Chapter 5 Variance

Higher moments

 \bullet

Defn: The **kth moment** of r.v. X is

$$
E[X^k] = \sum_{x} x^k \cdot P\{X = x\}
$$

Example:

 $X \sim Geometric(p)$.

Derive $E[X^2]$.

Can we say $\boldsymbol{E}[X^2] = \boldsymbol{E}[X] \cdot \boldsymbol{E}[X]$?

This doesn't work because X is not independent of X .

Higher moments

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$$

Example:

 $X \sim Geometric(p)$.

Derive $E[X^2]$.

$$
E[X^{2}] = \sum_{i=1}^{\infty} i^{2} p_{X}(i)
$$

=
$$
\sum_{i=1}^{\infty} i^{2} (1-p)^{i-1} \cdot p
$$

Not obvious how to compute this sum

2nd Moment of Geometric

2nd Moment of Geometric

Let's try

conditioning

 $X \sim Geometric(p)$.

Derive $E[X^2]$. Condition on value of 1st flip, Y.

 $E[X^2] = E[X^2 | Y = 1] \cdot P{Y = 1} + E[X^2 | Y = 0] \cdot P{Y = 0}$ $= 1 \cdot p + E[X^2 | Y = 0] \cdot (1-p)$ $[X | Y = 0] = X + 1$ $[X^2 | Y = 0] = (X + 1)^2$

 $= 1 \cdot p + E[(1+X)^2] \cdot (1-p)$

2nd Moment of Geometric

Derive $E[X^2]$. Condition on value of 1st flip, Y.

 $E[X^2] = E[X^2 | Y = 1] \cdot P{Y = 1} + E[X^2 | Y = 0] \cdot P{Y = 0}$

$$
= 1 \cdot p + E[X^2 | Y = 0] \cdot (1 - p)
$$

$$
=1\cdot p+E[(1+X)^2]\cdot(1-p)
$$

$$
=1\cdot p+E[1+2X+X^2]\cdot(1-p)
$$

$$
= 1 \cdot p + (1 + 2E[X] + E[X^2]) \cdot (1 - p)
$$

$$
E[X^2] = \frac{2-p}{p^2}
$$

Let's try

conditioning

Variance

Defn: The **variance** of r.v. X is the expected squared difference of X from its mean:

 $Var(X) = E[(X - E[X])^{2}]$

[&]quot;Introduction to Probability for Computing", Harchol-Balter '24

Choosing between Microsoft and a Startup

Variance of Bernoulli (p)

$$
X = \text{value of the coin flip } = \begin{cases} 1 & \text{w.p. } p \\ 0 & \text{o.w.} \end{cases}
$$

 $Var(X) = E[(X - p)^2]$ $= E[X^2 - 2Xp + p^2]$ $= E[X^2] - 2pE[X] + p^2$ $= p \cdot 1^2 - 2p \cdot p + p^2$ $= p - p^2 = p(1-p)$

Probability p of heads

Recall: $E[X] = p$

"Introduction to Probability for Computing", Harchol-Balter '24

Conditioning on Variance is NOT allowed

$$
X = \text{value of the coin flip } = \begin{cases} 1 & \text{w.p. } p \\ 0 & \text{o.w.} \end{cases}
$$

$$
Var(X) = Var
$$

= 0 · p + 0 ·
= 0

Probability p of heads

Recall:
$$
E[X] = p
$$

Alternative definitions of variance?

Standard deviation of X

Defn: The **standard deviation** of a r.v. X is:

$$
\sigma_X = \text{std}(X) = \sqrt{E[(X - E[X])^2]}
$$

We often write:

$$
Var(X)=\sigma_X^2
$$

The need for a different variation metric

Suppose we measure a quantity, first in cm (r.v. X) and then in mm (r.v. Y):

$$
X = \begin{bmatrix} 1 & \text{w.p.} & \frac{1}{3} \\ 2 & \text{w.p.} & \frac{1}{3} \\ 3 & \text{w.p.} & \frac{1}{3} \end{bmatrix} \qquad Y = \begin{bmatrix} 10 & \text{w.p.} & \frac{1}{3} \\ 20 & \text{w.p.} & \frac{1}{3} \\ 30 & \text{w.p.} & \frac{1}{3} \end{bmatrix}
$$

Feels like they should have same variance, since they're the same quantity, but they don't:

$$
Var(X) = \frac{2}{3}
$$
 $Var(Y) = \frac{200}{3}$ Need a new metric!

