

Large Graph Mining - Patterns, Explanations and ~~Cascade Analysis~~

Christos Faloutsos

CMU

Thank you!

- Foster Provost
- Sinan Aral
- Arun Sundararajan

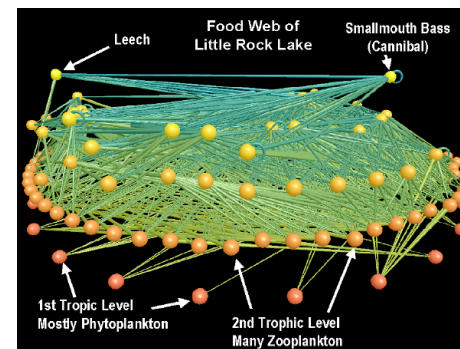
- Shirley Lau
- Sara Gorecki

Graphs - why should we care?

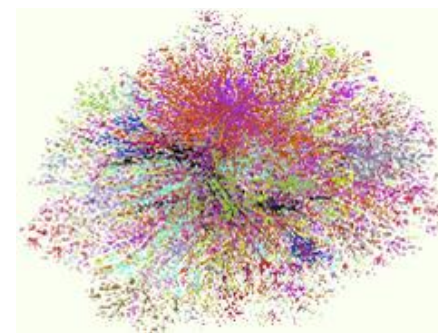


>\$10B revenue

>0.5B users



Food Web
[Martinez '91]



Internet Map
[lumeta.com]

Roadmap

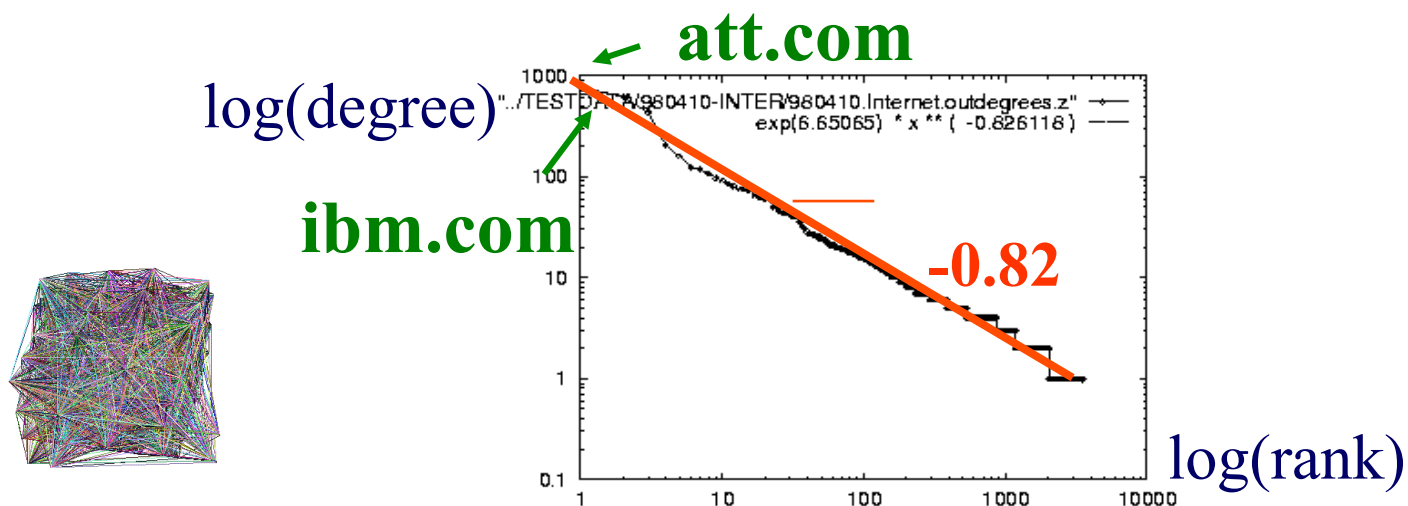


- Introduction – Motivation
- ➔ • Part#1: Patterns in graphs
 - Some (power) laws
 - The 'no good cuts' shock
 - A possible explanation: fractals
- [Part#2: Cascade analysis]
- Conclusions

Solution# S.1

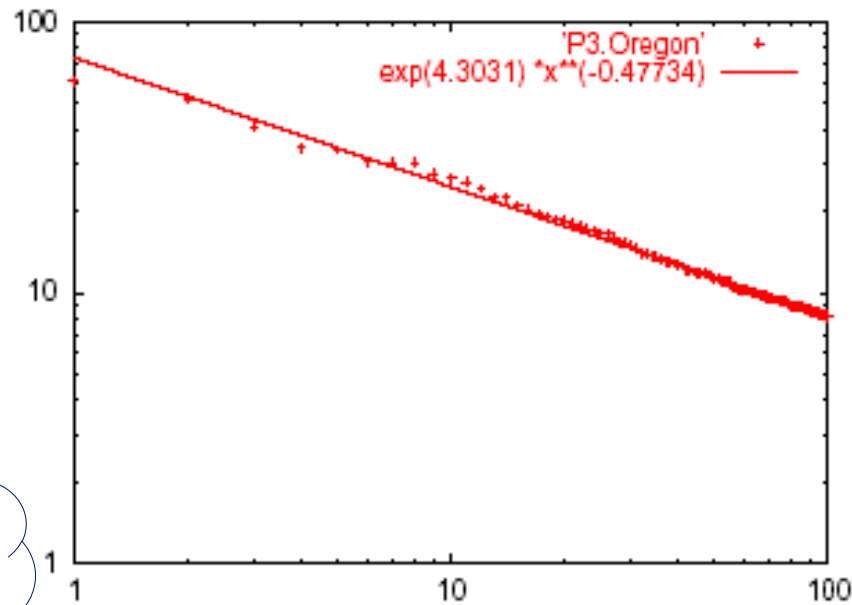
- Power law in the degree distribution [SIGCOMM99]

internet domains



Solution# S.2: Eigen Exponent E

Eigenvalue



Exponent = slope

$$E = -0.48$$

May 2001

$$\mathbf{A} \mathbf{x} = \lambda \mathbf{x}$$

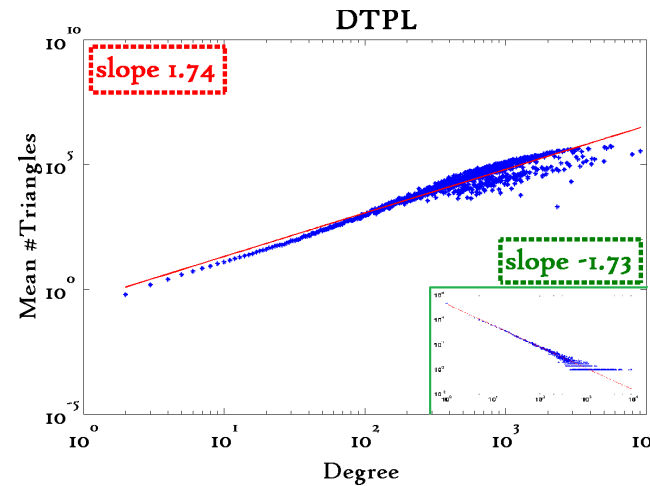
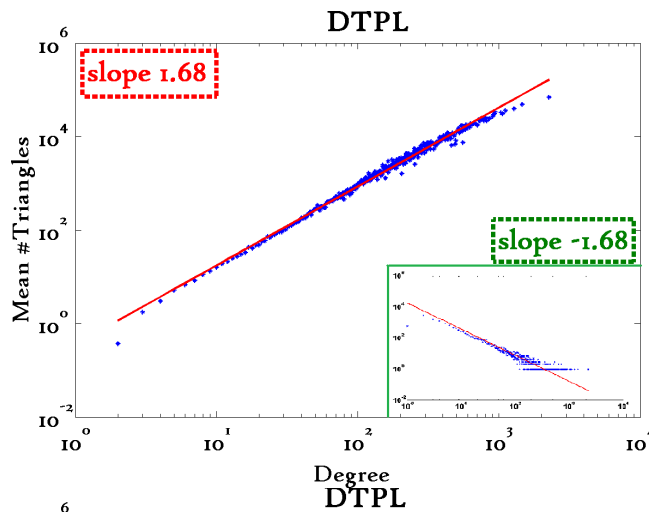
Rank of decreasing eigenvalue

- A2: power law in the eigenvalues of the adjacency matrix

Triangle Law: #S.3

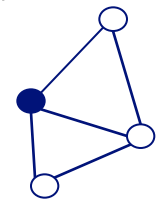
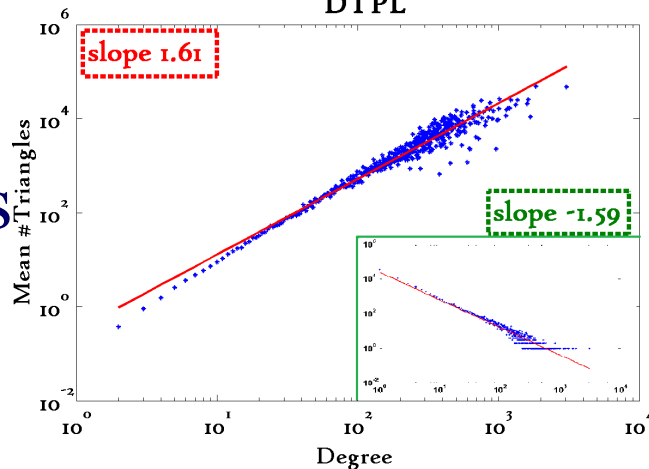
[Tsourakakis ICDM 2008]

Reuters



SN

Epinions



X-axis: degree
 Y-axis: mean # triangles
 n friends $\rightarrow \sim n^{1.6}$ triangles



MORE Graph Patterns

	Unweighted	Weighted
Static	<p>L01. Power-law degree distribution [Faloutsos et al. '99, Kleinberg et al. '99, Chakrabarti et al. '04, Newman '04]</p> <p>L02. Triangle Power Law (TPL) [Tsourakakis '08]</p> <p>L03. Eigenvalue Power Law (EPL) [Siganos et al. '03]</p> <p>L04. Community structure [Flake et al. '02, Girvan and Newman '02]</p>	<p>L10. Snapshot Power Law (SPL) [McGlohon et al. '08]</p>
Dynamic	<p>L05. Densification Power Law (DPL) [Leskovec et al. '05]</p> <p>L06. Small and shrinking diameter [Albert and Barabási '99, Leskovec et al. '05]</p> <p>L07. Constant size 2nd and 3rd connected components [McGlohon et al. '08]</p> <p>L08. Principal Eigenvalue Power Law (λ_1PL) [Akoglu et al. '08]</p> <p>L09. Bursty/self-similar edge/weight additions [Gomez and Santonja '98, Gribble et al. '98, Crovella and</p>	<p>L11. Weight Power Law (WPL) [McGlohon et al. '08]</p>

RTG: A Recursive Realistic Graph Generator using Random Typing Leman Akoglu and Christos Faloutsos. *PKDD'09*.



