

Lecture 8

Induction Variables and Strength Reduction

- I Overview of optimization
- II Algorithm to find induction variables

Reference: Muchnick 14.1

Example

```
FOR i = 0 to 100  
  A[i] = 0;
```

```
  i = 0  
L2: IF i>=100 GOTO L1  
    t1 = 4 * i  
    t2 = &A + t1  
    *t2 = 0  
    i = i+1  
    GOTO L2  
L1:
```

Definitions

1. A basic induction variable is a variable X
 - whose only definitions within the loop are assignments of the form $X = X+c$ or $X = X-c$, where c is either a constant or a loop-invariant variable.
2. An induction variable is
 - a basic induction variable
 - a variable defined once within the loop, whose value is a linear function of some basic induction variable at the time of the definition.
 $A = c_1 * B + c_2$
3. The FAMILY of a basic induction variable B
 - the set of induction variables A such that each time A is assigned in the loop, the value of A is a linear function of B .

Optimizations

1. Strength reduction:

Let A be an induction variable in family of basic induction variable B
($A = c_1 * B + c_2$)

- Create new variable: A'
- Initialization in preheader: $A' = c_1 * B + c_2 ;$
- Track value of B : add after $B=B+x$: $A' = A' + x * c_1 ;$
- Replace assignment to A : $A = A'$

Optimizations (cont.)

2. Optimizing non-basic induction variables

- copy propagation
- dead code elimination

3. Optimizing basic induction variables

Eliminate basic induction variables used only for

- calculating other induction variables and loop tests

Algorithm

- Select an induction variable A in the family of B, preferably with simple constants ($A = c_1 * B + c_2$).
- Replace a comparison such as
if B > X goto L1
by
if ($A' > c_1 X + c_2$) goto L1, assuming c_1 is positive
- if B is live at any exit from the loop, recompute it from A'
 - After the exit, $B = (A' - c_2) / c_1$

II. Basic Induction Variables

• A BASIC induction variable in a loop L

- a variable X whose only definitions within L are assignments of the form $X = X+c$ or $X = X-c$, where c is either a constant or a loop-invariant variable.

• **Algorithm:** can be detected by scanning L

• Example:

```
k = 0;
for (i = 0; i < n; i++) {
    k = k + 3;
    ... = m
    if (x < y)
        k = k + 4;
    if (a < b)
        m = 2 * k
    k = k - 2
    ... = m
```

Each iteration may execute a different number of increments/decrements!!

Strength Reduction Algorithm

- **Key idea**
 - For each induction variable A, ($A = c_1B + c_2$ at time of definition)
 - variable A' holds expression $c_1B + c_2$ at all times
 - replace definition of A with $A = A'$ only when executed
- **Result**
 - Program is correct
 - Definition of A does not need to refer to B

Finding Induction Variable Families

- **Let B be a basic induction variable**
 - Find all induction variables A in family of B:
 - $A = c_1 * B + c_2$
(where B refers to the value of B at time of definition)
- **Conditions**
 - If A has a single assignment in the loop L,
and assignment is one of:

$$A = B * c \quad A = c * B$$
$$A = B / c \quad (\text{assuming } A \text{ is real})$$
$$A = B + c \quad A = c + B$$
$$A = B - c$$
$$A = c - B$$

 - OR, ... (next page)

Finding Induction Variable Families (cont)

- Let D be an induction variable in the family of B
($D = c_1 * B + c_2$)
 - If A has a single assignment in the loop L,
and assignment is one of:
 $A = D * c$ $A = c * D$
 $A = D / c$ (assuming A is real)
 $A = D + c$ $A = c + D$
 $A = D - c$
 $A = c - D$
 - No definition of D outside L reaches the assignment to A
 - Between the lone point of assignment to D in L
and the assignment to A,
there are no definitions of B

Conclusions

- **Precise definitions of induction variables**
- **Systematic identification of induction variables**
- **Strength reduction**
- **Clean up:**
 - eliminating basic induction variables
 - used in other induction variable calculations
 - replacement of loop tests
 - eliminating other induction variables
 - standard optimizations