

Computation and Deduction

Lecture 15: Natural Deduction

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1. Natural Deduction
2. Local Reduction
3. Congruences

Predicate Logic

```
i : type. % individuals
o : type. % formulas
%name i T S
%name o A B C

and      : o -> o -> o. %infix right 11 and
imp      : o -> o -> o. %infix right 10 imp
or       : o -> o -> o. %infix right 11 or
not      : o -> o.      %prefix 12 not
true     : o.
false    : o.
forall   : (i -> o) -> o.
exists   : (i -> o) -> o.
```

Natural Deduction I

nd : o -> type. % deductions

%name nd D E

andi : nd A -> nd B -> nd (A and B).

andel : nd (A and B) -> nd A.

ander : nd (A and B) -> nd B.

impi : (nd A -> nd B) -> nd (A imp B).

impe : nd (A imp B) -> nd A -> nd B.

oril : nd A -> nd (A or B).

orir : nd B -> nd (A or B).

ore : nd (A or B) -> (nd A -> nd C) -> (nd B -> nd C)
-> nd C.

noti : ({p:o} nd A -> nd p) -> nd (not A).

note : nd (not A) -> {C:o} nd A -> nd C.

Natural Deduction II

`truei` : `nd (true)`.

`falsee` : `nd (false) -> nd C`.

`foralli` : `{a:i} nd (A a) -> nd (forall A)`.

`foralle` : `nd (forall A) -> {T:i} nd (A T)`.

`existsi` : `{T:i} nd (A T) -> nd (exists A)`.

`existse` : `nd (exists A) -> ({a:i} nd (A a) -> nd C) -> nd C`.

Local Reductions I

```
==>L : nd A -> nd A -> type.  %name ==>L L  
%infix none 14 ==>L
```

```
redl_andl   : (andel (andi D E)) ==>L D.
```

```
redl_andr   : (ander (andi D E)) ==>L E.
```

```
redl_imp    : (impe (impi D) E) ==>L (D E).
```

Local Reductions II

`redl_orl` : `(ore (oril D) E1 E2) ==>L (E1 D)`.

`redl_orr` : `(ore (orir D) E1 E2) ==>L (E2 D)`.

`redl_not` : `(note (noti D) C E) ==>L (D C E)`.

`redl_forall` : `(foralle (foralli D) T) ==>L (D T)`.

`redl_exists` : `(existse (existsi T D) E) ==>L (E T D)`.

Congruences, Conjunction

```
==> : nd A -> nd A -> type.           %name ==> R
```

```
%infix none 14 ==>
```

```
red_local      : D ==>L D'
                -> D ==> D'.
```

```
% Conjunction
```

```
red_andi1      : D1 ==> D1'
                -> (andi D1 D2) ==> (andi D1' D2).
```

```
red_andi2      : D2 ==> D2'
                -> (andi D1 D2) ==> (andi D1 D2').
```

```
red_ande1      : D ==> D'
                -> (and e1 D) ==> (and e1 D').
```

```
red_ande2      : D ==> D'
                -> (and e2 D) ==> (and e2 D').
```

Congruences, Implication

% Implication

red_impi : ($\{u:\text{nd } A\} u \implies u \rightarrow (D u) \implies (D' u)$)
 $\rightarrow (\text{impi } D) \implies (\text{impi } D')$.

red_impe1 : $D1 \implies D1'$
 $\rightarrow (\text{impe } D1 D2) \implies (\text{impe } D1' D2)$.

red_impe2 : $D2 \implies D2'$
 $\rightarrow (\text{impe } D1 D2) \implies (\text{impe } D1 D2')$.

Congruences, Quantification

% Universal Quantification

red_foralli : ($\{a:i\} (D a) \implies (D' a)$)
 \rightarrow (foralli D) \implies (foralli D').

red_foralle : D \implies D'
 \rightarrow (foralle D T) \implies (foralle D' T).

% Existential Quantification

red_existsi : D \implies D'
 \rightarrow (existsi T D) \implies (existsi T D').

red_existse1 : D1 \implies D1'
 \rightarrow (existse D1 D2) \implies (existse D1' D2).

red_existse2 : ($\{a:i\} \{u:nd (A a)\} (D2 a u) \implies (D2' a u)$)
 \rightarrow (existse D1 D2) \implies (existse D1 D2').