Chapter 5 Variance

### Higher moments

.

#### <u>Defn</u>: The *k***th moment** of r.v. *X* is

$$\boldsymbol{E}[X^k] = \sum_{x} x^k \cdot \boldsymbol{P}\{X = x\}$$

#### Example:

 $X \sim Geometric(p)$ .

Derive  $\boldsymbol{E}[X^2]$ .

Can we say  $E[X^{2}] = E[X] \cdot E[X]?$ 

# This doesn't work because X is not independent of X.

### Higher moments

E

#### <u>Defn</u>: The *k***th moment** of r.v. *X* is

$$\boldsymbol{E}[X^k] = \sum_{x} x^k \cdot \boldsymbol{P}\{X = x\}$$

#### Example:

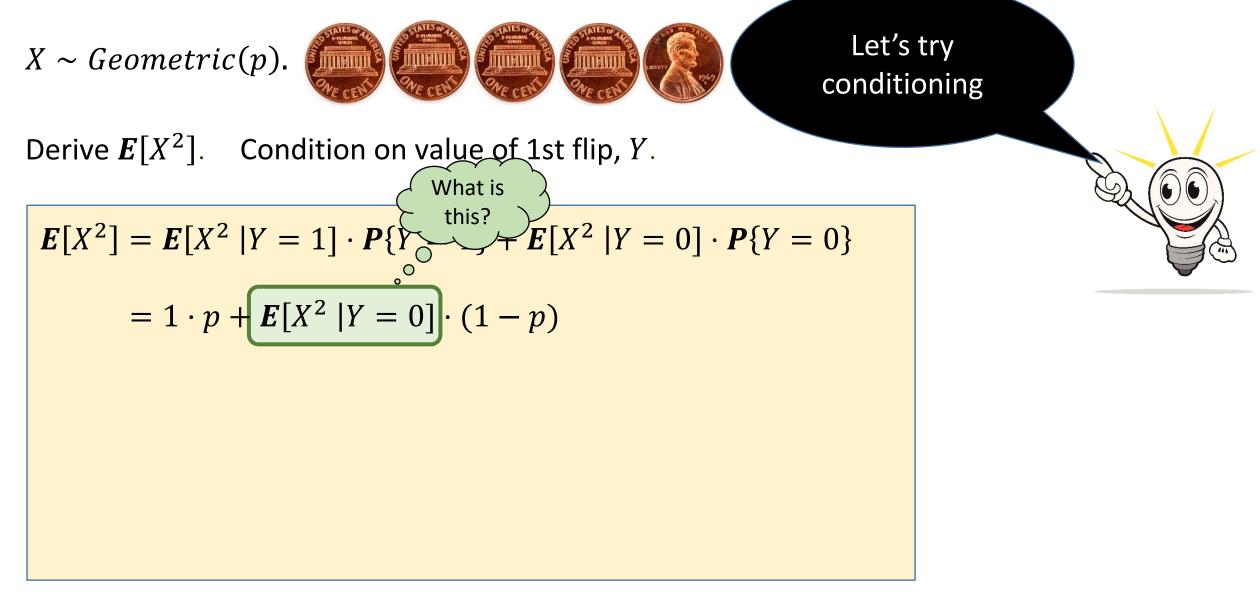
 $X \sim Geometric(p)$ .

Derive  $\boldsymbol{E}[X^2]$ .

$$[X^{2}] = \sum_{i=1}^{\infty} i^{2} p_{X}(i)$$
$$= \sum_{i=1}^{\infty} i^{2} (1-p)^{i-1} \cdot p$$

Not obvious how to compute this sum

## 2<sup>nd</sup> Moment of Geometric



## 2<sup>nd</sup> Moment of Geometric

Let's try

conditioning

 $X \sim Geometric(p)$ .



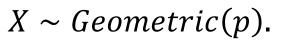
Derive  $E[X^2]$ . Condition on value of 1st flip, Y.

 $\boldsymbol{E}[X^2] = \boldsymbol{E}[X^2 | Y = 1] \cdot \boldsymbol{P}\{Y = 1\} + \boldsymbol{E}[X^2 | Y = 0] \cdot \boldsymbol{P}\{Y = 0\}$ 

$$= 1 \cdot p + E[X^{2} | Y = 0] \cdot (1 - p)$$
$$[X | Y = 0] \stackrel{d}{=} X + 1$$
$$[X^{2} | Y = 0] \stackrel{d}{=} (X + 1)^{2}$$

$$= 1 \cdot p + \boldsymbol{E}[(1+X)^2] \cdot (1-p)$$

# 2<sup>nd</sup> Moment of Geometric





Derive  $E[X^2]$ . Condition on value of 1st flip, Y.

$$\boldsymbol{E}[X^2] = \boldsymbol{E}[X^2 | Y = 1] \cdot \boldsymbol{P}\{Y = 1\} + \boldsymbol{E}[X^2 | Y = 0] \cdot \boldsymbol{P}\{Y = 0\}$$

$$= 1 \cdot p + E[X^2 | Y = 0] \cdot (1 - p)$$

$$= 1 \cdot p + \boldsymbol{E}[(1+X)^2] \cdot (1-p)$$

$$= 1 \cdot p + E[1 + 2X + X^{2}] \cdot (1 - p)$$

$$= 1 \cdot p + (1 + 2E[X] + E[X^2]) \cdot (1 - p)$$

$$\frac{\text{Result}}{\boldsymbol{E}[X^2]} = \frac{2-p}{p^2}$$

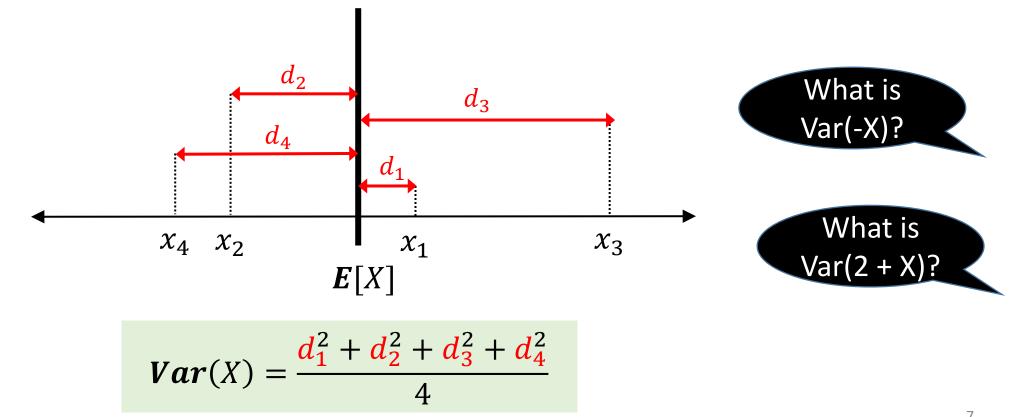
Let's try

conditioning

#### Variance

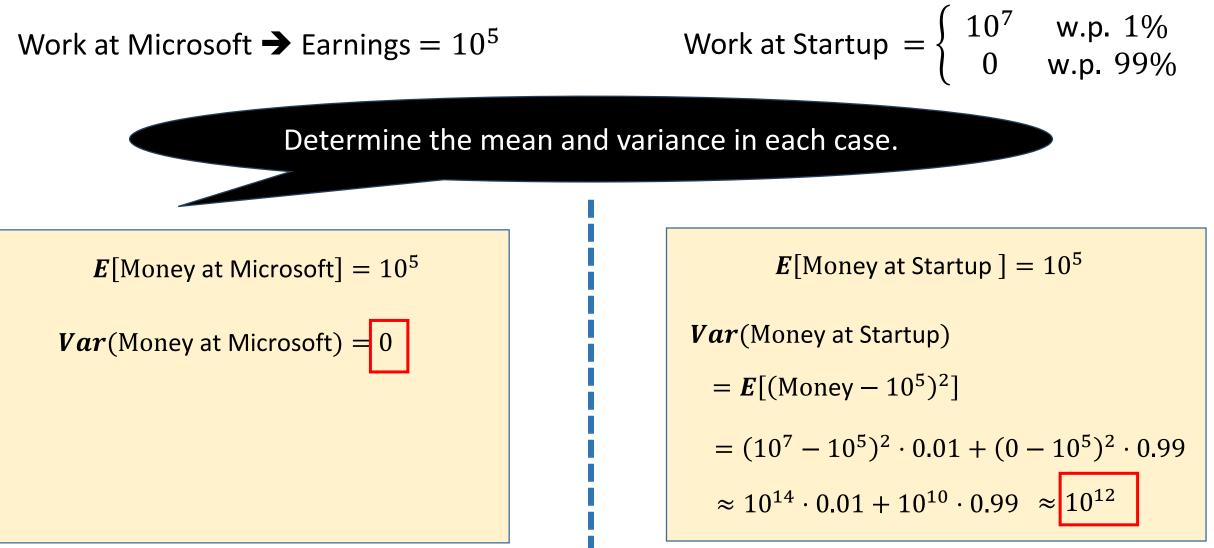
<u>Defn</u>: The **variance** of r.v. X is the expected squared difference of X from its mean:

 $Var(X) = E[(X - E[X])^2]$ 



<sup>&</sup>quot;Introduction to Probability for Computing", Harchol-Balter '24

## Choosing between Microsoft and a Startup



## Variance of Bernoulli(p)

$$X = \text{value of the coin flip} = \begin{cases} 1 & \text{w.p. } p \\ 0 & \text{o.w.} \end{cases}$$

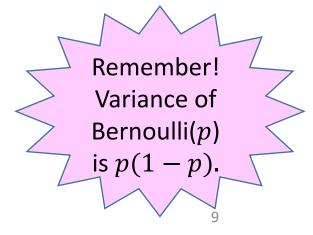
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 $Var(X) = E[(X - p)^{2}]$ =  $E[X^{2} - 2Xp + p^{2}]$ =  $E[X^{2}] - 2pE[X] + p^{2}$ =  $p \cdot 1^{2} - 2p \cdot p + p^{2}$ =  $p - p^{2} = p(1 - p)$ 



Probability p of heads

<u>Recall</u>:  $\boldsymbol{E}[X] = p$ 



"Introduction to Probability for Computing", Harchol-Balter '24

## Conditioning on Variance is NOT allowed

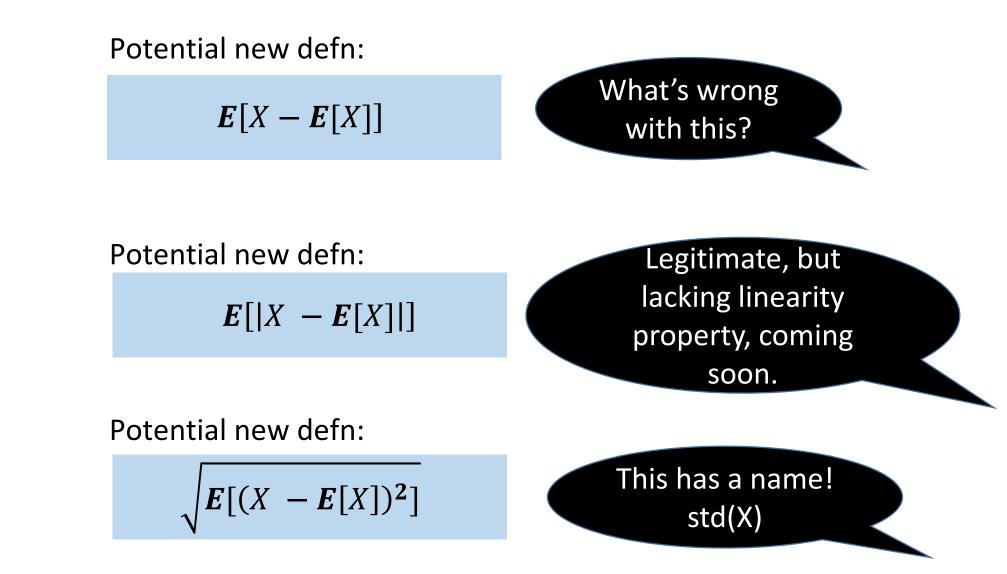
$$X = \text{value of the coin flip} = \begin{cases} 1 & \text{w.p. } p \\ 0 & \text{o.w.} \end{cases}$$

Probability p of heads

Recall: 
$$\boldsymbol{E}[X] = p$$

$$Var(X) = Var = 1) \cdot p + V \quad (X = 0) \cdot (1 - p)$$
$$= 0 \cdot p + 0 \cdot 1$$
$$= 0$$

