## INSTRUCTIONS

- Due: Monday, 29 January 2024 at 10:00 PM ET.
- Format: Write your answers in the yoursolution.tex file and compile a pdf (preferred) or you can type directly on the blank pdf. Make sure that your answers are within the dedicated regions for each question/part. If you do not follow this format, we may deduct points. You may use digital tools (e.g., an iPad) to handwrite your solutions, but make sure they are legible. We reserve the right to take points off if we can't read your solution. We will NOT accept handwritten solutions of any kind.
- Images: To insert pictures, we recommend drawing it on PowerPoint or Google Drawings, saving it as an image and including it in your latex source.
- How to submit: Submit a pdf with your answers on Gradescope. Log in and click on our class 15-281 and click on the submission titled HW2 and upload your pdf containing your answers.
- Policy: See the course website for homework policies and Academic Integrity.

| Name |  |
| :--- | :--- |
| Andrew ID |  |
| Hours to complete? |  |
|  |  |

For staff use only

| Q 1 | Q 2 | Q 3 | Q 4 | Q 5 | Q 6 | Total |
| ---: | ---: | ---: | ---: | ---: | ---: | ---: |
| $/ 19$ | $/ 22$ | $/ 15$ | $/ 24$ | $/ 17$ | $/ 3$ | $/ 100$ |

## Q1. [19 pts] Search and Heuristics

Recall the following pseudocode for graph search presented in lecture and recitation, and consider the graph below.

```
def graph_search(problem):
    initialize the explored set to be empty
    initialize frontier as specific work list # (stack, queue, priority queue)
    add intial state of problem to frontier
    while frontier not empty:
        choose a node and remove it from frontier
        if node contains goal state:
            return the corresponding solution
        add node to explored set
        for each resulting child from node:
            if child state not in frontier and child state not in explored set:
                add child to the frontier
        return failure
```


(a) [8 pts] Consider a depth-first graph search on the graph, where S is the start and G is the goal state. Assume that ties are broken alphabetically (so a partial plan $\mathrm{S} \rightarrow \mathrm{X} \rightarrow \mathrm{A}$ would be expanded before $\mathrm{S} \rightarrow \mathrm{X}$ $\rightarrow \mathrm{B}$ and $\mathrm{S} \rightarrow \mathrm{A} \rightarrow \mathrm{Z}$ would be expanded before $\mathrm{S} \rightarrow \mathrm{B} \rightarrow \mathrm{A})$.
(i) [4 pts] What final path would DFS return?

Answer:
(ii) $[4 \mathrm{pts}]$ Please list the explored nodes in the order they were explored.
Answer:

The pseudocode and graph have been reproduced for your convenience:

```
def graph_search(problem):
    initialize the explored set to be empty
    initialize frontier as specific work list # (stack, queue, priority queue)
    add intial state of problem to frontier
    while frontier not empty:
        choose a node and remove it from frontier
        if node contains goal state:
            return the corresponding solution
        add node to explored set
        for each resulting child from node:
            if child state not in frontier and child state not in explored set:
                add child to the frontier
        return failure
```


(b) [8 pts] Now consider a breadth-first graph search on the graph, where S is the start and G is the goal state. Assume the same tie-breaking scheme as the question above.
(i) [4 pts] What final path would BFS return?

Answer:
(ii) [4 pts] Please list the explored nodes in the order they were explored.
Answer:
(c) [3 pts] What final path would iterative deepening return?
Answer:

## Q2. [22 pts] Search and Heuristics



Imagine a car-like agent wishes to exit a maze like the one shown above. The agent is directional and at all times faces some direction $d \in(N, S, E, W)$. With a single action, the agent can either move forward at an adjustable velocity $v$ or turn.

The moving actions are faster, maintain and slower. For these actions, the agent then moves a number of squares equal to its new adjusted velocity. Let $v$ denote the agent's current velocity and let $v^{\prime}$ denote the agent's new adjusted velocity. Please note the agent cannot take the actions slower or maintain if its velocity is 0 .

- Faster: $v^{\prime}=v+1$
- Slower: $v^{\prime}=v-1$
- Maintain: $v^{\prime}=v$

The turning actions are left and right, which change the agent's direction by 90 degrees. Turning is only permitted when the velocity is zero (and leaves it at zero).

- Left: change the agent's direction by 90 degrees counterclockwise
- Right: change the agent's direction by 90 degrees clockwise

For example, if the agent is currently on $(0,0)$ facing north with velocity 0 (as pictured) and wants to get to $(2,0)$ facing east with velocity 0 , the sequence of actions will be: right, faster, maintain, slower.

