

P v.s. NP and Approximation Algorithms

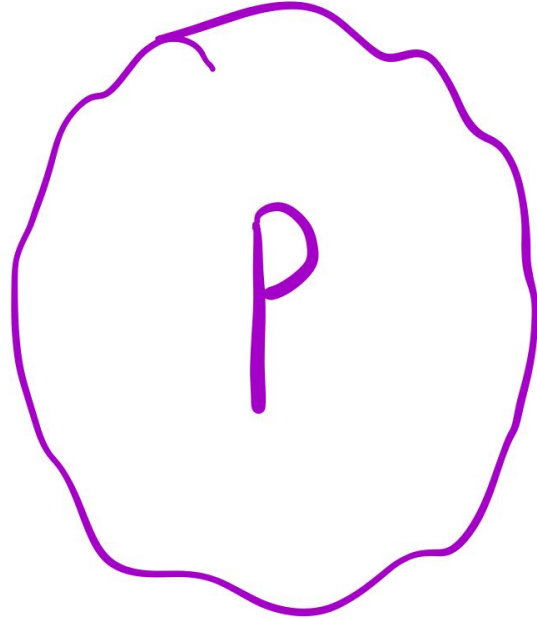
August 14th, 2021 (Class #6)

Our Best Results

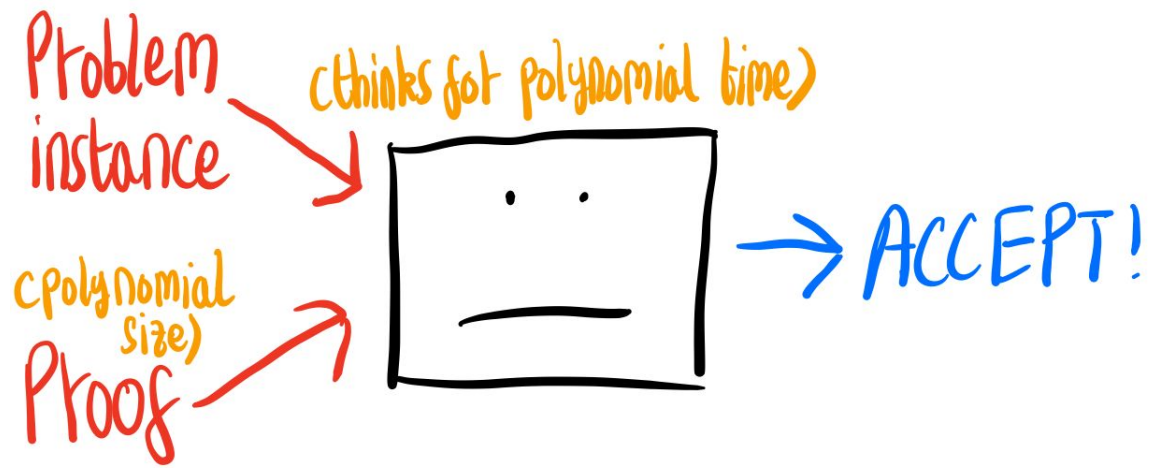
Problem	Running time
Multiplication	$O(n^{\log_2 3})$
Sorting	$O(n \log n)$
Unweighted SP	$O(n + m)$
Max Disj. Intervals	$O(n)$
Weighted SP	$O(m \log n)$


We've shown that all these problems are...

members of



What is NP?



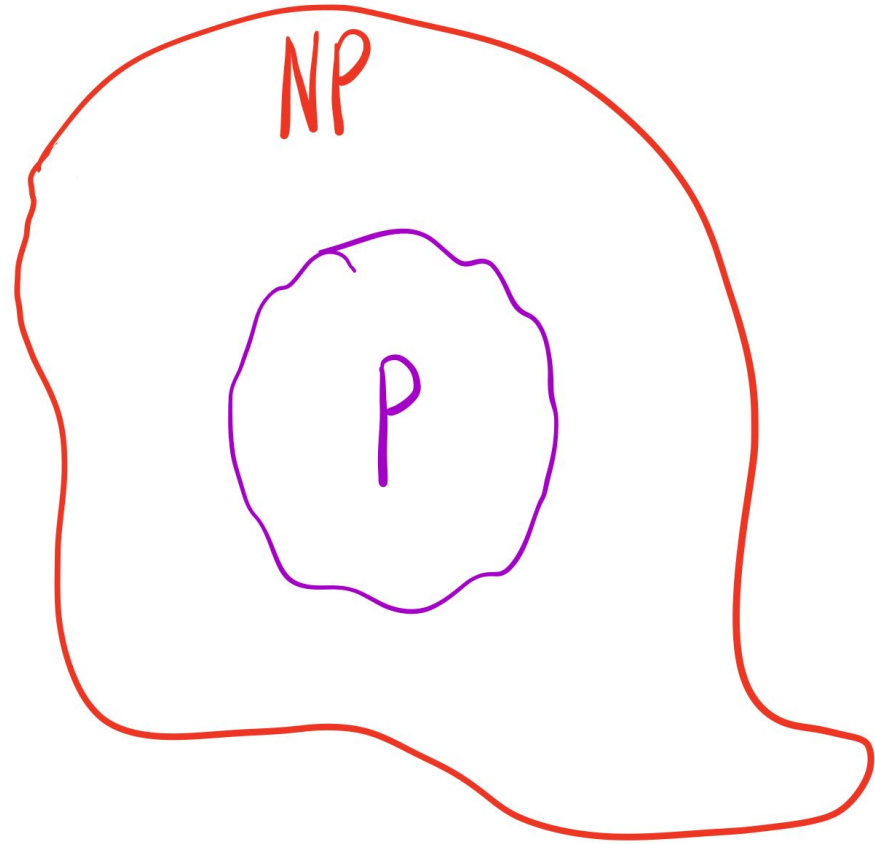
If the true answer is **YES**,
then  will **ACCEPT** for
some proof.

Otherwise, always **REJECT**.

The P v.s. NP Question

NP { If we can **verify** an
answer quickly, (check a proof)

P { Can we **find** it quickly?
(come up with a proof)