Squared coefficient of variation

Equivalent definition of variance

Theorem 5.7: $Var(X) = E[X^2] - E[X]^2$

 $Var(X) = E[(X - E[X])^{2}]$

$$
= E[X^2 - 2XE[X] + E[X]^2]
$$

 $= E[X^2] - 2 E[X]E[X] + E[X]^2$

$$
= E[X^2] - E[X]^2
$$

Linearity of Variance

Theorem 5.8: Let X and Y be random variables where $X \perp Y$. Then

 $Var(X + Y) = Var(X) + Var(Y)$

Variance of Binomial (n, p)

Experiment: Flip a coin, with probability p of Heads, n times

Random Variable $X =$ number of heads

Key Observation:

 $X = X_1 + X_2 + \cdots + X_n$, where $X_i \sim$ Bernoulli (p) \bullet

Applying Linearity of Variance:

$$
Var(X) = Var(X_1) + Var(X_2) + \cdots + Var(X_n)
$$

$$
= p(1-p) + p(1-p) + \dots + p(1-p) = np(1-p)
$$

What is $E[X_i]$?

 $\boldsymbol{\eta}$

What is $Var(X_i)$?

Sums versus copies

Let X_1 and X_2 be independent and identically distributed (i.i.d.) random variables, where $X_1 \sim X_2 \sim X$.

$$
Y = X_1 + X_2 \qquad \text{versus} \qquad \qquad Z = 2X
$$

more variance?

Covariance

Defn 5.11: The **covariance** of two random variables X and Y is:

$$
Cov(X,Y) = E[(X - E[X]) \cdot (Y - E[Y])]
$$

Intuition: If the large values of X tend to happen with the large values of Y, and the small values of X tend to happen with the small values of Y, then $(X - E[X]) \cdot (Y - E[Y])$ is positive on average, so $Cov(X, Y) > 0$, and we say that X and Y are **positively correlated**.

Likewise if $Cov(X, Y) < 0$, we say that X and Y are **negatively correlated**.