Illegal actions include

- Any action that would result in a collision with a wall
- Any action that would reduce $v$ below 0 or above a maximum speed $V_{\max }$
- Maintaining a velocity of 0

The agent's goal is to find a plan which parks it $(v=0)$ in the goal direction on the exit square using as few actions (time steps) as possible. Note that the cost of a path is defined by the number of actions the agent takes.
(a) [2 pts] Suppose the agent wants to take the leftmost path (i.e., the one that passes through $(1,2)$ ) from the start $(0,0)$ facing north to the goal $(0,4)$ facing west. Write down the shortest sequence of actions for it to take.
$\square$
(b) [2 pts] If the grid is M by N and the maximum speed is $V_{\max }$, what is the size of the state space? You should assume that all configurations are reachable from the start state.

## State Space Size:

(c) [2 pts] A "child" of a state $s$ is any other state $s^{\prime}$ reachable via a legal action from state $s$. Is it possible that a state in the state space has no children? If so, give an example of such a state. If not, briefly explain why every state must have at least one child.
$\bigcirc$ YesNo

## Example State or Explanation:

(d) [4 pts] What is the maximum branching factor of this problem? Draw an example state ( x , y , orientation, velocity, grid/walls) that has this branching factor, and list the set of available actions. For example, in the above picture, if the agent was in $(0,0)$ facing North with a velocity of $v=0$, the branching factor would be 2. The agent could turn left or right (but not go faster since it would hit a wall).

You may assume that illegal actions are simply not returned by the problem model and therefore not counted in the branching factor. You do not necessarily have to use the example grid above. If you need to include a drawing of your own, label properly and make sure it fits in the solution box.

## Maximum Branching Factor:

## Maximum Branching Example State and Available Actions:

(e) $[4 \mathrm{pts}]$ Is the Manhattan distance from the agent's location to the exit's location admissible?

If not, draw an example state ( $\mathrm{x}, \mathrm{y}$, orientation, velocity, grid/walls) where this heuristic overestimates at that state, and specify: 1) the heuristic value at that state and 2) the actual cost from that state to the goal.

You may assume that illegal actions are simply not returned by the problem model. You do not necessarily have to use the example grid above. Make sure to label your drawing, including the goal state and action sequence, and fit it into the solution box.

$$
\bigcirc \text { Yes } \quad \bigcirc \text { No }
$$

## Example State, Heuristic Value, Actual Cost:

(f) [4 pts] Is the following heuristic admissible? Manhattan distance / $V_{\max }$.

If not, draw an example state ( $x, y$, orientation, velocity, grid/walls) where this heuristic overestimates at that state, and specify: 1) the heuristic value at that state and 2) the actual cost from that state to the goal.
You may assume that illegal actions are simply not returned by the problem model. You do not necessarily have to use the example grid above. Make sure to label your drawing, including the goal state and action sequence, and fit it into the solution box.
$\bigcirc$ Yes $\bigcirc$ No

## Example State, Heuristic Value, Actual Cost:

(g) [1 pt] If we used an inadmissible heuristic in A* Tree search, could it change the completeness of the search? Assume the graph is finite and the heuristic is non-negative.YesNo
(h) [1 pt] If we used an inadmissible heuristic in A* Tree search, could it change the optimality of the search? Assume the graph is finite and the heuristic is non-negative.
$\bigcirc$ Yes $\bigcirc$ No
(i) [2 pts] What is a general advantage that an inadmissible heuristic might have over an admissible one. Select all that apply.An inadmissible heuristic may be easier to compute, leading to a faster state heuristic computation time.
$\square$ An inadmissible heuristic can be a closer estimate to the actual cost (even if it's an overestimate) than an admissible heuristic, thus exploring fewer nodes.An inadmissible heuristic will still find optimal paths when the actual costs are non-negative.

An inadmissible heuristic may be used to completely block off searching part of a graph in a search algorithm.

## Q3. [15 pts] Search Nodes

Consider the tree search (i.e. no explored set) of an arbitrary search problem with max branching factor $b$. Each search node $n$ has a backward (cumulative) cost of $g(n)$, an admissible heuristic of $h(n)$, and a depth of $d(n)$. Let $n_{c}$ be a minimum-cost goal node, and let $n_{s}$ be a shallowest goal node.

For each of the following, give an expression that characterizes the set of nodes that are explored before the search terminates. For instance, if we asked for the set of nodes with positive heuristic value, you could say: for all $n$, such that $h(n) \geq 0$. Don't worry about ties (so you won't need to worry about $>$ versus $\geq$ ). If there are no nodes for which the expression is true, you must write "none."

Note that you are not required to use all the functions given, $f, g$, and $h$, in your inequality.
(a) [5 pts] Give an inequality in terms of the functions $g$, $h$, and $d$, as well as the nodes $n_{c}$ and $n_{s}$ defined above to describe the nodes $n$ that are explored in a breadth-first search before terminating.

## Inequality: All $n$, such that:

(b) [5 pts] Give an inequality in terms of the functions $g$, $h$, and $d$, as well as the nodes $n_{c}$ and $n_{s}$ defined above to describe the nodes $n$ that are explored in a uniform cost search before terminating.
$\square$
(c) [5 pts] Now for this question, assume the heuristic $h$ is consistent. Give an inequality in terms of the functions $g$, $h$, and $d$, as well as the nodes $n_{c}$ and $n_{s}$ defined above to describe the nodes $n$ that are explored in an $\mathbf{A}^{*}$ search before terminating.