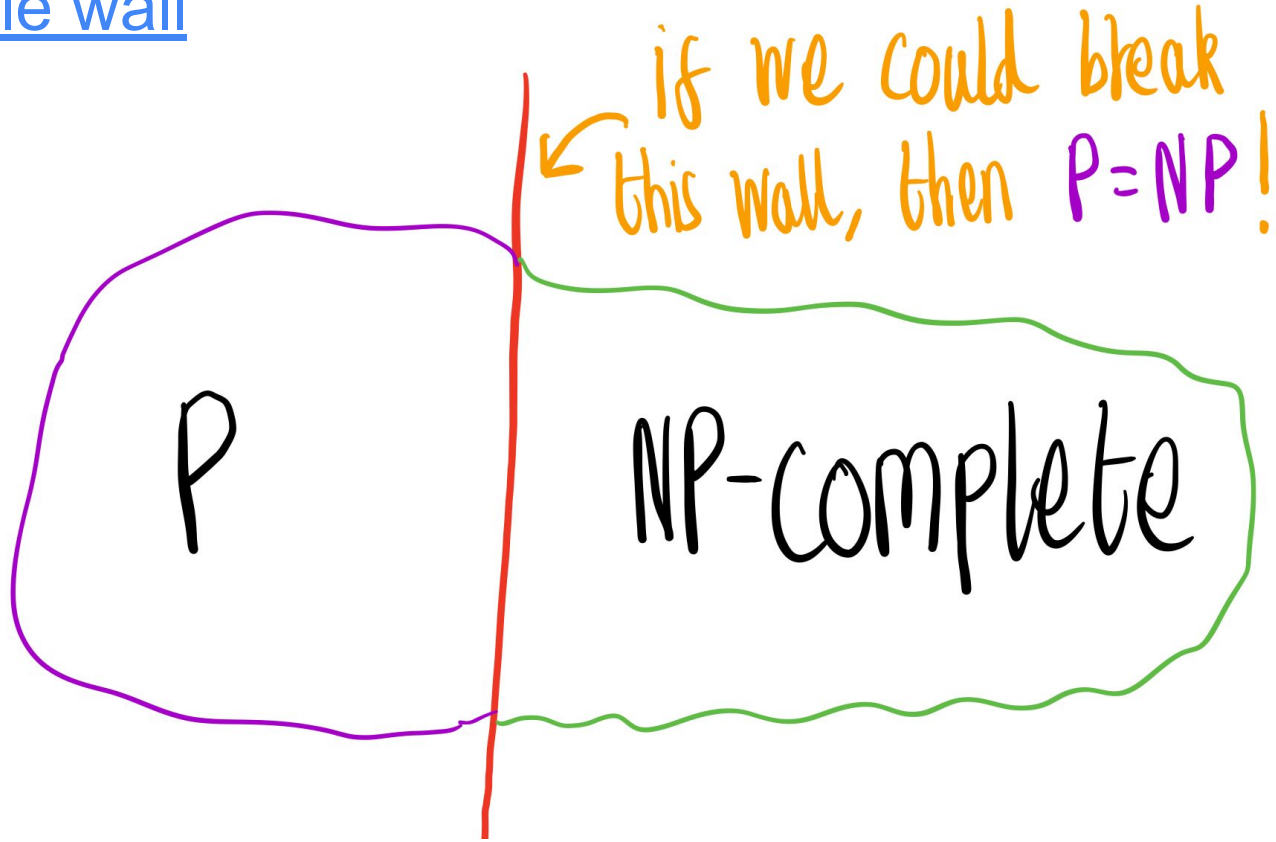


NP-complete problems

If we could ^{quickly} solve an NP-Complete problem, we could ^{quickly} solve every problem in NP.

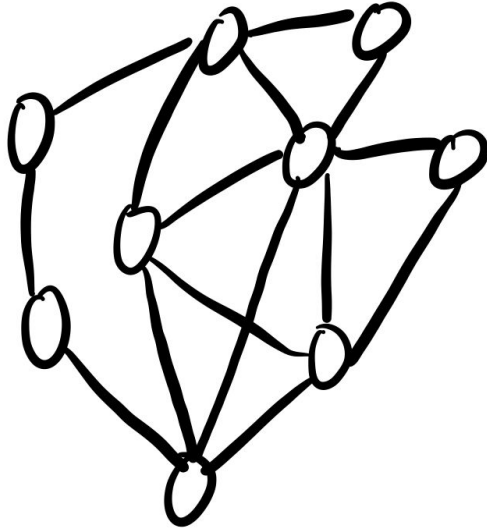
The surprise: There exist very reasonable sounding problems that are provably NP-Complete.

The invisible wall



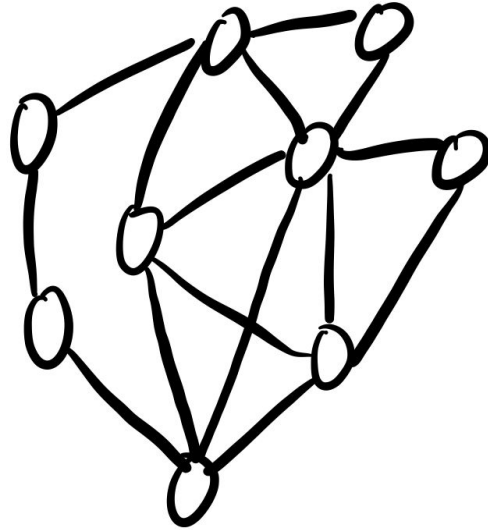
Eulerian Path v.s. Hamiltonian Path

polynomial time



Use every **edge**
once.

exponential time



visit every **node**
once.

2-SAT v.s. 3-SAT

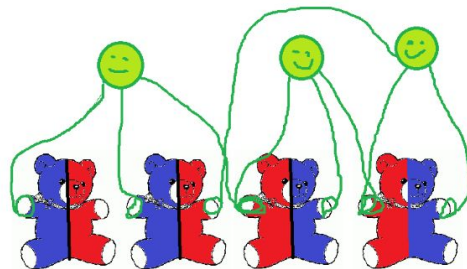
polynomial time

(A or B) and
(A or not C) and
(not B or not A) and
(not B or C)

2 variables per
clause

exponential time

(A or B or C) and
(C or not B or not D) and
...

