

OK, now to HDXs.

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$$X = (U, \mathcal{F})$$

- Simplicial Complex = downward closed set system  $\leftarrow$  eg. matroid ind sets  
(SC)

- dim of a set  $S$  =  $|S|-1$ .  $\nearrow$  "face"
- dim of SC = ~~dim~~ dim of largest set in it.  $\nearrow$  "facet"
- pure SC = all facets have same size (also as in matroid).

- So matroids are pure SCs with exchange prop

$$X(i) = \text{sets of size } i+1. \quad (\dim i)$$

$$X = X(0) \cup X(1) \cup \dots \cup X(d)$$

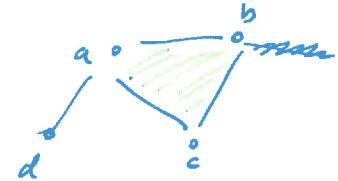
"vertices"  $\uparrow$  "edges"

"facets"

$$\begin{matrix} & ab & bc & ac & ad \\ a & & b & c & d \\ & \emptyset & \leftarrow \dim -1. \end{matrix}$$

- 1-skeleton of  $X$  = graph  $(X(0), X(1))$

called  $G(X)$



- Links of  $X$ .

given  $\alpha \in \mathbb{F}$ ,  $X_\alpha = \{\beta \setminus \alpha \mid \beta \supseteq \alpha, \beta \in \mathcal{F}\}$ .  $\leftarrow$  "contraction" of  $\alpha$

If  $X$  is a pure  $d$ -dim SC  $\Rightarrow X_\alpha$  is pure  $(d-1)$ -dim SC.

Local Spectral Expanders:  $X$  is pure 1-dim SC.  $\leftarrow$   $X$  is  $\delta$ -(local)-spectral-expander

If  $\ell$  faces  $\alpha$  of  $\dim \leq 2$ ,  $G(X_\alpha)$  has  $\alpha_2$  if the random walk matrix  $\leq \delta$ .

Moreover we can define a probability distribution over facets, and over each layer, as follows:-

$X$  be a pure SC of dim  $d$ . Let  $\pi_d$  be a probability over  $X(d)$ .

Now can define a distribution  $X(d-1)$  by setting for each  $\alpha \in X(d-1)$  :-

$$\pi_{d-1}(\alpha) = \frac{1}{\text{size of sets } \beta \text{ in } X(d)} \sum_{\substack{\beta \in X(d) \\ \beta \supseteq \alpha}} \pi_d(\beta).$$

this gives a distribution  $X(d-1)$ . Now induct.

$$\longrightarrow X \longrightarrow$$

Given  $(X, \pi_d)$  ~~be~~ d-dim pure SC with weights. For  $\alpha$  and  $\tau \in X_\alpha$

$$\text{get } \pi^\alpha(\tau) = P[\beta = \alpha \cup \tau \mid \beta \supseteq \alpha] = \frac{\pi(\alpha \cup \tau)}{\binom{|\alpha \cup \tau|}{|\alpha|} \cdot \pi(\alpha)}.$$

just the natural thing.

$$\longrightarrow X \longrightarrow$$

Now: gives weights on the 1-skeleton (edges and vertices) of  $X$ .

Call  $G_\alpha$  the ~~1-skeleton~~ skeleton of the link,  $\pi_\alpha^\alpha$  be distributed over edges.

$$D_\alpha = \text{diagonal matrix with } D_\alpha(x, x) = \sum_{y \in X_\alpha^{(0)}} \pi_\alpha^y(\{x, y\}) = 2\pi_\alpha^d(x).$$

Now the walk is not unweighted : the weights decide the prob of going places.

$$W_\alpha = D_\alpha^{-1} A_\alpha. \quad \text{has } \lambda_{\max} = 1, \text{ with } f_{\max} = 1.$$

[ALOV]

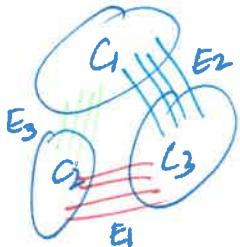
Thm: The SC of any matroid with the uniform distribution on facets is a 0-local expander.

Pf: Show that ① every link <sup>has</sup> connected skeleton, and  
② every top link has skeleton which is 0-local-expander.

① Say contract  $\alpha$ . So left with matroid. Can get from any vrtx to any other b/c of exchange axiom.

