Principles of Software Construction: Objects, Design, and Concurrency

Introduction to concurrency, part 4
In the trenches of parallelism

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#### Administrivia

- Homework 5b due Wednesday night
  - Turn in by 9 a.m. Thursday to be considered for Best Framework award
  - Homework 5 Best Frameworks available Friday
- Homework 5c due Tuesday, 11:59 p.m.



# Key concepts from last Thursday



#### Policies for thread safety

- 1. Thread-confined state mutate but don't share
- 2. Shared read-only state share but don't mutate
- 3. Shared thread-safe object synchronizes itself internally
- 4. Shared guarded client synchronizes object(s) externally

#### Advice for building thread-safe objects

- Do as little as possible in synchronized region: get in, get out
  - Obtain lock
  - Examine shared data
  - Transform as necessary
  - Drop the lock
- If you must do something slow, move it outside the synchronized region



#### 3. Shared thread-safe state

- Thread-safe objects that perform internal synchronization
- You can build your own, but not for the faint of heart
- You're better off using ones from java.util.concurrent
- j.u.c also provides skeletal implementations

## java.util.concurrent is BIG (1)

- 1. Atomic variables: java.util.concurrent.atomic
  - Support various atomic read-modify-write ops
- 2. Concurrent collections
  - Shared maps, sets, lists
- 3. Data exchange collections
  - Blocking queues, deques, etc.
- 4. Executor framework
  - Tasks, futures, thread pools, completion service, etc.
- 5. Synchronizers
  - Semaphores, cyclic barriers, countdown latches, etc.
- 6. Locks: java.util.concurrent.locks
  - Read-write locks, conditions, etc.



## java.util.concurrent is BIG (2)

- Pre-packaged functionality: java.util.Arrays
  - Parallel sort, parallel prefix
- Completable futures
  - Multistage asynchronous concurrent computations
- Flows
  - Publish/subscribe service
- And more
  - It just keeps growing



## Today

- j.u.c. Executor framework overview
- Concurrency in practice: In the trenches of parallelism

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#### 4. Executor framework overview

- Flexible interface-based task execution facility
- Key abstractions
  - Runnable basic task
  - Callable<T> task that returns a value (and can throw an exception)
  - Future<T> a promise to give you a T
  - Executor machine that executes tasks
  - Executor service Executor on steroids
    - Lets you manage termination
    - Can produce Future instances

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#### Executors – your one-stop shop for executor services

- Executors.newSingleThreadExecutor()
  - A single background thread
- Executors.newFixedThreadPool(int nThreads)
  - A fixed number of background threads
- Executors.newCachedThreadPool()
  - Grows in response to demand



### A very simple (but useful) executor service example

- Background execution in a long-lived worker thread
  - To start the worker thread:

```
ExecutorService executor =
   Executors.newSingleThreadExecutor();
```

— To submit a task for execution:

```
executor.execute(runnable);
```

— To terminate gracefully:

```
executor.shutdown(); // Allows tasks to finish
```

#### Other things you can do with an executor service

- Wait for a task to completeFoo foo = executorSvc.submit(callable).get();
- Wait for any or all of a collection of tasks to complete invoke{Any,All}(Collection<Callable<T>> tasks)
- Retrieve results as tasks complete
   ExecutorCompletionService
- Schedule tasks for execution at a time in the future
   ScheduledThreadPoolExecutor
- etc., ad infinitum

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#### Concurrency at the language level

• Consider: Collection<Integer> collection = ...; int sum = 0; for (int i : collection) { sum += i;• In python: collection = ... sum = 0for item in collection: sum += item

#### Parallel quicksort in Nesl

```
function quicksort(a) =
  if (#a < 2) then a
  else
  let pivot = a[#a/2];
    lesser = {e in a| e < pivot};
    equal = {e in a| e == pivot};
    greater = {e in a| e > pivot};
    result = {quicksort(v): v in [lesser,greater]};
  in result[0] ++ equal ++ result[1];
```

- Operations in {} occur in parallel
- 210-esque questions: What is total work? What is span?

### Prefix sums (a.k.a. inclusive scan, a.k.a. scan)

 Goal: given array x[0...n-1], compute array of the sum of each prefix of x

```
[ sum(x[0...0]),
   sum(x[0...1]),
   sum(x[0...2]),
   ...
  sum(x[0...n-1]) ]
```

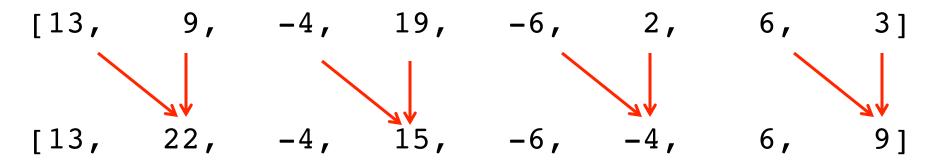
• e.g., x = [13, 9, -4, 19, -6, 2, 6, 3]prefix sums: [13, 22, 18, 37, 31, 33, 39, 42]

#### Parallel prefix sums

- Intuition: Partial sums can be efficiently combined to form much larger partial sums. E.g., if we know sum(x[0...3]) and sum(x[4...7]), then we can easily compute sum(x[0...7])
- e.g., x = [13, 9, -4, 19, -6, 2, 6, 3]

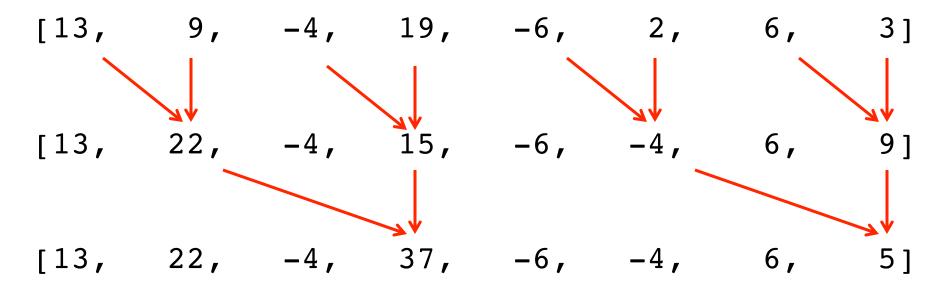
### Parallel prefix sums algorithm, upsweep

Compute the partial sums in a more useful manner



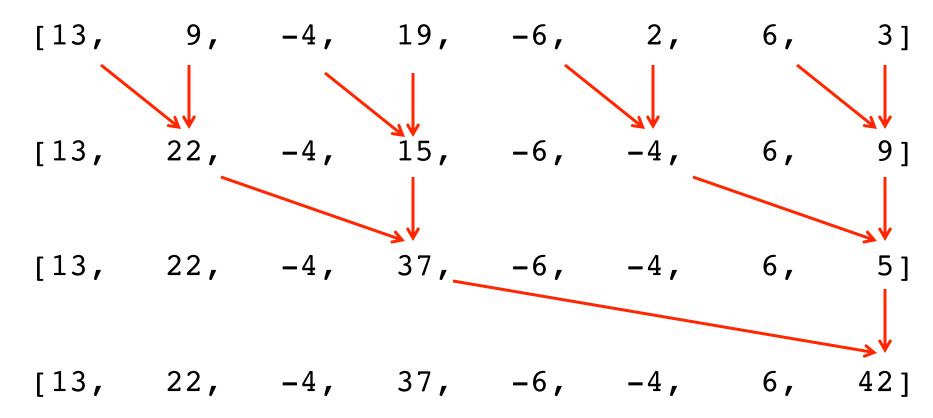
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#### Parallel prefix sums algorithm, downsweep

Now unwind to calculate the other sums

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Now unwind to calculate the other sums

• Recall, we started with:

$$[13, 9, -4, 19, -6, 2, 6, 3]$$

In-class example for parallel prefix sums

[7, 5, 8, -36, 17, 2, 21, 18]

In-class example for parallel prefix sums (downsweep)

[7, 12, 8, -16, 17, 19, 21, 42]

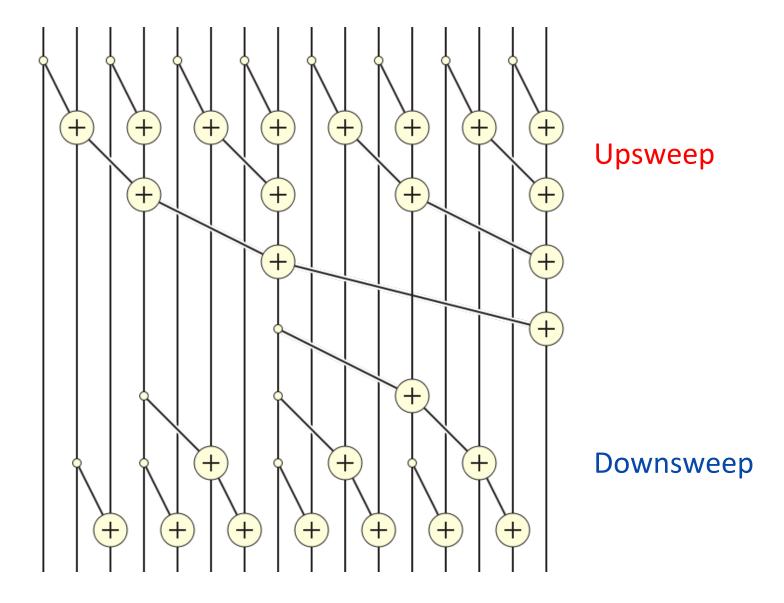
In-class example for parallel prefix sums (downsweep)

$$[7, 12, 8, -16, 17, 19, 21, 42]$$

$$[7, 12, 20, -16, 1, 3, 24, 42]$$

[7, 5, 8, -36, 17, 2, 21, 18]