## Alternative definitions of variance?



#### Standard deviation of X

<u>Defn</u>: The **standard deviation** of a r.v. X is:

$$\sigma_X = std(X) = \sqrt{E[(X - E[X])^2]}$$

We often write:

$$Var(X) = \sigma_X^2$$

## The need for a different variation metric

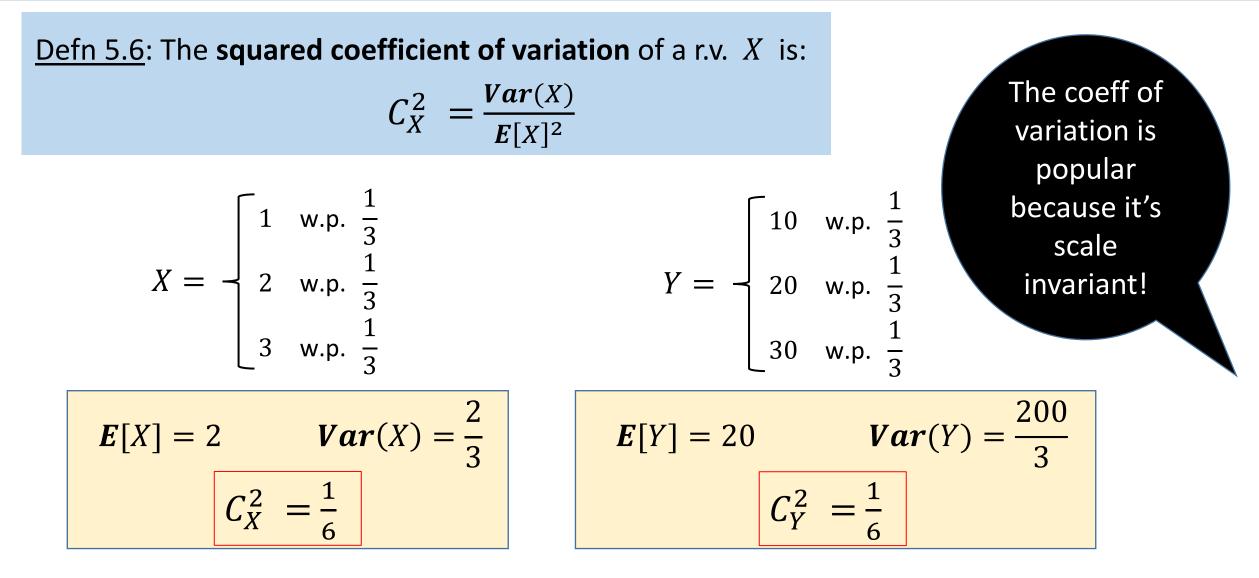
Suppose we measure a quantity, first in cm (r.v. X) and then in mm (r.v. Y):

$$X = - \begin{cases} 1 & \text{w.p.} & \frac{1}{3} \\ 2 & \text{w.p.} & \frac{1}{3} \\ 3 & \text{w.p.} & \frac{1}{3} \end{cases} \qquad Y = - \begin{cases} 10 & \text{w.p.} & \frac{1}{3} \\ 20 & \text{w.p.} & \frac{1}{3} \\ 30 & \text{w.p.} & \frac{1}{3} \end{cases}$$

Feels like they should have same variance, since they're the same quantity, but they don't:

$$Var(X) = \frac{2}{3}$$
   
  $Var(Y) = \frac{200}{3}$    
 Need a new metric!

#### Squared coefficient of variation



### Equivalent definition of variance

Theorem 5.7:  $Var(X) = E[X^2] - E[X]^2$ 

 $Var(X) = E[(X - E[X])^2]$ 

$$= \boldsymbol{E}[X^2 - 2X\boldsymbol{E}[X] + \boldsymbol{E}[X]^2]$$

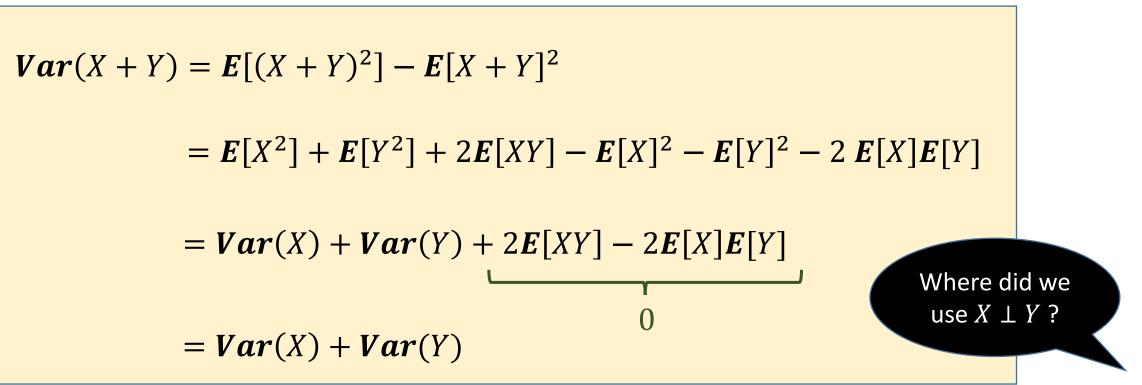
 $= E[X^{2}] - 2 E[X]E[X] + E[X]^{2}$ 

$$= \boldsymbol{E}[X^2] - \boldsymbol{E}[X]^2$$

## Linearity of Variance

**Theorem 5.8:** Let *X* and *Y* be random variables where  $X \perp Y$ . Then

Var(X + Y) = Var(X) + Var(Y)



