

# 15-745 Lecture 2

## Programming in SUIF

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## Background (SUIF1)

- **SUIF**: Stanford University Intermediate Format
- **Goal**: Develop a **research** compiler infrastructure
  - Easy to modify
  - Easy to augment
  - (subgoal) Enforce modularity
  - Speed not an issue
- **Reality**:
  - Sometimes easy, sometimes hard
    - Extending the IR was VERY hard!
  - Even for research, speed is an issue

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## SUIF2

- **Goals**:
  - Better support for modifying the IR
  - Faster
- **Reflective IR**
- Tools for **adding to IR**
- Support **modules & performance**
  - i.e., no disk accesses between passes

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## The SUIF System

```

    graph TD
      PGI[PGI Fortran] --> OSUIF[OSUIF]
      EDG_C[EDG C] --> OSUIF
      EDG_Cplusplus[EDG C++] --> OSUIF
      Java --> OSUIF
      OSUIF -- "*" --> SUIF2[SUIF2]
      SUIF2 --> C[C]
      SUIF2 --> MachSUIF[MachSUIF]
      MachSUIF --> Alpha[Alpha]
      MachSUIF --> x86[x86]
      SUIF2 --- IA[Interprocedural Analysis]
      SUIF2 --- P[Parallelization]
      SUIF2 --- LO[Locality Opt]
      MachSUIF --- SO[Scalar opt]
      MachSUIF --- IS[Inst. Scheduling]
      MachSUIF --- RA[Register Allocation]
    
```

\* C++ OSUIF to SUIF is incomplete From Lam, PLDI workshop

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## MACHSUIF

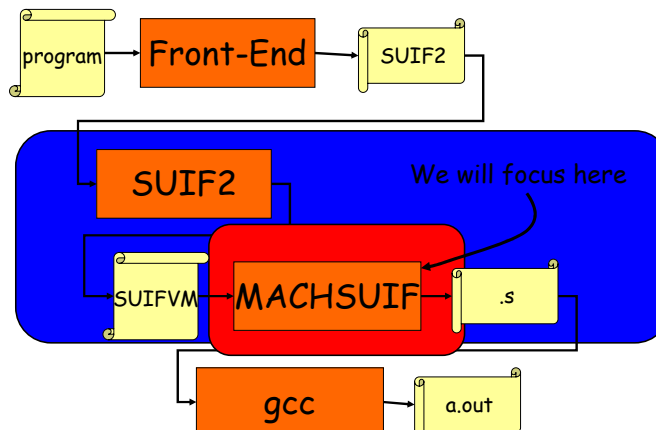
- Machine SUIF developed at Harvard
- A backend for SUIF2
- Goals:
  - Machine specific optimizations
  - Support architecture research
  - Easy to extend
  - Support multiple targets
  - 1-to-1 correspondence
  - Portable across different compilation environments

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## Compiler Phases

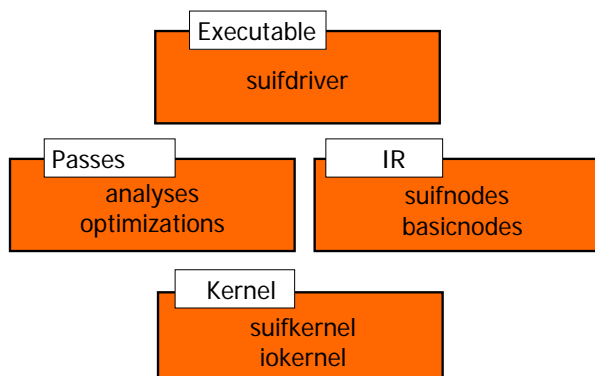


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## SUIF2 - Modular



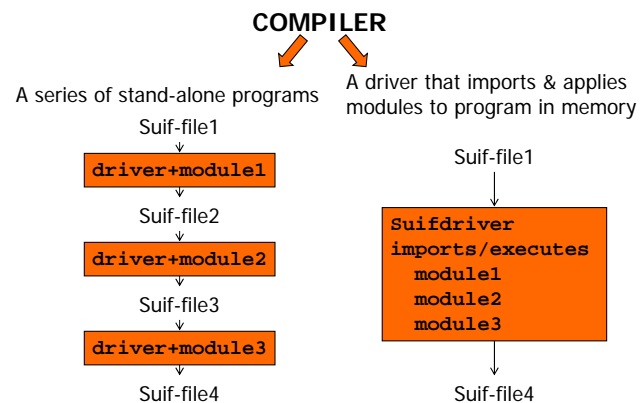
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## SUIF2 - 2 Methods



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## IR Structure

- **FileSetBlock**: the main container for files and external symbol tables
  - **FileBlock**: represents a source file: file scope symbol tables
    - **ProcedureDefinition**