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Theorem 5.12: Cov(X, Y) = E[XY] - E[X]E[Y]
```
Central moments

Defn 5.13: The **kth central moment** of r.v. X is

$$
E[(X - E[X])^{k}] = \sum_{x} (x - E[X])^{k} \cdot P\{X = x\}
$$

The 2nd central moment is the variance, representing how much the distribution varies from its mean.

Q: What's the difference between the 2nd and 4th central moments?

The 4th central moment is similar to variance, but outliers (those far from the mean) count a lot more!

Third central moment and skew

The **3rd central moment** of r.v. *X* is $E[(X - E[X])^3]$. Roughly, the 3rd moment captures the **skew** of the distribution.

 $X_{\scriptscriptstyle 4}$

Sum of random number of random variables

In many applications, we need to add a number of i.i.d. r.v.s, where the total number of r.v.s added is itself a r.v.

> Get new prize every day, until wheel says stop.

> > X_3

 X_2

 X_1

In

 $\overline{\mathbf{v}}$

 $S = \sum$

 $l=1$

 X_i

<u>N</u>

$$
Total earnings = \sum_{i=1}^{N} X_i
$$

 $N \perp \{X_1, X_2, X_3, \dots\}$

where
$$
N \sim \text{Geometric}\left(\frac{1}{6}\right)
$$

Let $X_1, X_2, X_3, ...$ be i.i.d. r.v.s, where $X_i \sim X$. Let $N \perp \{X_1, X_2, X_3, ... \}$

"Introduction to Probability for Computing", Harchol-Balter '24

Let $X_1, X_2, X_3, ...$ be i.i.d. r.v.s, where $X_i \sim X$. Let $N \perp \{X_1, X_2, X_3, ... \}$

Let
$$
S = \sum_{i=1}^{N} X_i
$$

$$
E[S] = \sum_{n=1}^{\infty} E[S \mid N = n] \cdot P\{N = n\}
$$

$$
= \sum_{n=1}^{\infty} E\left[\sum_{i=1}^{n} X_i \mid N = n\right] \cdot P\{N = n\}
$$

$$
= \sum_{n=1}^{\infty} nE[X] \cdot P\{N = n\} = E[X] \cdot E[N]
$$

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Summary Theorem 5.14:

Let
$$
X_1, X_2, X_3, ...
$$
 be i.i.d. r.v.s, where $X_i \sim X$.
Let $S = \sum_{i=1}^{N} X_i$, where $N \perp \{X_1, X_2, X_3, ... \}$

Then

 $E[S^2] = E[N] \cdot Var(X) + E[N^2] \cdot E[X]^2$ $E[S] = E[N] \cdot E[X]$ $Var(S) = E[N] \cdot Var(X) + Var(N) \cdot E[X]^2$

Example: Epidemic growth modeling

At each time step, every leaf independently either:

- forks off 2 children, w.p. $\frac{1}{2}$ 2
- stays inert w.p. $\frac{1}{2}$ 2

 X_t is number of leaves in tree after t steps.

Example: Epidemic growth modeling

At each time step, every leaf independently either:

- forks off 2 children, w.p. $\frac{1}{2}$ 2
- stays inert w.p. $\frac{1}{2}$

 X_t is number of leaves in tree after t steps.

Hint:

Tail bounds

Defn: The **tail** of random variable *X* is $P\{X > x\}$.

Example: T denotes response time at a web service. Want to ensure the fraction of people with response time $> 0.5s$ is not too high.

Want an **upper bound** on $P{T > 0.5}$. This is called a **tail bound**.

Tail bounds

Another Example: *n* items are hashed into a table of size n . Assume each item ends up in a random bucket. Ideally, we have 1 item per bucket. What is the fraction of time that your search time $> k$? (i.e., what's the probability your bucket has $> k$ items?)

Let
$$
N = \text{Hitems in bucket 1}
$$

\n
$$
P\{N > k\} = \sum_{i=k+1}^{n} P\{N = i\} = \sum_{i=k+1}^{n} {n \choose i} \left(\frac{1}{n}\right)^i \left(1 - \frac{1}{n}\right)^{n-i}
$$
\nWe don't know how to compute such bounds in general.

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Point: We'll see that just knowing the mean and variance suffices for a tail **bound**. In some cases, the mean alone suffices (although this bound is quite weak).

Markov's inequality

Theorem: (Markov's inequality) If r.v. X is non-negative, then $\forall a > 0$,

$$
P\{X \ge a\} \le \frac{E[X]}{a}
$$

$$
E[X] = \sum_{x=0}^{\infty} x \cdot p_X(x) \ge \sum_{x=a}^{\infty} x \cdot p_X(x)
$$

$$
\ge \sum_{x=a}^{\infty} a \cdot p_X(x)
$$

$$
= a \sum_{x=a}^{\infty} p_X(x) = a \cdot P\{X \ge a\}
$$

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Chebyshev's inequality