# Variance of Binomial(n, p)

**Experiment**: Flip a coin, with probability p of Heads, n times

**Random Variable** *X* = number of heads

**Key Observation:** 

 $X = X_1 + X_2 + \dots + X_n$ , where  $X_i \sim \text{Bernoulli}(p)$  •

**Applying Linearity of Variance:** 

$$Var(X) = Var(X_1) + Var(X_2) + \dots + Var(X_n)$$

$$= p(1-p) + p(1-p) + \dots + p(1-p) = np(1-p)$$

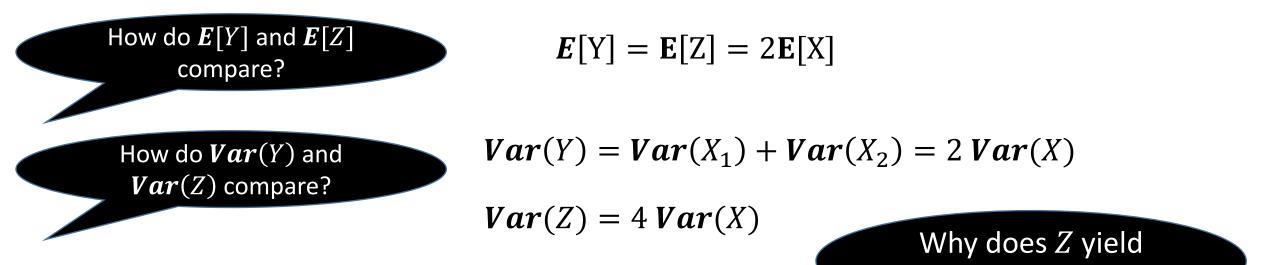
Remember! Variance of Binomial(n, p)is np(1-p).

What is  $\boldsymbol{E}[X_i]$ ? What is  $\boldsymbol{Var}(X_i)$ ?

#### Sums versus copies

Let  $X_1$  and  $X_2$  be independent and identically distributed (i.i.d.) random variables, where  $X_1 \sim X_2 \sim X$ .

$$Y = X_1 + X_2$$
 versus  $Z = 2X$ 



more variance?

#### Covariance

<u>Defn 5.11</u>: The **covariance** of two random variables *X* and *Y* is:

$$Cov(X,Y) = E[(X - E[X]) \cdot (Y - E[Y])]$$

<u>Intuition</u>: If the large values of X tend to happen with the large values of Y, and the small values of X tend to happen with the small values of Y, then  $(X - E[X]) \cdot (Y - E[Y])$  is positive on average, so Cov(X, Y) > 0, and we say that X and Y are **positively correlated**.

Likewise if Cov(X, Y) < 0, we say that X and Y are **negatively correlated**.

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Theorem 5.12: Cov(X, Y) = E[XY] - E[X]E[Y]
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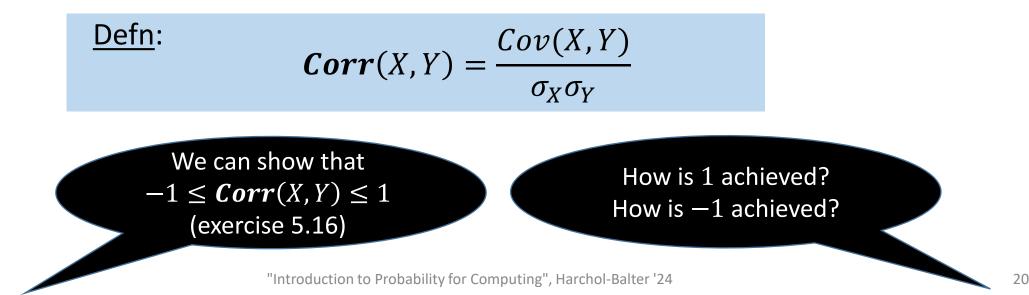
## **Correlation Coefficient**

<u>Defn 5.11</u>: The **covariance** of two random variables *X* and *Y* is:

$$Cov(X,Y) = E[(X - E[X]) \cdot (Y - E[Y])]$$

<u>*Problem*</u>: Covariance is sensitive to scale. If  $X \rightarrow 2X$ , the covariance doubles.

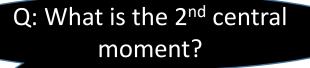
<u>Solution</u>: The correlation coefficient is a normalization that is insensitive to scale.



#### Central moments

<u>Defn 5.13</u>: The *k*th central moment of r.v. *X* is

$$E[(X - E[X])^{k}] = \sum_{x} (x - E[X])^{k} \cdot P\{X = x\}$$

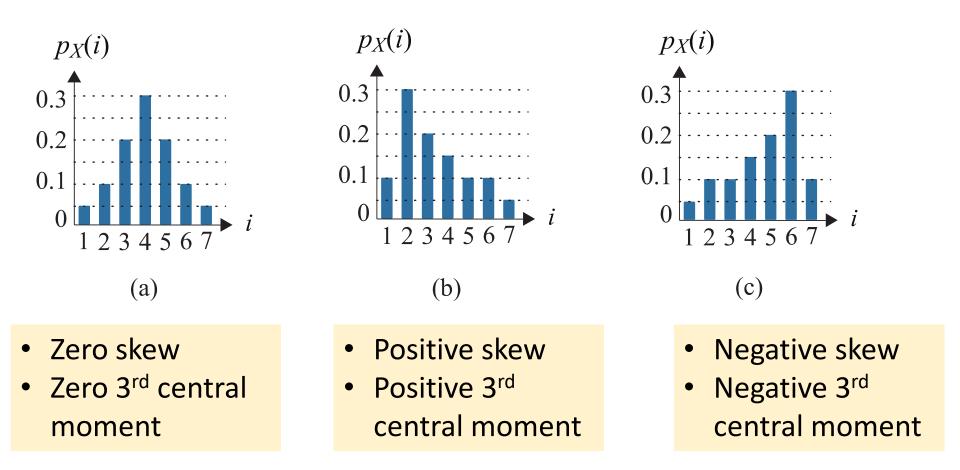


The 2<sup>nd</sup> central moment is the variance, representing how much the distribution varies from its mean.

Q: What's the difference between the 2<sup>nd</sup> and 4<sup>th</sup> central moments? The 4<sup>th</sup> central moment is similar to variance, but outliers (those far from the mean) count a lot more!

## Third central moment and skew

The **3<sup>rd</sup> central moment** of r.v. X is  $E[(X - E[X])^3]$ . Roughly, the 3<sup>rd</sup> moment captures the **skew** of the distribution.



In many applications, we need to add a number of i.i.d. r.v.s, where the total number of r.v.s added is itself a r.v.