## Doubling array size adds two more levels



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#### Parallel prefix sums

#### pseudocode

```
// Upsweep
prefix_sums(x):
  for d in 0 to (\lg n)-1:
                           // d is depth
    parallelfor i in 2^{d}-1 to n-1, by 2^{d+1}:
      x[i+2^d] = x[i] + x[i+2^d]
// Downsweep
for d in (\lg n)-1 to 0:
  parallelfor i in 2^{d}-1 to n-1-2^{d}, by 2^{d+1}:
    if (i-2^d >= 0):
      x[i] = x[i] + x[i-2^d]
```

#### Parallel prefix sums algorithm, in code

• An iterative Java-esque implementation:

```
void iterativePrefixSums(long[] a) {
  int gap = 1;
  for (; gap < a.length; gap *= 2) {
    parfor(int i=gap-1; i+gap < a.length; i += 2*gap) {</pre>
      a[i+gap] = a[i] + a[i+gap];
  for (; gap > 0; gap /= 2) {
    parfor(int i=gap-1; i < a.length; i += 2*gap) {</pre>
      a[i] = a[i] + ((i-gap >= 0) ? a[i-gap] : 0);
```

#### Parallel prefix sums algorithm, in code

A recursive Java-esque implementation:

```
void recursivePrefixSums(long[] a, int gap) {
  if (2*gap - 1 >= a.length) {
    return;
  parfor(int i=gap-1; i+gap < a.length; i += 2*gap) {</pre>
    a[i+gap] = a[i] + a[i+gap];
  recursivePrefixSums(a, gap*2);
  parfor(int i=gap-1; i < a.length; i += 2*gap) {</pre>
    a[i] = a[i] + ((i-gap >= 0) ? a[i-gap] : 0);
```

How good is this?



How good is this?

– Work: O(n)

– Span: O(lg n)

See PrefixSums.java,
 PrefixSumsSequentialWithParallelWork.java

### Goal: parallelize the PrefixSums implementation

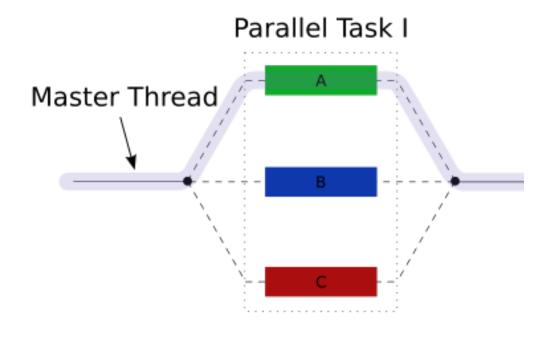
Specifically, parallelize the parallelizable loops

```
parfor(int i = gap-1; i+gap < a.length; i += 2*gap) {
   a[i+gap] = a[i] + a[i+gap];
}</pre>
```

Partition into multiple segments, run in different threads

```
for(int i = left+gap-1; i+gap < right; i += 2*gap) {
   a[i+gap] = a[i] + a[i+gap];
}</pre>
```

## The fork-join pattern



```
if (my portion of the work is small)
   do the work directly
else
   split my work into pieces
   recursively process the pieces
```

#### Fork/join in Java

- The java.util.concurrent.ForkJoinPool class
  - Implements ExecutorService
  - Executes java.util.concurrent.ForkJoinTask<V> or java.util.concurrent.RecursiveTask<V> or java.util.concurrent.RecursiveAction
- In a long computation:
  - Fork a thread (or more) to do some work
  - Join the thread(s) to obtain the result of the work

#### The RecursiveAction abstract class

```
public class MyActionFoo extends RecursiveAction {
    public MyActionFoo(...) {
        store the data fields we need
    }
    @Override
    public void compute() {
        if (the task is small) {
            do the work here;
            return;
        invokeAll(new MyActionFoo(...), // smaller
                   new MyActionFoo(...), // subtasks
                   ...);
                                         // ...
    }
```

### A ForkJoin example

- See PrefixSumsParallelForkJoin.java
- See the processor go, go go!



How good is this?

– Work: O(n)

– Span: O(lg n)

• See PrefixSumsParallelArrays.java

- How good is this?
  - Work: O(n)
  - Span: O(lg n)
- See PrefixSumsParallelArrays.java
- See PrefixSumsSequential.java

- How good is this?
  - Work: O(n)
  - Span: O(lg n)
- See PrefixSumsParallelArrays.java
- See PrefixSumsSequential.java
  - n-1 additions
  - Memory access is sequential
- For PrefixSumsSequentialWithParallelWork.java
  - About 2n useful additions, plus extra additions for the loop indexes
  - Memory access is non-sequential
- The punchline:
  - Don't roll your own. Know the libraries
  - Cache and constants matter



In-class example for parallel prefix sums

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