② Let's sketch the proof for graphic matroids, where  $d = n - 2$  (since bases one of size  $n - 1$ ). Now any link  $X_\alpha$  corresponds to an ~~acyclic~~ edges of a forest  $F \subseteq E$  with  $|F| = n - 3$  edges). This gives  $\Rightarrow 3$  components

the following



Hence the link has skeleton whose elements correspond to these colored edges ( $E_1, E_2, E_3$  say), and whose pairs correspond to pairs of colored edges that can belong to some base. Note these are edges in a complete bipartite graph, since we can add any pair of



It is a fact that the complete  $k$ -partite graphs have at most one positive eigenvalue (and this is a characterization of complete  $k$ -partite graphs).

We discussed this on Praza so lets give a proof (another to be added later).

Pf1: Cauchy's Interlacing thm says that if

$$A = B + \alpha \alpha^T \quad \text{then} \quad \lambda_1(A) \geq \lambda_1(B) \geq \lambda_2(A) \geq \dots$$

rank 1 (PSD)

Nw: the adjacency matrix of a complete  $k$ -partite graph is

$$A = J + \underbrace{(-B_1 - B_2 - \dots - B_k)}_B. \text{ where } B_i = \chi_{E_i} \chi_{E_i}^T \text{ and } J = \mathbb{1} \mathbb{1}^T$$
$$= B + \mathbb{1} \mathbb{1}^T.$$

Nw  $B$  is negative semidefinite so  $\lambda_{\max}(B) \leq 0$ .

Nw  $\lambda_2(A) \leq \lambda_1(B) \leq 0$ .

□

Pf2:

OK: given a distribution over facets (usually uniform), we have a natural way to get distributions over links, and over each 1-skeleton. (and so over random walks)

Each has a spectral gap (for the matrix  $D_\alpha^{-1} A_\alpha$ ).

$(X, \pi)$  is an  $\delta$ -local-spectral expander if  $\lambda_2$  for each of these matrices is at most  $\delta$ . (1st largest eigenvalue  $\lambda_1 = 1$  for each, with  $f_1 \propto \mathbf{1}$ ).

Do we need to check the expansion for all links?

Oppenheim gave a nice "frickle-down" theorem that says: if the top links are good expanders then the lower links are, too (with some loss).

Thm [Opp18]:  $(X, \pi)$  be pure SC of dim d (weighted by  $\pi$ ).

Sps.  $G_\phi = (X(0), X(1), \pi_\phi)$  is connected and  $\lambda_2(D_\phi^{-1} A_\phi) \leq \delta$  for all  $v \in X(0)$

$$\Rightarrow \lambda_2(D_\phi^{-1} A_\phi^{-1}) \leq \frac{\delta}{1-\delta}$$

Corollary: if top links  $\lambda_2(D_\alpha^{-1} A_\alpha) \leq \delta \forall \alpha \in X(d-2)$

and  $G_\beta$  is connected for all  $\beta \in X(k) \forall k \leq d-2$

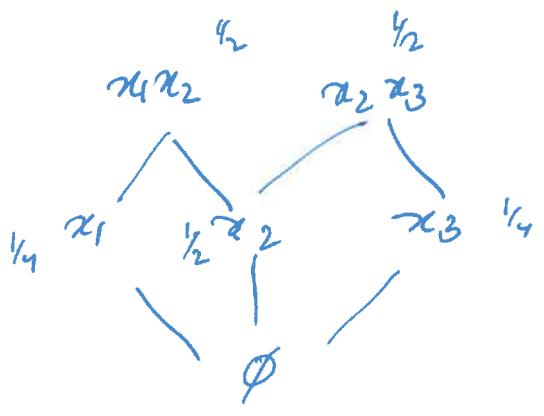
$$\text{then } \lambda_2(D_\phi^{-1} A_\phi^{-1}) \leq \frac{\delta}{1-\cancel{\delta}} \quad (\cancel{d-1})$$

Pf: omitted

But note: if each  $G_\alpha$  is a 0-local expander

$\Rightarrow G_\phi$  is also a 0-local expander!

