

## Mining Billion-node Graphs: Patterns and Tools

Christos Faloutsos

**CMU** 

(on sabbatical at google)



## Thank you!

• Prof. Sharad Mehrotra



Mark Cartnal



Carolyn Simpson



### Our goal:

Open source system for mining huge graphs:

PEGASUS project (PEta GrAph mining System)

- www.cs.cmu.edu/~pegasus
- code and papers



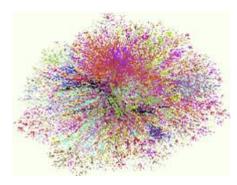
#### **Outline**



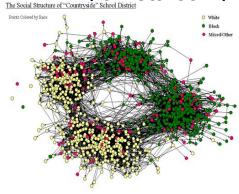
- Introduction Motivation
  - Problem#1: Patterns in graphs
  - Problem#2: Tools
  - Problem#3: Scalability
  - Conclusions



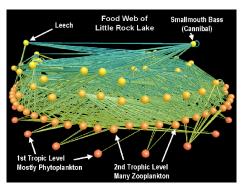
## Graphs - why should we care?



Internet Map [lumeta.com]



Friendship Network [Moody '01]



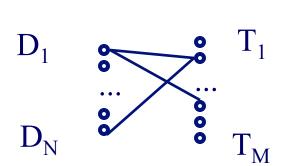
Food Web [Martinez '91]

- Social networks
  - (orkut, linkedIn ...)
- twitter



## Graphs - why should we care?

• IR: bi-partite graphs (doc-terms)



web: hyper-text graph

• ... and more:



## Graphs - why should we care?

- 'viral' marketing
- web-log ('blog') news propagation
- computer network security: email/IP traffic and anomaly detection

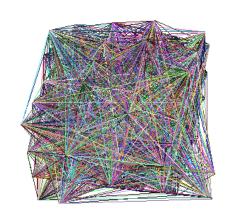
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  - Static graphs
  - Weighted graphs
  - Time evolving graphs
  - Problem#2: Tools
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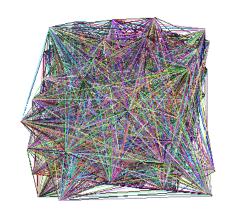
# Problem #1 - network and graph mining



- What does the Internet look like?
- What does FaceBook look like?
  - What is 'normal'/'abnormal'?
- which patterns/laws hold?

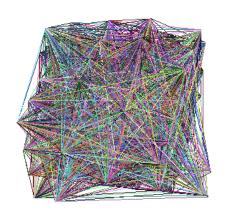


# Problem #1 - network and graph mining



- What does the Internet look like?
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  - To spot anomalies (rarities), we have to discover patterns

# Problem #1 - network and graph mining



- What does the Internet look like?
- What does FaceBook look like?
- What is 'normal'/'abnormal'?
- which patterns/laws hold?
  - To spot anomalies (rarities), we have to discover patterns
  - Large datasets reveal patterns/anomalies that may be invisible otherwise...

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## Graph mining

• Are real graphs random?



## Laws and patterns

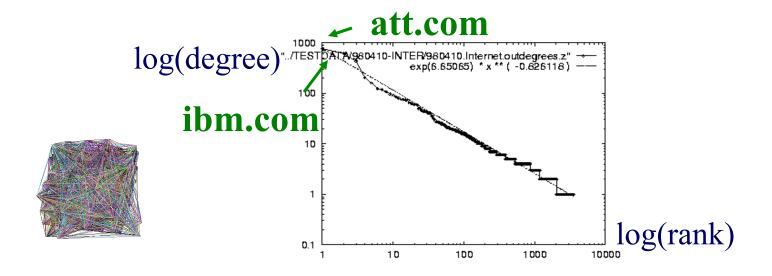
- Are real graphs random?
- A: NO!!
  - Diameter
  - in- and out- degree distributions
  - other (surprising) patterns
- So, let's look at the data



### **Solution# S.1**

• Power law in the degree distribution [SIGCOMM99]

#### internet domains

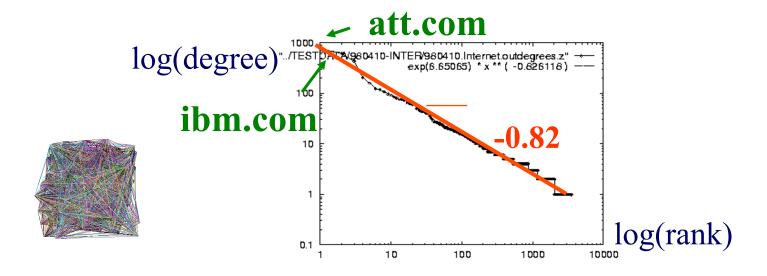




### Solution# S.1

• Power law in the degree distribution [SIGCOMM99]

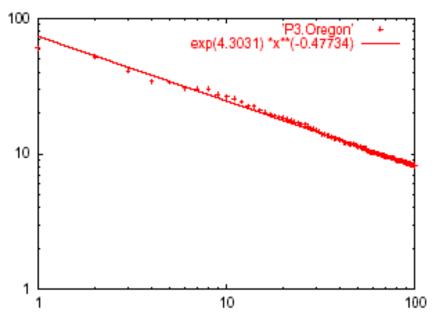
#### internet domains





## Solution# S.2: Eigen Exponent *E*

#### Eigenvalue



Exponent = slope

$$E = -0.48$$

May 2001

Rank of decreasing eigenvalue

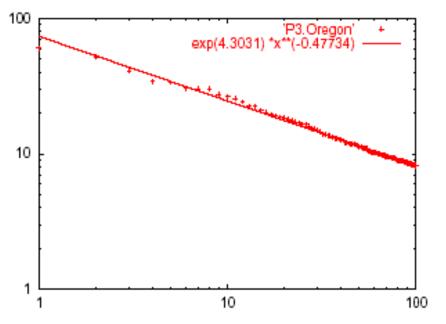
• A2: power law in the eigenvalues of the adjacency matrix

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## Solution# S.2: Eigen Exponent *E*

#### Eigenvalue



Exponent = slope

$$E = -0.48$$

May 2001

Rank of decreasing eigenvalue

• [Mihail, Papadimitriou '02]: slope is ½ of rank exponent

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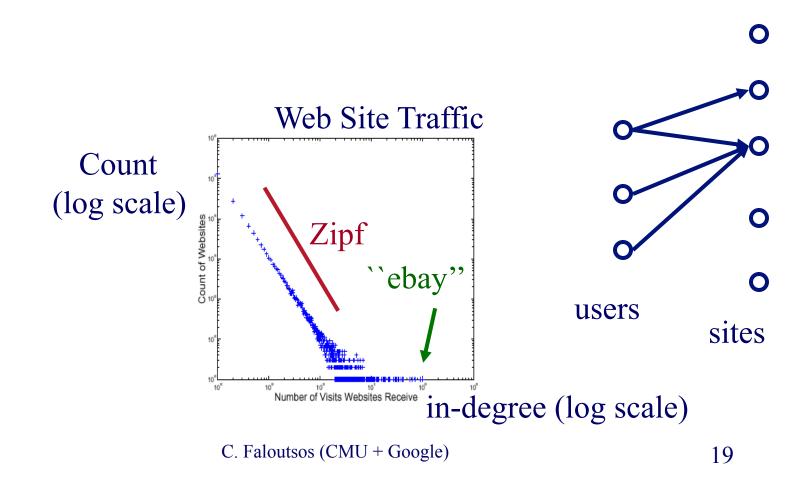
### **But:**

How about graphs from other domains?



## More power laws:

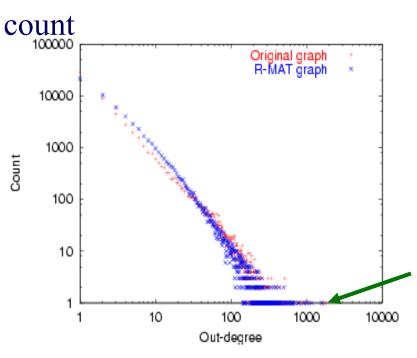
• web hit counts [w/ A. Montgomery]



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## epinions.com



who-trusts-whom
 [Richardson +
 Domingos, KDD
 2001]

trusts-2000-people user

(out) degree

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### And numerous more

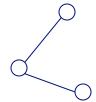
- # of sexual contacts
- Income [Pareto] –'80-20 distribution'
- Duration of downloads [Bestavros+]
- Duration of UNIX jobs ('mice and elephants')
- Size of files of a user
- •
- 'Black swans'

#### **Outline**

- Introduction Motivation
- Problem#1: Patterns in graphs
  - Static graphs
    - degree, diameter, eigen,
- triangles
- cliques
- Weighted graphs
- Time evolving graphs
- Problem#2: Tools



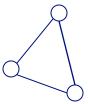
## Solution# S.3: Triangle 'Laws'



Real social networks have a lot of triangles



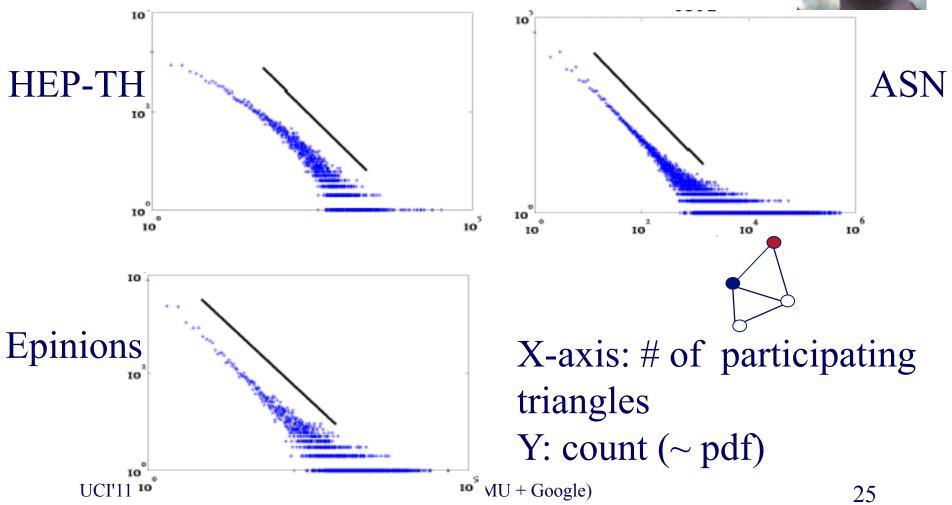
## Solution# S.3: Triangle 'Laws'



- Real social networks have a lot of triangles
  - Friends of friends are friends
- Any patterns?

## Triangle Law: #S.3 [Tsourakakis ICDM 2008]



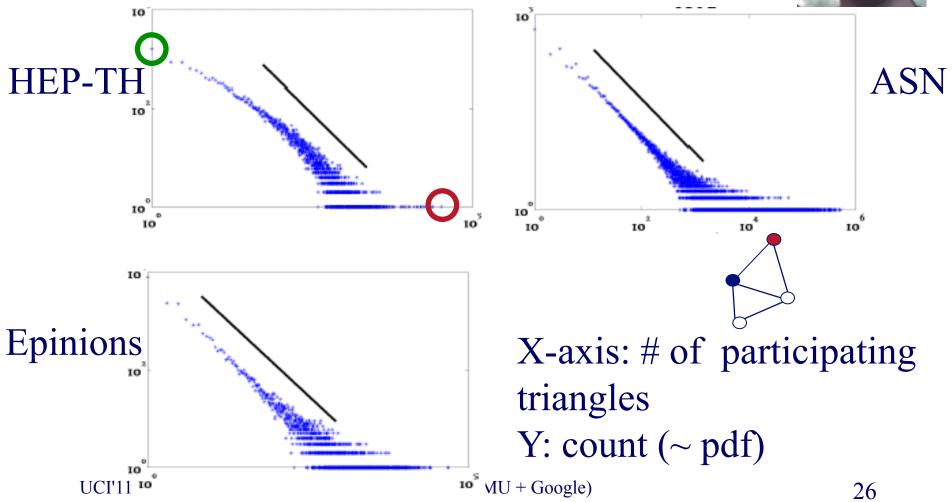


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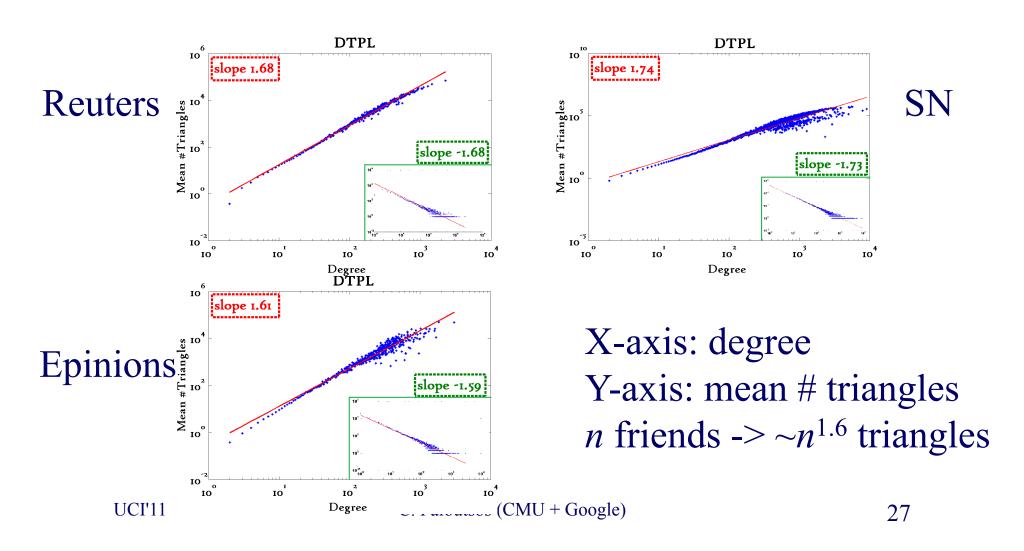
## Triangle Law: #S.3 [Tsourakakis ICDM 2008]





26

# Triangle Law: #S.4 [Tsourakakis ICDM 2008]





# Triangle Law: Computations [Tsourakakis ICDM 2008]

But: triangles are expensive to compute (3-way join; several approx. algos) Q: Can we do that quickly?

details



## Triangle Law: Computations

[Tsourakakis ICDM 2008]

```
But: triangles are expensive to compute (3-way join; several approx. algos)
Q: Can we do that quickly?
A: Yes!

#triangles = 1/6 Sum (\lambda_i^3)

(and, because of skewness (S2),

we only need the top few eigenvalues!
```

details

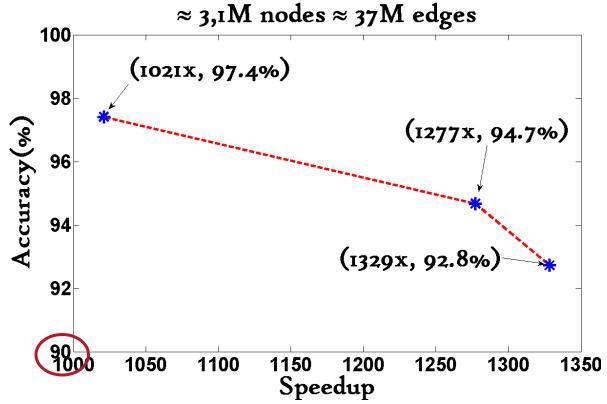




## Triangle Law: Computations

#### [Tsourakakis ICDM 2008]

Wikipedia graph 2006-Nov-04 ≈ 3.1M nodes ≈ 37M edges



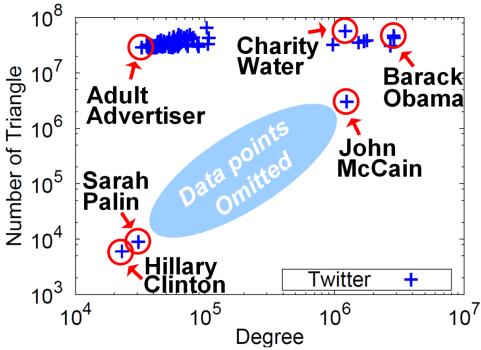


## Triangle counting for large graphs?

Anomalous nodes in Twitter(~ 3 billion edges)
[U Kang, Brendan Meeder, +, PAKDD'11]



## Triangle counting for large graphs?



Anomalous nodes in Twitter(~ 3 billion edges)
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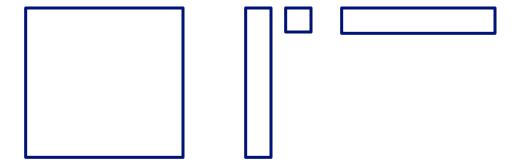
B. Aditya Prakash, Mukund Seshadri, Ashwin Sridharan, Sridhar Machiraju and Christos Faloutsos: *EigenSpokes: Surprising Patterns and Scalable Community Chipping in Large Graphs*, PAKDD 2010, Hyderabad, India, 21-24 June 2010.



## **EigenSpokes**

- Eigenvectors of adjacency matrix
  - equivalent to singular vectors (symmetric, undirected graph)

$$A = U\Sigma U^T$$

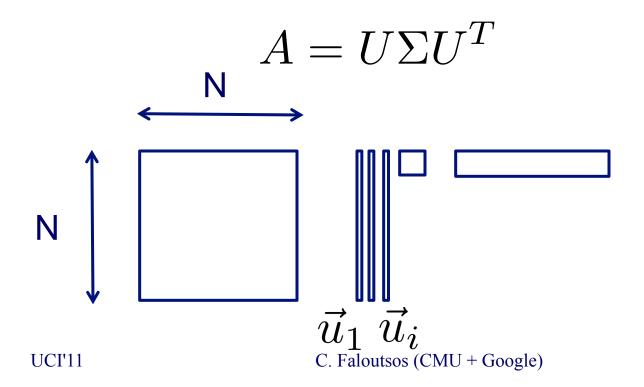






## **EigenSpokes**

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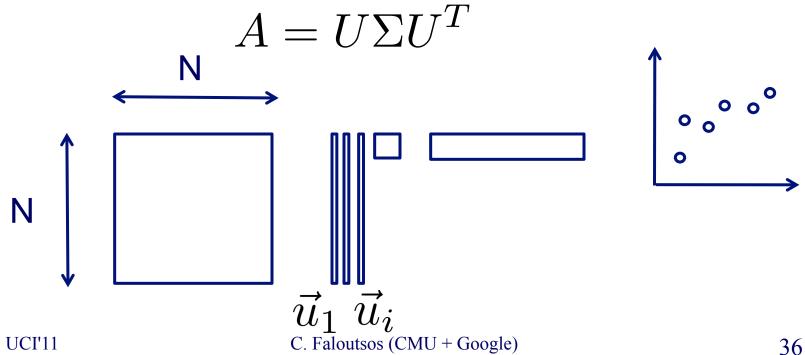






## **EigenSpokes**

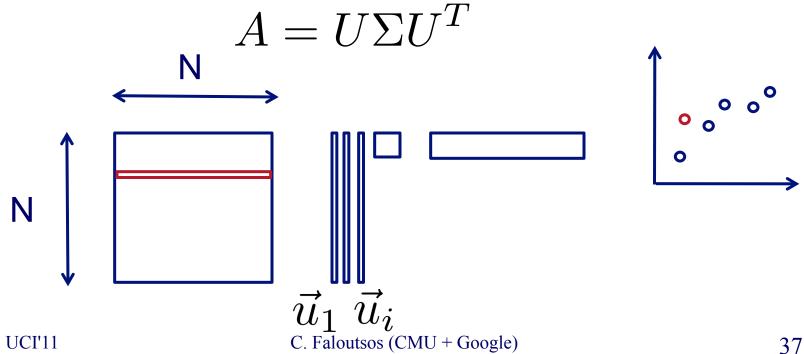
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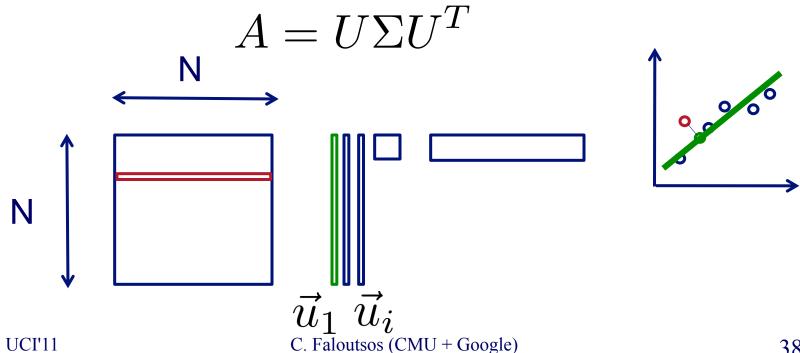
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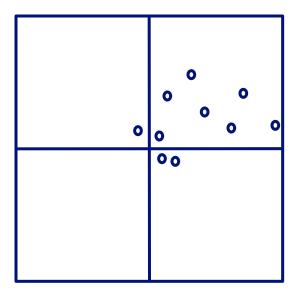




• EE plot:

2<sup>nd</sup> Principal component u2

- Scatter plot of scores of u1 vs u2
- One would expect
  - Many points @origin
  - A few scattered~randomly



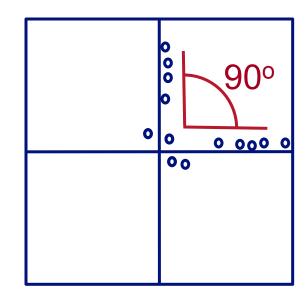
**u**1

1<sup>st</sup> Principal component



**u**2

- EE plot:
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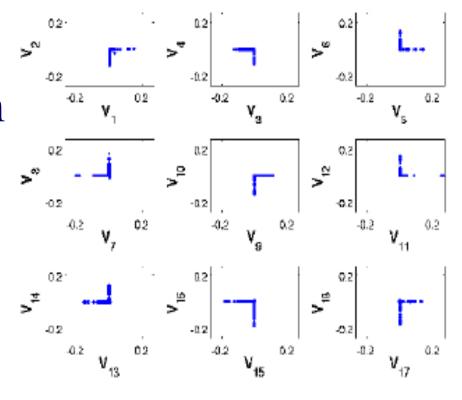
**u**1



## EigenSpokes - pervasiveness

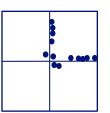
- Present in mobile social graph
  - across time and space

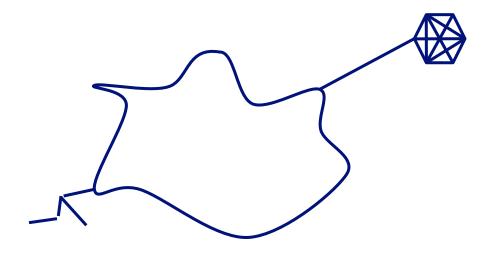
Patent citation graph





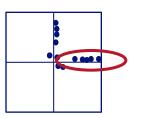
Near-cliques, or nearbipartite-cores, loosely connected

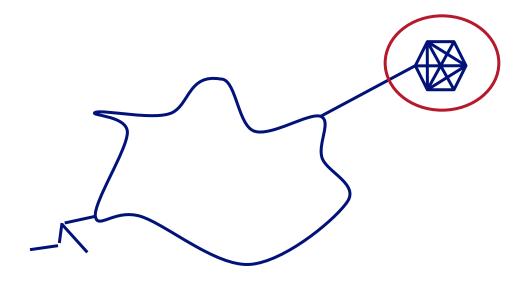






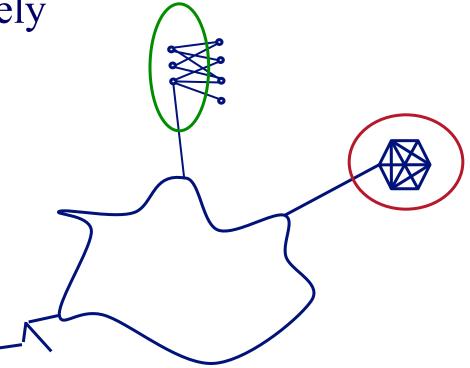
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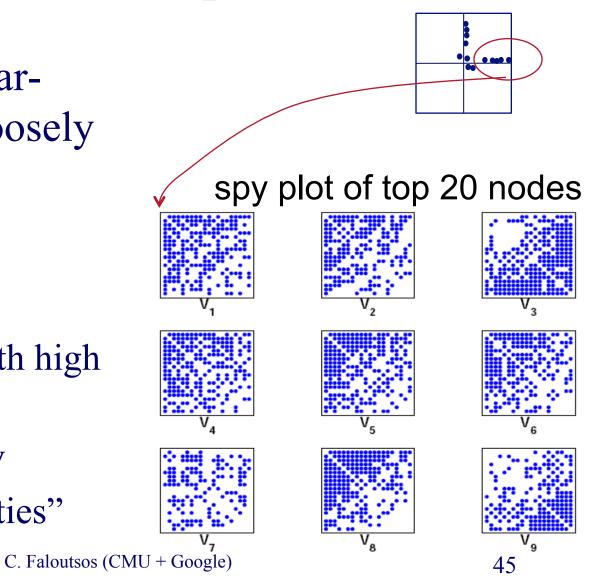




Near-cliques, or nearbipartite-cores, loosely connected

#### So what?

- Extract nodes with high scores
- high connectivity
- Good "communities"



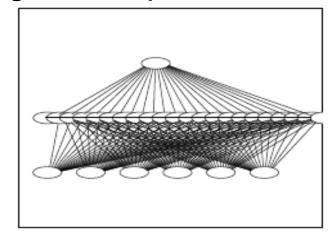


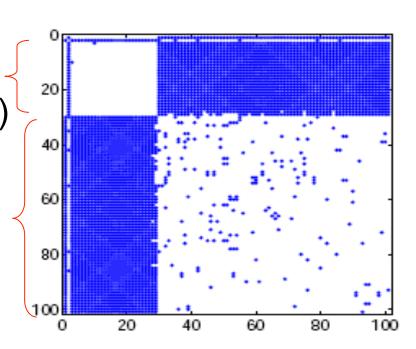
## **Bipartite Communities!**

patents from same inventor(s)

`cut-and-paste' bibliography!

magnified bipartite community





#### **Outline**

- Introduction Motivation
- Problem#1: Patterns in graphs
  - Static graphs
    - degree, diameter, eigen,
    - triangles
    - cliques



- Weighted graphs
- Time evolving graphs
- Problem#2: Tools



# Observations on weighted graphs?

A: yes - even more 'laws'!





M. McGlohon, L. Akoglu, and C. Faloutsos Weighted Graphs and Disconnected Components: Patterns and a Generator. SIG-KDD 2008



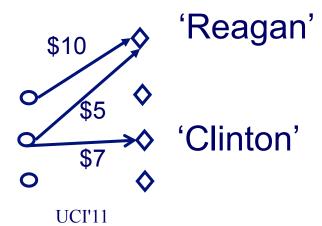
#### **Observation W.1: Fortification**

Q: How do the weights of nodes relate to degree?



#### **Observation W.1: Fortification**

## More donors, more \$ ?

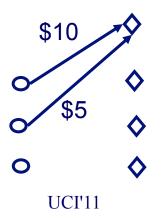




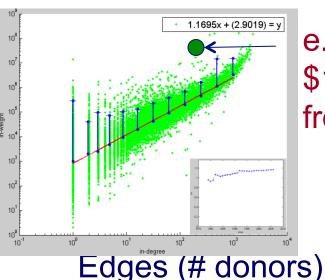
# Observation W.1: fortification: Snapshot Power Law

- Weight: super-linear on in-degree
- exponent 'iw': 1.01 < iw < 1.26

## More donors, even more \$



In-weights (\$)



e.g. John Kerry, \$10M received, from 1K donors

**Orgs-Candidates** 

C. Faloutsos (CMU + Google)

51

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•



#### **Problem: Time evolution**

 with Jure Leskovec (CMU -> Stanford)



and Jon Kleinberg (Cornell – sabb. @ CMU)





#### T.1 Evolution of the Diameter

- Prior work on Power Law graphs hints at slowly growing diameter:
  - diameter  $\sim$  O(log N)
  - diameter  $\sim$  O(log log N)





What is happening in real data?



#### T.1 Evolution of the Diameter

- Prior work on Power Law graphs hints at slowly growing diameter:

  - diameter ~ (log N)diameter ~ O(log log N)

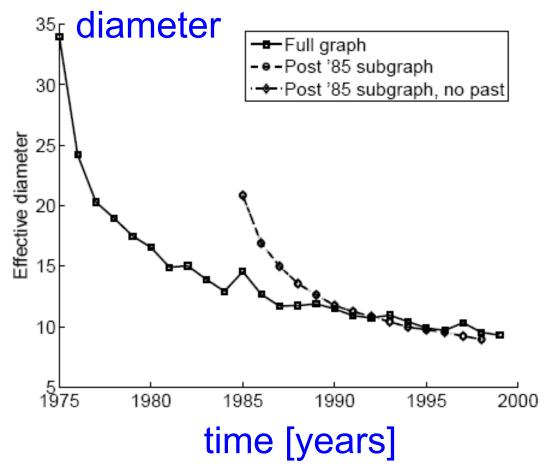


- What is happening in real data?
- Diameter shrinks over time



#### T.1 Diameter – "Patents"

- Patent citation network
- 25 years of data
- @1999
  - 2.9 M nodes
  - 16.5 M edges



# T.2 Temporal Evolution of the Graphs

- N(t) ... nodes at time t
- E(t) ... edges at time t
- Suppose that

$$N(t+1) = 2 * N(t)$$

• Q: what is your guess for

$$E(t+1) = ?2 * E(t)$$

# T.2 Temporal Evolution of the Graphs

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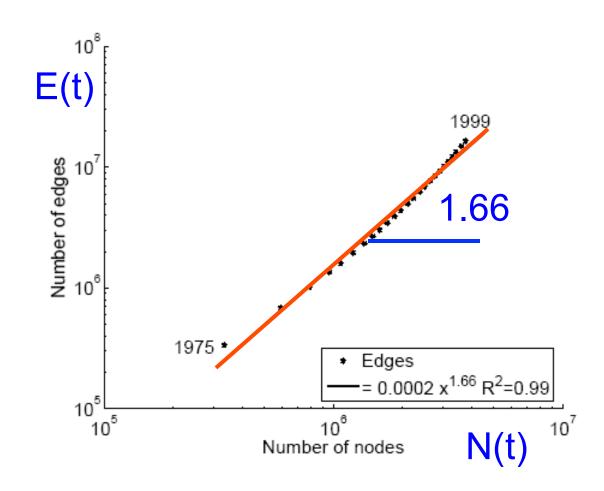
$$N(t+1) = 2 * N(t)$$

- Q: what is your guess for  $E(t+1) = (2)^* E(t)$
- A: over-doubled!
  - But obeying the ``Densification Power Law''



## T.2 Densification – Patent Citations

- Citations among patents granted
- (a) 1999
  - -2.9 M nodes
  - 16.5 M edges
- Each year is a datapoint



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  - Weighted graphs



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•



### More on Time-evolving graphs

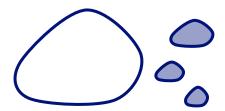
M. McGlohon, L. Akoglu, and C. Faloutsos Weighted Graphs and Disconnected Components: Patterns and a Generator. SIG-KDD 2008



Q: How do NLCC's emerge and join with the GCC?

(``NLCC'' = non-largest conn. components)

- −Do they continue to grow in size?
- or do they shrink?
- or stabilize?

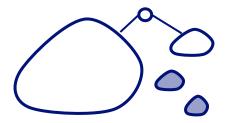




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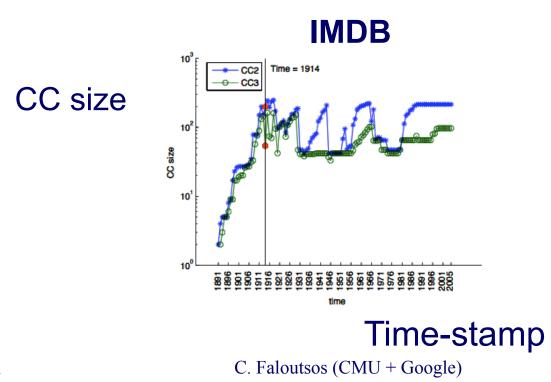
YES – Do they continue to grow in size?

YES – or do they shrink?

YES – or stabilize?



• After the gelling point, the GCC takes off, but NLCC's remain ~constant (actually, **oscillate**).



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#### **Timing for Blogs**

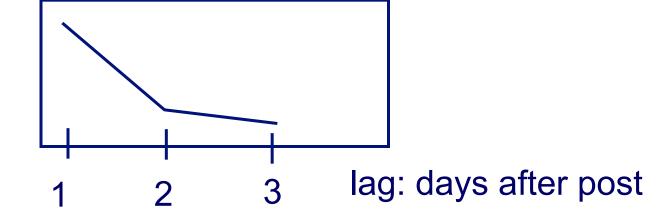
- with Mary McGlohon (CMU->Google)
- Jure Leskovec (CMU->Stanford)
- Natalie Glance (now at Google)
- Mat Hurst (now at MSR)

[SDM'07]

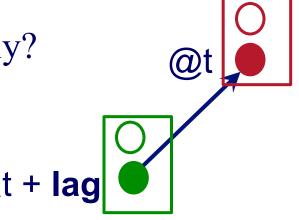


## T.4: popularity over time

# in links



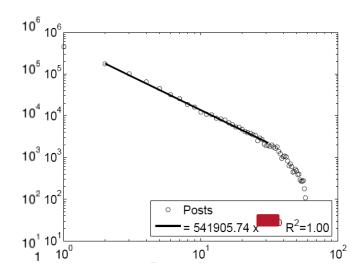
Post popularity drops-off – exponentially?





## T.4: popularity over time

# in links (log)



days after post (log)

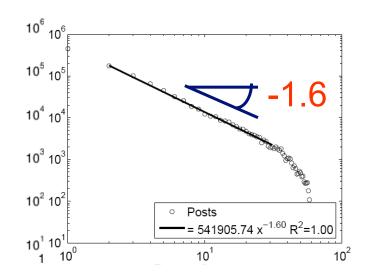
Post popularity drops-off – exponentally? POWER LAW!

Exponent?



## T.4: popularity over time

# in links (log)

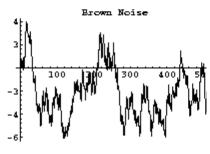


days after post (log)

Post popularity drops-off – exported ally? POWER LAW!

Exponent? -1.6

- close to -1.5: Barabasi's stack model
- and like the zero-crossings of a random walk
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DFT of Brown Noise

69



### -1.5 slope

J. G. Oliveira & A.-L. Barabási Human Dynamics: The Correspondence Patterns of Darwin and Einstein. *Nature* **437**, 1251 (2005) . [PDF]



## T.5: duration of phonecalls

Surprising Patterns for the Call Duration Distribution of Mobile Phone Users

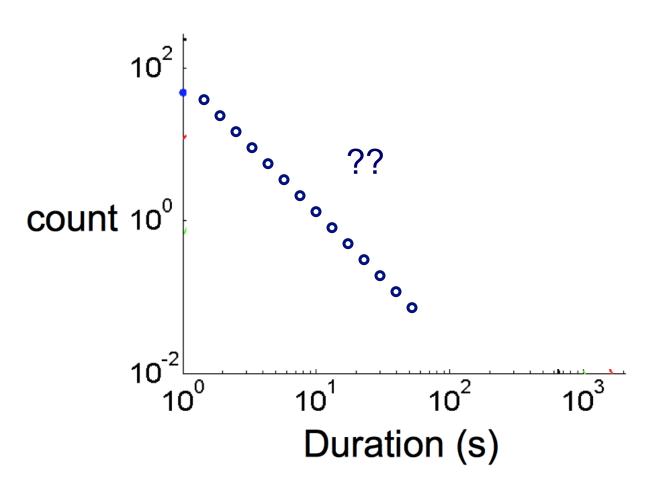


Pedro O. S. Vaz de Melo, Leman Akoglu, Christos Faloutsos, Antonio A. F. Loureiro

PKDD 2010

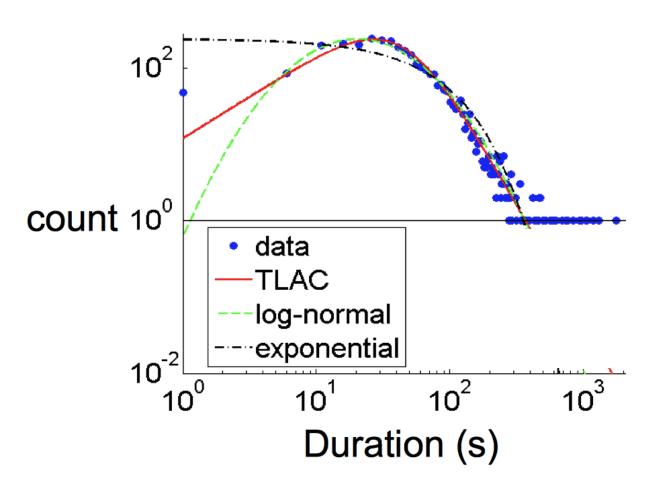


## Probably, power law (?)





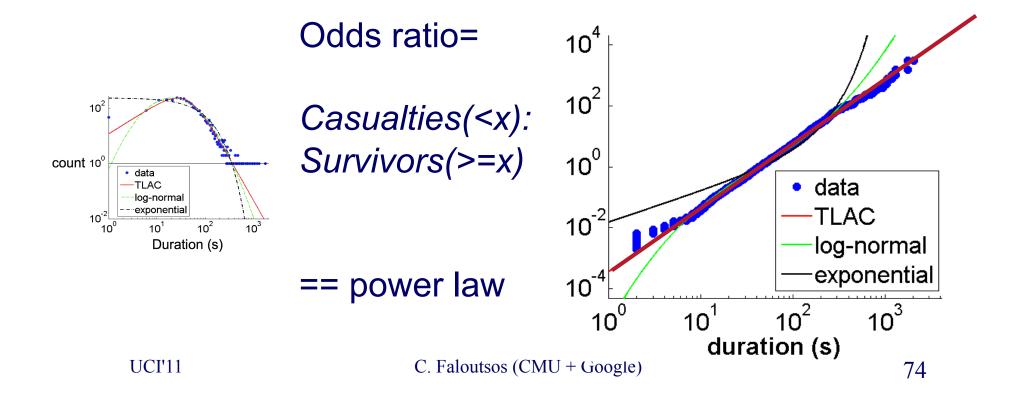
#### No Power Law!





## 'TLaC: Lazy Contractor'

- The longer a task (phonecall) has taken,
- The even longer it will take





## **Data Description**

- Data from a private mobile operator of a large city
  - 4 months of data
  - 3.1 million users
  - more than 1 billion phone records
- Over 96% of 'talkative' users obeyed a TLAC distribution ('talkative': >30 calls)

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- OddBall (anomaly detection)
- Belief Propagation
- Immunization
- Problem#3: Scalability
- Conclusions



# OddBall: Spotting Anomalies in Weighted Graphs

Leman Akoglu, Mary McGlohon, Christos Faloutsos

> Carnegie Mellon University School of Computer Science

PAKDD 2010, Hyderabad, India



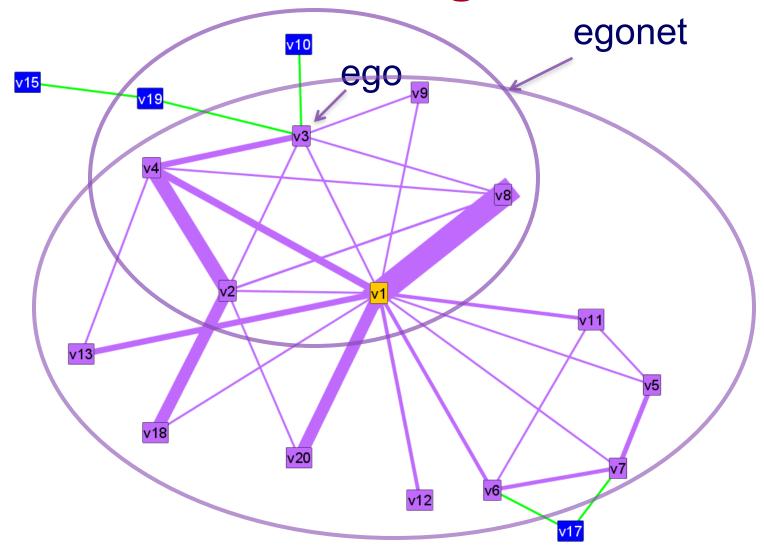
#### Main idea

For each node,

- extract 'ego-net' (=1-step-away neighbors)
- Extract features (#edges, total weight, etc etc)
- Compare with the rest of the population

#### Carnegie Mellon

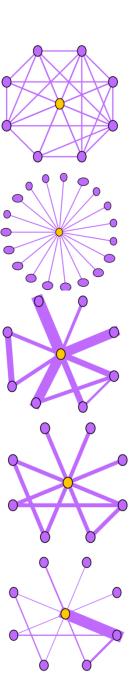
## What is an egonet?



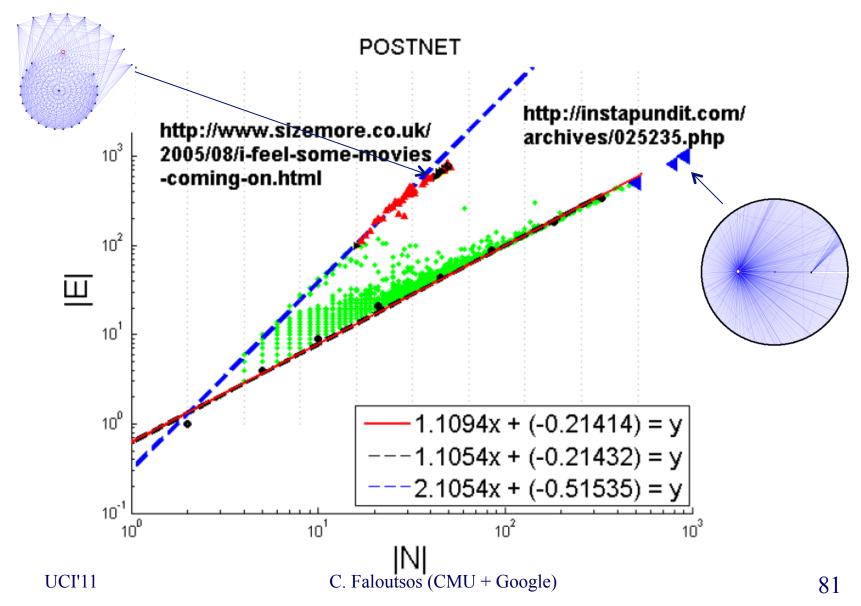


#### **Selected Features**

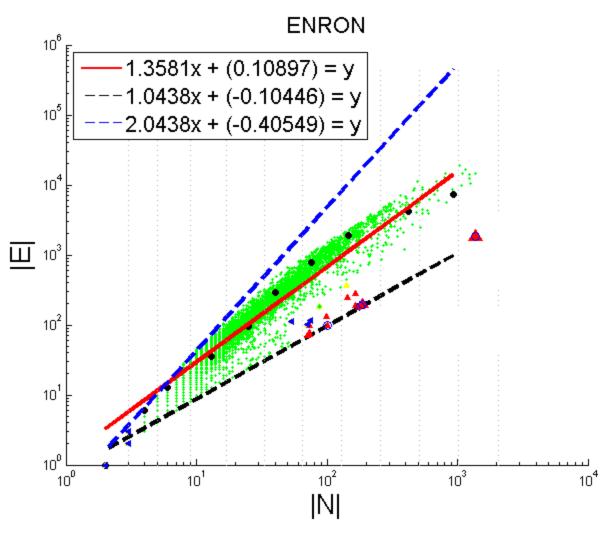
- $N_i$ : number of neighbors (degree) of ego i
- $E_i$ : number of edges in egonet i
- W<sub>i</sub>: total weight of egonet i
- $\lambda_{w,i}$ : principal eigenvalue of the weighted adjacency matrix of egonet I







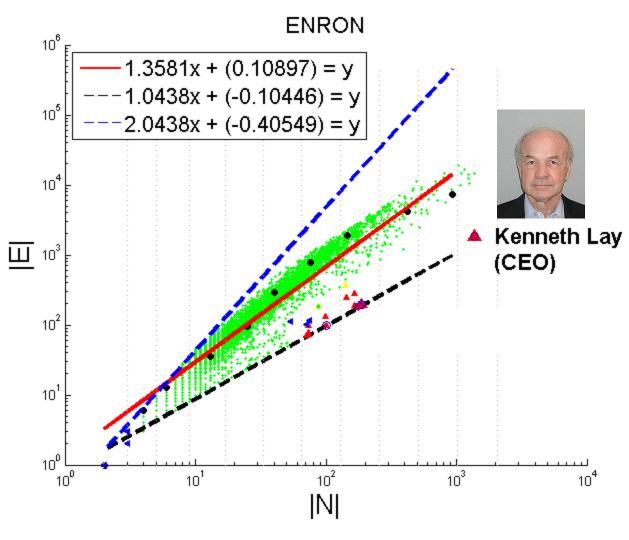




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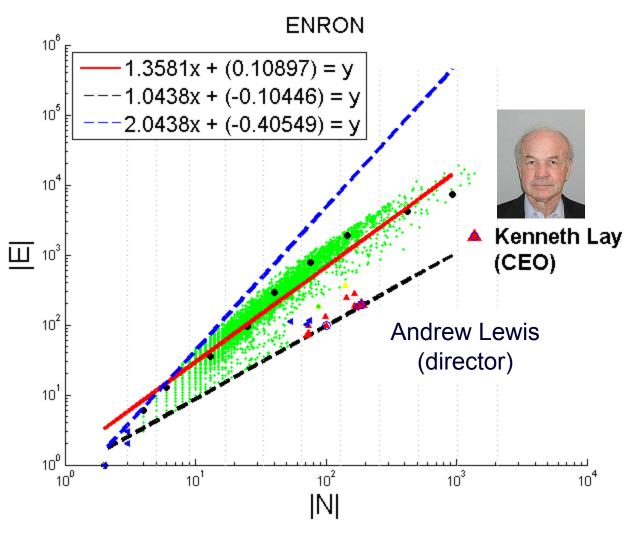




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#### **Outline**

- Introduction Motivation
- Problem#1: Patterns in graphs
- Problem#2: Tools
  - OddBall (anomaly detection)



- Belief Propagation
- Immunization
- Problem#3: Scalability
- Conclusions

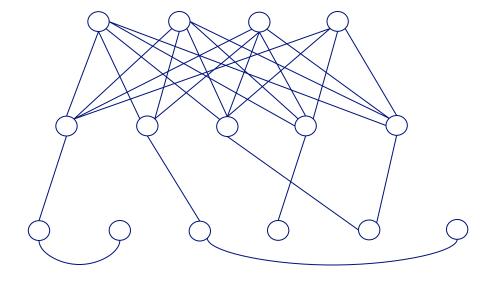


## **E-bay Fraud detection**



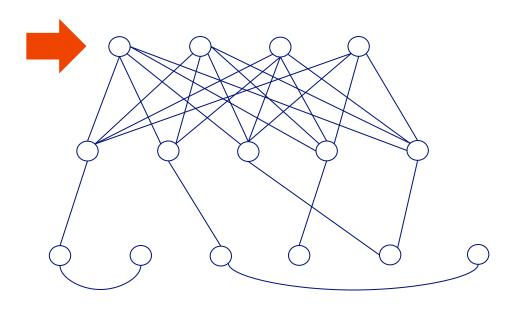


w/ Polo Chau & Shashank Pandit, CMU [www'07]



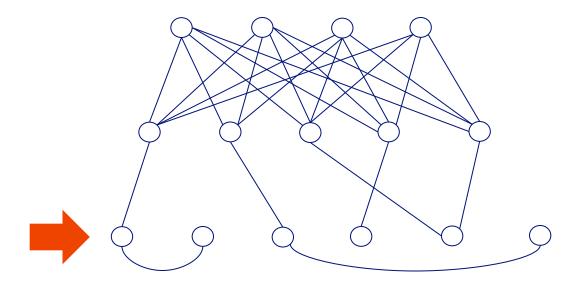


## **E-bay Fraud detection**



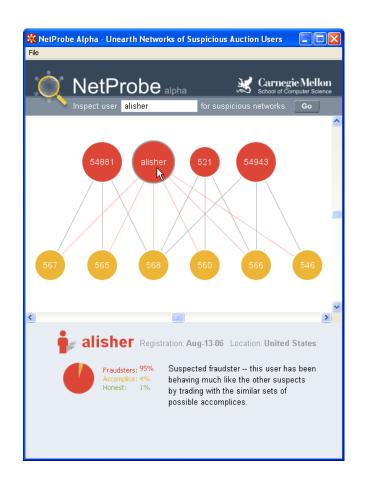


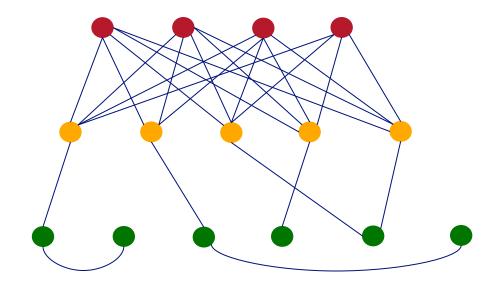
## **E-bay Fraud detection**





#### **E-bay Fraud detection - NetProbe**







## Popular press



The Washington Post

Los Angeles Times

And less desirable attention:

• E-mail from 'Belgium police' ('copy of your code?')



## **Polonium:** Tera-Scale Graph Mining and Inference for Malware Detection

SDM 2011, Mesa, Arizona



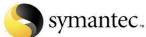
**Polo Chau** Machine Learning Dept Vice President & Fellow



symantec...

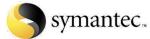
**Carey Nachenberg** 





**Jeffrey Wilhelm Principal Software Engineer** 





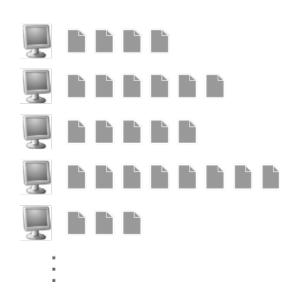
**Adam Wright** Software Engineer

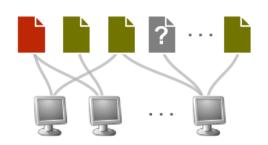


**Prof. Christos Faloutsos** Computer Science Dept



#### Carnegie Mellon





#### The Data

60+ terabytes of data *anonymously* contributed by participants of worldwide *Norton Community Watch* program

50+ million machines

900+ million executable files

Constructed a machine-file bipartite graph (0.2 TB+)

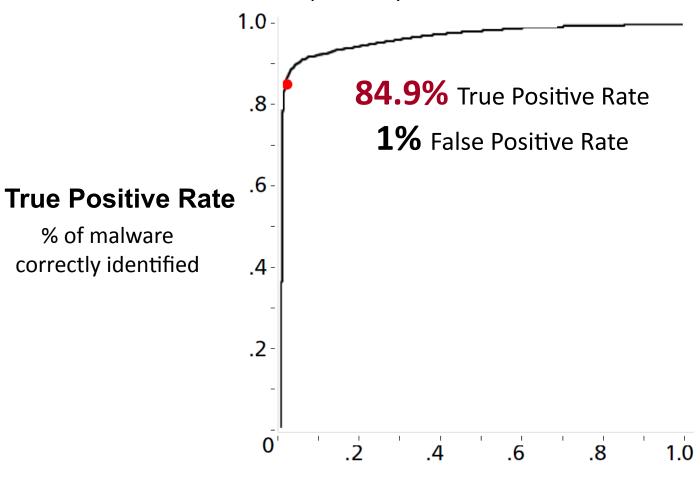
1 billion nodes (machines and files)

37 billion edges



#### **One-Iteration Results**

for files reported by four or more machines



**False Positive Rate** 

#### **Outline**

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- Problem#2: Tools
  - OddBall (anomaly detection)
  - Belief propagation



- Immunization
- Problem#3: Scalability -PEGASUS
- Conclusions

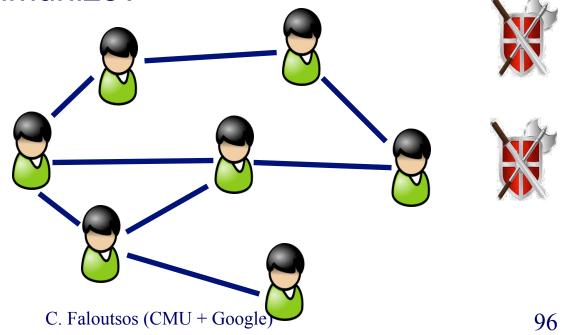


## Immunization and epidemic thresholds

- Q1: which nodes to immunize?
- Q2: will a virus vanish, or will it create an epidemic?

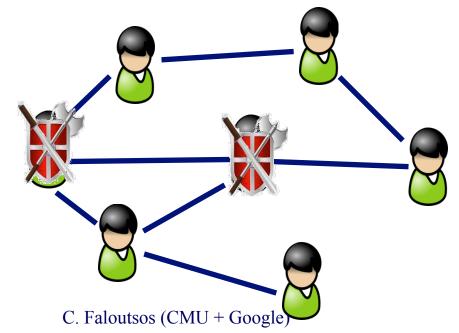


- Given
  - a network,
  - k vaccines, and
  - the virus details
- •Which nodes to immunize?



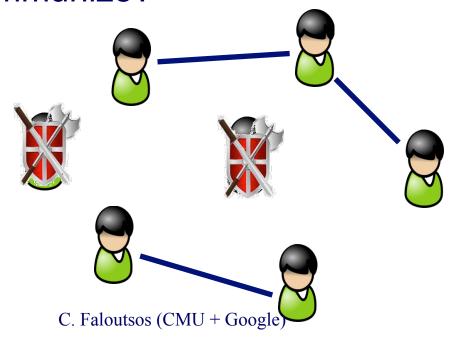


- Given
  - a network,
  - k vaccines, and
  - the virus details
- •Which nodes to immunize?





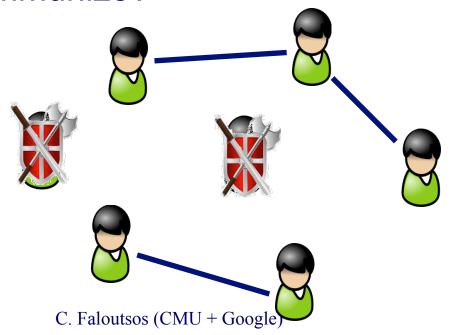
- Given
  - a network,
  - k vaccines, and
  - the virus details
- •Which nodes to immunize?





- Given
  - a network,
  - k vaccines, and
  - the virus details
- •Which nodes to immunize?

A: immunize the ones that maximally raise the `epidemic threshold' [Tong+, ICDM'10]



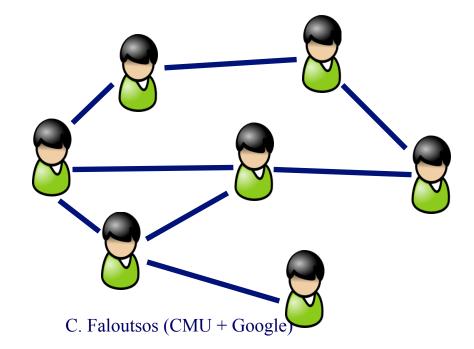


## Q2: will a virus take over?

- Flu-like virus (no immunity, 'SIS')
- Mumps (life-time immunity, 'SIR')
- Pertussis (finite-length immunity, 'SIRS')

β: attack prob

δ: heal prob





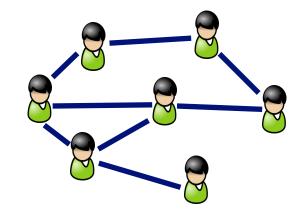
## Q2: will a virus take over?

- Flu-like virus (no immunity, 'SIS')
- Mumps (life-time immunity, 'SIR')
- Pertussis (finite-length immunity, 'SIRS')

β: attack prob

δ: heal prob

A: depends on connectivity (avg degree? Max degree? variance? Something else?

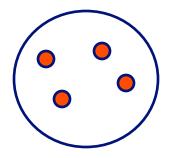




## Epidemic threshold τ

What should  $\tau$  depend on?

- avg. degree? and/or highest degree?
- and/or variance of degree?
- and/or third moment of degree?
- and/or diameter?











## **Epidemic threshold**

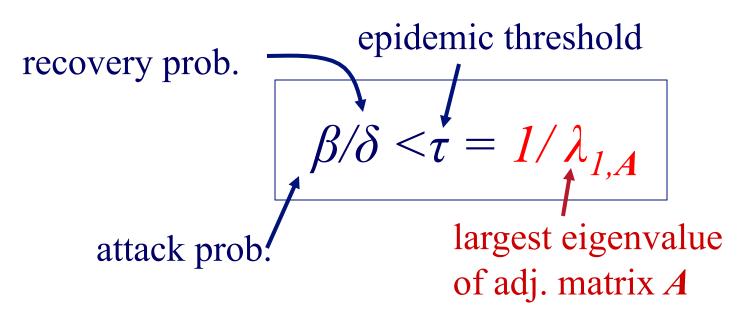
• [Theorem] We have no epidemic, if

$$\beta/\delta < \tau = 1/\lambda_{1,A}$$



## **Epidemic threshold**

• [Theorem] We have no epidemic, if



Proof: [Wang+03] (for SIS=flu only)



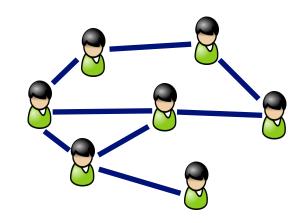
#### A2: will a virus take over?

- For all typical virus propagation models (flu, mumps, pertussis, HIV, etc)
- The **only** connectivity measure that matters, is

$$1/\lambda_1$$

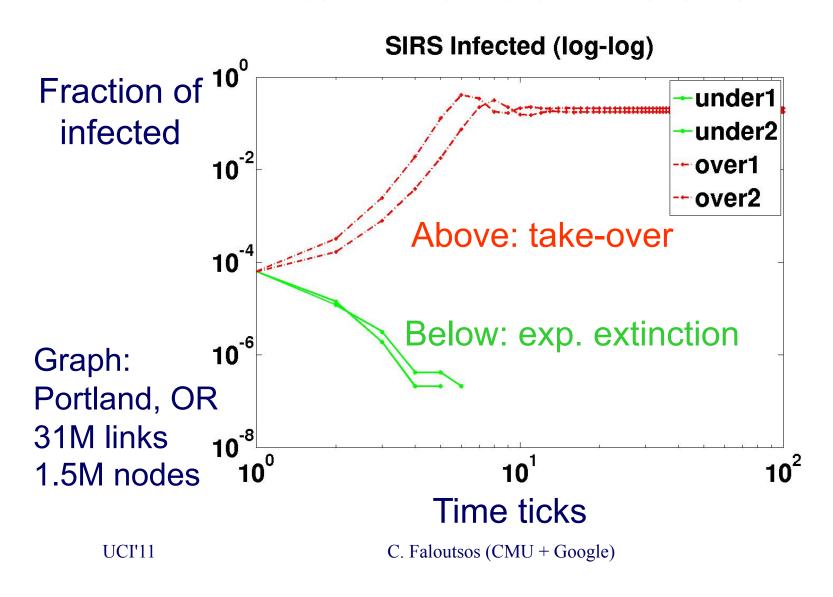
the first eigenvalue of the adj. matrix

[Prakash+, '10, arxiv]





#### A2: will a virus take over?



106

#### **Outline**

- Introduction Motivation
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  - Belief propagation
  - Immunization
- Problem#3: Scalability -PEGASUS
  - Conclusions





## **Scalability**

- Google: > 450,000 processors in clusters of ~2000 processors each [Barroso, Dean, Hölzle, "Web Search for a Planet: The Google Cluster Architecture" IEEE Micro 2003]
- Yahoo: 5Pb of data [Fayyad, KDD'07]
- Problem: machine failures, on a daily basis
- How to parallelize data mining tasks, then?
- A: map/reduce hadoop (open-source clone)
   <a href="http://hadoop.apache.org/">http://hadoop.apache.org/</a>





## Outline – Algorithms & results

	Centralized	Hadoop/ PEGASUS
Degree Distr.	old	old
Pagerank	old	old
Diameter/ANF	old	HERE
Conn. Comp	old	HERE
Triangles	done	HERE
Visualization	started	

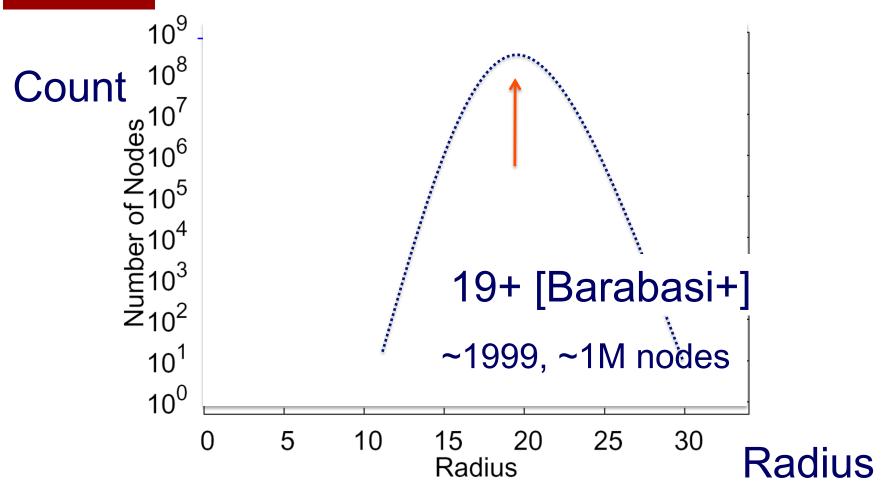


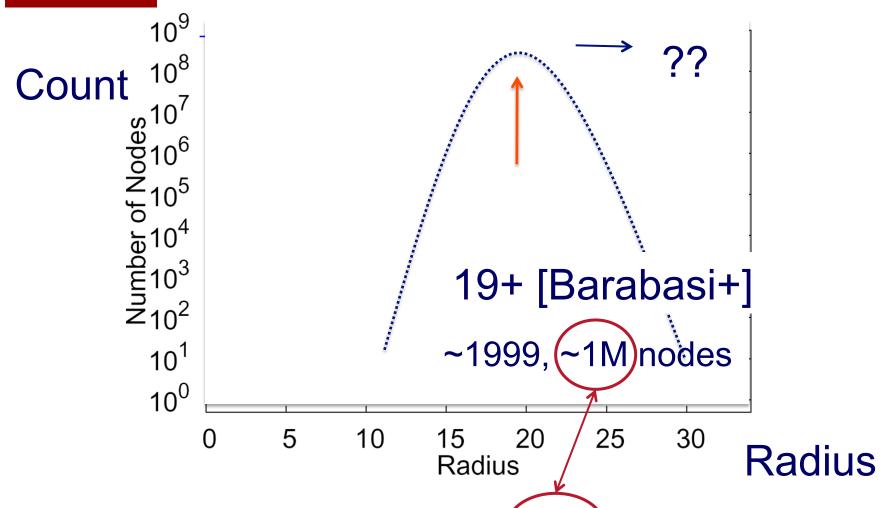




#### HADI for diameter estimation

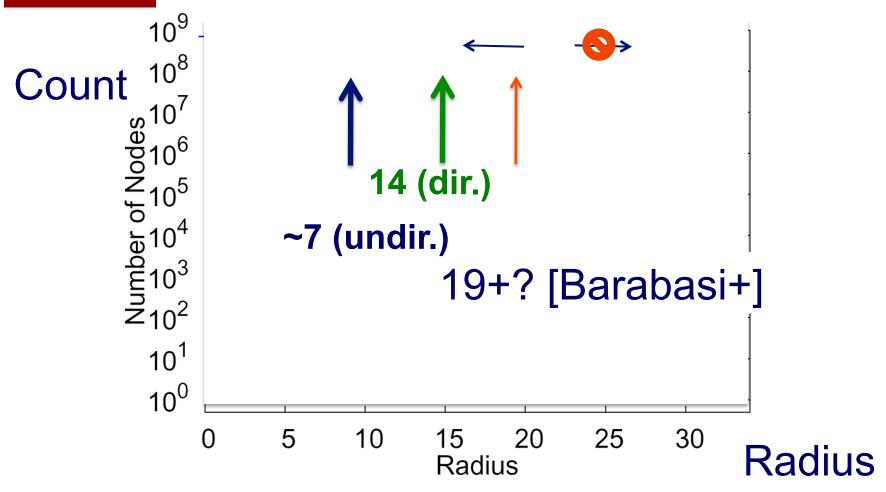
- Radius Plots for Mining Tera-byte Scale
   Graphs U Kang, Charalampos Tsourakakis,
   Ana Paula Appel, Christos Faloutsos, Jure
   Leskovec, SDM'10
- Naively: diameter needs O(N\*\*2) space and up to O(N\*\*3) time – prohibitive (N~1B)
- Our HADI: linear on E (~10B)
  - Near-linear scalability wrt # machines
  - Several optimizations -> 5x faster





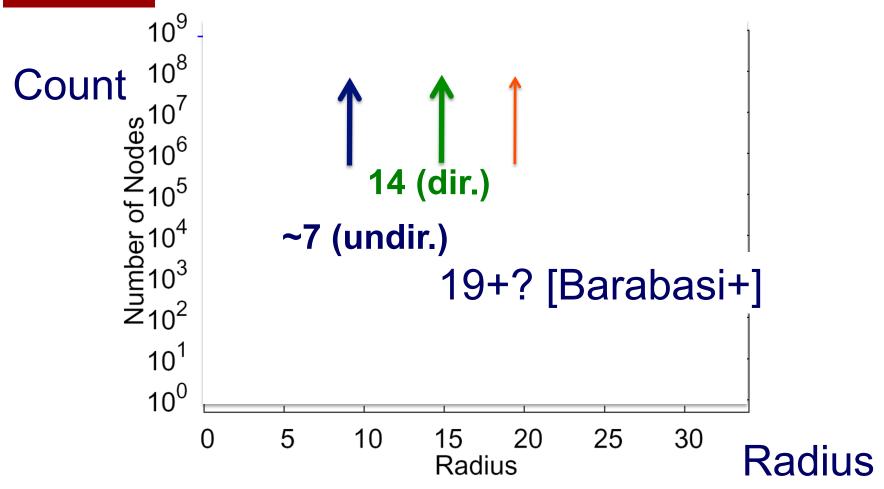
YahooWeb graph (120Gb, 1.4B hodes, 6.6 B edges)

Largest publicly available graph ever studied.



YahooWeb graph (120Gb, 1.4B nodes, 6.6 B edges)

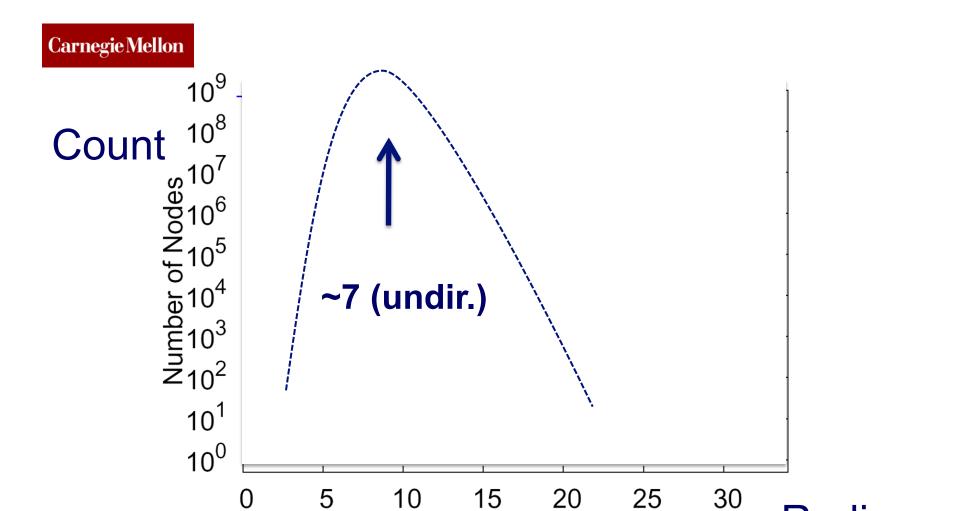
Largest publicly available graph ever studied.



YahooWeb graph (120Gb, 1.4B nodes, 6.6 B edges)

- •7 degrees of separation (!)
- Diameter: shrunk

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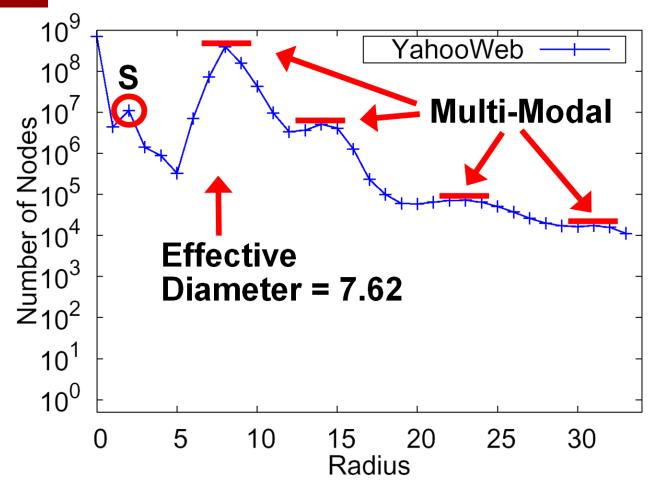


YahooWeb graph (120Gb, 1.4B nodes, 6.6 B edges) Q: Shape?

Radius

0

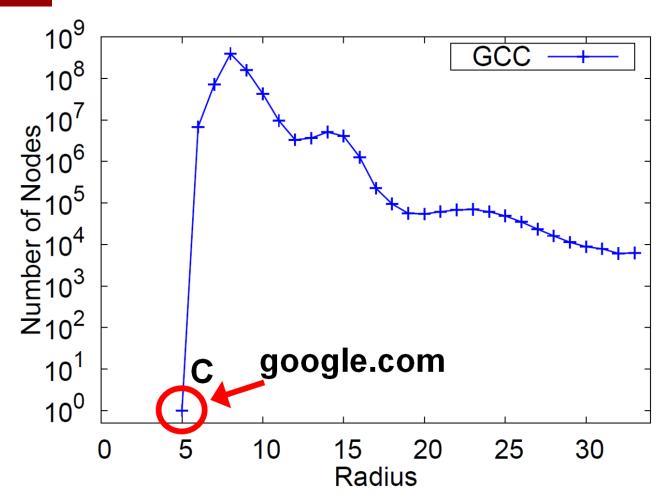
Radius



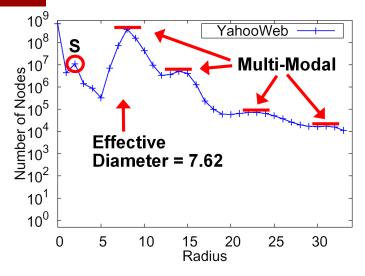
YahooWeb graph (120Gb, 1.4B nodes, 6.6 B edges)

- effective diameter: surprisingly small.
- Multi-modality (?!)

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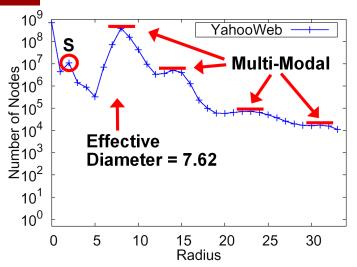
#### Radius Plot of GCC of YahooWeb.



YahooWeb graph (120Gb, 1.4B nodes, 6.6 B edges)

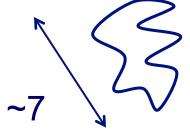
- effective diameter: surprisingly small.
- Multi-modality: probably mixture of cores.

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### Conjecture:

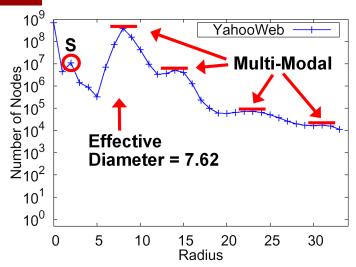


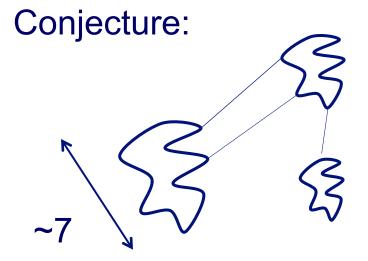


S BR

YahooWeb graph (120Gb, 1.4B nodes, 6.6 B edges)

- effective diameter: surprisingly small.
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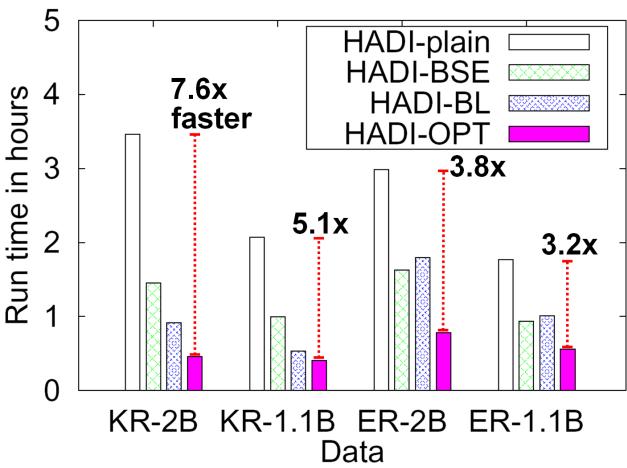


YahooWeb graph (120Gb, 1.4B nodes, 6.6 B edges)

- effective diameter: surprisingly small.
- Multi-modality: probably mixture of cores .







Running time - Kronecker and Erdos-Renyi Graphs with billions edges.



## Outline – Algorithms & results

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Diameter/ANF	old	HERE
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Triangles		HERE
Visualization	started	



## Generalized Iterated Matrix Vector Multiplication (GIMV)

<u>PEGASUS: A Peta-Scale Graph Mining</u> System - Implementation and Observations.

U Kang, Charalampos E. Tsourakakis, and Christos Faloutsos.

(ICDM) 2009, Miami, Florida, USA. Best Application Paper (runner-up).



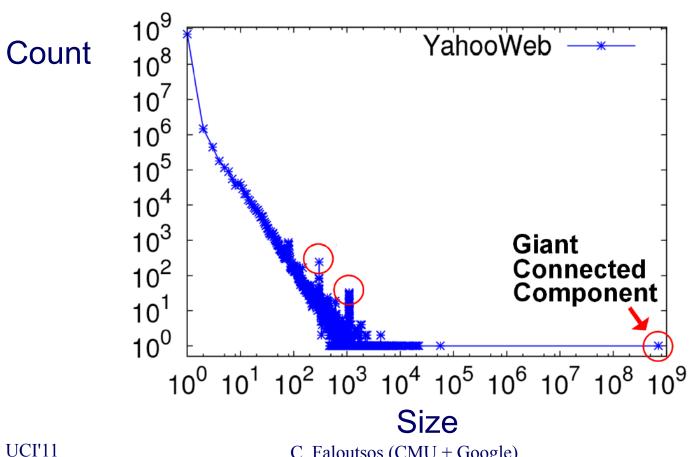
# Generalized Iterated Matrix Vector Multiplication (GIMV)

- PageRank
- proximity (RWR)
- Diameter
- Connected components
- (eigenvectors,
- Belief Prop.
- ...)

Matrix – vector Multiplication (iterated)

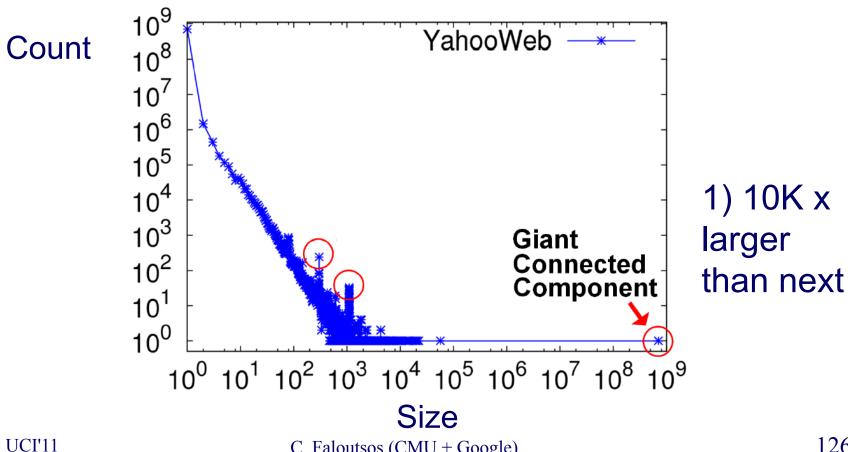


• Connected Components – 4 observations:



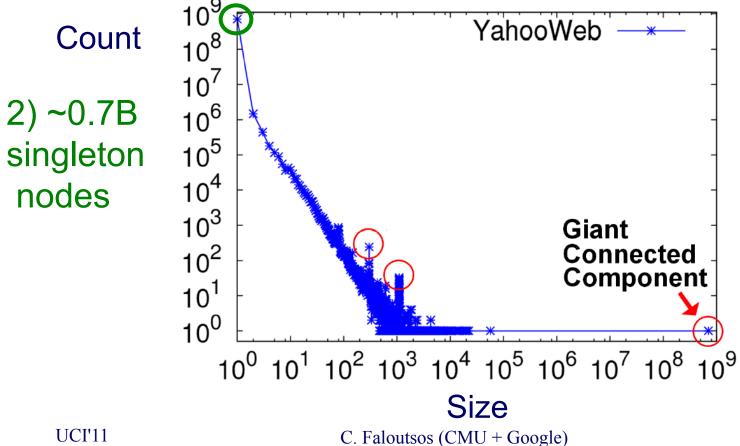


Connected Components



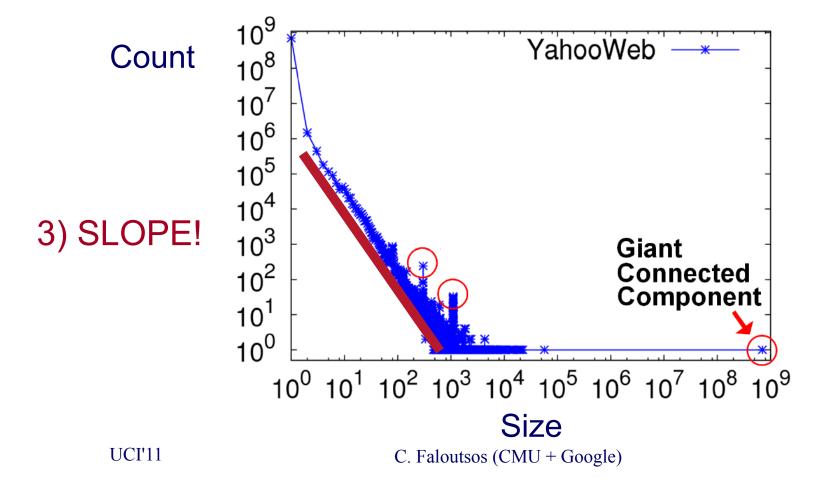


Connected Components



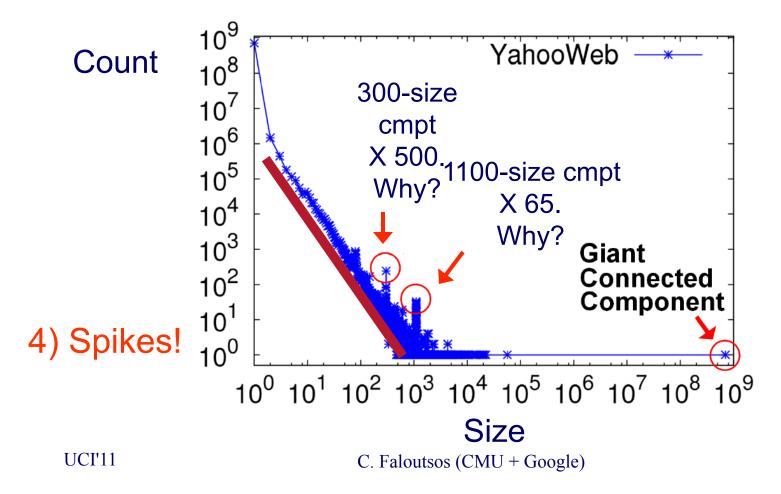


Connected Components



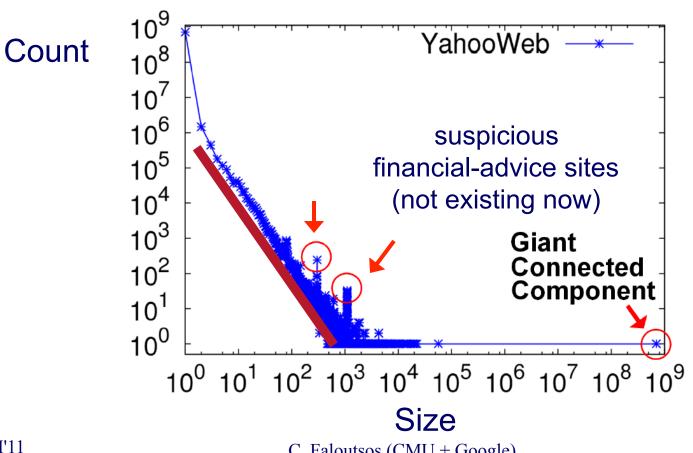


Connected Components





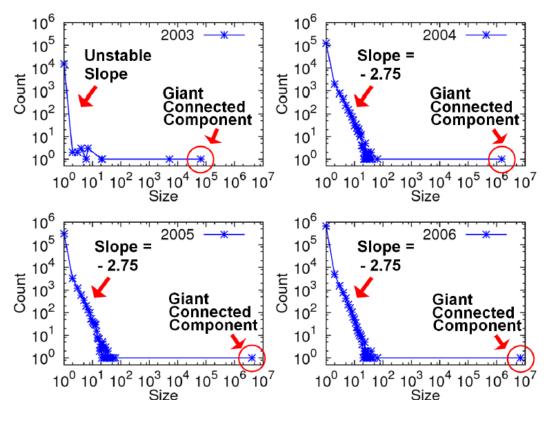
Connected Components





#### **GIM-V At Work**

- Connected Components over Time
- LinkedIn: 7.5M nodes and 58M edges



Stable tail slope after the gelling point

#### **Outline**

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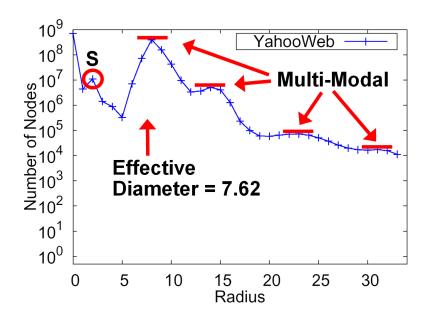
## OVERALL CONCLUSIONS – low level:

- Several new **patterns** (fortification, triangle-laws, conn. components, etc)
- New tools:
  - anomaly detection (OddBall), belief propagation, immunization
- Scalability: PEGASUS / hadoop



## OVERALL CONCLUSIONS – high level

• **BIG DATA:** Large datasets reveal patterns/ outliers that are invisible otherwise



## OVERALL CONCLUSIONS –

high level



Databases, Map/reduce

Data center monitoring

Big data / analytics

Graphs/ Social net

Cyber-security

Fraud detection

Environmental data monitoring

Health db



C. Faloutsos (CMU + Google)



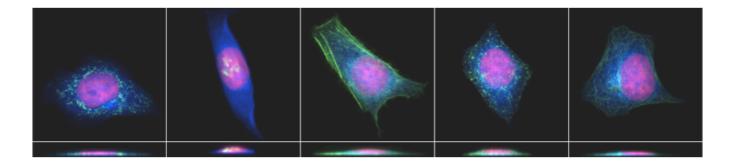






## Biological DB (w/ Bob Murphy, CMU)





#### Protein localization images:

- find similar ones
- do classification
- do clustering; anomaly detection

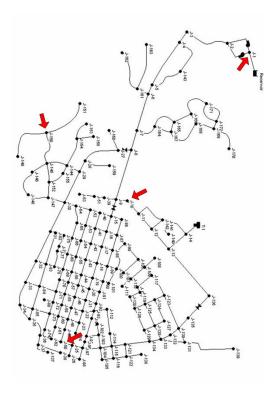
Environmental

data monitoring

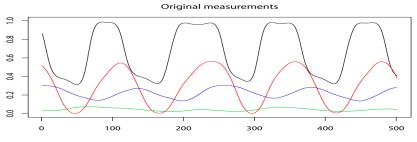
Drinking water monitoring (w/ Jeanne Vanbriesen,

Civ. Eng./CMU)





- Q1: How can we summarize the chlorine-level sequences?
- Q2: Where to put sensors, to spot attacks?



C. Faloutsos (CMU + Google)



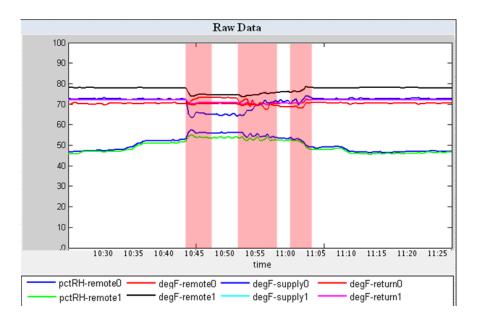


### monitoring Data center mon. & control



• (w/ Greg Ganger, CMU; A. Terzis, JHU)





 Abnormal dehumidification and reheating cycle is identified



• Leman Akoglu, Christos Faloutsos: *RTG: A Recursive Realistic Graph Generator Using Random Typing*. ECML/PKDD (1) 2009: 13-28

• Deepayan Chakrabarti, Christos Faloutsos: *Graph mining: Laws, generators, and algorithms*. ACM Comput. Surv. 38(1): (2006)



- Deepayan Chakrabarti, Yang Wang, Chenxi Wang, Jure Leskovec, Christos Faloutsos: *Epidemic thresholds in real networks*. ACM Trans. Inf. Syst. Secur. 10(4): (2008)
- Deepayan Chakrabarti, Jure Leskovec, Christos
   Faloutsos, Samuel Madden, Carlos Guestrin, Michalis
   Faloutsos: *Information Survival Threshold in Sensor* and P2P Networks. INFOCOM 2007: 1316-1324



• Christos Faloutsos, Tamara G. Kolda, Jimeng Sun: *Mining large graphs and streams using matrix and tensor tools*. Tutorial, SIGMOD Conference 2007: 1174



• T. G. Kolda and J. Sun. *Scalable Tensor Decompositions for Multi-aspect Data Mining*. In: ICDM 2008, pp. 363-372, December 2008.



- Jure Leskovec, Jon Kleinberg and Christos Faloutsos *Graphs over Time: Densification Laws, Shrinking Diameters and Possible Explanations*, KDD 2005 (Best Research paper award).
- Jure Leskovec, Deepayan Chakrabarti, Jon M. Kleinberg, Christos Faloutsos: *Realistic, Mathematically Tractable Graph Generation and Evolution, Using Kronecker Multiplication*. PKDD 2005: 133-145



- Jimeng Sun, Yinglian Xie, Hui Zhang, Christos Faloutsos. *Less is More: Compact Matrix Decomposition for Large Sparse Graphs*, SDM, Minneapolis, Minnesota, Apr 2007.
- Jimeng Sun, Spiros Papadimitriou, Philip S. Yu, and Christos Faloutsos, *GraphScope: Parameter-free Mining of Large Time-evolving Graphs* ACM SIGKDD Conference, San Jose, CA, August 2007



• Jimeng Sun, Dacheng Tao, Christos Faloutsos: *Beyond streams and graphs: dynamic tensor analysis*. KDD 2006: 374-383



- Hanghang Tong, Christos Faloutsos, and Jia-Yu Pan, *Fast Random Walk with Restart and Its Applications*, ICDM 2006, Hong Kong.
- Hanghang Tong, Christos Faloutsos,
   Center-Piece Subgraphs: Problem
   Definition and Fast Solutions, KDD 2006,
   Philadelphia, PA



 Hanghang Tong, Christos Faloutsos, Brian Gallagher, Tina Eliassi-Rad: Fast best-effort pattern matching in large attributed graphs.
 KDD 2007: 737-746



## **Project info**

www.cs.cmu.edu/~pegasus



Chau, ∡Polo



Akoglu,

Leman





Koutra, Danae



McGlohon, Mary

Prakash, Aditya





Tong, Hanghang

**★** Out, next year

Thanks to: NSF IIS-0705359, IIS-0534205, CTA-INARC; Yahoo (M45), LLNL, IBM, SPRINT, Google, INTEL, HP, iLab

## OVERALL CONCLUSIONS –

high level



Databases, Map/reduce

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C. Faloutsos (CMU + Google)