MORE Graph Patterns

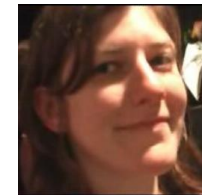
	Unweighted	Weighted
Static	<p>L02. Power-law degree distribution [Faloutsos et al. '99, Kleinberg et al. '99, Chakrabarti et al. '04, Newman '04]</p> <p>L03. Triangle Power Law (TPL) [Tsourakakis '08]</p> <p>L03. Eigenvalue Power Law (EPL) [Siganos et al. '03]</p> <p>L04. Community structure [Flake et al. '02, Girvan and Newman '02]</p>	<p>L05. Snapshot Power Law (SPL) [McGlohon et al. '08]</p>
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MORE Graph Patterns

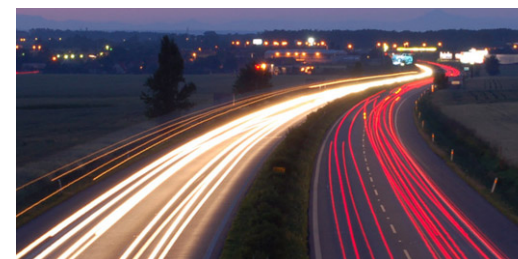
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- Mary McGlohon, Leman Akoglu, Christos Faloutsos. *Statistical Properties of Social Networks*. in "Social Network Data Analytics" (Ed.: Charu Aggarwal)
- Deepayan Chakrabarti and Christos Faloutsos, [*Graph Mining: Laws, Tools, and Case Studies*](#) Oct. 2012, Morgan Claypool.



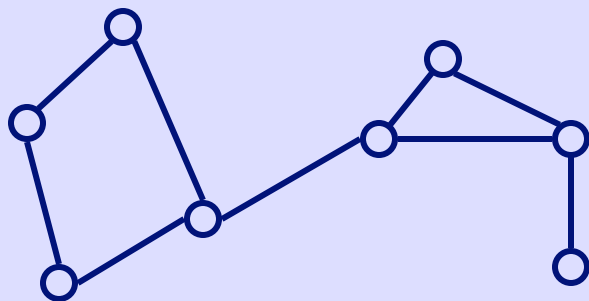
Roadmap

- Introduction – Motivation
- Part#1: Patterns in graphs
 - Some (power) laws
 - ➔ – The 'no good cuts' shock
 - A possible explanation: fractals
- Part#2: Cascade analysis
- Conclusions



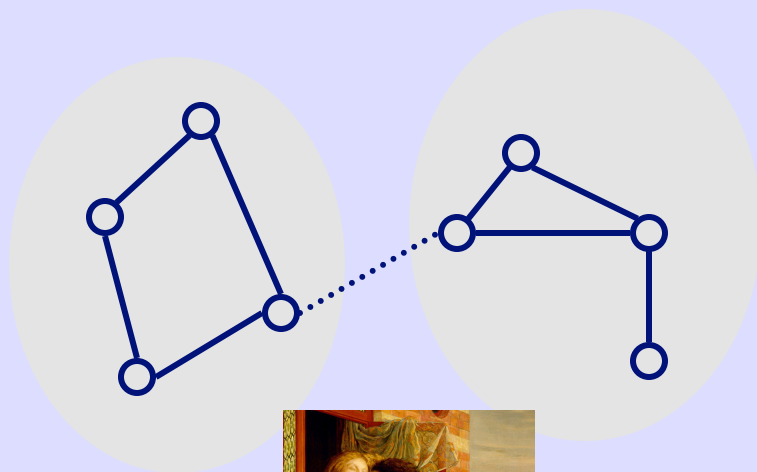
Background: Graph cut problem

- Given a graph, and k
- Break it into k (disjoint) communities

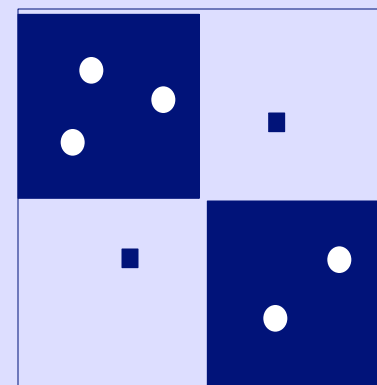


Graph cut problem

- Given a graph, and k
- Break it into k (disjoint) communities
- (assume: block diagonal = ‘cavemen’ graph)



$$k = 2$$



Many algo's for graph partitioning

- METIS [Karypis, Kumar +]
- 2nd eigenvector of Laplacian
- Modularity-based [Girwan+Newman]
- Max flow [Flake+]
- ...
- ...
- ...



Strange behavior of min cuts

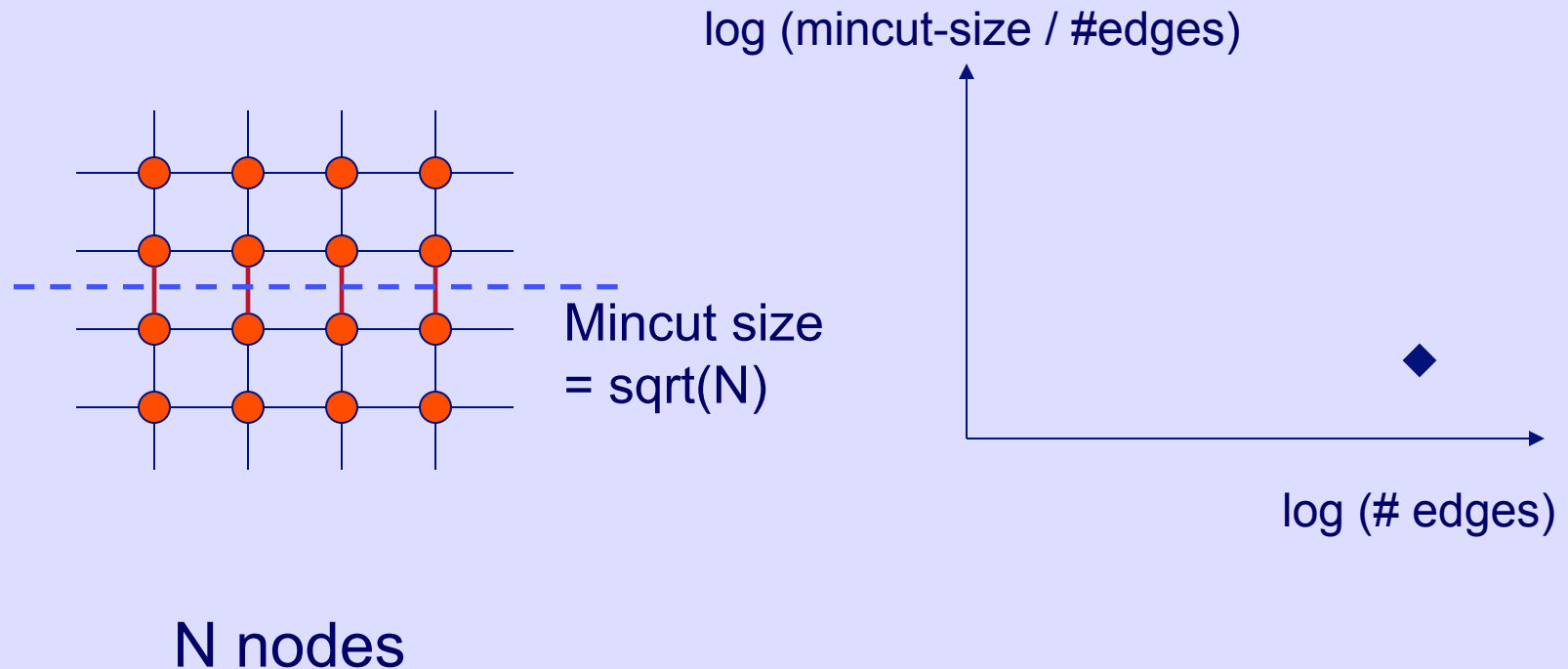
- Subtle details: next
 - Preliminaries: min-cut plots of ‘usual’ graphs

NetMine: New Mining Tools for Large Graphs, by D. Chakrabarti, Y. Zhan, D. Blandford, C. Faloutsos and G. Blelloch, in the SDM 2004 Workshop on Link Analysis, Counter-terrorism and Privacy

Statistical Properties of Community Structure in Large Social and Information Networks, J. Leskovec, K. Lang, A. Dasgupta, M. Mahoney. WWW 2008.

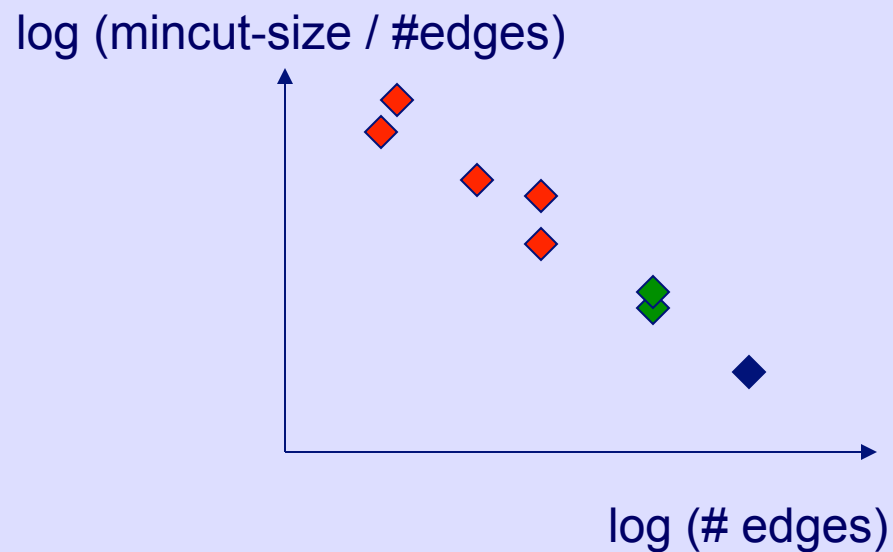
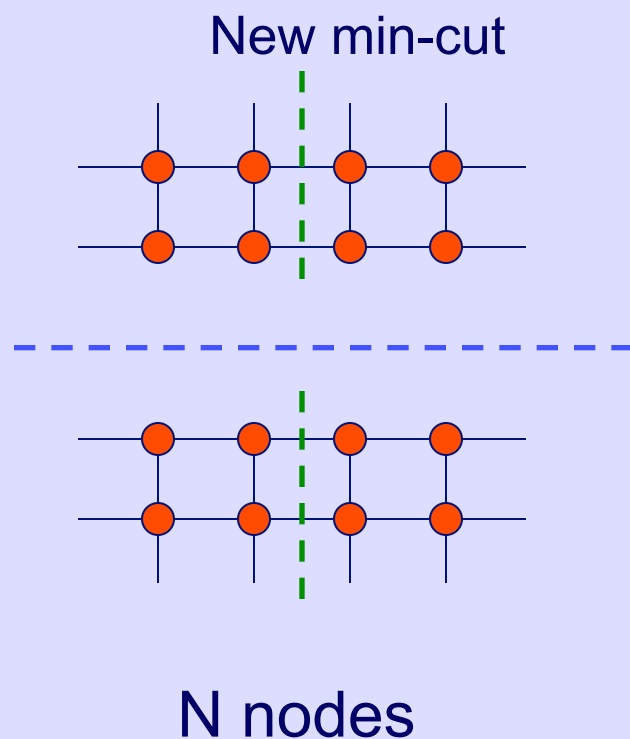
“Min-cut” plot

- Do min-cuts recursively.



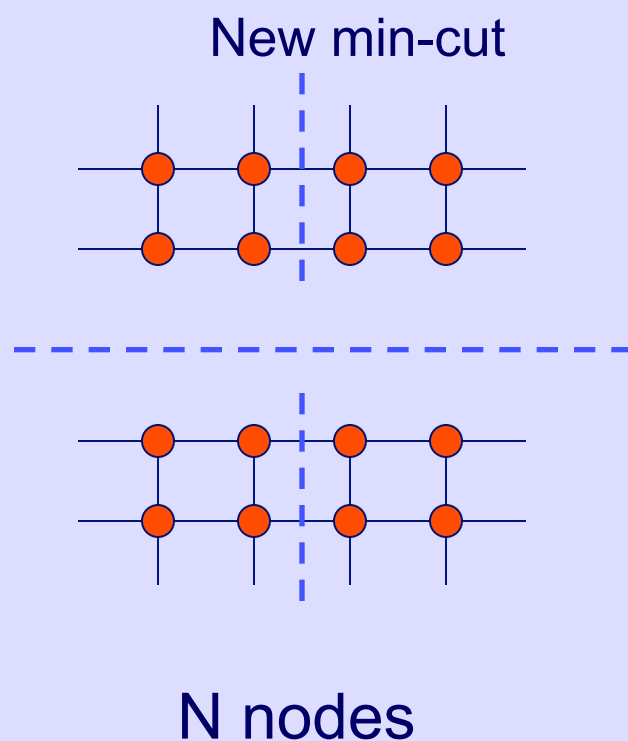
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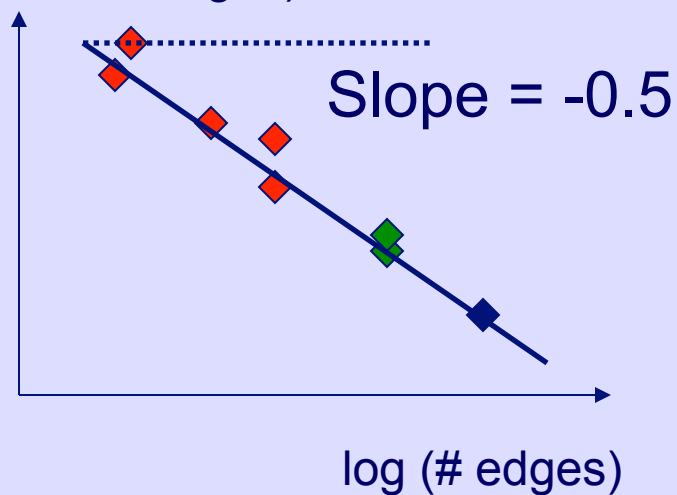
“Min-cut” plot

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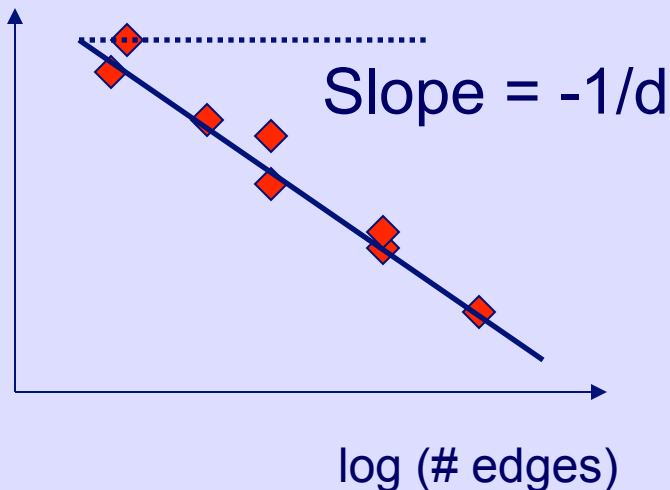
$\log(\text{mincut-size} / \#\text{edges})$

↓
Better
cut

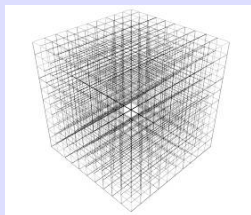


“Min-cut” plot

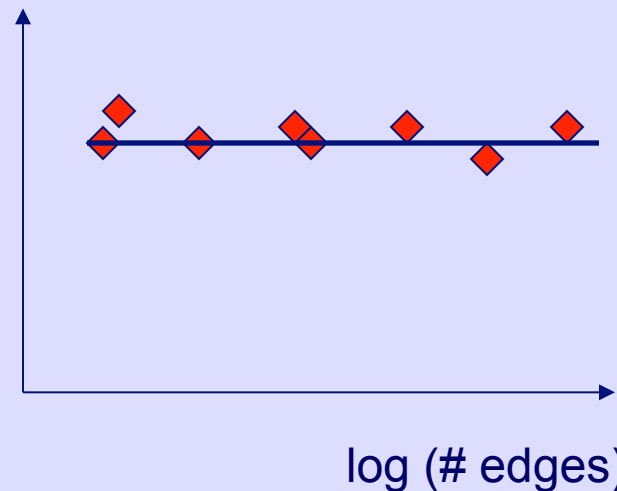
log (mincut-size / #edges)



For a d -dimensional grid, the slope is $-1/d$



log (mincut-size / #edges)



For a random graph
(and clique),
the slope is 0

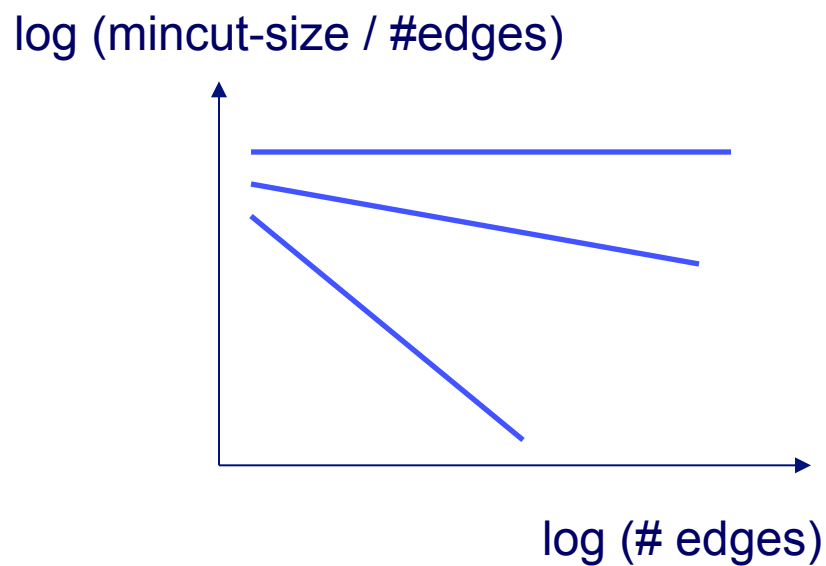
Experiments

- Datasets:
 - **Google Web Graph**: 916,428 nodes and 5,105,039 edges
 - **Lucent Router Graph**: Undirected graph of network routers from www.isi.edu/scan/mercator/maps.html; 112,969 nodes and 181,639 edges
 - **User → Website Clickstream Graph**: 222,704 nodes and 952,580 edges

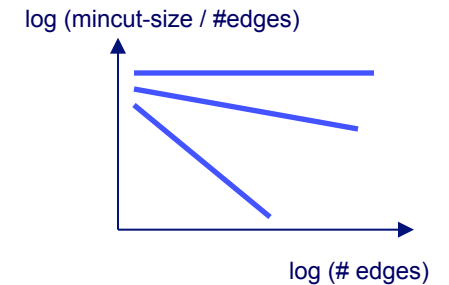
NetMine: New Mining Tools for Large Graphs, by D. Chakrabarti, Y. Zhan, D. Blandford, C. Faloutsos and G. Blelloch, in the SDM 2004 Workshop on Link Analysis, Counter-terrorism and Privacy

“Min-cut” plot

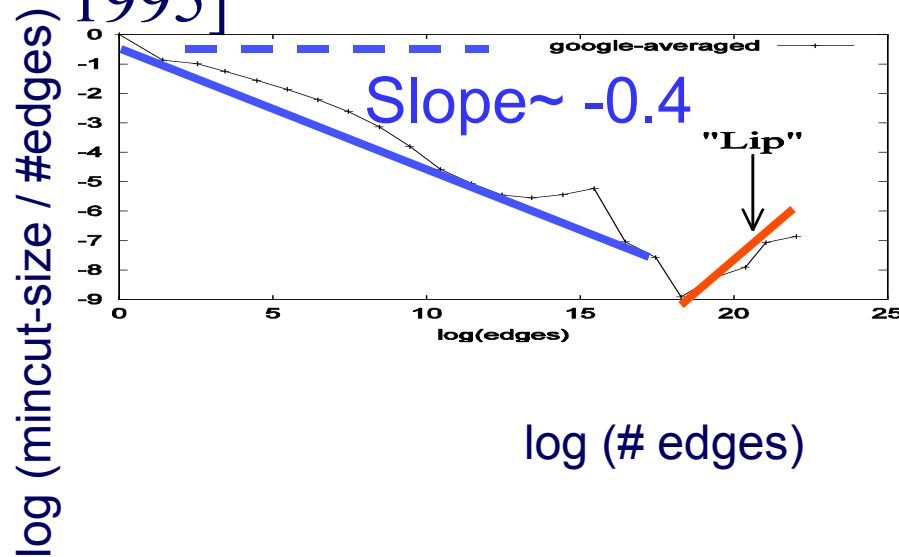
- What does it look like for a real-world graph?



Experiments

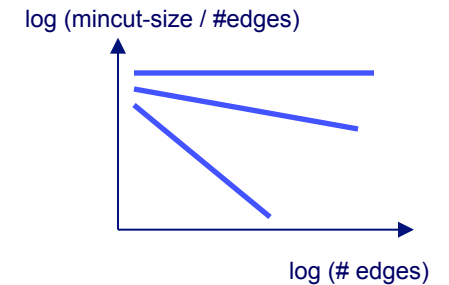


- Used the METIS algorithm [Karypis, Kumar, 1995]

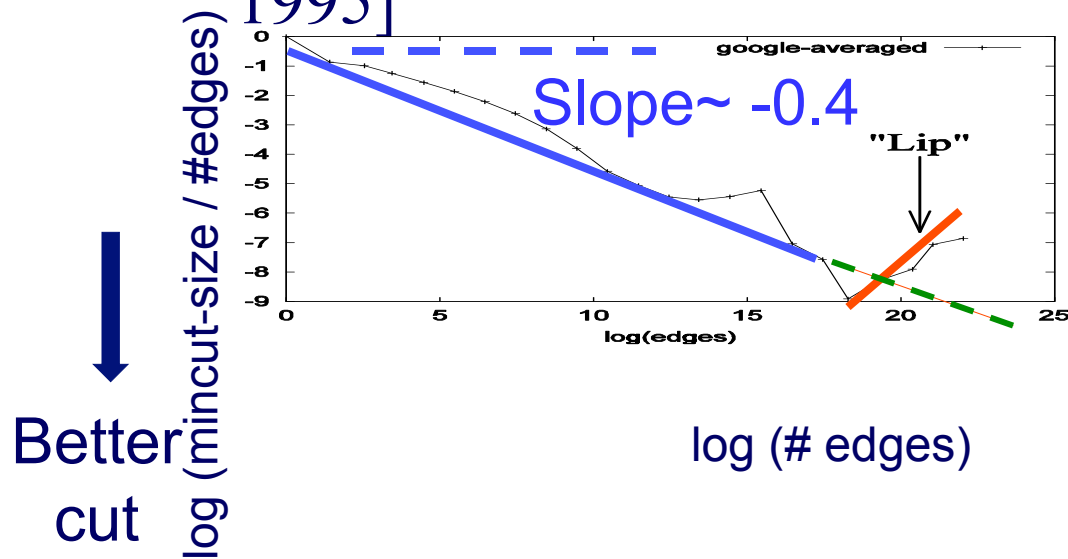


- Google Web graph
- Values along the y-axis are averaged
- “lip” for large # edges
- Slope of -0.4, corresponds to a 2.5-dimensional grid!

Experiments



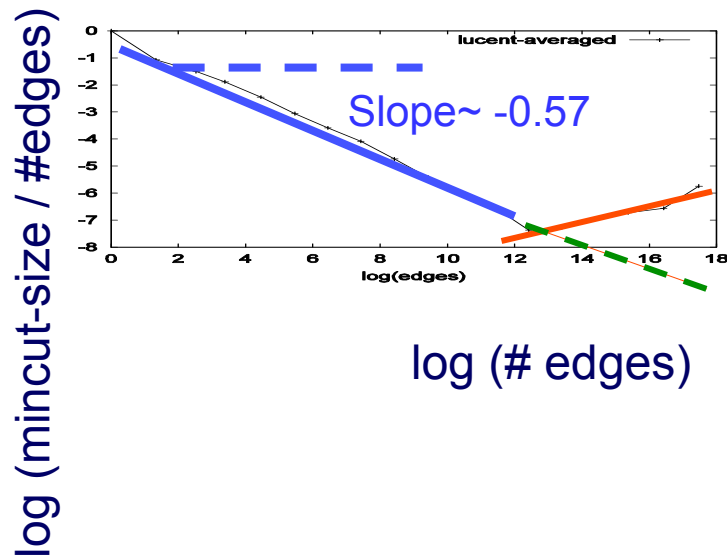
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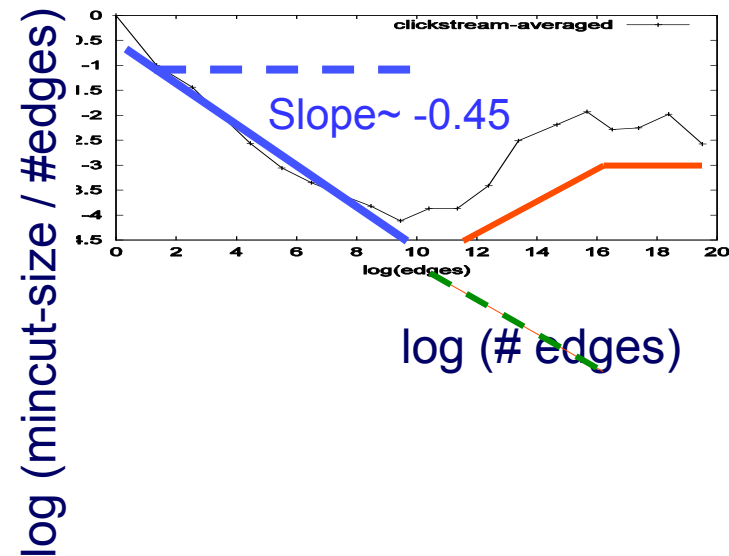
- Google Web graph
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Experiments

- Same results for other graphs too...



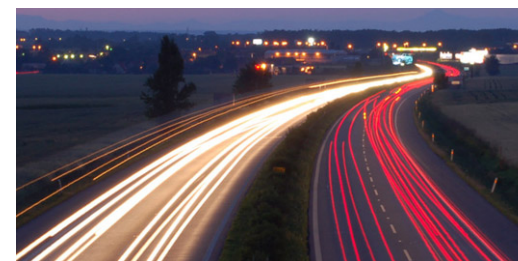
Lucent Router graph



Clickstream graph

Roadmap

- Introduction – Motivation
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 - ➔ – A possible explanation: fractals
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2 Questions, one answer

- Q1: why so many power laws
- Q2: why no ‘good cuts’?

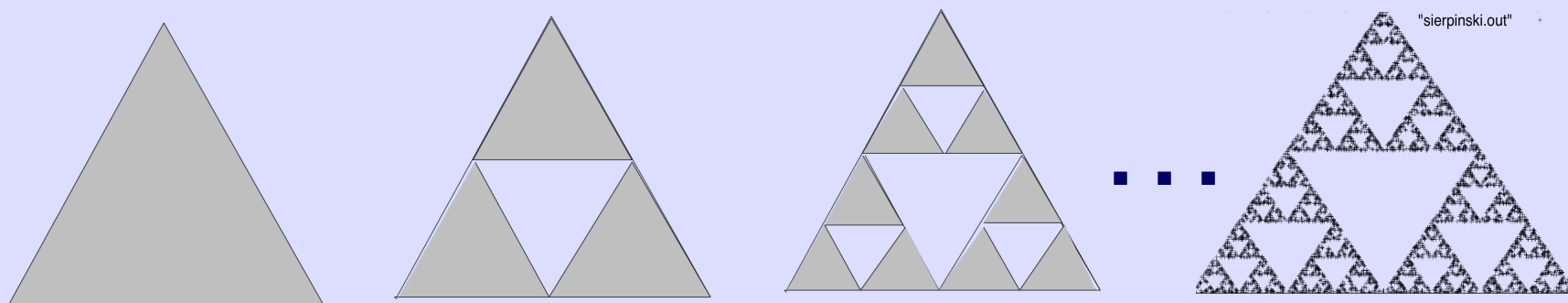
possible

2 Questions, one answer

- Q1: why so many power laws
- Q2: why no ‘good cuts’?
- A: Self-similarity = fractals = ‘RMAT’ ~ ‘Kronecker graphs’

20'' intro to fractals

- Remove the middle triangle; repeat
- -> Sierpinski triangle
- (Bonus question - dimensionality?)
 - >1 (inf. perimeter – $(4/3)^\infty$)
 - <2 (zero area – $(3/4)^\infty$)



20'' intro to fractals

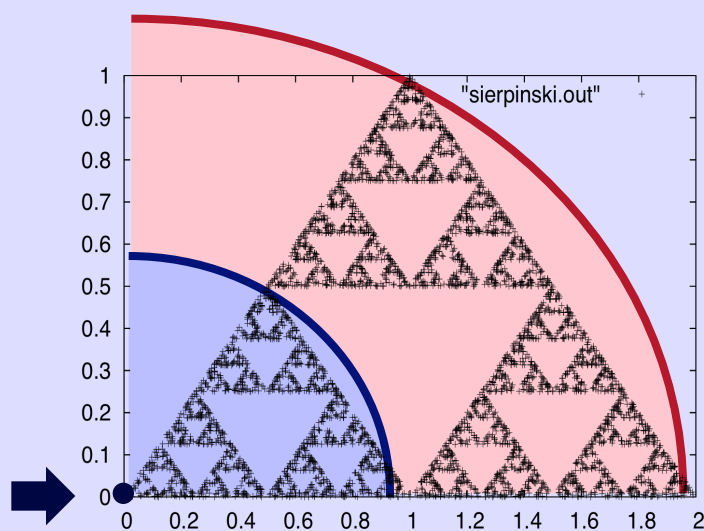
Self-similarity -> no char. scale

-> power laws, eg:

2x the radius,

3x the #neighbors $nn(r)$

$$nn(r) = C r^{\log 3 / \log 2}$$



20'' intro to fractals

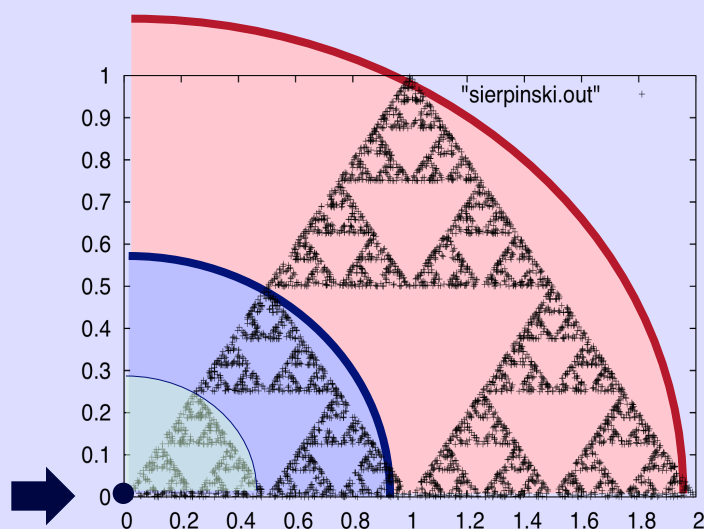
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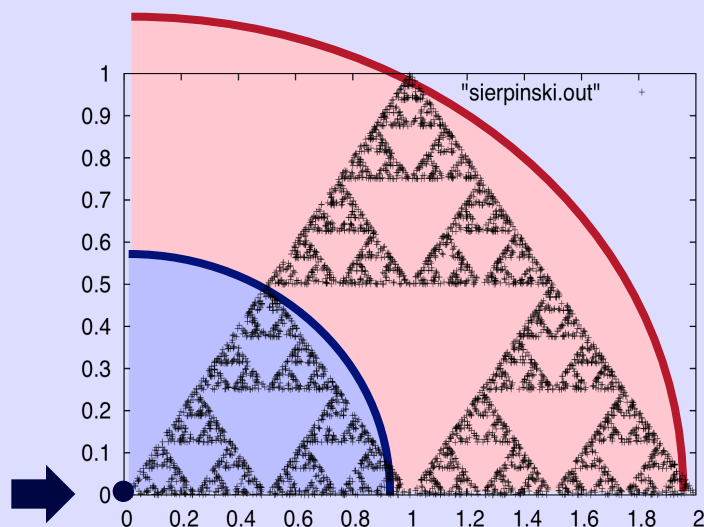
3x the #neighbors

$$nn = C r^{\log 3 / \log 2}$$

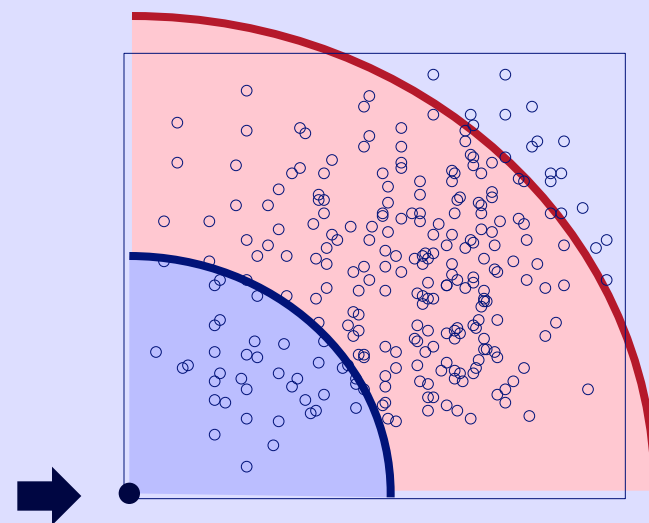
2x the radius,

4x neighbors

$$nn = C r^{\log 4 / \log 2} = C r^2$$



WIN workshop, NYU



(c) 2013, C. Faloutsos

20'' intro to fractals

Self-similarity -> no char. scale

-> power laws, eg:

2x the radius,

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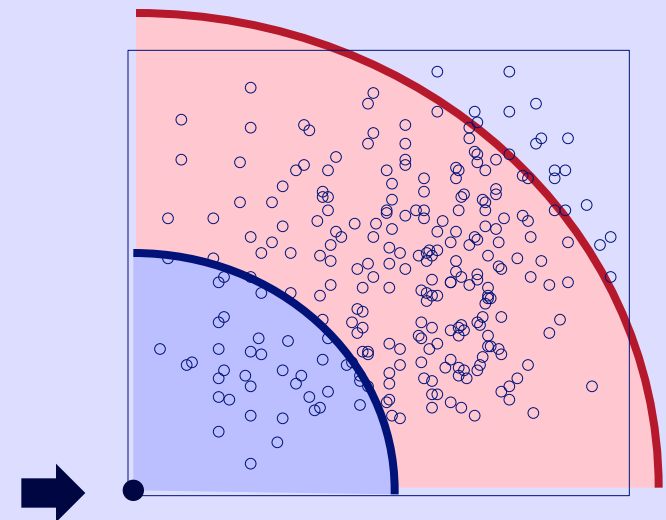
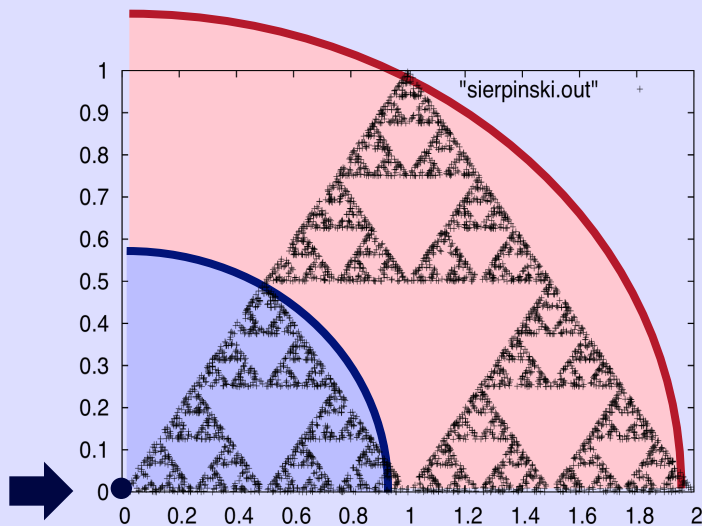
$$n_n = C r^{\log 3 / \log 2} \leftarrow = 1.58$$

2x the radius,

4x neighbors

$$n_n = C r^{\log 4 / \log 2} = C r^2$$

Fractal dim.



20'' intro to fractals

Self-similarity -> no char. scale

-> **power laws**, eg:

2x the radius,

3x the #neighbors

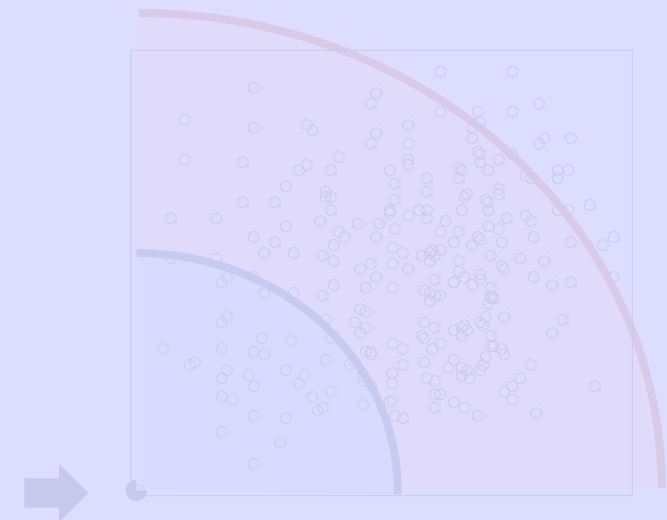
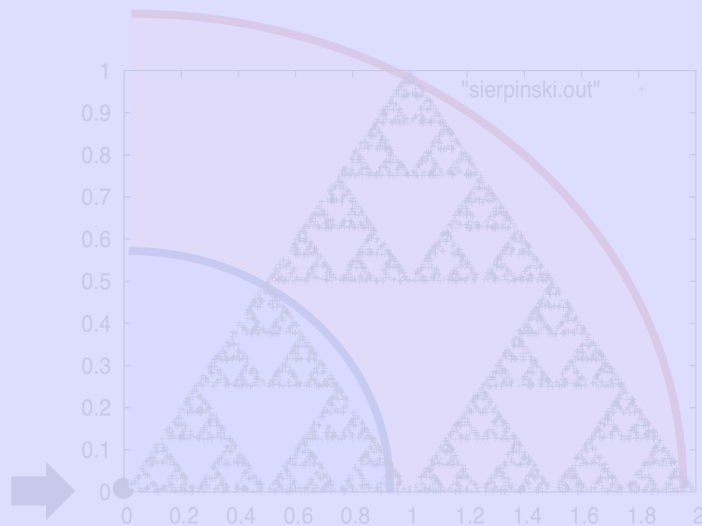
$$n_n = C r^{\log 3 / \log 2}$$

2x the radius,

4x neighbors

$$n_n = C r^{\log 4 / \log 2} = C r^2$$

Fractal dim.



How does self-similarity help in graphs?