This is what we show for  
case ~~where~~ for top links ~~of~~  
of matroidal SC.



$$A = \begin{bmatrix} x_1 & y_2 & 1/2 & 0 \\ x_2 & 1/4 & y_2 & 1/4 \\ x_3 & 0 & 1/2 & 1/2 \\ x_4 & - & - & - \end{bmatrix}$$

not symmetric

$$P_0^\Delta = \begin{bmatrix} 1 & 1/2 & 0 & 0 \\ 1/2 & 1 & 1/2 & 0 \\ 0 & 1/2 & 1 & 1/2 \\ 0 & 0 & 1/2 & 1 \end{bmatrix}$$

$$P_1^\nabla = \begin{bmatrix} 1 & 1/2 & 0 & 0 \\ 1/2 & 1 & 1/2 & 0 \\ 0 & 1/2 & 1 & 1/2 \\ 0 & 0 & 1/2 & 1 \end{bmatrix}$$

eigenvalues of  $P_0^\Delta$  =  $1 = \lambda_1, \frac{1}{2}, 0$

evecs.  $\begin{pmatrix} 1 & 1 & 1 \\ -1 & 0 & 1 \\ 1 & -1 & 1 \end{pmatrix}$

evecs.  $\begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix}$

↑ not orthonormal, since  
 $P_0^\Delta \neq P_1^\nabla$  mean  
 necc. symmetric

Lemma: Eigenvalues of  $P_{Kt}^\Delta$  and  $P_{Kt}^\nabla$  are the same.

And so we can talk about the spectral gap of a pair of layers  $X(k), X(k+t)$   
 (one for the down-up walk on  $X(k+t)$ )  
 another for up-down walk on  $X(k)$ .)

~~mixing time~~

Last time we saw that if  $G$  was a regular  $n$  vertex graph with  
 $\max(\lambda_2, |\lambda_n|) \leq \varepsilon \Rightarrow$  called  $(n, d, \varepsilon)$  graph.

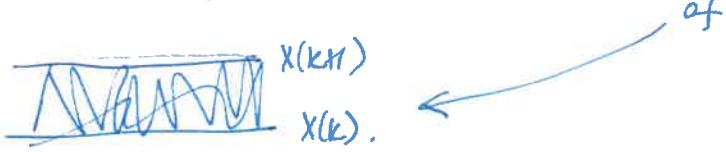
then random walk "mixes" in ~~time~~  $O\left(\frac{\log(n/\varepsilon)}{\varepsilon}\right)$  time  
 (TV distance to uniform)  
 $\leq \varepsilon$

If graph not regular  $\Rightarrow$  cannot use same proof, b/c matrices not symmetric.  
 but still walks on undirected graphs (perhaps with weights).

Now the mixing time (to the stationary distribution  $\pi$ )  $\leq \frac{1}{\varepsilon} \log\left(\frac{n}{\delta} \cdot \frac{d_{\max}}{d_{\min}}\right)$

max/min degree ratio

Hence still interested to bound the spectral radius, by  $\epsilon$ .



In fact we care about the spectral radius of  $P_d^\nabla$

But that's the same as  $P_{d-1}^\Delta$ .

Now if we could relate  $P_{d-1}^\Delta \rightarrow P_{d-1}^\nabla$  etc.... Actually relate  $P_{d-1}^\Delta \rightarrow P_{d-1}^\nabla$  lose  $\delta$

which is the proof

non lazy walk.  
↓  
and then relate  $P_{d-1}^\Delta \rightarrow P_{d-1}^\nabla$  lose a little more.  
.....

$$1 - \frac{1}{\alpha} \log \frac{1}{\delta}$$

Thm [Kaufmann-Oppenheim] if  $(X, \pi)$  is a  $\delta$ -local-spectral expander then  $\forall k \leq d$ .

$$\lambda_2(P_k^\nabla) \leq 1 - \frac{1}{k+1} + k\delta$$

$$\text{For matroidal sc } \delta=0 \Rightarrow \lambda_2(P_k^\nabla) \leq 1 - \frac{1}{k+1} \Rightarrow \lambda_2(P_d^\nabla) \leq \frac{d}{d+1}.$$

$$\Rightarrow \text{spectral gap} \leq \Theta(\frac{1}{d})$$

$$\text{Now plugging into mixing bounds: \# steps} \leq O\left(\frac{1}{\epsilon} \log\left(\frac{N}{\delta}\right)\right) \quad \begin{matrix} \leftarrow \# \text{ times} \\ \text{\# steps} \end{matrix}$$

$$\text{but } \epsilon = \frac{1}{d}$$

$$N = n^d \quad \begin{matrix} \leftarrow \# \text{ bases} \\ \text{\# bases} \end{matrix}$$

$$\Rightarrow O(d^2 \log(n/\delta)) \text{ time!}$$

$$\begin{cases} d = \text{rank of matroid}. \end{cases}$$

Intuition:

