Chapter 13 Generating Random Variables for Simulation

We've seen many examples of distributions. How do we generate instances of a given distribution?

Two methods:

Inverse Transform method

Accept-Reject method

Both assume we already have generator for U(0,1)

Inverse Transform Method

Requirement: To generate instances of r.v. *X*:

- 1. Need to know c.d.f. of X, that is, $F_X(x) = P\{X \le x\}$
- 2. Need to be able to easily invert $F_X(x)$, i.e. get x from $F_X(x)$

Inverse Transform Method

Let u be our random instance of U(0,1). Want to map u to an instance x of the r.v. X.



Q: What properties should *g* have?

> A value in [0, x] should be output with probability: $F_X(x)$

 \succ A value in [0, x] is output with probability: <u>u</u>

$$\implies u = F_X(x) \qquad \implies x = F_X^{-1}(u) \qquad \implies g(\cdot) = F_X(\cdot)$$

Inverse Transform Method

Let u be our random instance of U(0,1). Want to map u to an instance x of the r.v. X.



Inverse Transform method to generate continuous r.v. X:

- 1. Generate random $u \in U(0,1)$.
- 2. Return x such that $F_X(x) = u$.

Inverse Transform Method: Example

Q: Generate $X \sim Exp(\lambda)$, given $u \in U(0,1)$.

 $F_X(x) = u$ $1 - e^{-\lambda x} = u$ $e^{-\lambda x} = 1 - u$ $-\lambda x = \ln(1-u)$ Return $x = -\frac{1}{\lambda}\ln(1-u)$

Same idea applies to **discrete** r.v.s.

Let u be our random instance of U(0,1). Want to map u to an instance x of the r.v. X.



Inverse Transform method to generate discrete r.v. X:

1. Generate random $u \in U(0,1)$.

2. If $0 < u \le p_0$, output $\underline{x_0}$. If $p_0 < u \le p_0 + p_1$, output $\underline{x_1}$. If $p_0 + p_1 < u \le p_0 + p_1 + p_2$, output $\underline{x_2}$.

If $\sum_{i=0}^{\ell-1} p_i < \mathbf{u} \le \sum_{i=0}^{\ell} p_i$, then output x_{ℓ} , where $0 \le \ell \le k$

To be practical, requires closed-form and invertible expression for $F_X(x)$.



Example: Generate instance x of X given instance u of U(0,1).

$$X = \begin{cases} 1 & \text{w.p. } 0.1 \\ 2 & \text{w.p. } 0.1 \\ \vdots \\ 10 & \text{w.p. } 0.1 \end{cases}$$

Solution: We want to determine x as a function of u.

$$p_X(i) = 0.1$$

$$F_X(x) = \mathbf{P}\{X \le x\} = \sum_{i=1}^{x} p_X(i) = x \cdot (0.1)$$
Set $u = F_X(x) = x \cdot (0.1)$
Hence $x = \lceil 10u \rceil \circ \bigcirc$
Note we need ceiling.

Example:

Generate instances of $X \sim Geometric(p)$

Solution:

See Exercise 13.3 in your textbook.

Interview Question

Generate instance of *X* where:

$$X = \begin{cases} 1 & \text{w.p. } p_1 = 10^{-6} \cdot k_1 = 10^{-6} \cdot 0.33 \\ 2 & \text{w.p. } p_2 = 10^{-6} \cdot k_2 = 10^{-6} \cdot 1.63 \\ 3 & \text{w.p. } p_3 = 10^{-6} \cdot k_3 = 10^{-6} \cdot 0.75 \\ \dots \\ 10^6 & \text{w.p. } p_{10^6} = 10^{-6} \cdot k_{10^6} = 10^{-6} \cdot 1.02 \end{cases} \xrightarrow{\text{C}}$$

<u>Given</u>: All of the constants, $k_1, k_2, k_3, \dots, k_{10^6}$ are between (0, 2).

Question: Suppose you know all the p_i 's. How do you generate an instance of X?

Inverse Transform:

- 1. Compute all the partial sums. Store in an array A[1] to $A[10^6]$.
- 2. Given u, find *approximate* bin in array: $x \approx [u \cdot 10^6]$, then search around there for partial sum interval that contains u.

Interview Question

Generate instance of *X* where:

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<u>Given</u>: All of the constants, $k_1, k_2, k_3, \dots, k_{10^6}$ are between (0, 2).

Question: Now suppose you don't know the p_i 's. Can look up one at a time, but takes time to look up each one, and they change over time.

Can't get partial sums → Can't use Inverse Transform Method → Need new approach!

Interview Question

Generate instance of *X* where:

$$X = \begin{cases} 1 & \text{w.p. } p_1 = 10^{-6} \cdot k_1 = 10^{-6} \cdot 0.33 \\ 2 & \text{w.p. } p_2 = 10^{-6} \cdot k_2 = 10^{-6} \cdot 1.63 \\ 3 & \text{w.p. } p_3 = 10^{-6} \cdot k_3 = 10^{-6} \cdot 0.75 \\ \dots \\ 10^6 & \text{w.p. } p_{10^6} = 10^{-6} \cdot k_{10^6} = 10^{-6} \cdot 1.02 \end{cases} \xrightarrow{\text{Given}}_{All \ are}$$

<u>Given</u>: All of the constants, $k_1, k_2, k_3, \dots, k_{10^6}$ are between (0, 2).

Question: Now suppose you don't know the p_i 's. Can look up one at a time, but takes time to look up each one, and they change over time.



> Imagine r.v. Y, where
$$p_i = 10^{-6}$$
, ∀i
> Generate instance i of Y (easy to do).
> Now look up $p_X(i)$. Return i as instance of X if $p_X(i)$ is high compared to $p_Y(i)$?

hehind

A method with fewer requirements ...