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## Logistics

- Can run on **Linux (2.2)** or **Alpha**
- To get setup you should execute:
  - `cs745=/afs/cs/academic/classes/15745-f03`
  - `eval ` $cs745/public/bin/setup-suif-env -sh ``
- Now a ton of environment vars are set
- You can cross-compile
  - which may be useful when debugging new passes

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## Logistics

- Running the Front-end/Generating SUIF2 IR
  - **Linux**: `c2s foo.c`
  - **Alpha**: `c2sby1 foo.c`
 } Creates `foo.suif`
  - Run this on the target machine for cpp
- Converting **High-SUIF** to **Low-SUIF**
  - `do_lower foo.suif foo.lsf`
- Converting to **MACHSUIF IR/SUIFVM**
  - `do_svm foo.lsf foo.svm`
- ...
- Eventually,
  - `gcc -o foo foo.s`

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## Different Targets

- Different target machines have
  - different **opcodes**
  - different **reg files, conventions**, etc.
- **MACHSUIF** approach
  - early code generation
  - parameterize passes
  - `suifvm`
- **Logistics**
  - `do_gen -target_lib X`
  - `MACHSUIF_TARGET_LIB=X`
  - X can currently be: **alpha** or **x86**

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## Different Compiler Environments

- Support different compiler substrates without having to rewrite code
  - SUIF **static** compilation
  - DECO **dynamic** compilation
- Optimization Programming Interface (OPI)
  - Defines an interface that can manipulate IR in a target-machine specific manner.
  - Defines containers:
    - lists of instructions
    - control flow graph
  - Relies on substrate for I/O, etc.

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## Existing Passes

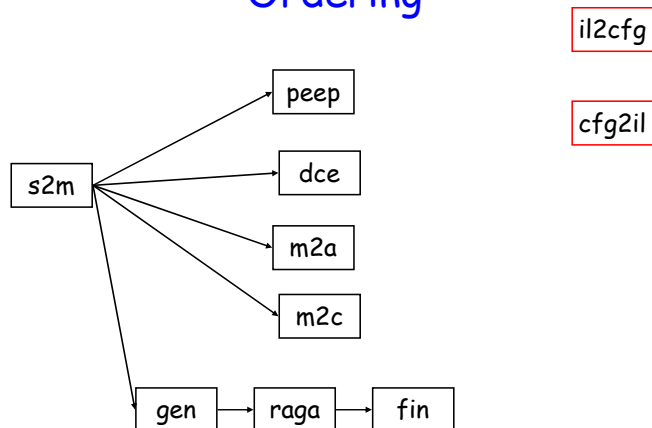
- **s2m**: convert suif to suifvm
- **gen**: convert suifvm to target dialect
- **raga**: register allocation
- **dce**: dead code elimination
- **fin**: finalize (frame layout), proc entry/exit
- **il2cfg**: instruction lists to cfg
- **cfg2il**: cfg to instruction lists
- **m2a**: create .s file
- **m2c**: create .c file

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## Ordering



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## Using OPI to Create a Pass

- **Two parts** to every pass:
  - Substrate independent part
  - Wrapper
- **Independent part** performs the optimization
- **Wrapper**
  - Binds the pass to the target (if necessary)
  - Links pass to substrate

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## Independent Part

- Each substrate-independent part is a class
- Class defines **at least three methods**:
  - **initialize**: code run **before** the pass
  - **do\_opt\_unit**: perform the pass
  - **finalize**: code to run **afterwards**
- Utilizes OPI functions and libraries
  - **CFG**: control flow graphs
  - **CFA**: control flow analysis
  - **BVD**: bit vector based dataflow

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## Wrapper

- Wraps the stand-alone class in a subclass of **PipelineablePass**
- **PipelineablePass** deals with:
  - command line parsing
  - file I/O
  - iterating on files, procedures, etc.
- If the pass is **parameterized** (i.e., must treat different targets differently)
  - call **focus()** for each **FileBlock** and **ProcedureDefinition**

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```

void
PeepSuifPass::do_file_set_block(FileSetBlock *fsb) {
    set_opi_predefined_types(fsb);
}

PeepSuifPass::do_file_block(FileBlock *fb) {
    claim(has_note(fb, k_target_lib),
          "expected target_lib annotation on file block");

    focus(fb);
    peep.initialize();
}

PeepSuifPass::do_procedure_definition(ProcedureDefinition *pd) {
    focus(pd);
    peep.do_opt_unit(pd);
    defocus(pd);
}

PeepSuifPass::finalize() {
    peep.finalize();
}

```

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## Defining Pass Peep

```

class Peep {
public:
    Peep() { }

    void initialize();
    void do_opt_unit(OptUnit*);
    void finalize();