 $S = \sum_{i} X_i$ 

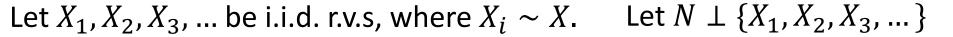


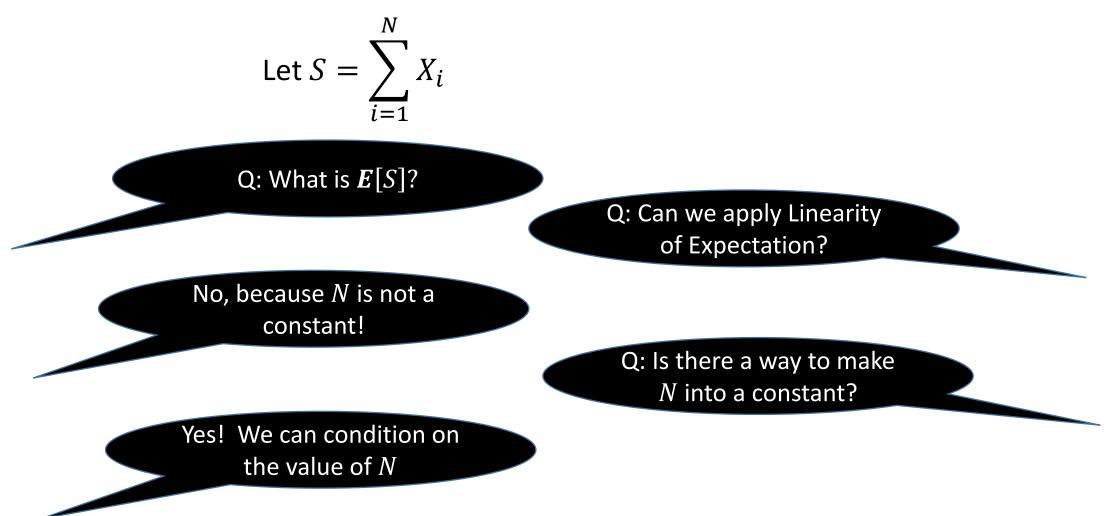
N

$$N \perp \{X_1, X_2, X_3, \dots\}$$

Total earnings  $=\sum_{i=1}^{N} X_i$ 

where 
$$N \sim Geometric\left(\frac{1}{6}\right)$$





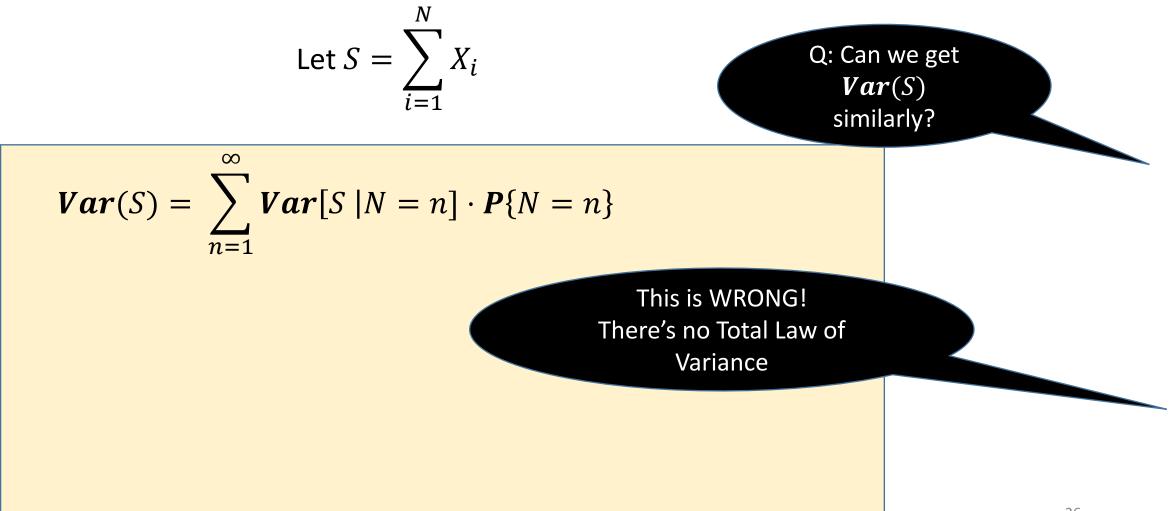
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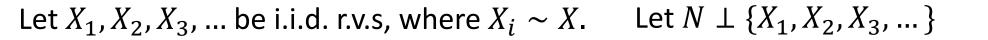
Let  $X_1, X_2, X_3, \dots$  be i.i.d. r.v.s, where  $X_i \sim X$ . Let  $N \perp \{X_1, X_2, X_3, \dots\}$ 

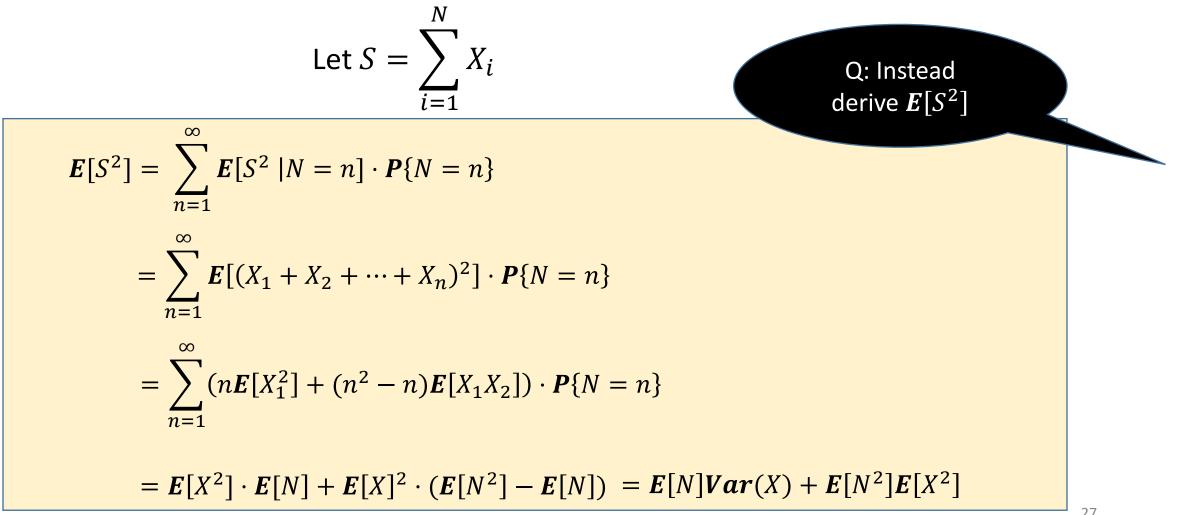
Let 
$$S = \sum_{i=1}^{N} X_i$$

$$E[S] = \sum_{n=1}^{\infty} E[S \mid N = n] \cdot P\{N = n\}$$
$$= \sum_{n=1}^{\infty} E\left[\sum_{i=1}^{n} X_i \mid N = n\right] \cdot P\{N = n\}$$
$$= \sum_{n=1}^{\infty} nE[X] \cdot P\{N = n\} = E[X] \cdot E[N]$$

Let  $X_1, X_2, X_3, \dots$  be i.i.d. r.v.s, where  $X_i \sim X$ . Let  $N \perp \{X_1, X_2, X_3, \dots\}$ 







Summary Theorem 5.14:

Let 
$$X_1, X_2, X_3, ...$$
 be i.i.d. r.v.s, where  $X_i \sim X$ .  
Let  $S = \sum_{i=1}^{N} X_i$ , where  $N \perp \{X_1, X_2, X_3, ...\}$ 



Then

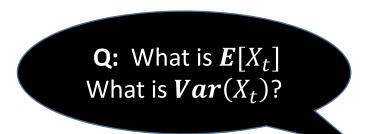
 $E[S] = E[N] \cdot E[X]$  $E[S^{2}] = E[N] \cdot Var(X) + E[N^{2}] \cdot E[X]^{2}$  $Var(S) = E[N] \cdot Var(X) + Var(N) \cdot E[X]^{2}$ 

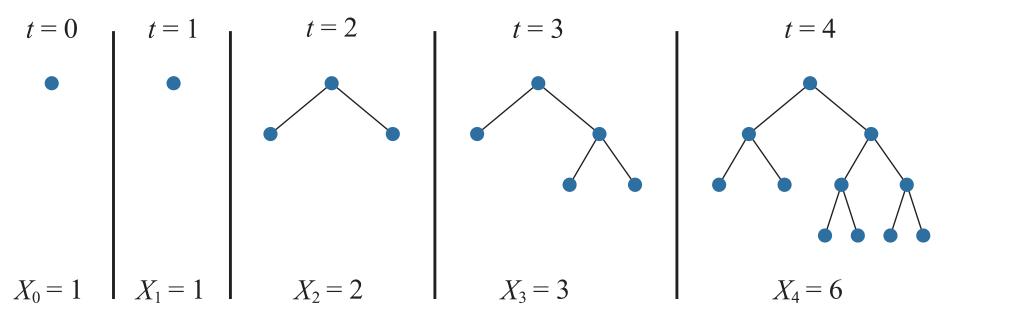
## Example: Epidemic growth modeling

At each time step, every leaf independently either:

- forks off 2 children, w.p.  $\frac{1}{2}$
- stays inert w.p.  $\frac{1}{2}$

 $X_t$  is number of leaves in tree after t steps.





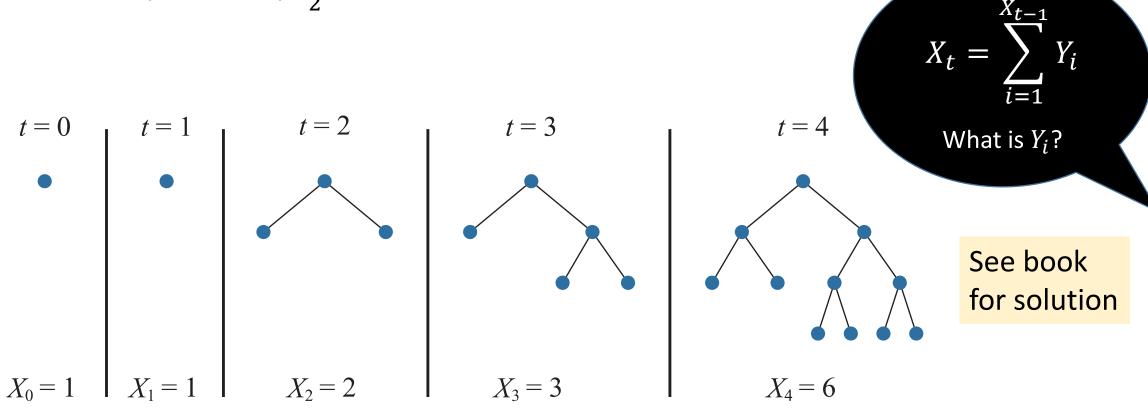
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Hint:



### Tail bounds

<u>Defn</u>: The **tail** of random variable X is  $P{X > x}$ .

Example: T denotes response time at a web service. Want to ensure the fraction of people with response time > 0.5s is not too high.

Want an **upper bound** on  $P{T > 0.5}$ . This is called a **tail bound**.

## Tail bounds

<u>Another Example</u>: *n* items are hashed into a table of size *n*. Assume each item ends up in a random bucket. Ideally, we have 1 item per bucket. What is the fraction of time that your search time > k? (i.e., what's the probability your bucket has > k items?)

Let 
$$N =$$
#items in bucket 1 How is  $N$  distributed?  $N \sim Binomial\left(n, \frac{1}{n}\right)$   
 $P\{N > k\} = \sum_{i=k+1}^{n} P\{N = i\} = \sum_{i=k+1}^{n} {n \choose i} \left(\frac{1}{n}\right)^{i} \left(1 - \frac{1}{n}\right)^{n-i}$  We don't know how to compute such bounds in general.

\

<u>Point</u>: We'll see that just knowing the mean and variance suffices for a tail **bound**. In some cases, the mean alone suffices (although this bound is quite weak).