Goal: To generate instances of r.v. X

Inverse Transform Method

- 1. Need to know closed-form $F_X(x)$
- 2. Need to know be able to easily invert $F_X(x)$, i.e. get x from $F_X(x) = u$

Accept-Reject Method

- 1. Need the p.m.f. of X (or p.d.f., if continuous)
- Need to know how to generate some other r.v. Y, where X and Y take on the same set of values, i.e.,

$$p_X(i) > 0 \iff p_Y(i) > 0$$

A method with fewer requirements ...

Goal: To generate instances of r.v. X



High-level idea:

We generate an instance of Y. Then we some probability we return that value as our instance of X (accept). Otherwise, we reject that value and try again. **Accept-Reject Method**

- 1. Need the p.m.f. of *X* (or p.d.f., if continuous)
- Need to know how to generate some other r.v. Y, where X and Y take on the same set of values, i.e.,

$$p_X(i) > 0 \iff p_Y(i) > 0$$

Accept-Reject Method

Accept-Reject method to generate discrete r.v. X:

1. Find a discrete r.v. Y which we already know how to generate, where $p_X(i) > 0 \iff p_Y(i) > 0$



Accept-Reject Method: Correctness

Claim: $P{X \text{ is set to } i} = p_X(i)$

Proof: Frac of time *i* is generated & accepted = $P{i$ is generated} · $P{i$ is accepted | *i* is generated}

$$= p_Y(i) \cdot \frac{p_X(i)}{c \ p_Y(i)} = \boxed{\frac{p_X(i)}{c}}$$

Frac of time any value is accepted = \sum_{i} Frac of time *i* is generated & accepted = $\sum_{i} \frac{p_X(i)}{c} = \begin{bmatrix} \frac{1}{c} \end{bmatrix}$



Accept-Reject Method

Theorem: On average, only need to generate *c* values of *Y* before one is accepted.



Back to the Interview Question

Generate instance of *X* where:

$$X = \begin{cases} 1 & \text{w.p. } p_1 = 10^{-6} \cdot k_1 = 10^{-6} \cdot 0.33 \\ 2 & \text{w.p. } p_2 = 10^{-6} \cdot k_2 = 10^{-6} \cdot 1.63 \\ 3 & \text{w.p. } p_3 = 10^{-6} \cdot k_3 = 10^{-6} \cdot 0.75 \\ \dots \\ 10^6 & \text{w.p. } p_{10^6} = 10^{-6} \cdot k_{10^6} = 10^{-6} \cdot 1.02 \end{cases}$$

<u>Given</u>: All the constants, $k_1, k_2, k_3, \dots, k_{10^6}$ are between (0, 2).

Question: You don't know the p_i 's. Can look up one at a time, but takes time to look up each one, and they change over time.

Solution: Accept-Reject! 1. Let $Y = [1, 2, 3, ..., 10^6]$, each with probability 10^{-6} 2. Observe $\frac{p_X(i)}{p_Y(i)} < c$, $\forall i$ 3. Generate *i*: instance of *Y*. Look up $p_X(i)$ only. 4. With probability $\frac{p_X(i)}{c \cdot p_Y(i)}$: ACCEPT *i*, RETURN X = i. Else, Return to Step 3.

Accept-Reject Method: Continuous

Accept-Reject method to generate continuous r.v. X with p.d.f. $f_X(t)$:

- 1. Find a discrete r.v. Y which we already know how to generate, where $f_X(t) > 0 \iff f_Y(t) > 0$
- 2. Let c > 1 be the smallest constant s.t. $\frac{f_X(t)}{f_Y(t)} \le c, \qquad \forall t \text{ s.t. } f_X(t) > 0.$
- 3. Generate an instance *t* of *Y*.
- 4. With probability $AcceptRatio(t) = \frac{f_X(t)}{cf_Y(t)}$, accept t and return X = t. Else, reject t and return to step 3.

GOAL: Generate r.v. *X* where $f_X(t) = 20t(1-t)^3$, 0 < t < 1



Q: Given this image of $f_X(t)$, what's a good choice for Y?

A: $f_Y(t) = 1$, 0 < t < 1 has same range and is easy to generate.

Q: Looking at the picture, how high is *c* approximately ?

A: $c \approx 2$. So only need 2 guesses to attempts on averge to generate X.

GOAL: Generate r.v. *X* where $f_X(t) = 20t(1-t)^3$, 0 < t < 1



[&]quot;Introduction to Probability for Computing", Harchol-Balter '24

GOAL: Generate r.v. *N* where $N \sim Normal(0,1)$



Suffices to generate X = |N| and then multiply X by -1 with probability $\frac{1}{2}$

Q: What's a good choice for r.v. Y with range 0 to ∞ ?

A: Let
$$Y \sim Exp(1)$$
. $f_Y(t) = e^{-t}, 0 < t < \infty$

Q: From the picture, how high is *c* approximately?

A: $c \approx 1.3$. So only need 1.3 guesses to attempts on averge to generate X.



More precisely:

$$c = \max_{t} \left\{ \frac{f_X(t)}{f_Y(t)} \right\} = \max_{t} \sqrt{\frac{2}{\pi}} e^{t - \frac{t^2}{2}}$$

Suffices to maximize $t - \frac{t^2}{2}$

$$0 = \frac{d}{dt} \left(t - \frac{t^2}{2} \right) = 1 - t \quad \Leftrightarrow \quad t = 1$$

$$c = \frac{f_X(1)}{f_Y(1)} = \sqrt{\frac{2e}{\pi}} \approx 1.3$$

