

$$\mathcal{F}_0 \subseteq \mathcal{F}_1 \subseteq \dots \subseteq \mathcal{F}_n = \Sigma_1 \quad (\Omega, \Sigma, P)$$

now X_i is measurable in \mathcal{F}_i (X is adapted to \mathcal{F}).

and sps. $E[X_i | \mathcal{F}_{i-1}] = X_{i-1}$

↑ take X_i and average over regions of \mathcal{F}_{i-1}

then (X_0, X_1, \dots) is a martingale.

by induction $E[X_t] = \text{~~XXXX~~} = E[X_0]$. typically set this to 0.

Eg. random walk. on line.

Example: $X_{i+1} = X_i + Y_{i+1}$ ← indep random var. (0 mean).

but could also be $X_{i+1} = X_i + Y_{i+1}/X_i$ ← depends on past, but $E[Y_{i+1} | X_1, \dots, X_i] = 0$.
Conditionally still 0 mean.

Eg. random walk with absorbing borders.

Stopping time measurable wrt filtration.

sps. $ET < \infty$ and all step sizes are bounded a.s.

then $E[X_T] = E[X_0]$. (Doob's Optional Stopping Thm)

E.g.: $E[S_T] = E[S_0] = 0$ in random walk. $S_t = X_1 + X_2 + \dots + X_t$

so. $p_a(-a) + p_b(b) = 0 \Rightarrow p_a = \frac{b}{a+b}$.

but also. $S_t^2 - t$ is a martingale for the increments being 0 mean var 1.

$\Rightarrow E[S_t^2] = E[t] = \frac{b}{a+b} a^2 + \frac{a}{a+b} b^2 = ab \Rightarrow n^2$ time to hit barriers @ $\pm n$.

Ville's thm: if S_n is a supermartingale, nonnegative then $Pr(\sup_n S_n > x) \leq \frac{E[S_0]}{x}$.

Like Markov but now for supermartingales.

And set supremum for free.

Concentration bounds (Azuma-Hoeffding)

(2)

$0 = X_0, X_1, X_2, \dots$ martingale with $|X_i - X_{i-1}| \leq C_i$.

then $P[|X_n| \geq \lambda] \leq \exp(-\lambda^2/2\sigma^2)$ where $\sigma^2 = \sum C_i^2$.

Andriy bides like Ville's

$P[\max_{s \leq t} |X_s| \geq \lambda] \leq \exp(-\lambda^2/2\sigma^2)$ as well.

Example: sps. $f(\bar{x}) \in \mathbb{R}$ $\bar{x} = (x_1, x_2, \dots, x_n)$. X_i indep r.v.s.

then def $X_0 = E[f]$

$$X_1 = E[f | X_1]$$

⋮

$$X_n = E[f | X_1, X_2, \dots, X_n] = f.$$

Moreover: sps. changing coordinates changes f by $\leq C_i \Rightarrow |X_i - X_{i-1}| \leq C_i$ ← verify!!

Now apply Azuma: $P[|f - E[f]| \geq \lambda] \leq \exp(-\lambda^2/2\sigma^2)$. [McDiarmid]

Example: chromatic number of random graph $G_{n,p}$ is concentrated on $\leq O(\sqrt{n})$ values. [Shannon - Spencer]

[Bollobas: at most 4 values.]

[Alon Krivelevich, Achlioptas and Hass: 2 values] $\uparrow p = d/n$.

$\uparrow p = n^{1/2 - \epsilon}$

Btw: instead of requiring that $|X_i - X_{i-1}| \leq C_i$ as.

we could use that $[X_i | X_{i-1}]$ has ~~exponential tails~~ subgaussian tails.

and hence can extend to some of

subgaussian r.v.s.

Discrepancy

$$S = S_1, S_2, \dots, S_m \subseteq \Omega$$

$$\chi: \Omega \rightarrow \{-1, 1\}$$

$$\text{disc}_\chi(S) = \sum_{i \in S} \chi(i)$$

$$\text{disc}_\chi(S) = \max_{S \subseteq \Omega} \chi(S)$$

Want χ to get low discrepancy $\text{disc}_\chi(S)$.