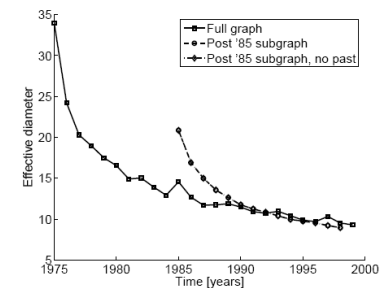
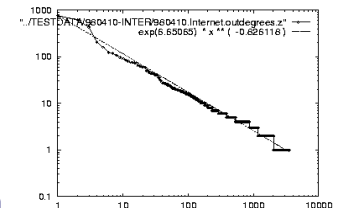
- A: RMAT/Kronecker generators
 - With self-similarity, we get all power-laws, automatically,
 - And small/shrinking diameter
 - And ‘no good cuts’

R-MAT: A Recursive Model for Graph Mining,
by D. Chakrabarti, Y. Zhan and C. Faloutsos,
SDM 2004, Orlando, Florida, USA

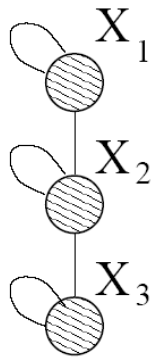
Realistic, Mathematically Tractable Graph Generation and Evolution, Using Kronecker Multiplication,
by J. Leskovec, D. Chakrabarti, J. Kleinberg,
and C. Faloutsos, in PKDD 2005, Porto, Portugal

Graph gen.: Problem defn

- Given a growing graph with count of nodes N_1 , N_2 , ...
- Generate a realistic sequence of graphs that will obey all the patterns
 - Static Patterns
 - S1 Power Law Degree Distribution
 - S2 Power Law eigenvalue and eigenvector distribution
 - Small Diameter
 - Dynamic Patterns
 - T2 Growth Power Law (2x nodes; 3x edges)
 - T1 Shrinking/Stabilizing Diameters



Kronecker Graphs

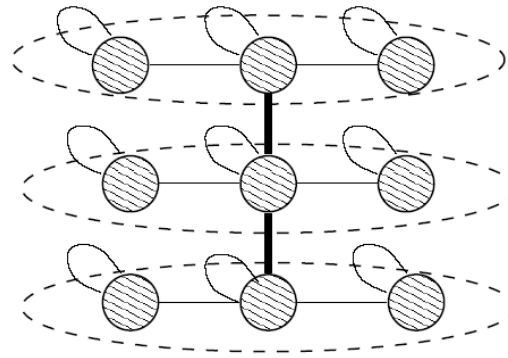
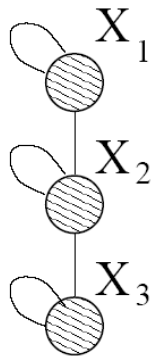


1	1	0
1	1	1
0	1	1

G_1

Adjacency matrix

Kronecker Graphs



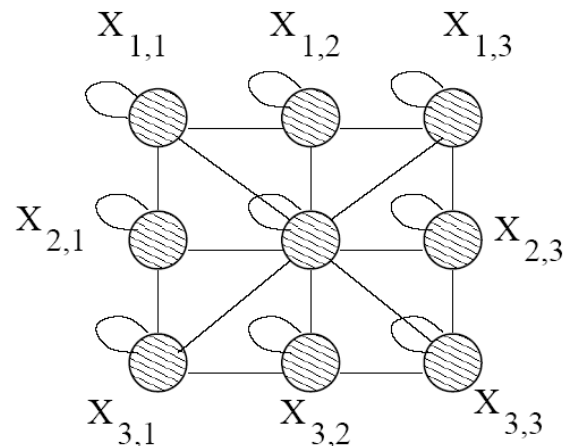
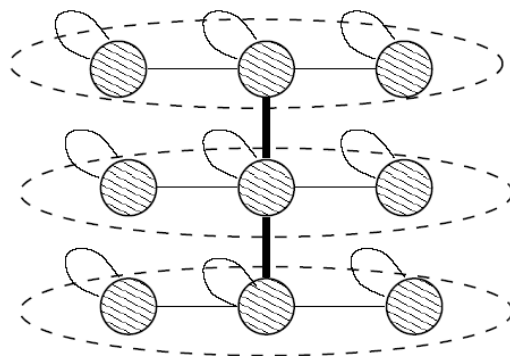
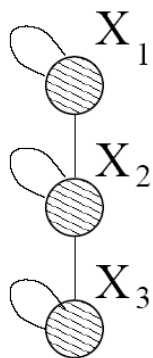
Intermediate stage

1	1	0
1	1	1
0	1	1

G_1

Adjacency matrix

Kronecker Graphs



Intermediate stage

1	1	0
1	1	1
0	1	1

G_1

Adjacency matrix

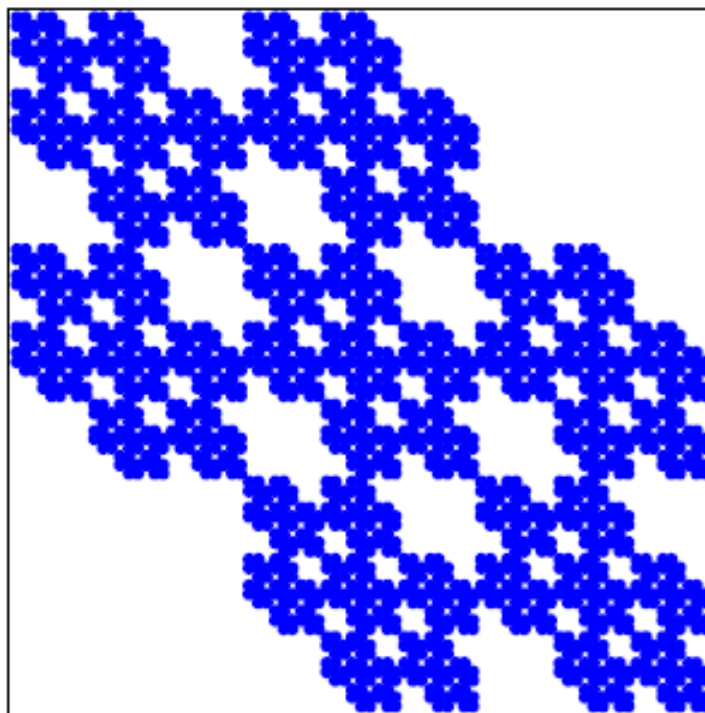
G_1	G_1	0
G_1	G_1	G_1
0	G_1	G_1

$G_2 = G_1 \otimes G_1$

Adjacency matrix

Kronecker Graphs

- Continuing multiplying with G_1 we obtain G_4 and so on ...

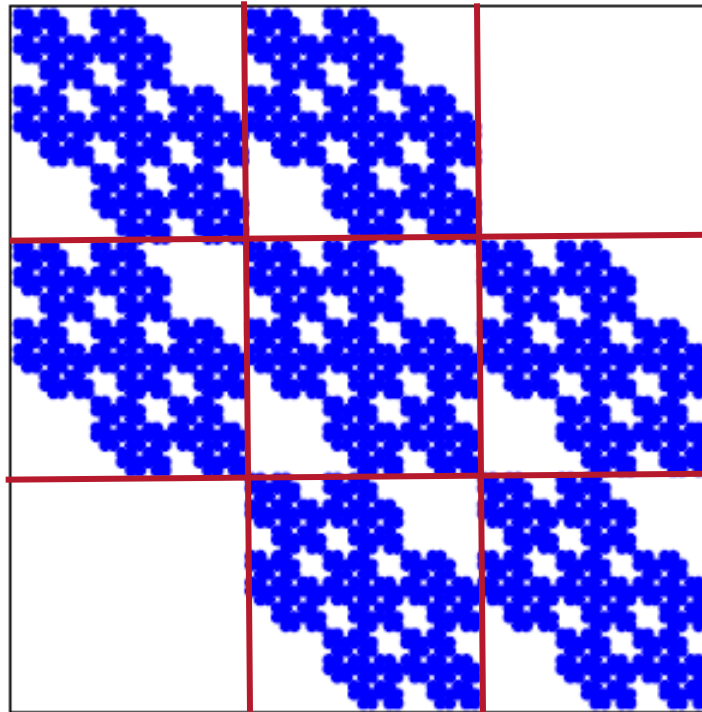


G_4 adjacency matrix

(c) 2013, C. Faloutsos

Kronecker Graphs

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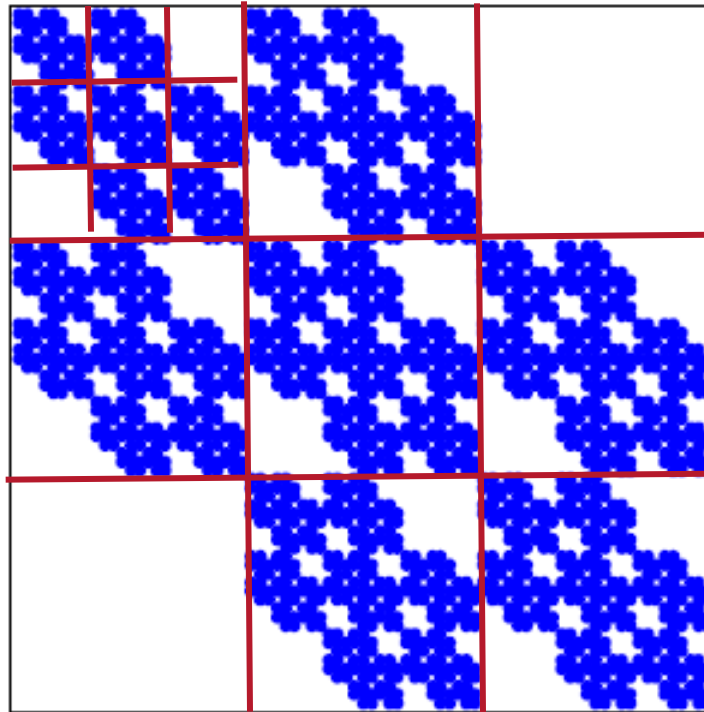


