 87 04 08
804877c: a0 87 04 08 a8
8048781: 87 04 08
8048784: b0 87
8048786: 04 08

Address	Entry
8048770:	08048788
8048774:	08048790
8048778:	08048798
804877c:	080487a0
8048780:	080487a8
8048784:	080487b0

Alternate Decoding Technique

Use GDB

```
gdb code-examples
```

```
(gdb) x/6xw 0x8048770
```

- Examine 6 hexadecimal format "words" (4-bytes each)
- Use command "**help x**" to get format documentation

```
0x8048770 <unparse_symbol+20>:
```

```
0x08048788
```

```
0x08048790
```

```
0x08048798
```

```
0x080487a0
```

```
0x8048780 <unparse_symbol+36>:
```

```
0x080487a8
```

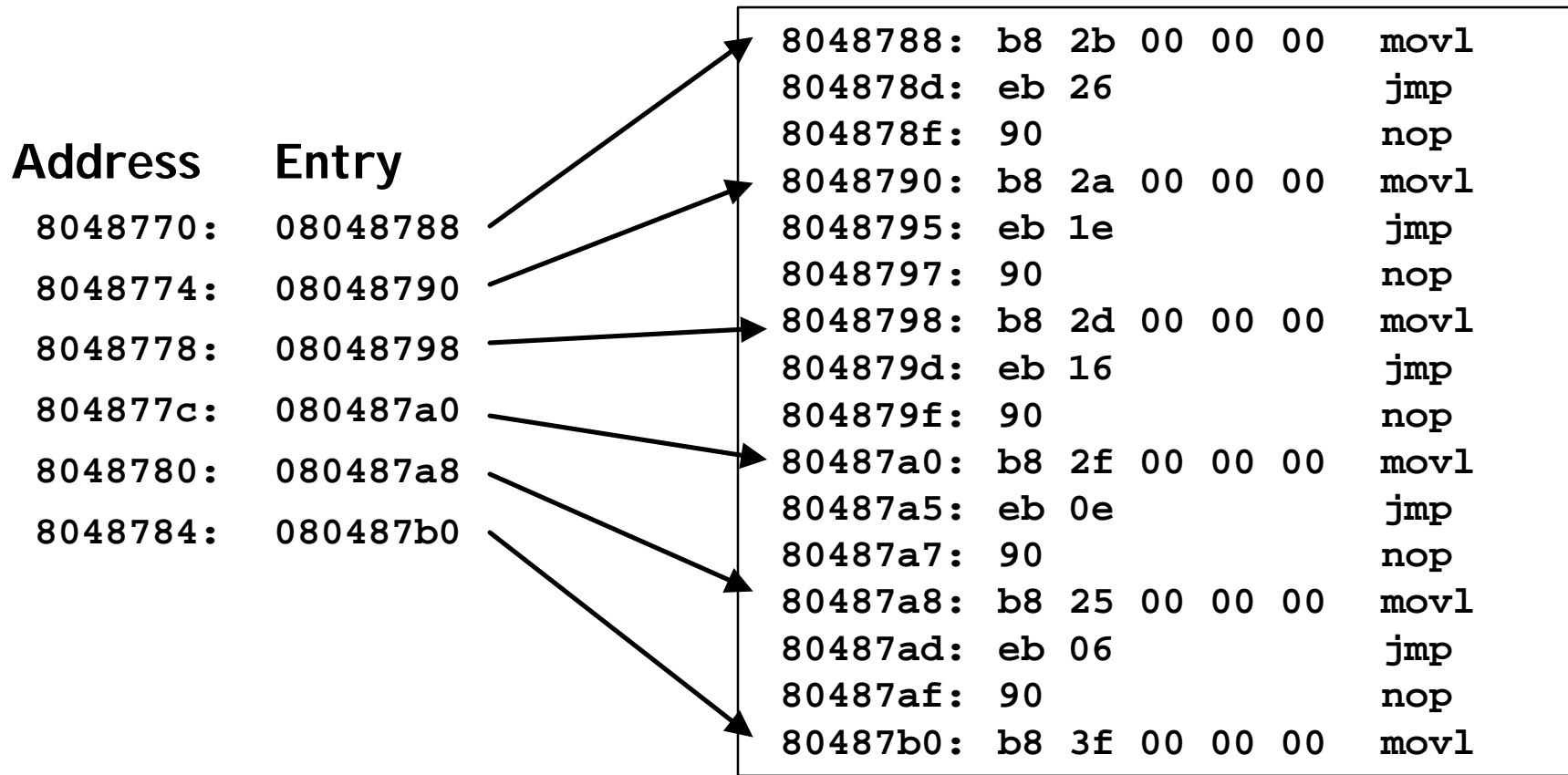
```
0x080487b0
```

Disassembled Targets

- No-operations (nop) inserted to align target addresses

```
8048788: b8 2b 00 00 00  movl  $0x2b,%eax
804878d: eb 26           jmp   80487b5 <unparse_symbol+0x59>
804878f: 90            nop
8048790: b8 2a 00 00 00  movl  $0x2a,%eax
8048795: eb 1e           jmp   80487b5 <unparse_symbol+0x59>
8048797: 90            nop
8048798: b8 2d 00 00 00  movl  $0x2d,%eax
804879d: eb 16           jmp   80487b5 <unparse_symbol+0x59>
804879f: 90            nop
80487a0: b8 2f 00 00 00  movl  $0x2f,%eax
80487a5: eb 0e           jmp   80487b5 <unparse_symbol+0x59>
80487a7: 90            nop
80487a8: b8 25 00 00 00  movl  $0x25,%eax
80487ad: eb 06           jmp   80487b5 <unparse_symbol+0x59>
80487af: 90            nop
80487b0: b8 3f 00 00 00  movl  $0x3f,%eax
```


Matching Disassembled Targets



Summary

C Control

- if-then-else
- do-while
- while
- switch

Assembler Control

- jump
- Conditional jump

Compiler

- Must generate assembly code to implement more complex control

Standard Techniques

- All loops converted to do-while form
- Large switch statements use jump tables

Conditions in CISC

- CISC machines generally have condition code registers

Conditions in RISC

- Use general registers to store condition information
- Special comparison instructions
- E.g., on Alpha:

```
cmple $16,1,$1
```

- Sets register \$1 to 1 when Register \$16 <= 1