Principles of Software Construction: Objects, Design, and Concurrency

Part 3: Concurrency

Introduction to concurrency, part 4

In the trenches of parallelism

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Administrivia

- Homework 5b due 11:59 p.m. tonight
  - Turn in by tomorrow 9 a.m. to be considered as a Best Framework
- Optional reading due today:
  - Java Concurrency in Practice, Chapter 12

NEW COURSE: LANGUAGE DESIGN & PROTOTYPING

17-396/17-696 – SPRING 2020

Little languages are everywhere! Would you like to – or do you need to – design your own?

In this course, you will:

- Learn how to **critique a language design**
- Practice several **language prototyping** approaches (interpreters, transpilers, fluent APIs)
- Apply techniques for **evaluating language designs with users**
- **Design and prototype your own language** in the final project

Prof. Jonathan Aldrich – T/Th 3-4:20

http://www.cs.cmu.edu/~aldrich/courses/17-396/

Software Engineering (SE) at CMU

- 17-214: Code-level design
  - Extensibility, reuse, concurrency, functional correctness
- 17-313: Human aspects of software development
  - Requirements, teamwork, scalability, security, scheduling, costs, risks, business models
- 17-413 Practicum, 17-415 Seminar, Internship
- Various courses on requirements, architecture, software analysis, SE for startups, API design, etc.
- SE Minor: http://isri.cmu.edu/education/undergrad

Key concepts from last Thursday

Policies for thread safety

- Thread-confined
- Shared read-only
- Shared thread-safe
  - Objects that perform internal synchronization
- Guarded
  - Objects that must be synchronized externally

Shared thread-safe

- "Thread-safe" objects that perform internal synchronization
- Build your own, or know the Java concurrency libraries

Advice for building thread-safe objects

- Do as little as possible in synchronized region: get in, get out
  - Obtain lock
  - Examine shared data
  - Transform as necessary
  - Drop the lock
- If you must do something slow, move it outside the synchronized region

Example: adding concurrency to the observer pattern

```
private final List<Observer<E>> observers = new ArrayList<>();
public void addObserver(Observer<E> observer) {
   synchronized(observers) { observers.add(observer); }
}
public boolean removeObserver(Observer<E> observer) {
   synchronized(observers) { return observers.remove(observer); }
}
private void notifyOf(E element) {
   synchronized(observers) {
       for (Observer<E> observer : observers)
           observer.notify(this, element);  // Risks liveness and
   }                                        // safety failures!
}
```

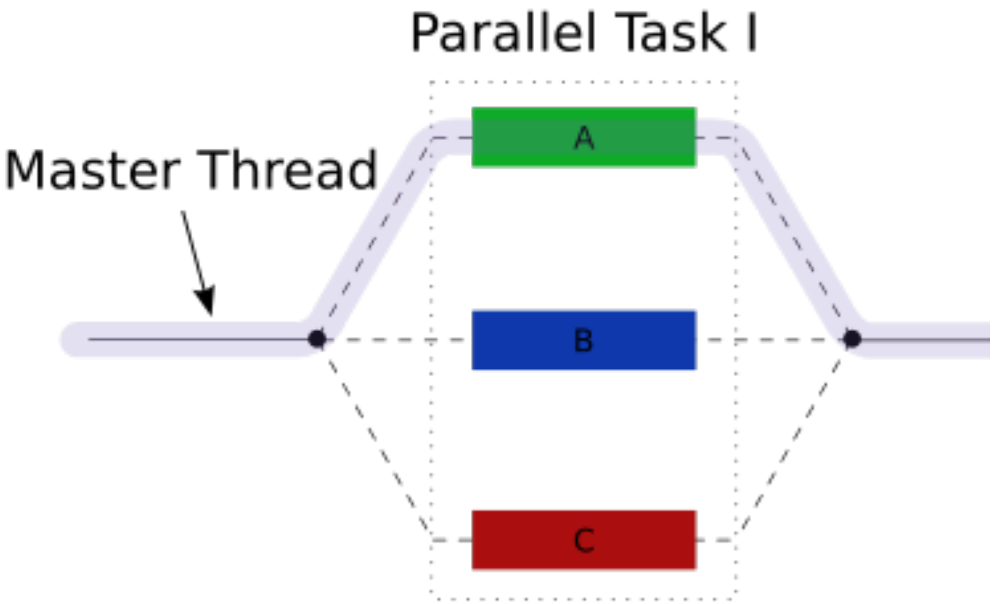
One solution: *snapshot iteration*

```
private void notifyOf(E element) {
    List<Observer<E>> snapshot = null;

    synchronized(observers) {
        snapshot = new ArrayList<>(observers);
    }

    for (Observer<E> observer : snapshot) {
        observer.notify(this, element); // Safe
    }
}
```

A better solution: CopyOnWriteArrayList

```
private final List<Observer<E>> observers =
       new CopyOnWriteArrayList<>();

public void addObserver(Observer<E> observer) {
   observers.add(observer);
}
public boolean removeObserver(Observer<E> observer) {
   return observers.remove(observer);
}
private void notifyOf(E element) {
   for (Observer<E> observer : observers)
       observer.notify(this, element);
}
```

The fork-join pattern




```
if (my portion of the work is small)
    do the work directly
else
    split my work into pieces
    invoke the pieces and wait for the results
```

A framework for asynchronous computation

- The `java.util.concurrent.Future<V>` interface

```
V       get();
V       get(long timeout, TimeUnit unit);
boolean isDone();
boolean cancel(boolean mayInterruptIfRunning);
boolean isCancelled();
```

- The `java.util.concurrent.ExecutorService` interface

```
void              execute(Runnable task);
Future            submit(Runnable task);
Future<V>         submit(Callable<V> task);
List<Future<V>>   invokeAll(Collection<Callable<V>> tasks);
Future<V>         invokeAny(Collection<Callable<V>> tasks);
```

Today

- Concurrency in practice: In the trenches of parallelism

Concurrency at the language level

- Consider:

```
Collection<Integer> collection = …;
int sum = 0;
for (int i : collection) {
    sum += i;
}
```

- In python:

```
collection = …
sum = 0
for item in collection:
    sum += item
```

Parallel quicksort in Nesl

```
function quicksort(a) =
  if (#a < 2) then a
  else
   let pivot   = a[#a/2];
       lesser  = {e in a| e < pivot};
       equal   = {e in a| e == pivot};
       greater = {e in a| e > pivot};
       result  = {quicksort(v): v in [lesser,greater]};
   in result[0] ++ equal ++ result[1];
```

- Operations in {} occur in parallel
- 210-esque questions: What is total work? What is depth?

Prefix sums (a.k.a. inclusive scan, a.k.a. scan)

- Goal: given array `x[0…n-1]`, compute array of the sum of each prefix of x

```
[ sum(x[0…0]),
  sum(x[0…1]),
  sum(x[0…2]),
  …
  sum(x[0…n-1]) ]
```

- e.g., `x =` `[13, 9, -4, 19, -6, 2, 6, 3]`

  prefix sums: `[13, 22, 18, 37, 31, 33, 39, 42]`

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Parallel prefix sums

- Intuition: Partial sums can be efficiently combined to form much larger partial sums. E.g., if we know `sum(x[0…3])` and `sum(x[4…7])`, then we can easily compute `sum(x[0…7])`
- e.g., `x =     [13,  9, -4, 19, -6,  2,  6,  3]`

Parallel prefix sums algorithm, upsweep

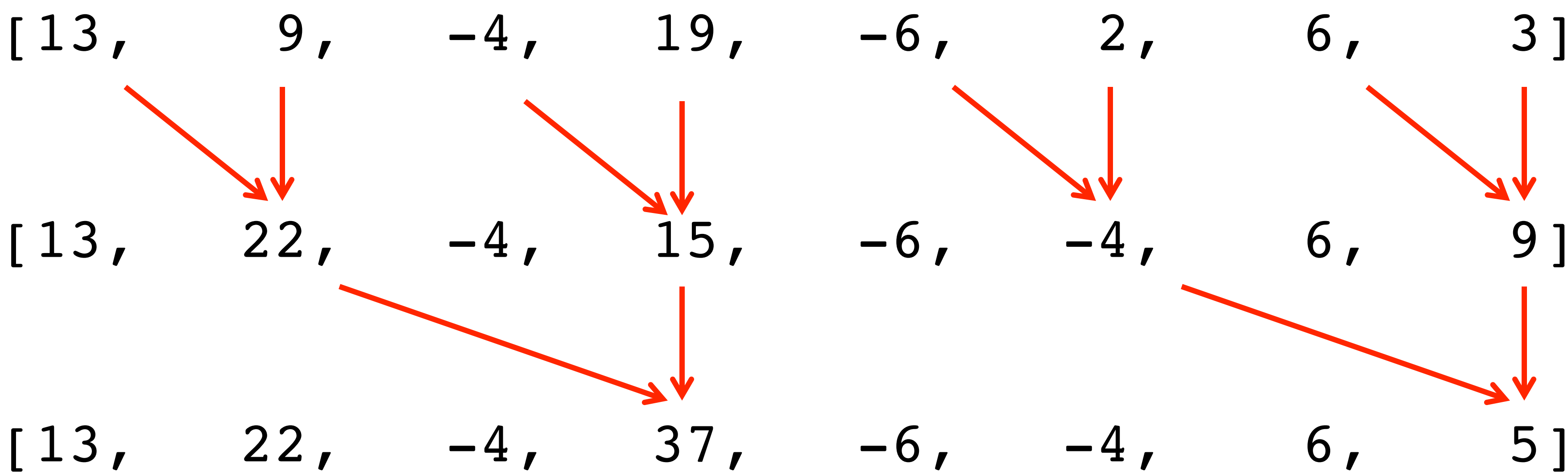
Compute the partial sums in a more useful manner

```
[13,    9,    -4,    19,    -6,    2,    6,    3]
```

```
[13,    22,    -4,    15,    -6,    -4,    6,    9]
```

Parallel prefix sums algorithm, upsweep

Compute the partial sums in a more useful manner



Parallel prefix sums algorithm, upsweep

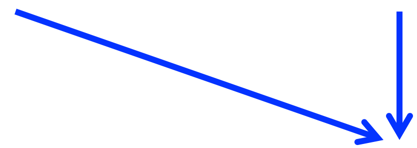
Compute the partial sums in a more useful manner

```
[13,  9, -4, 19, -6,  2, 6,  3]
[13, 22, -4, 15, -6, -4, 6,  9]
[13, 22, -4, 37, -6, -4, 6,  5]
[13, 22, -4, 37, -6, -4, 6, 42]
```

Parallel prefix sums algorithm, downsweep

Now unwind to calculate the other sums

```
[13,   22,   -4,   37,   -6,   -4,    6,   42]
```



```
[13,   22,   -4,   37,   -6,   33,    6,   42]
```

Parallel prefix sums algorithm, downsweep

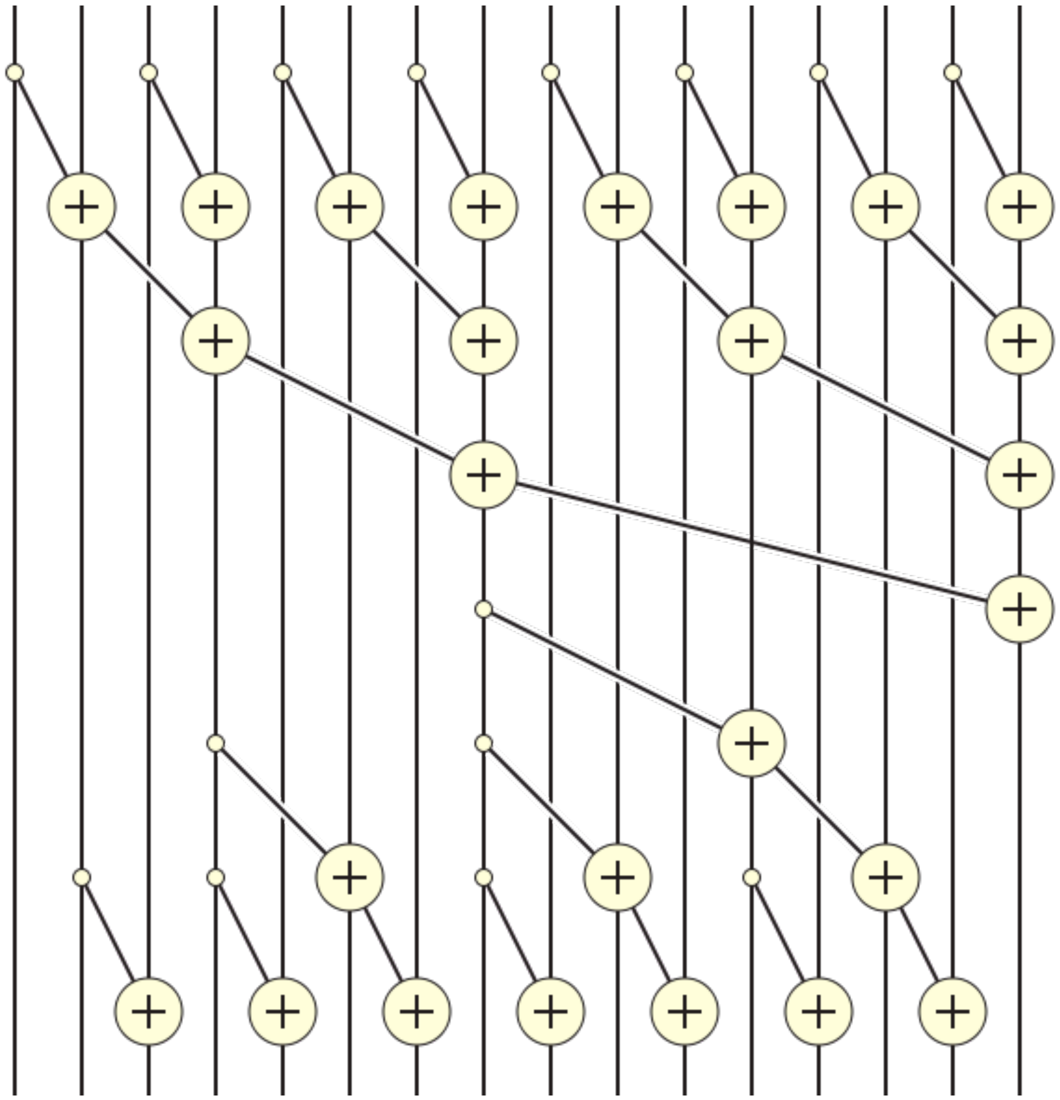
Now unwind to calculate the other sums

```
[13,   22,   -4,   37,   -6,   -4,    6,   42]

[13,   22,   -4,   37,   -6,   33,    6,   42]

[13,   22,   18,   37,   31,   33,   39,   42]
```

- Recall, we started with:

```
[13,    9,   -4,   19,   -6,    2,    6,    3]
```

Doubling array size adds two more levels



Upsweep

Downsweep

Parallel prefix sums

pseudocode

```
// Upsweep
prefix_sums(x):
  for d in 0 to (lg n)-1:          // d is depth
    parallelfor i in 2^d-1 to n-1, by 2^(d+1):
      x[i+2^d] = x[i] + x[i+2^d]

// Downsweep
for d in (lg n)-1 to 0:
  parallelfor i in 2^d-1 to n-1-2^d, by 2^(d+1):
    if (i-2^d >= 0):
      x[i] = x[i] + x[i-2^d]
```