    ...
}

```

OptUnit (in SUIF environment) is a  
ProcedureDefinition

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## OPI Data Structures

- Reference or Value Semantics?
  - It is made explicit
    - (Be safe and do explicit copies)
- Basics:
  - constants, symbols, and types
- Instructions
- Operands
- Containers
- Annotations

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## Instructions

- `Instr*` is type of every instruction.
- All contain an `Opcode`.
  - Opcode value is target specific
- Classes of instructions:
  - `Active`
    - `alm` (arith, logical, memory)
    - `cti` (control transfer)
  - `Inactive`
    - `label`
    - `dot` (pseudo-ops)
- Constructors (`new_instr_alm`)
- Predicates (`is_alm`)

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## Operands

- Instructions can contain operands of type `Opnd`.
- All Operands
  - have a type (`get_type(opnd)`)
  - a kind (`get_kind(opnd)`, `is_reg(opnd)`)
  - can be copied (`clone(opnd)`)
  - support `==` and `!=`
  - can be hashed (`hash(opnd)`)
  - can be printed (`fprint(opnd)`)
- Operands come in many flavors:
 

- Null	<code>is_null</code>
- Variable symbols	<code>is_var</code>
- Registers	<code>is_reg</code>
- Immediate	<code>is_immed</code>
- Address	<code>is_addr_sym</code>
- Address expression	<code>is_addr_exp</code>

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## Register Operands

- Hard or Virtual
- Virtual Register:
  - Creation: `opnd_reg(type)`
  - Testing: `is_virtual_reg(opnd)`
- Hard Register
  - Creation: `opnd_reg(num, type)`
  - Testing: `is_hard_ref(opnd)`
  - Getting Number: `get_reg(r)`

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## Variable Symbols

- `get_var(v)` returns the `VarSym*` for this operand
- `VarSym*` contains information about the symbol:
  - scope
  - definition
  - type
  - whether address is taken, etc.

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## Example

```
bool
is_mortal(Opnd opnd)
{
    if (is_reg(opnd))
        return true;
    if (is_var(opnd)) {
        VarSym *vs = get_var(opnd);
        return (!is_addr_taken(vs) && is_auto(vs));
    }
    return false;
}
```

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## Sequences

- `InstrList`, `CfgNode` (contain instructions)
- `Instr` (contains operands)
- Many functions based on position and handles
  - `_size`
  - `_start`, `_last`, `_end`
  - `get_x(container, int)`
  - `get_x(container, handle)`
  - `++`, `--`
  - `prepend`, `append`, ...

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## CFG

- `nodes_size(cfg)`
- `nodes_start(cfg)`
- `nodes_last(cfg)`
- `nodes_end(cfg)`
- `get_node(cfg, integer)`
- `get_node(cfg, handle)`
- `get_entry_node(cfg)`
- `get_exit_node(cfg)`
- Adding nodes is handled separately
  - e.g., `insert_empty_node(cfg, tail, head)`

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## Iterating Example

```
void
Peep::do_opt_unit(OptUnit *unit)
{
    claim(is_kind_of<Cfg>(get_body(unit)),
          "Body is not in CFG form");
    unit_cfg = static_cast<Cfg*>(get_body(unit));
    for (int i = 0; i < nodes_size(unit_cfg); ++i) {
        CfgNode *b = get_node(unit_cfg, i);
        ...
    }
}
```

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## What's important about a CFG Node?

- Instructions it contains
  - e.g., `instrs_size(node)`
- Predecessors and Successors
  - e.g., `preds(node)`
- How it ends?
  - `get_cti_handle(node)`
  - `ends_in_ubr(node)`

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## Iterating Example

```
void
Peep::do_opt_unit(OptUnit *unit)
{
    claim(is_kind_of<Cfg>(get_body(cur_unit)),
          "Body is not in CFG form");
    unit_cfg = static_cast<Cfg*>(get_body(unit));
    for (int i = 0; i < nodes_size(unit_cfg); ++i) {
        CfgNode *b = get_node(unit_cfg, i);

        Instr *inst;
        InstrHandle ih = last(b);
        int bsize = size(b);
        for (int i = 0; i < bsize; ) {
            inst = *ih;
            ++i, --ih;      // scan instructions backwards
        }
    }
}
```

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## Some Predefined Optimizations

- `remove_unreachable_nodes`
- `optimize_jumps`

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## Annotations

- Allow any information to be added to (almost) any node.
- They are **persistent**
- **key, value** pair
  - key is of type **NoteKey**
- Kinds of annotations:
  - **flag** annotations
    - existence is what matters; no key
  - **singleton** annotations
    - value is of type: **long, Integer, IdString, IrObject\***  
(**IrObject** are all ptr types. NOT: **opnd!**)
  - **list** annotations, **custom** annotations

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## There Is More

- Read: Overview and OPI users guide.
- Do HW1

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