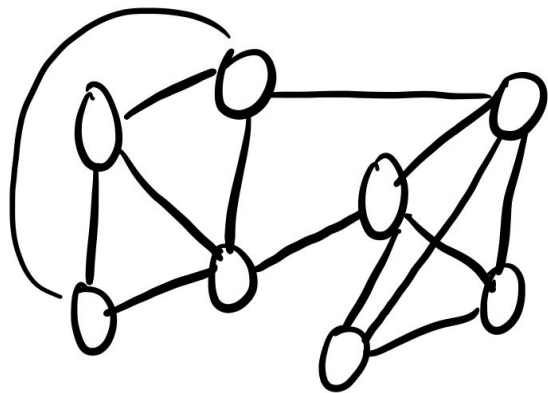


3 variables per
clause

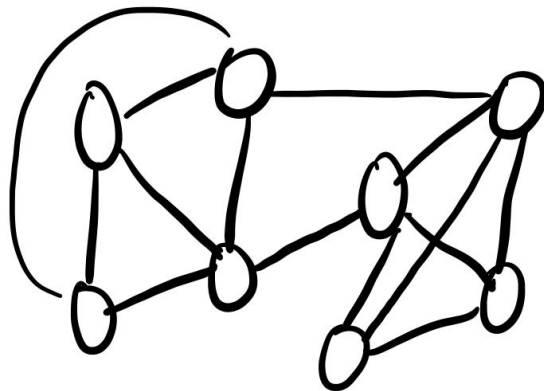
Min Cut v.s. Max Cut

polynomial time

exponential time



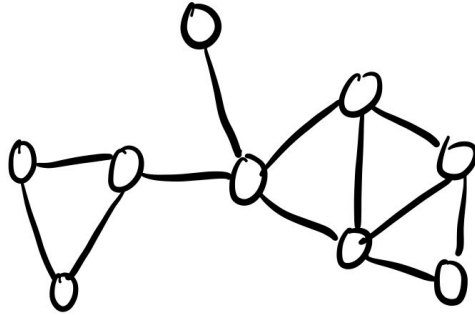
Cut with **fewest**
crossing edges.



Cut with **most**
crossing edges.

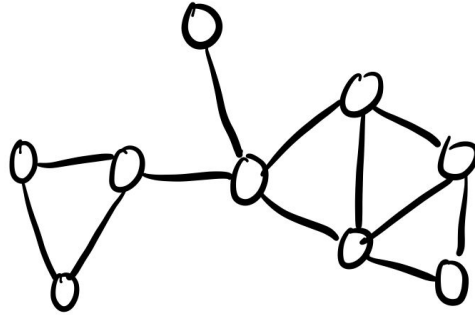
Maximum Matching v.s. Minimum Vertex Cover

polynomial time



Max number of
non-adjacent
edges.

exponential time



Min number of
vertices that
cover every edge.