## Markov's inequality

**Theorem:** (Markov's inequality) If r.v. X is non-negative, then  $\forall a > 0$ ,

$$\boldsymbol{P}\{X \ge a\} \le \frac{\boldsymbol{E}[X]}{a}$$

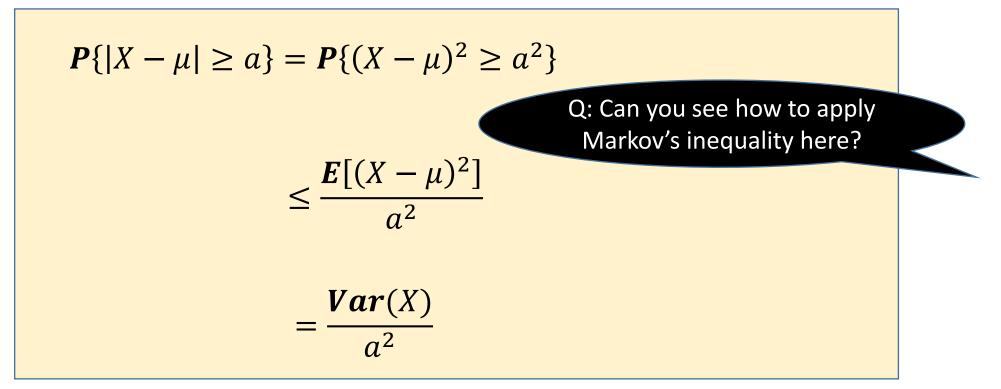
$$E[X] = \sum_{x=0}^{\infty} x \cdot p_X(x) \ge \sum_{x=a}^{\infty} x \cdot p_X(x)$$
$$\ge \sum_{x=a}^{\infty} a \cdot p_X(x)$$
$$= a \sum_{x=a}^{\infty} p_X(x) = a \cdot P\{X \ge a\}$$

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## Chebyshev's inequality

**Theorem:** (Chebyshev's inequality) Let *X* be any r.v. with finite mean,  $\mu$ , and finite variance. Then  $\forall a > 0$ ,

$$\mathbf{P}\{|X-\mu| \ge a\} \le \frac{\mathbf{Var}(X)}{a^2}$$



## Chebyshev's inequality

**Theorem:** (Chebyshev's inequality) Let *X* be any r.v. with finite mean,  $\mu$ , and finite variance. Then  $\forall a > 0$ ,

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#### Example:

$$N \sim Binomial\left(n, \frac{1}{n}\right)$$

Provide upper bound on:  $P{N \ge 6}$ 

$$P\{N \ge 6\} \le P\{|N-1| \ge 5\}$$
$$\le \frac{Var(N)}{25}$$
$$\le \frac{1}{25}$$

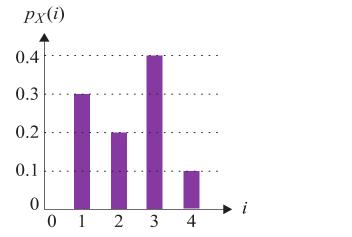
#### Stochastic dominance

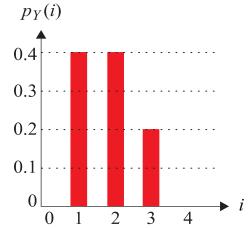
<u>Defn 5.18</u>: Given two random variables X and Y, if

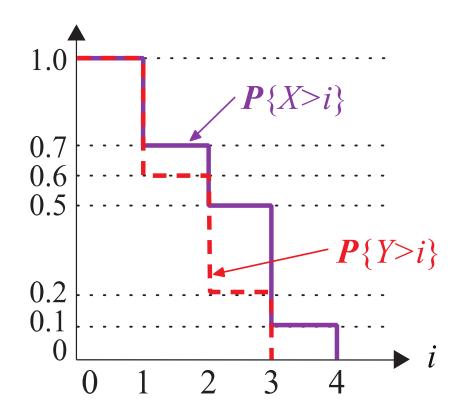
 $\boldsymbol{P}\{X > i\} \ge \boldsymbol{P}\{Y > i\}, \qquad \forall i$ 

we say that X stochastically dominates Y:

 $X \geq_{st} Y$ 





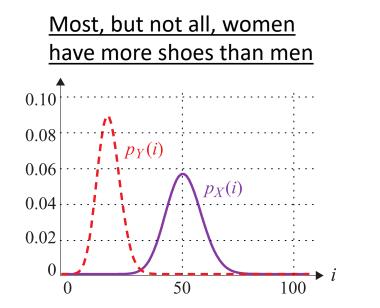


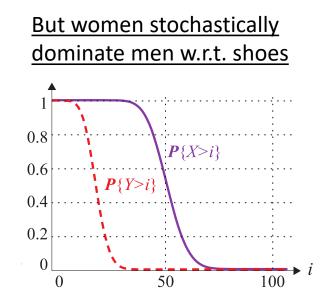
### Stochastic dominance

X = Number pairs of shoes owned by women ~  $Poisson(\lambda = 27)$ 



Y = Number pairs of shoes owned by men ~  $Poisson(\lambda = 12)$ 





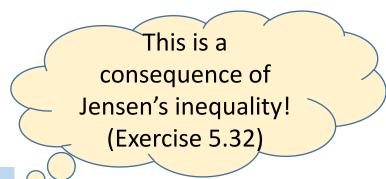
#### Jensen's inequality: motivation

We already know that

$$\boldsymbol{E}[X^2] \geq \boldsymbol{E}[X]^2$$

Is it also the case that

$$E[X^3] \ge E[X]^3$$
?  
 $E[X^4] \ge E[X]^4$ ?  
 $E[X^{4.5}] \ge E[X]^{4.5}$ ?



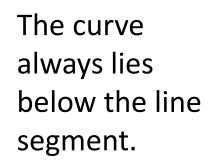
**Theorem:** Let *X* be any positive r.v. Then  $\forall a \in Reals$ ,

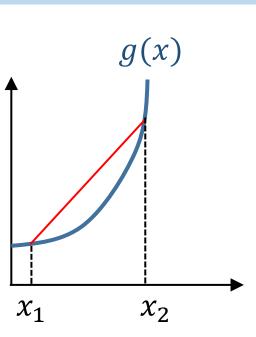
#### $\boldsymbol{E}[X^a] \geq \boldsymbol{E}[X]^a$

## Jensen's inequality

<u>Defn 5.21</u>: A function g(x) is **convex** on interval *S* if, for any  $x_1, x_2 \in S$ , and any  $\alpha \in [0,1]$ , we have:

$$g(\alpha x_1 + (1 - \alpha)x_2) \le \alpha g(x_1) + (1 - \alpha)g(x_2)$$





g(x) is convex on *S* iff  $g''(x) \ge 0$ ,  $\forall x \in S$ .