G_4 adjacency matrix

(c) 2013, C. Faloutsos

Kronecker Graphs

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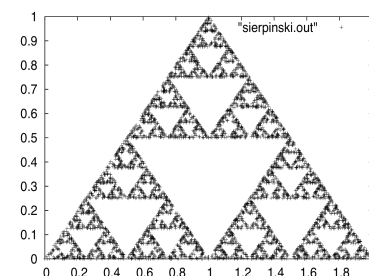
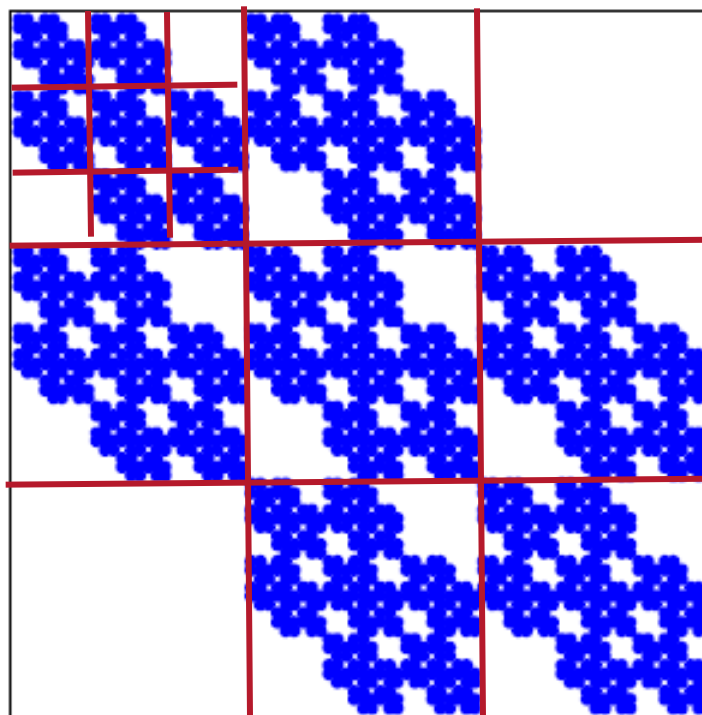
G_4 adjacency matrix

(c) 2013, C. Faloutsos

Kronecker Graphs

- Continuing multiplying with G_1 we obtain G_4 and so on ...

Holes within holes;
Communities
within communities



G_4 adjacency matrix

(c) 2013, C. Faloutsos

Problem Definition

- Given a growing graph with nodes N_1, N_2, \dots
- Generate a realistic sequence of graphs that will obey all the patterns
 - Static Patterns
 - ✓ Power Law Degree Distribution
 - ✓ Power Law eigenvalue and eigenvector distribution
 - ✓ Small Diameter
 - Dynamic Patterns
 - ✓ Growth Power Law
 - ✓ Shrinking/Stabilizing Diameters
- First generator for which we can **prove** all these properties

Impact: Graph500

- Based on RMAT (= 2x2 Kronecker)
- Standard for graph benchmarks
- <http://www.graph500.org/>
- Competitions 2x year, with all major entities: LLNL, Argonne, ITC-U. Tokyo, Riken, ORNL, Sandia, PSC, ...

To iterate is human, to recurse is devine

R-MAT: A Recursive Model for Graph Mining,
by D. Chakrabarti, Y. Zhan and C. Faloutsos,
SDM 2004, Orlando, Florida, USA

Roadmap

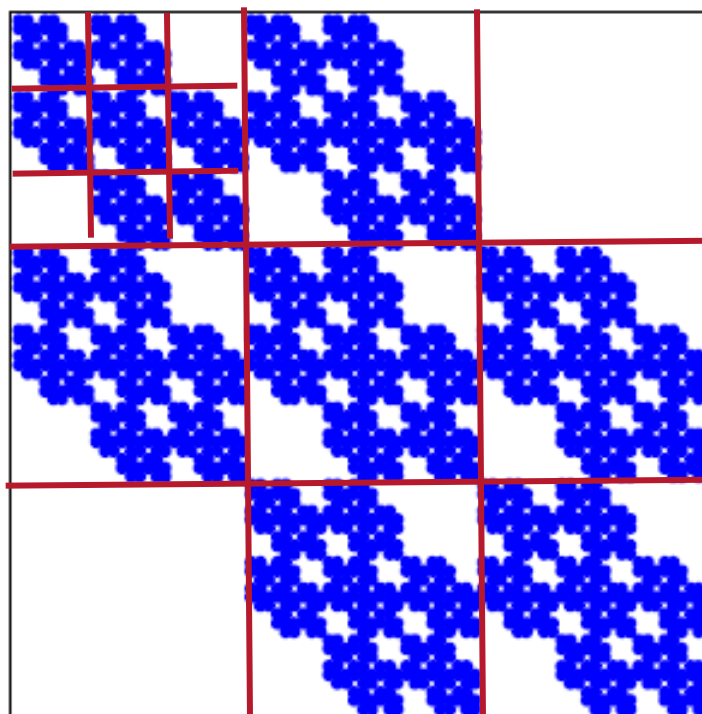
- Introduction – Motivation
- Part#1: Patterns in graphs
 - ...
 - Q1: Why so many power laws?
 - ➔ – Q2: Why no ‘good cuts’?
- Part#2: Cascade analysis
- Conclusions



A: real graphs ->
self similar ->
power laws

Kronecker Product – a Graph

- Continuing multiplying with G_1 we obtain G_4 and so on ...



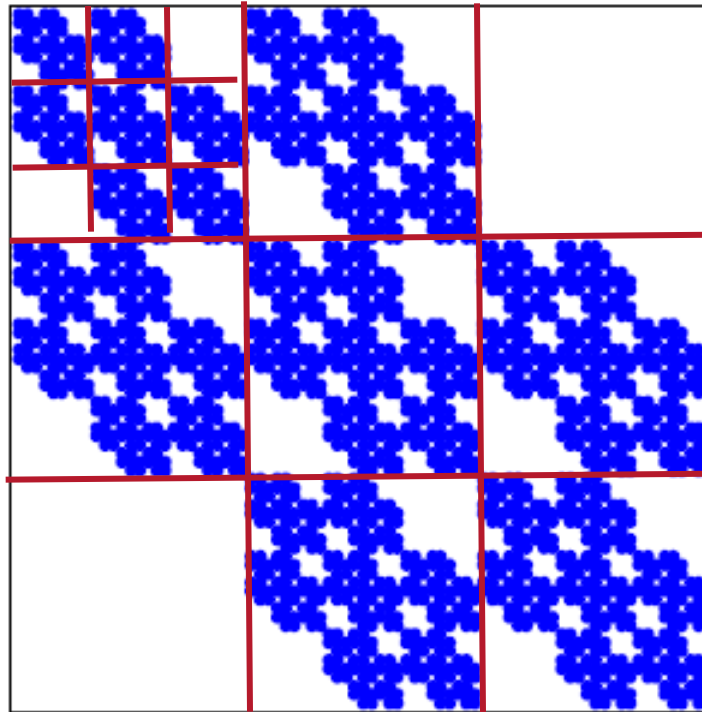
G_4 adjacency matrix

(c) 2013, C. Faloutsos

Kronecker Product – a Graph

- Continuing multiplying with G_1 we obtain G_4 and so on ...

Communities within communities within communities ...



How many
Communities?
3?
9?
27?

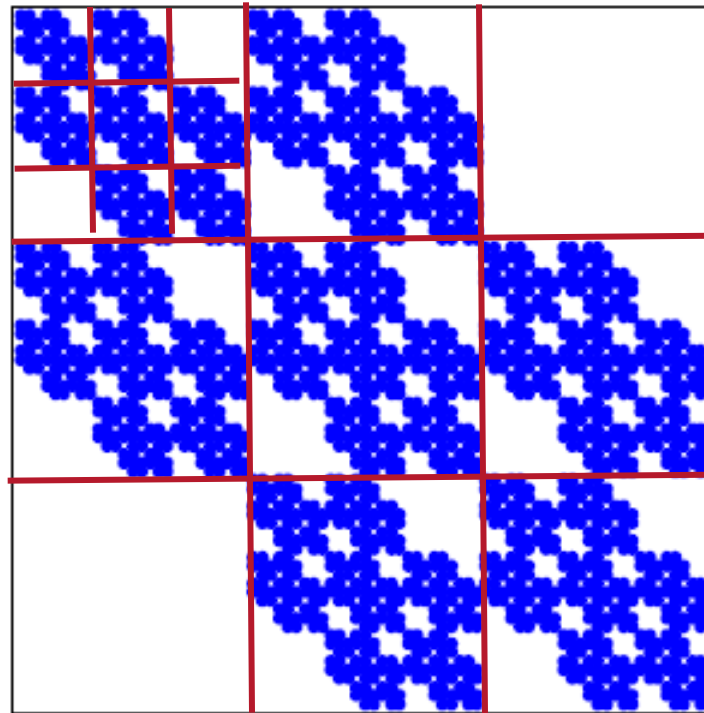
G_4 adjacency matrix

(c) 2013, C. Faloutsos

Kronecker Product – a Graph

- Continuing multiplying with G_1 we obtain G_4 and so on ...

Communities within communities within communities ...



G_4 adjacency matrix

(c) 2013, C. Faloutsos

How many
Communities?

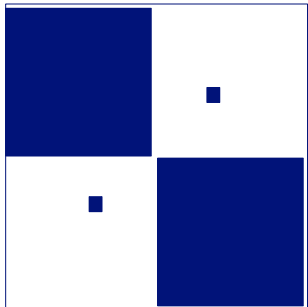
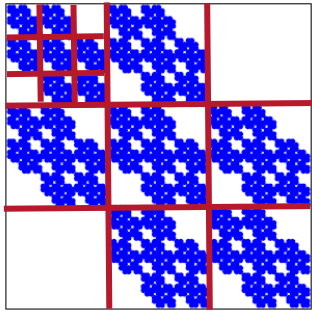
3?

9?

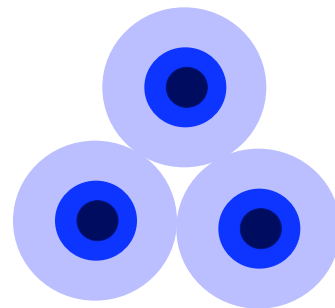
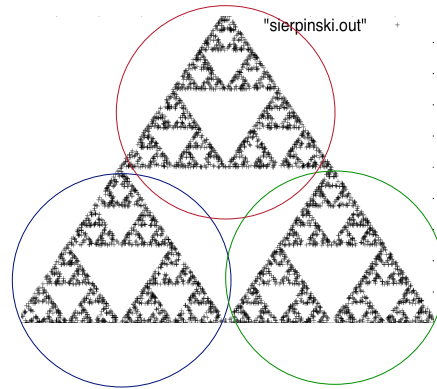
27?

A: one – but
not a typical,
block-like
community...