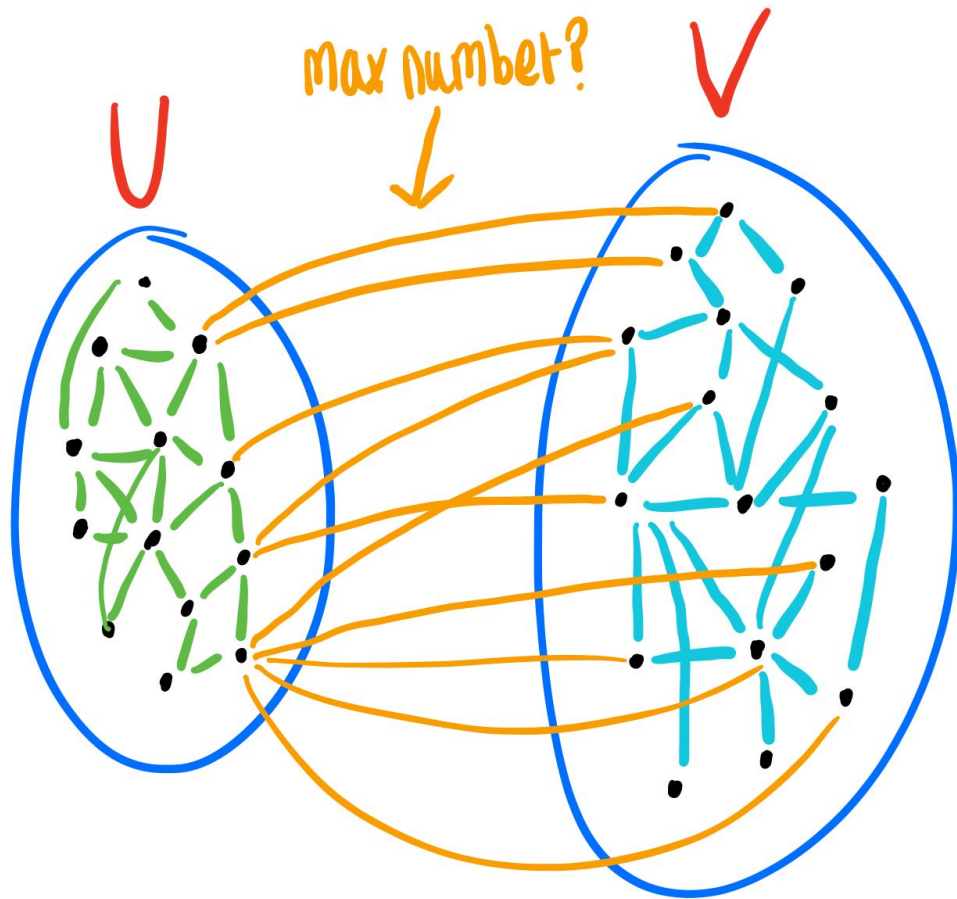
Break for 5 minutes

Approximation Algorithms

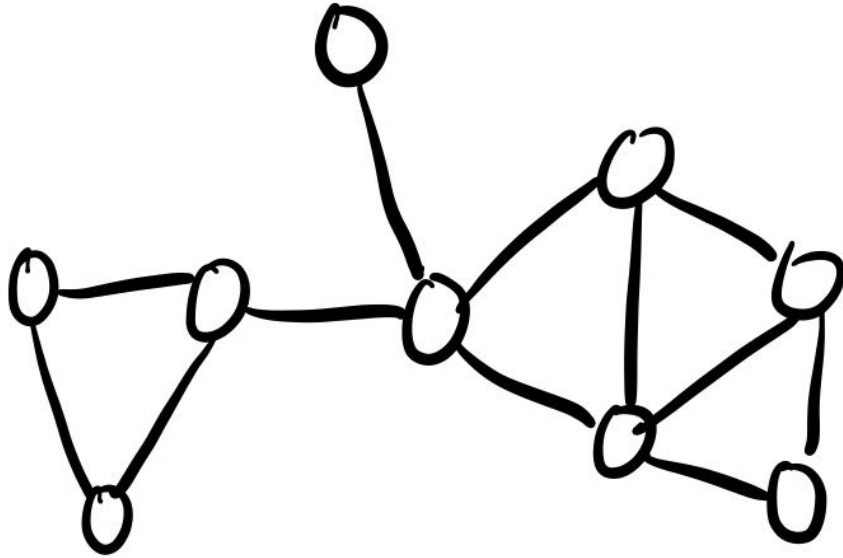
For "optimisation problems":

Get within a **multiplicative factor** of the optimal.

Approximation Algorithm for Max Cut?



Approximation Algorithm for Vertex Cover?



Min number of
vertices that
cover every edge.

The big ideas of algorithm design

- **Correctness**

Correct answers to well-defined problems.

- **Asymptotic complexity** (measure of efficiency)

How does number of steps SCALE with input size?

- **Worst-case analysis**

Want to do well on every possible input!

→ Leads to an incredibly rich design space...

as we've seen!

Further Exploration

Exploring more algorithms...

MIT OCW (Profs. Devadas & Demaine)

Coursera (Prof. Roughgarden)

Further Exploration

Implementing algorithms!

- Just use your favorite language and play around.
 - * Python
- Competitive Programming
 - * Codeforces, Atcoder, etc.
 - * CP Handbook by Laaksonen
 - * USACO

Hope you've enjoyed the classes!