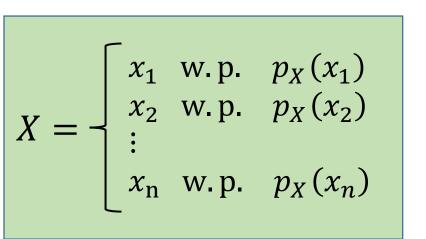
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## Jensen's inequality

<u>Defn 5.22</u>: A function g(x) is **convex** on interval *S* if, for any  $x_1, x_2, ..., x_n \in S$ , and any  $\alpha_1, \alpha_2, ..., \alpha_n \in [0,1]$ , where  $\sum_i \alpha_i = 1$ , we have:

 $g(\alpha_1 x_1 + \alpha_2 x_2 + \dots + \alpha_n x_n) \le \alpha_1 g(x_1) + \alpha_2 g(x_2) + \dots + \alpha_n g(x_n)$ 

 $g(p_X(x_1)x_1 + \dots + p_X(x_n)x_n) \le p_X(x_1)g(x_1) + \dots + p_X(x_n)g(x_n)$ 



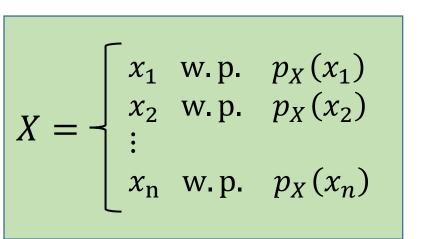
 $\implies g(\boldsymbol{E}[X]) \leq \boldsymbol{E}[g(X)]$ 

## Jensen's inequality

**Theorem 5.23:** (Jensen's inequality) If g(x) is convex on interval S and X takes on values on interval S, then:

 $g(\boldsymbol{E}[X]) \leq \boldsymbol{E}[g(X)]$ 

 $g(p_X(x_1)x_1 + \dots + p_X(x_n)x_n) \le p_X(x_1)g(x_1) + \dots + p_X(x_n)g(x_n)$ 

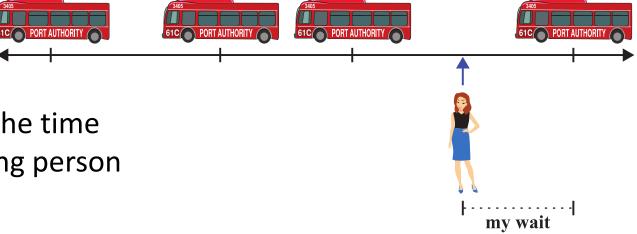


 $\implies g(\boldsymbol{E}[X]) \leq \boldsymbol{E}[g(X)]$ 

<u>Defn</u>: The **inspection paradox** says that, in high-variability settings, the mean seen by a random observer can be very different from the true mean.

Mean time between buses is 10 minutes.

However if there is some variability in the time between buses, then a randomly arriving person will wait more than 5 minutes.



Expected wait can even be >10 minutes!



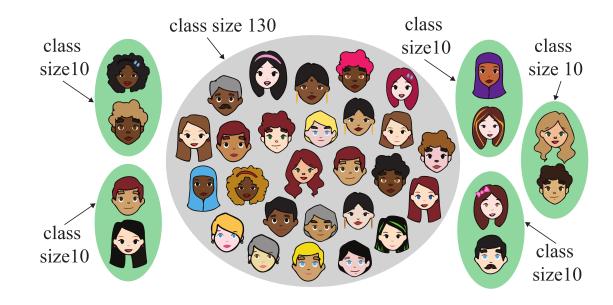
<u>Defn</u>: The **inspection paradox** says that, in high-variability settings, the mean seen by a random observer can be very different from the true mean.

Average class size reported by students is 100.

But the dean claims average class size is 30.

No one is lying.



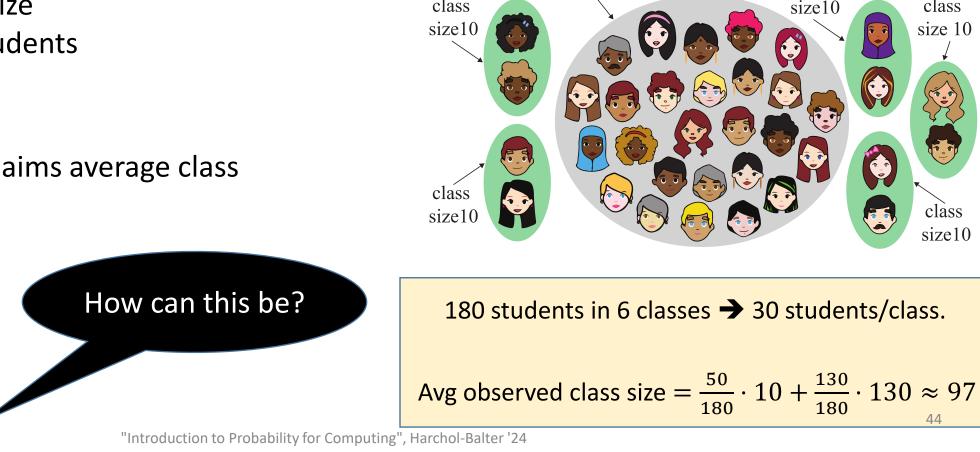


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class

class size 130

class

class

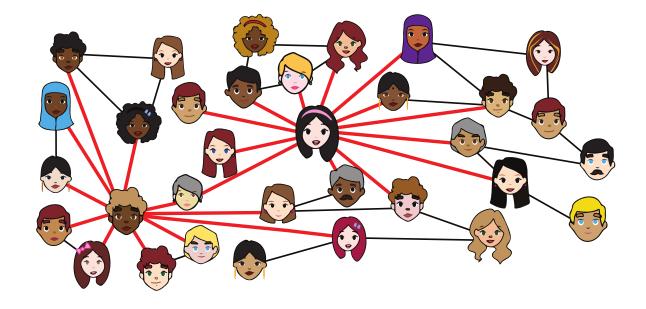
<u>Defn</u>: The **inspection paradox** says that, in high-variability settings, the mean seen by a random observer can be very different from the true mean.

The average Facebook user has 44 friends.

But the average friend of a Facebook user has 104 friends.

In fact, with probability 76%, your friend is more popular than you are.

#### How can this be?



Most people have few friends.

A few people are very popular with many friends. Which classification most likely describes you? Which most likely describes your friend?