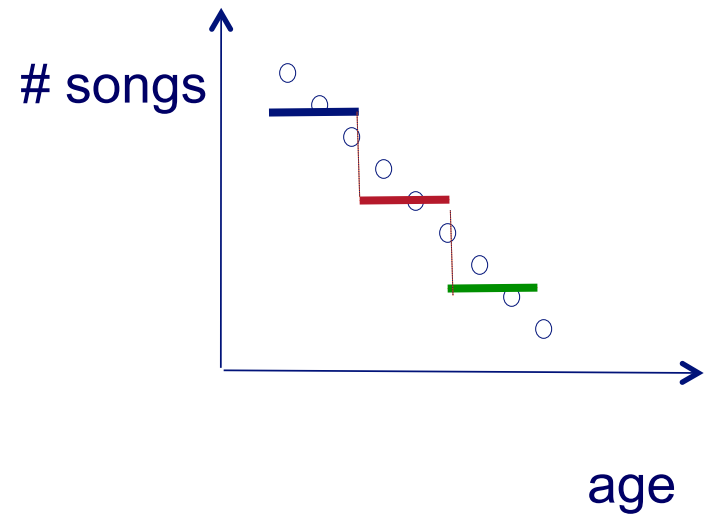
Communities?

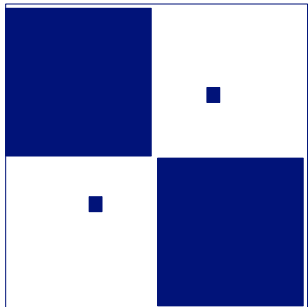
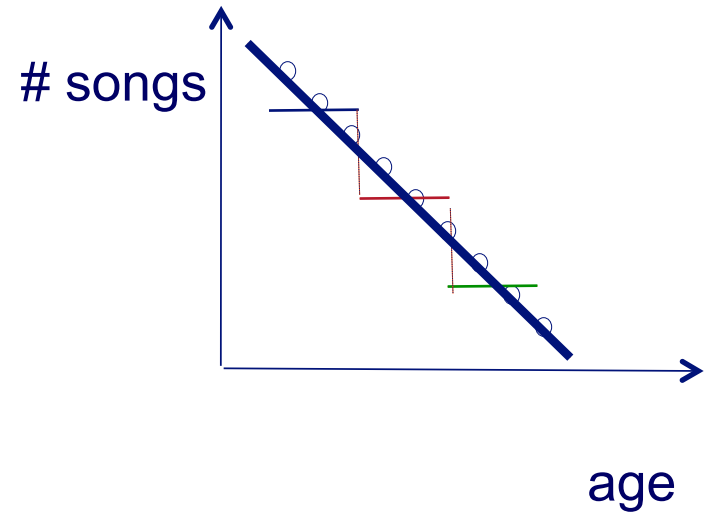
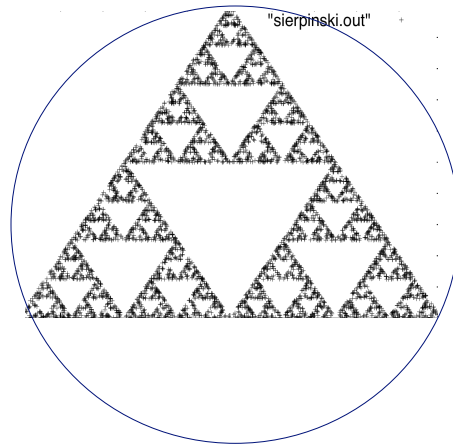
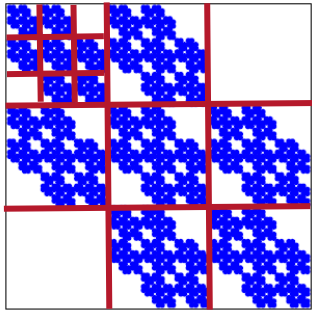


(Gaussian) Clusters?



Piece-wise flat parts?

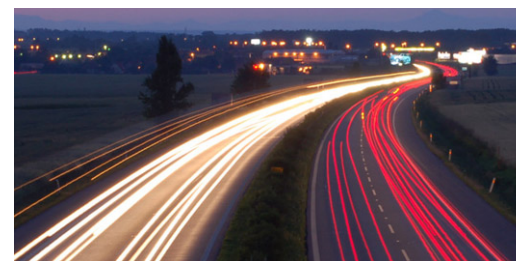




Wrong questions to ask!

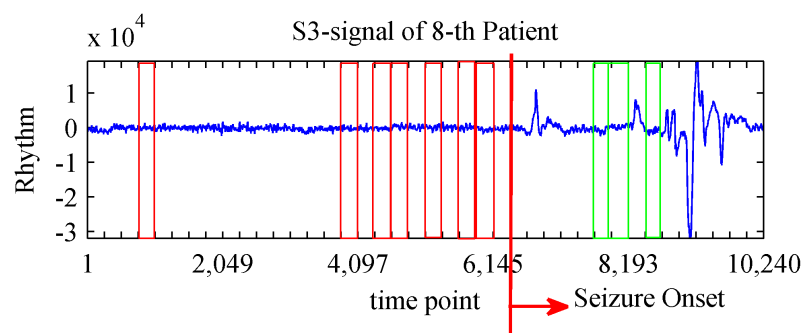
Roadmap

- Introduction – Motivation
- Part#1: Patterns in graphs
 - ...
 - Q1: The 'no good cuts' shock
 - Q2: Why no 'good cuts'?
- ➔ • What next?
- Conclusions



Challenge #1: ‘Connectome’ – brain wiring

- Which neurons get activated by ‘tomato’
- How wiring evolves
- Modeling epilepsy



Tom Mitchell



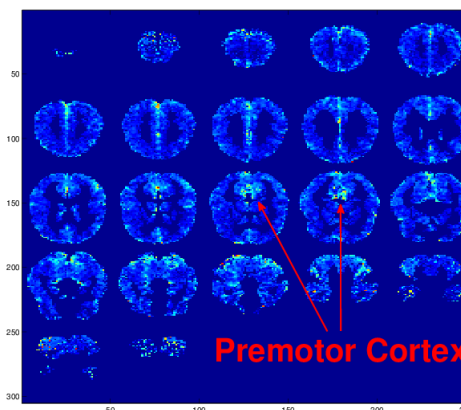
George Karypis



N. Sidiropoulos



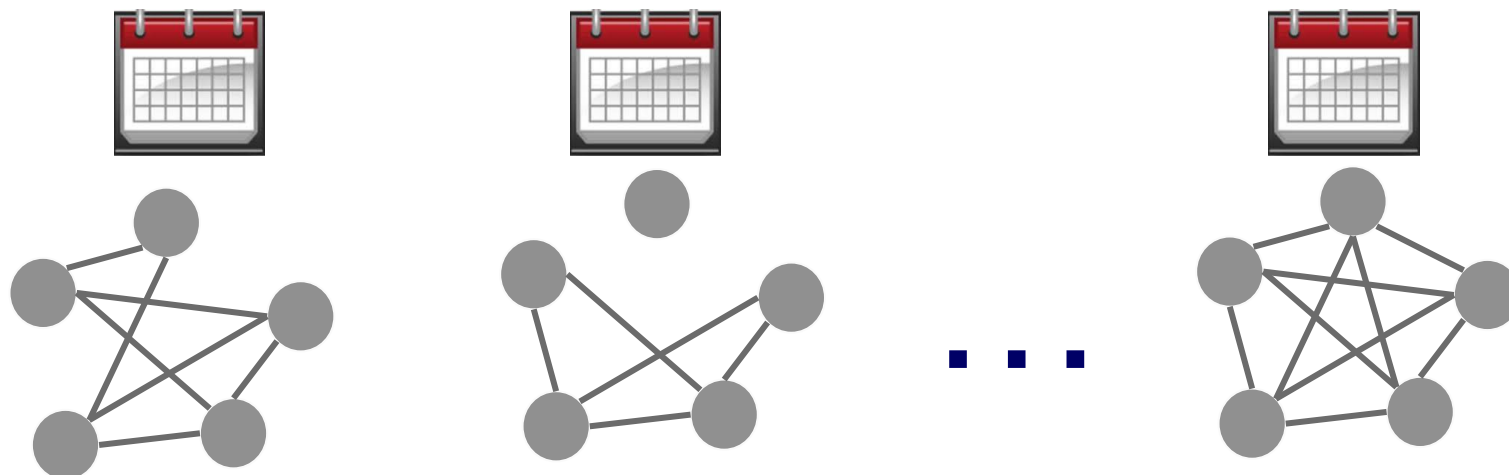
V. Papalexakis



‘glass’
‘tomato’
‘bell’

Challenge#2: Time evolving networks / tensors

- Periodicities? Burstiness?
- What is ‘typical’ behavior of a node, over time
- Heterogeneous graphs (= nodes w/ attributes)



Summary

- ***many*** patterns in real graphs
 - Power-laws everywhere
 - ‘no good cuts’
- Self-similarity (RMAT/Kronecker): good model

Thanks



Disclaimer: All opinions are mine; not necessarily reflecting the opinions of the funding agencies

Thanks to: NSF IIS-0705359, IIS-0534205, CTA-INARC; Yahoo (M45), LLNL, IBM, SPRINT, Google, INTEL, HP, iLab

Project info: PEGASUS



www.cs.cmu.edu/~pegasus

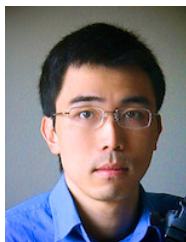
Results on large graphs: with Pegasus +
hadoop + M45

Apache license

Code, papers, manual, video



Prof. U Kang



Prof. Polo Chau

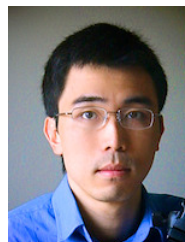
Cast



Akoglu,
Leman



Beutel,
Alex



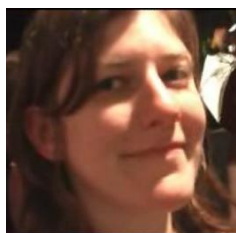
Chau,
Polo



Kang, U



Koutra,
Danai



McGlohon,
Mary



Prakash,
Aditya



Papalexakis,
Vagelis



Tong,
Hanghang

TAKE HOME MESSAGE:

Cross-disciplinary