Theorem: (Chebyshev's inequality) Let X be any r.v. with finite mean, μ , and finite variance. Then $\forall a > 0$,

$$
P\{|X-\mu|\geq a\}\leq \frac{Var(X)}{a^2}
$$

Chebyshev's inequality

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$$
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$$

Example:

$$
N \sim Binomial\left(n,\frac{1}{n}\right)
$$

Provide upper bound on: $P{N \ge 6}$

$$
P{N \ge 6} \le P{|N - 1| \ge 5}
$$

$$
\le \frac{Var(N)}{25}
$$

$$
\le \frac{1}{25}
$$

Stochastic dominance

Defn 5.18: Given two random variables X and Y , if

 $P\{X > i\} \geq P\{Y > i\}, \quad \forall i$

we say that X stochastically dominates Y :

 $X \geq_{st} Y$

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Stochastic dominance

 $X =$ Number pairs of shoes owned by women ~ $Poisson(\lambda = 27)$

 $Y =$ Number pairs of shoes owned by men ~ $Poisson(\lambda = 12)$

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Jensen's inequality: motivation

We already know that

 $E[X^2] \ge E[X]^2$

Is it also the case that

 $E[X^3] \geq E[X]^3$? $E[X^4] \geq E[X]^4$? $E[X^{4.5}] \geq E[X]^{4.5}$?

Theorem: Let X be any positive r.v. Then $\forall a \in \text{Reals}$,

$$
E[X^a] \ge E[X]^a
$$

Jensen's inequality

Defn 5.21: A function $g(x)$ is **convex** on interval S if, for any $x_1, x_2 \in S$, and any $\alpha \in [0,1]$, we have:

$$
g(\alpha x_1 + (1 - \alpha)x_2) \le \alpha g(x_1) + (1 - \alpha)g(x_2)
$$

 $g(x)$ is convex on S iff $g''(x) \geq 0$, $\forall x \in S$.

Jensen's inequality

<u>Defn 5.22</u>: A function $g(x)$ is **convex** on interval S if, for any $x_1, x_2, ..., x_n \in S$, and any $\alpha_1, \alpha_2, ..., \alpha_n \in [0,1]$, where $\sum_i \alpha_i = 1$, we have:

 $g(\alpha_1 x_1 + \alpha_2 x_2 + \cdots + \alpha_n x_n) \leq \alpha_1 g(x_1) + \alpha_2 g(x_2) + \cdots + \alpha_n g(x_n)$

 $g(p_X(x_1)x_1 + \cdots + p_X(x_n)x_n) \leq p_X(x_1)g(x_1) + \cdots + p_X(x_n)g(x_n)$

 $\implies g(E[X]) \leq E[g(X)]$

Jensen's inequality

Theorem 5.23: (Jensen's inequality) If $g(x)$ is **convex** on interval S and X takes on values on interval S , then:

 $g(E[X]) \leq E[g(X)]$

 $g(p_X(x_1)x_1 + \cdots + p_X(x_n)x_n) \leq p_X(x_1)g(x_1) + \cdots + p_X(x_n)g(x_n)$

$\implies g(E[X]) \leq E[g(X)]$

Defn: The **inspection paradox** says that, in high-variability settings, the mean seen by a random observer can be very different from the true mean.

Mean time between buses is 10 minutes.

However if there is some variability in the time between buses, then a randomly arriving person will wait more than 5 minutes.

Expected wait can even be >10 minutes!

my wait

Defn: The **inspection paradox** says that, in high-variability settings, the mean seen by a random observer can be very different from the true mean.

Average class size reported by students is 100.

But the dean claims average class size is 30.

Defn: The **inspection paradox** says that, in high-variability settings, the mean seen by a random observer can be very different from the true mean.

class size 130 class Average class size class class size10 size10 size 10 reported by students is 100. But the dean claims average class class size is 30. class size1 $size10$ No one is lying.

How can this be? 180 students in 6 classes > 30 students/class. Avg observed class size $= \frac{50}{180} \cdot 10 + \frac{130}{180} \cdot 130 \approx 97$ "Introduction to Probability for Computing", Harchol-Balter '24

Defn: The **inspection paradox** says that, in high-variability settings, the mean seen by a random observer can be very different from the true mean.

The average Facebook user has 44 friends.

But the average friend of a Facebook user has 104 friends.

In fact, with probability 76%, your friend is more popular than you are.

How can this be?

Most people have few friends.

A few people are very popular with many friends. Which classification most likely describes you? Which most likely describes your friend?