Parallel prefix sums algorithm, in code

- An iterative Java-esque implementation:

```
void iterativePrefixSums(long[] a) {
  int gap = 1;
  for ( ; gap < a.length; gap *= 2) {
    parfor(int i=gap-1; i+gap < a.length; i += 2*gap) {
      a[i+gap] = a[i] + a[i+gap];
    }
  }
  for ( ; gap > 0; gap /= 2) {
    parfor(int i=gap-1; i < a.length; i += 2*gap) {
      a[i] = a[i] + ((i-gap >= 0) ? a[i-gap] : 0);
    }
  }
```

Parallel prefix sums algorithm, in code

- A recursive Java-esque implementation:

```
void recursivePrefixSums(long[] a, int gap) {
  if (2*gap – 1 >= a.length) {
    return;
  }

  parfor(int i=gap-1; i+gap < a.length; i += 2*gap) {
    a[i+gap] = a[i] + a[i+gap];
  }

  recursivePrefixSums(a, gap*2);

  parfor(int i=gap-1; i < a.length; i += 2*gap) {
    a[i] = a[i] + ((i-gap >= 0) ? a[i-gap] : 0);
  }
}
```

Parallel prefix sums algorithm

- How good is this?

Parallel prefix sums algorithm

- How good is this?
  - Work: O(n)
  - Depth: O(lg n)
- See `PrefixSums.java`, `PrefixSumsSequentialWithParallelWork.java`

Goal: parallelize the PrefixSums implementation

- Specifically, parallelize the parallelizable loops

```
parfor(int i = gap-1;  i+gap < a.length;  i += 2*gap) {
  a[i+gap] = a[i] + a[i+gap];
}
```

- Partition into multiple segments, run in different threads

```
for(int i = left+gap-1;  i+gap < right;  i += 2*gap) {
  a[i+gap] = a[i] + a[i+gap];
}
```

Recall from Thursday: Fork/join in Java

- The `java.util.concurrent.ForkJoinPool` class
  - Implements `ExecutorService`
  - Executes `java.util.concurrent.ForkJoinTask<V>` or `java.util.concurrent.RecursiveTask<V>` or `java.util.concurrent.RecursiveAction`
- In a long computation:
  - Fork a thread (or more) to do some work
  - Join the thread(s) to obtain the result of the work

The `RecursiveAction` abstract class

```
public class MyActionFoo extends RecursiveAction {
    public MyActionFoo(…) {
        store the data fields we need
    }

    @Override
    public void compute() {
        if (the task is small) {
            do the work here;
            return;
        }

        invokeAll(new MyActionFoo(…),  // smaller
                  new MyActionFoo(…),  // subtasks
                  …);                  // …
    }
}
```

A ForkJoin example

- See `PrefixSumsParallelForkJoin.java`
- See the processor go, go go!

Parallel prefix sums algorithm

- How good is this?
  - Work: O(n)
  - Depth: O(lg n)
- See `PrefixSumsParallelArrays.java`

Parallel prefix sums algorithm

- How good is this?
  - Work: O(n)
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- See `PrefixSumsParallelArrays.java`
- See `PrefixSumsSequential.java`

Parallel prefix sums algorithm

- How good is this?
  - Work: O(n)
  - Depth: O(lg n)
- See `PrefixSumsParallelArrays.java`
- See `PrefixSumsSequential.java`
  - n-1 additions
  - Memory access is sequential
- For `PrefixSumsSequentialWithParallelWork.java`
  - About 2n useful additions, plus extra additions for the loop indexes
  - Memory access is non-sequential
- The punchline:
  - Don't roll your own. Know the libraries
  - Cache and constants matter

In-class example for parallel prefix sums

```
[7,    5,    8,  -36,   17,    2,   21,   18]
```