Thm: \exists set-systems such that $\text{disc}_\chi(S) \geq \Omega(\sqrt{n})$.
 $m = O(n)$ sets n elts

take Hadamard matrix $H_1 = \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix}$ $H_k = \begin{pmatrix} H_{k-1} & H_{k-1} \\ H_{k-1} & -H_{k-1} \end{pmatrix}$.

now $m = 2 \cdot 2^k$ sets for each row i , $A_i = +$ entries $\bar{A}_i = -$ entries.

$$\text{then } \chi(A_i) - \chi(\bar{A}_i) = \langle \chi, h_i \rangle$$

$$\Rightarrow |\langle \chi, h_i \rangle| \leq |\chi(A_i)| + |\chi(\bar{A}_i)| \quad (*)$$

But h 's form an orthogonal basis, with vectors of length $\|h_i\|^2 = 2^k$.

$$\Rightarrow \sum_i (\langle \chi, h_i \rangle)^2 = \sum_i \frac{\langle \chi, h_i \rangle^2}{2^k} = \|\chi\|^2 = 2^k$$

$$\Rightarrow \exists i \text{ st } \langle \chi, h_i \rangle^2 \geq 2^{k/2}$$

$$\Rightarrow |\langle \chi, h_i \rangle| \geq 2^{k/4}$$

using (*) get $\exists i$, either A_i or \bar{A}_i satisfy $\chi(S) \geq \frac{1}{2} \sqrt{n}$.

Spencer Showed: set-system with $m = O(n)$ sets,
 Factory with $\text{disc}_x(\mathcal{S}) \leq O(\sqrt{n})$.

Recent proofs using CS-type ideas (based on making the results algorithmic)

Bansal, Lovett-Meka, Rothvoss, ~~El~~ ^{El} El Karoui-Singh, Bansal-Ladha-Vempala, ...

Give the Lovett-Meka proof

Want $x(S_i) \leq c\sqrt{n \log(m/n)} \quad \forall i$

$x_i \in \{-1, 1\}$.

↗ relax to $x_i \in [-1, 1]$.

☹ $x_i = 0$ is a solution with discrepancy zero!!

So want to round.

try 1: independent rounding:- $P_0[\text{disc}(S_i) \geq c\sqrt{n \log m}] \leq 1/\text{poly}(m)$
 then union bound.

But: get $\text{disc} \sqrt{n \log m}$ for m sets, and hence $\sqrt{n \log m}$ for n sets
 "penalty for union bound"

Smarter idea: small steps, and fix problems as they arise

lets make this geometric:-

$x(S_j) \leq c\sqrt{n \log(m/n)}$

$\rightarrow \langle a_j, x \rangle \leq \|a_j\| \cdot c_i \quad \forall i \in [m]$

$x_j \in [-1, 1]$

$\forall j \in [n]$

↖ scale $a_i \in \mathbb{R}^n$ st. $\|a_i\| = 1$
 (wlog)

↑ think of this as

$\Theta(\sqrt{\log(m/n)})$

start @ x_0 (think as 0).

Call constraint a_i dangerous if $\langle a_i, x \rangle \geq c_i - \delta$
 variable x_j frozen if $|x_j| \geq 1 - \delta$. } *apx satisfied*

Each time t Find subspace orthogonal to all dangerous constraints
 & all frozen variables.

until subspace dimension smaller than $n/4$.
 V_t (re. will have $3n/4$ of these constraints/vars).

Idea: want to ensure that many of these are frozen vars.
 not ~~high~~ dangerous constraints.

[Pick $g_t \sim N(V_t)$ gaussian from that subspace.] $\delta \geq \delta \cdot \sqrt{\log \frac{mn}{\delta}}$
 $X_0^0 = x_0, \quad X_{\frac{t}{\delta}}^t \leftarrow X_{\frac{t}{\delta}}^{t-1} + \delta g_t$ *↑ ting*
 martingale.

Claim 1: $\dim(V_{t+\delta}) \leq \dim(V_t)$.

Claim 2: if we are not dangerous/frozen, it will not be violated w.p.

why? for a gaussian $P(|g| \geq \lambda) \leq 2 \exp(-\lambda^2/2)$
 $\text{now } |\delta g| > \delta \leq 2 \exp(-(\delta/\delta)^2/2) = 2 \exp(-\frac{\log(mn/\delta)}{2})$
 $= (\frac{\delta}{mn})^{o(1)}$

\Rightarrow unimbounded over all steps.

Claim 3: $E[\|x^t\|^2] = E[\|x^{t-1}\|^2] + \underbrace{2E[x^t, \delta g^t]}_{0 \text{ indep}} + \underbrace{\delta^2 E[\|g^t\|^2]}_{\delta^2 \dim(V_t)}$

$= \sum_{i=1}^t \dim(V_i) \Rightarrow$ as long as $\dim(V_i) \geq n/4$, progress.