## Inequality: All $n$, such that:

## Q4. [24 pts] Searching a Graph



| Node | $h_{1}$ | $h_{2}$ |
| :--- | :--- | :--- |
| A | 12 | 11 |
| B | 6 | 7 |
| C | 9 | 6 |
| D | 3 | 4 |
| E | 3 | 5 |
| F | 2 | 1 |
| G | 0 | 0 |

Consider the graph shown above. A is the start state and G is the goal state. The costs for each edge are shown on the graph. The graph is bi-directional so each edge can be traversed from either direction. Please refer to the search algorithms exactly as presented on the lecture slides as the ordering of the actions matters.
(a) [15 pts] For each of the following graph search strategies, mark with an X which (if any) of the listed paths it could return. Note that for some search strategies the specific path returned might depend on tie-breaking behavior. In any such cases, make sure to mark all paths that could be returned under alphabetical tie-breaking schemes. If a graph search strategy returns a path not listed, write out the correct path in the Other column.

| Algorithm | A-C-E-G | A-C-E-F-G | A-B-D-E-F-G | Other |
| :--- | :--- | :--- | :--- | :--- |
| UCS | (i) | (ii) | (iii) | (iv) |
| Greedy with heuristic $h_{1}$ | (v) | (vi) | (vii) | (viii) |
| Greedy with heuristic $h_{2}$ | (ix) | (x) | (xi) | (xii) |
| A* with heuristic $h_{1}$ | (xiii) | (xiv) | (xv) | (xvi) |
| A* with heuristic $h_{2}$ | (xvii) | (xviii) | (xix) | (xx) |

(b) [1 pt] What is the cost of the optimal path for uniform cost search from A to G?

Answer:
(c) $[4 \mathrm{pts}]$ Is $h_{1}$ admissible? Is it consistent?

Admissible:
$\bigcirc \mathrm{Yes}$
○ No
Consistent: $\bigcirc$ Yes $\bigcirc$ No
(d) [4 pts] Is $h_{2}$ admissible? Is it consistent?

Admissible:
$\bigcirc$ Yes
$\bigcirc$ No
Consistent: 〇 Yes 〇 No

## Q5. [17 pts] Search: Multiple Choice and Short Answer Questions

(a) [12 pts] Consider the following true/false questions with each question worth 2 points. For the following search problems, assume every action has a cost of at least $\epsilon$, with $\epsilon>0$. Assume any heuristics used are consistent.

Depth-first tree-search on a finite graph is guaranteed to be complete.
False

Breadth-first tree-search on a finite graph is guaranteed to be complete.

## OTrue

FalseIterative deepening tree-search on a finite graph is guaranteed to be complete.TrueFalse

For all graphs without cycles, graph-search contains a larger frontier than tree-search.TrueFalse

Iterative deepening graph-search has the time complexity of BFS and the space complexity of DFS.
$\qquad$ TrueFalse

If $h_{1}(\mathrm{~s})$ is a consistent heuristic and $h_{2}(\mathrm{~s})$ is a consistent heuristic, then $\min \left(h_{1}(\mathrm{~s}), h_{2}(\mathrm{~s})\right)$ must be consistent. $\bigcirc$ True $\bigcirc$ False
(b) [5 pts] Consider the state space graph shown below. S is the start state and G is the goal state. The costs for each edge are shown on the graph. For the following table below, fill in potential heuristic values such that the heuristic is admissible but not consistent.


| Heuristic Function |  |
| :--- | :--- |
| State | $h(s)$ |
| $S$ |  |
| $A$ |  |
| $C$ | 0 |
| $G$ |  |

## Q6. [3 pts] Ethics

(a) [3 pts] Please review this article discussed in recitation: This Navigation App Keeps You Out of High-crime Areas in Real Time, https://www.androidauthority.com/redzone-navigation-app-avoids-high-crime-areas-686894/.
(i) [1 pt] Let's assume this navigation app uses A* search with an admissible and consistent heuristic to find and return paths. Suppose you are working on this app, and your boss proposes a new heuristic. There is some crime rate metric that ranges from 0 to 5 , where 5 is high crime. Your boss proposes that you add this crime rate metric to the application's current heuristic, such that the heuristic is still admissible and consistent.
Will the path necessarily change with the new heuristic? If so, will the new path be optimal? If not, under what conditions would the path change?
Answer:
(ii) [1 pt] If for some reason you choose not to adjust the heuristic, what is one way you could modify the graph structure or edge costs so that your implementation could return paths that avoid high crime areas? Give your answer in one sentence.
Answer:
(iii) [1 pt] In what ways could people and stores residing in high crime neighborhoods be impacted by changes in navigation algorithms, particularly focusing on the economic and/or societal impact? Please provide 2-3 sentences.

## Answer:

