CoMingle: Distributed Logic Programming for Decentralized Mobile Ensemble

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Outline

1 Introduction

2 Example

3 Semantics

4 Status

5 Conclusion & Future Work

Distributed Programming

- *Computations that run at more than one place at once*
  - A 40 year old paradigm
  - Now more popular than ever
    - Cloud computing
    - Modern webapps
    - ***Mobile device applications***
- Hard to get right
  - Concurrency bugs (race conditions, deadlocks, . . .)
  - Communication bugs
  - "Normal" bugs
- Two views
  - *Node-centric* — program each node separately
  - *System-centric* — program the distributed system as a whole
    - Compiled to node-centric code
    - Used in limited settings (Google Web Toolkit, MapReduce)

What is CoMingle?

A programming language for distributed mobile apps

- Declarative, concise, based on linear logic
- Enables high-level *system-centric* abstraction
  - **specifies** distributed computations as *ONE* declarative program
  - **compiles** into node-centric fragments, executed by each node
- Designed to implement mobile apps that run across Android devices
- Inspired by CHR [Frühwirth and Raiser, 2011], extended with
  - Decentralization [Lam and Cervesato, 2013]
  - Comprehension patterns [Lam and Cervesato, 2014]
- Also inspired by Linear Meld [Cruz et al., 2014]

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CoMingle by Example

```
module comingle.lib.ExtLib import {
     size   :: A -> int.
}

predicate swap    :: (loc,int) -> trigger.
predicate item    :: int -> fact.
predicate display :: (string,A) -> actuator.

rule pivotSwap :: [X]swap(Y,P),
                  {[X]item(D)|D->Xs. D >= P},
                  {[Y]item(D)|D->Ys. D <= P}
                    --o [X]display(Msg,size(Ys),Y), {[X]item(D)|D<-Ys},
                        [Y]display(Msg,size(Xs),X), {[Y]item(D)|D<-Xs}
                        where Msg = "Received %s items from %s".
```

CoMingle by Example: Decentralized Multiset Rewriting

```
[X]swap(Y,P)
{[X]item(D)|D->Xs.D>=P} --o [X]display(Msg,size(Ys),Y), {[X]item(D)|D<-Ys}
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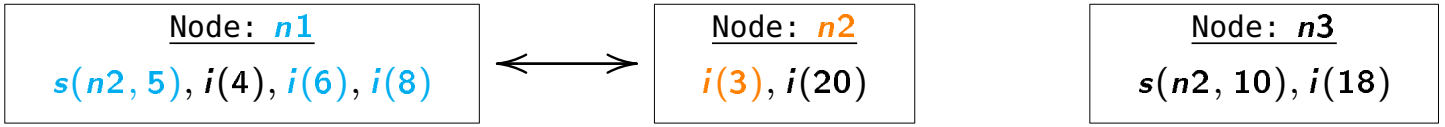
```

Let $s = \texttt{swap}$, $i = \texttt{item}$ and $d = \texttt{display}$



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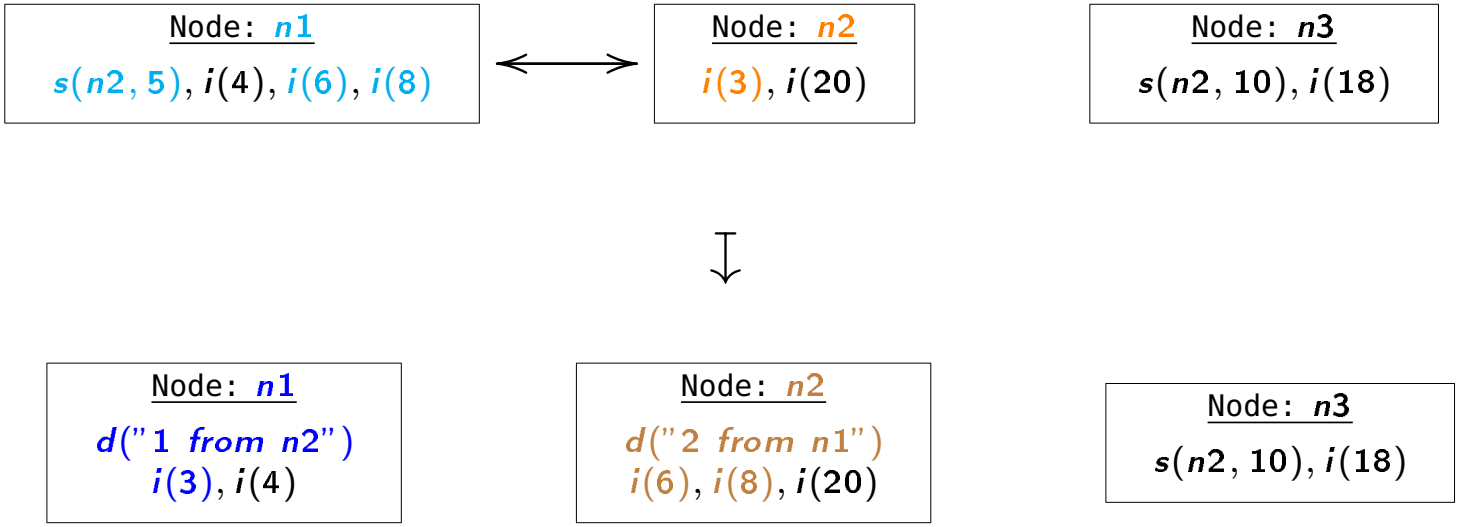

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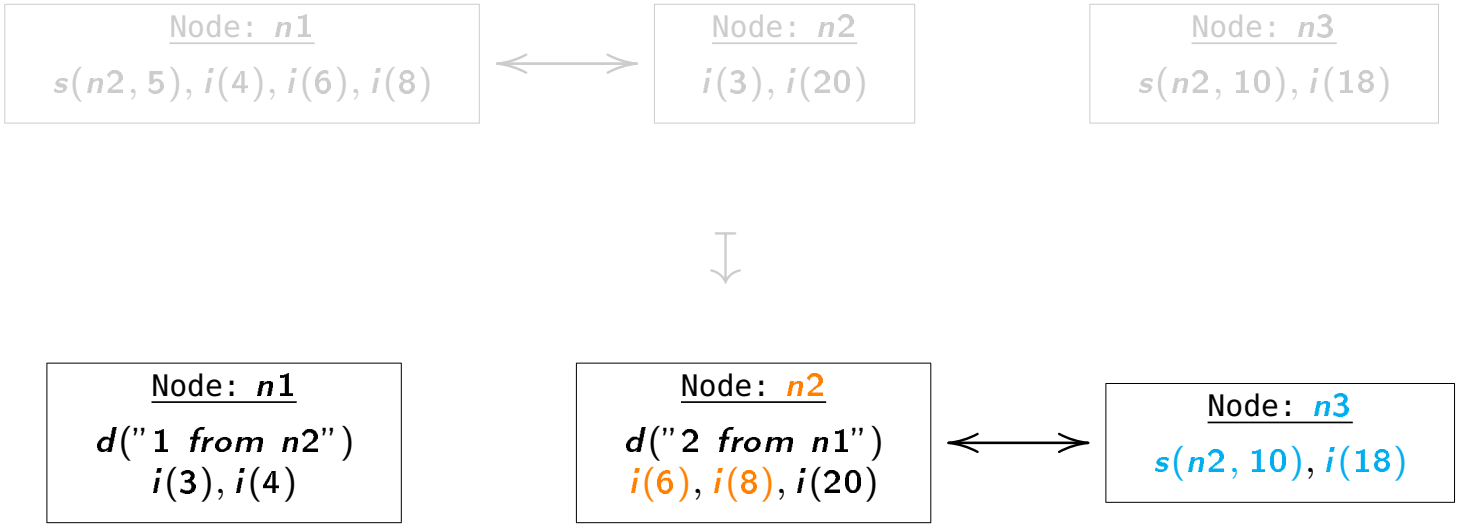

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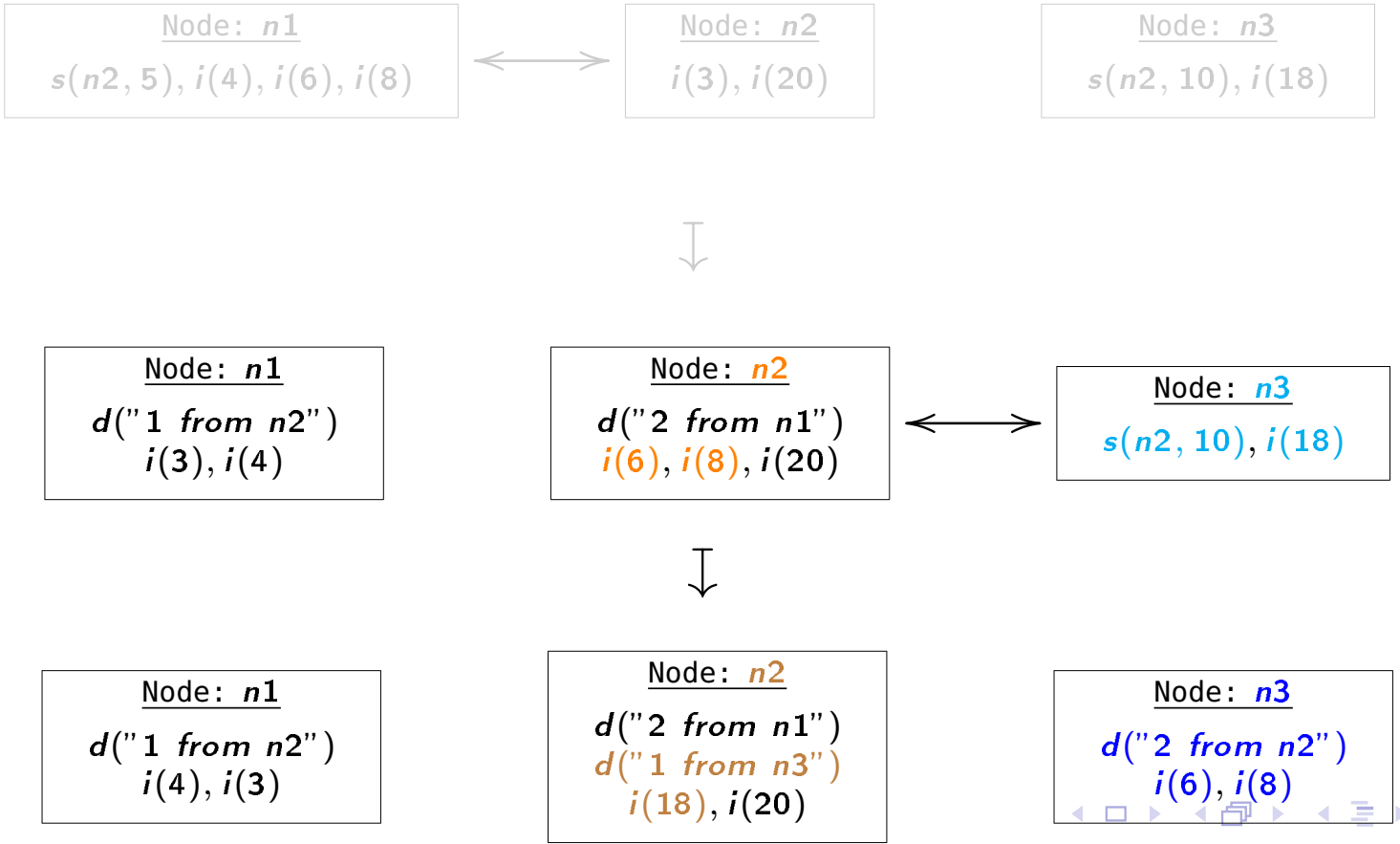
```

Let s = swap, i = item and d = display



CoMingle: Distributed Android Applications

- Traditional Distributed Graph processing problems:
  - Hyper-quicksort, Minimal Spanning Tree (GHS), Page Rank
  - Terminal/Quiescence states are the "results" (typically).
  - Emphasis on performance.
- CoMingle is targeted at *interactive mobile applications*:
  - Intermediate states are *observable* the "results".
  - Users can *interact* with the rewriting runtime.

Example Application: Drag Racing

- Inspired by Chrome Racer (`www.chrome.com/racer`)
- Race across a group of mobile devices
- Decentralized communication (over Wifi-Direct or LAN)

Implementing Drag Racing in CoMingle

```
rule init :: [I]initRace(Ls)
   --o {[A]next(B)|(A,B)<-Cs}, [E]last(),
       {[I]has(P), [P]all(Ps), [P]at(I), [P]rendTrack(Ls) | P<-Ps}
       where (Cs,E) = makeChain(I,Ls), Ps = list2mset(Ls).

rule start :: [X]all(Ps) \ [X]startRace() --o {[P]release()|P<-Ps}.

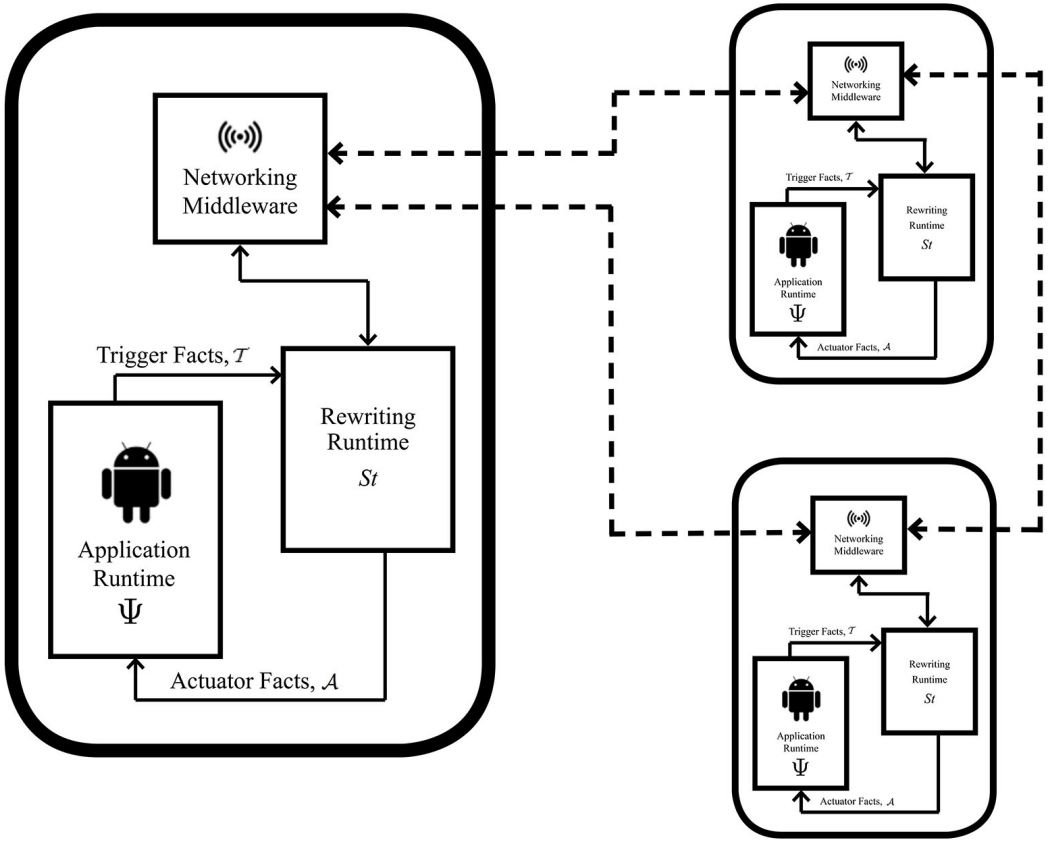
rule tap   :: [X]at(Y) \ [X]sendTap() --o [Y]recvTap(X).

rule trans :: [X]next(Z) \ [X]exiting(Y), [Y]at(X) --o [Z]has(Y), [Y]at(Z).

rule win   :: [X]last() \ [X]all(Ps), [X]exiting(Y) --o {[P]decWinner(Y) | P <- Ps}.
```

- + 862 lines of properly indented Java code
  - 700++ lines of local operations (e.g., display and UI operations)
  - $<$ 100 lines for initializing CoMingle runtime

CoMingle Architecture



CoMingle by Example: Triggers and Actuators

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                        [Y]display(Msg,size(Xs),X), {[Y]item(D)|D<-Xs}
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```

- **Abstracts** communications between node (i.e., X, Y)
- Executed by a **rewriting runtime** on each node
- Interacts with a local **application runtime** on each node
- Triggers: **inputs** from the application runtime
- Actuators: **outputs** into the application runtime

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                        [Y]display(Msg,size(Xs),X), {[Y]item(D)|D<-Xs}
                        where Msg = "Received %s items from %s".
```

- Predicate swap is a trigger
  - An input interface into the rewriting runtime
  - Only in rule heads
  - swap(Y,P) is added to rewriting state when button on device X is pressed

CoMingle by Example: Triggers and Actuators

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```

- Predicate `display` is an `actuator`
  - An output interface from the rewriting runtime
  - Only in rule body
  - `display("2 from n1")` executes a screen display callback function

CoMingle by Example: Triggers and Actuators

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```

- Predicate `item` is a standard `fact`
  - Can appear in rule head or body
  - Atoms of the rewriting state

CoMingle by Example

```
[X]swap(Y,P)
{[X]item(D)|D->Xs.D>=P} --o [X]display(Msg,size(Ys),Y), {[X]item(D)|D<-Ys}
{[Y]item(D)|D->Ys.D<=P}     [Y]display(Msg,size(Xs),X), {[Y]item(D)|D<-Xs}
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- High-level specification of distributed triggers/actuators

Distributed Triggers + *Distributed* State Patterns	--o	*Distributed* Actuators + *Distributed* State Patterns

- Declarative, concise and executable!
- Abstracts away
  - Low-level message passing
  - Synchronization
- Ensures atomicity and isolation

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Abstract Syntax

- A CoMingle program $\mathcal{P}$ is a set of rules of the form

$$r : H_p \setminus H_s \mid g \multimap B$$

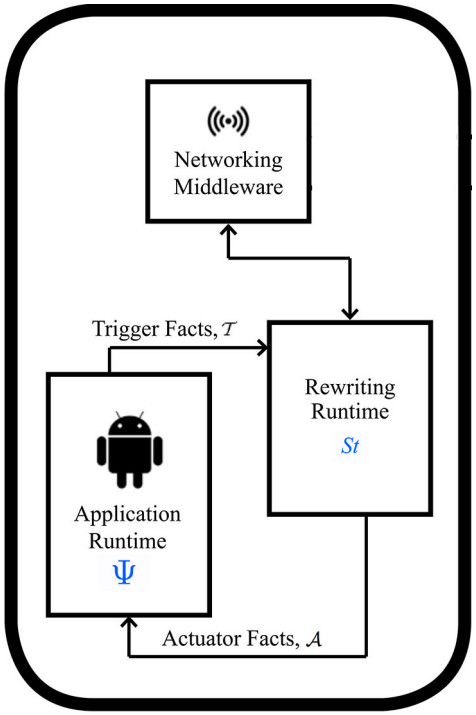
  - H_p, H_s and B: Multisets of patterns
  - g: Guard conditions

- A pattern is either
  - a fact: $[\ell]p(\vec{t})$
  - a comprehension: $\wr[\ell]p(\vec{t}) \mid g \int_{\vec{x} \in t}$

- Three kinds of facts
  - Triggers (only in H_p or H_s): Inputs from the "Android world"
  - Actuators (only in B): Outputs to the "Android world"
  - Standard facts: Atoms of rewriting state

Semantics of CoMingle: Abstract State Transitions

- *CoMingle state* $\langle St; \Psi \rangle$ represents the mobile ensemble
  - St is the *rewriting state*, a multiset of ground facts $[\ell]f$
  - Ψ is the *application state*, a set of local states $[\ell]\psi$
  - A location ℓ is a computing node



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  - St is the *rewriting state*, a multiset of ground facts $[\ell]f$
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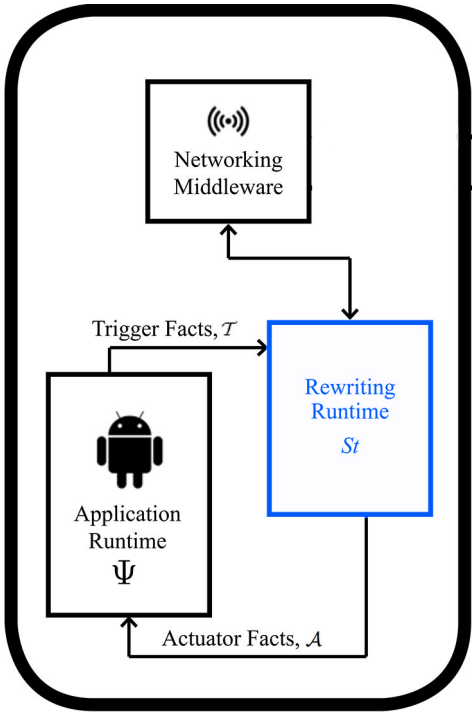
  - A location ℓ is a computing node
- The rewrite runtime: $\mathcal{P} \triangleright \langle St; \Psi \rangle \mapsto \langle St'; \Psi \rangle$
  - Applies a rule in $\mathcal{P}$
  - Several locations may participate
  - Decentralized multiset rewriting



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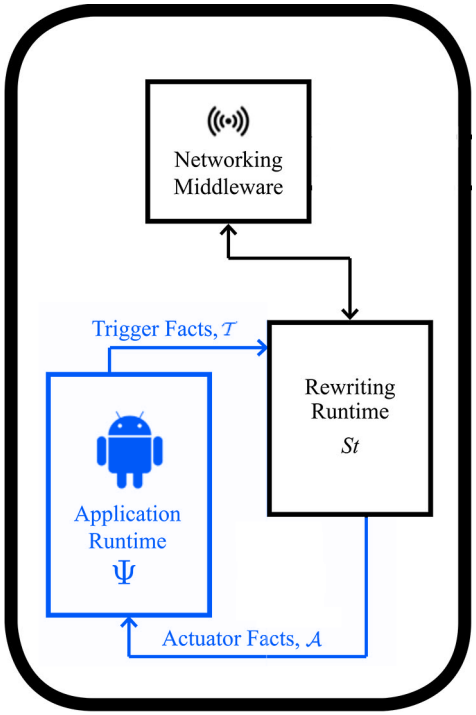
  - Applies a rule in $\mathcal{P}$
  - Several locations may participate
  - Decentralized multiset rewriting
- The application runtime: $\langle \mathcal{A}; \psi \rangle \mapsto_\ell \langle \mathcal{T}; \psi' \rangle$
  - Models local computation within a node
  - All within location ℓ



Rewriting Runtime: Rewriting Semantics

- Rewriting runtime transition: $\mathcal{P} \rhd \langle St; \Psi \rangle \mapsto \langle St'; \Psi \rangle$
  - Applies a rule in $\mathcal{P}$ to transform St into St'

$$\frac{(\overline{H}_p \setminus \overline{H}_s \mid g \multimap \overline{B}) \in \mathcal{P} \qquad \models \theta g \qquad \theta\overline{H}_p \triangleq^{\triangle}_{\mathbf{lhs}} St_p \quad \theta\overline{H}_s \triangleq^{\triangle}_{\mathbf{lhs}} St_s \quad \theta(\overline{H}_p, \overline{H}_s) \triangleq^{\neg}_{\mathbf{lhs}} St \quad \theta\overline{B} \ggg_{\mathbf{rhs}} St_b}{\mathcal{P} \rhd \langle St_p, St_s, St; \Psi \rangle \mapsto \langle St_p, St_b, St; \Psi \rangle}$$

- Matching: $\theta\overline{H}_p \triangleq^{\triangle}_{\mathbf{lhs}} St_p$ and $\theta\overline{H}_s \triangleq^{\triangle}_{\mathbf{lhs}} St_s$
- Maximality of Comprehensions: $\theta(\overline{H}_p, \overline{H}_s) \triangleq^{\neg}_{\mathbf{lhs}} St$
- Unfold: $\theta\overline{B} \ggg_{\mathbf{rhs}} St_b$
- See the paper/tech report, or ask me offline. =)

Application Runtime: Triggers and Actuators

- A local computation at location ℓ: $\langle \mathcal{A}; \psi \rangle \mapsto_\ell \langle \mathcal{T}; \psi' \rangle$
  - $\mathcal{A}$ is a set of actuator facts, introduced by the rewrite state St
  - $\mathcal{T}$ is a set of trigger facts, produced by the above local computation

$$\frac{\langle \mathcal{A}; \psi \rangle \mapsto_\ell \langle \mathcal{T}; \psi' \rangle}{\mathcal{P} \rhd \langle St, [\ell]\mathcal{A}; \Psi, [\ell]\psi \rangle \mapsto \langle St, [\ell]\mathcal{T}; \Psi, [\ell]\psi' \rangle}$$

- Entire computation must be happen at ℓ

Compilation of CoMingle Programs

System-centric specification
- High-level, concise
- Allows distributed events

```
rule pSwap :: [X]swap(Y,Z),
              {[X]item(I)|I->Is},
              {[Y]item(J)|J->Js},
              {[Z]item(K)|K->Ks}  --o [X]display(Msg,size(Js),Y), {[X]item(J)|J<-Js},
                                      [Y]display(Msg,size(Ks),Z), {[Y]item(K)|K<-Ks},
                                      [Z]display(Msg,size(Is),X), {[Z]item(I)|I<-Is}
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Choreographic Transformation ↓ [Lam and Cervesato, 2013]

Node-centric specification
- Match facts within a node
- Handles lower-level concurrency
  - Synchronization
  - Progress
  - Atomicity and Isolation

```
rule pSwapTest :: -{ [X]inTrans__(_) }, [X]swap(Y,Z), {[X]item(I)|I -> Is} \ 1
                  --o exists T__. [X]inTrans__(T__), [Y]pSwapProbeY(T__,Y,Is,Z), [Z]pSwapProbeZ(T__,Y,Is,Z).

rule pSwapProbeY :: { [Y]inTrans__(P__)|P__->Ps__ }, {[Y]item(J)|J -> Js}
                    \ [Y]pSwapProbeY(T__,Y,Is,Z) | strongest(T__,Ps__) --o [X]pSwapReadyY(T__,Y).


rule pSwapProbeZ :: { [Z]inTrans__(P__)|P__->Ps__ }, {[Z]item(K)|K -> Ks}
                    \ [Z]pSwapProbeZ(T__,Y,Is,Z) | strongest(T__,Ps__) --o [X]pSwapReadyZ(T__,Z).

rule pSwapEngage :: [X]inTrans__(T__) \ [X]pSwapReadyY(T__,Y), [X]pSwapReadyZ(T__,Z) --o [X]pSwapInit(T__,Y,Z).

rule pSwapInit :: [X]pSwapInit(T__,Y,Z), [X]itemLock(), [X]swap(Y,Z), {[X]item(I)|I -> Is}
                  --o [X]pSwapLHSX(T__,Y,Is,Z), [Y]pSwapReqY(T__,X,Is,Z), [Z]pSwapReqZ(T__,X,Y,Is).

rule pSwapReqYSucc :: [Y]pSwapReqY(T__,X,Is,Z), [Y]itemLock(), {[Y]item(J)|J -> Js}
                      --o [Y]pSwapLHSY(T__,Y,Js).


rule pSwapReqZSucc :: [Z]pSwapReqZ(T__,X,Y,Is), [Z]itemLock(), {[Z]item(K)|K -> Ks}
                      --o [Z]pSwapLHSZ(T__,Z,Ks).

rule pSwapCommit :: [X]inTrans__(T__), [X]pSwapLHSX(T__,Y,Is,Z), [X]pSwapLHSY(T__,Y,Js), [X]pSwapLHSZ(T__,Z,Ks)
                    --o [X]display(Msg,lvar0,Y), {[X]item(J)|J -> Js},
                        [Y]display(Msg,lvar1,Z), {[Y]item(K)|K -> Ks},
                        [Z]display(Msg,lvar2,X), {[Z]item(I)|I -> Is},
                        [Y]itemLock(), [X]itemLock(), [Z]itemLock()
                        Msg = "%s from %s", lvar0 = (sizeJs), lvar1 = (sizeKs), lvar2 = (sizeIs)

rule pSwapReqYFail :: [Y]pSwapReqY(T__,X,Is,Z) --o [X]pSwapAbort(T__).


rule pSwapReqZFail :: [Z]pSwapReqZ(T__,X,Y,Is) --o [X]pSwapAbort(T__).

rule pSwapAbortX :: [X]pSwapAbort(T__) \ [X]inTrans__(T__), [X]pSwapLHSX(T__,Y,Is,Z)
                    --o [X]itemLock(), [X]swap(Y,Z), {[X]item(I)|I -> Is}.

rule pSwapAbortY :: [X]pSwapAbort(T__) \ [X]pSwapLHSY(T__,Y,Js)
                    --o [Y]itemLock(), {[Y]item(J)|J -> Js}.


rule pSwapAbortZ :: [X]pSwapAbort(T__) \ [X]pSwapLHSZ(T__,Z,Ks)
                    --o [Z]itemLock(), {[Z]item(K)|K -> Ks}.
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rule pSwapReqYSucc :: [Y]pSwapReqY(T__,X,Is,Z), [Y]itemLock(), {[Y]item(J)|J -> Js}
                   --o [Y]pSwapLHSY(T__,Y,Js).

rule pSwapReqZSucc :: [Z]pSwapReqZ(T__,X,Y,Is), [Z]itemLock(), {[Z]item(K)|K -> Ks}
                   --o [Z]pSwapLHSZ(T__,Z,Ks).

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                 --o [X]display(Msg,lvar0,Y), {[X]item(J)|J -> Js},
                     [Y]display(Msg,lvar1,Z), {[Y]item(K)|K -> Ks},
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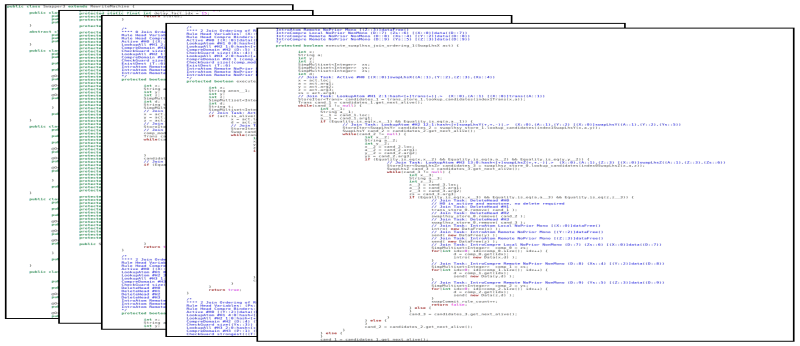
                 --o [X]itemLock(), [X]swap(Y,Z), {[X]item(I)|I -> Is}.

rule pSwapAbortY :: [X]pSwapAbort(T__) \ [X]pSwapLHSY(T__,Y,Js)
                 --o [Y]itemLock(), {[Y]item(J)|J -> Js}.

rule pSwapAbortZ :: [X]pSwapAbort(T__) \ [X]pSwapLHSZ(T__,Z,Ks)
                 --o [Z]itemLock(), {[Z]item(K)|K -> Ks}.
```

Imperative Compilation ↓ [Lam and Cervesato, 2014]

Low-level imperative compilation
- Java code
- Low-level network calls
- Operationalize multiset rewriting
- Trigger and actuator interfaces



Outline

1 Introduction

2 Example

3 Semantics

4 Status

5 Conclusion & Future Work

Implementation

- Prototype Available at

  https://github.com/sllam/comingle

- Networking via Wifi-Direct or Local-Area-Networks
- Proof-of-concept Apps
  - *Drag Racing* - Racing cars across mobile devices
  - *Battleships* - Traditional maritime war game, free-for-all style
  - *Wifi-Direct IP Directory* - Maintaining IP table for Wifi-Direct
  - *Musical Shares* - Bouncing a musical piece between mobile devices
  - *Swarbble* - Real-time team-based scrabble (In progress)
  - *Mafia* - The traditional party game, with a mobile twist (In progress)
- See tech.report [Lam and Cervesato, 2015] for details!

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Conclusion

- CoMingle: Distributed logic programming language
  - For programming distributed mobile applications
  - Based on decentralized multiset rewriting with comprehension patterns
- Prototype implementation
  - Available at `https://github.com/sllam/comingle`
  - Example apps available for download as well
  - Show your support, please STAR CoMingle GitHub repository!

Current and Future Work

- Front end refinements
  - Additional primitive types
  - More syntactic sugar
  - Refine Java interfaces
- Incremental extensions
  - Additional networking middlewares (Bluetooth, NFC, Wifi)
  - Orchresting Time Sensitive Events (e.g., Musical Shares, Mafia)
  - Sensor abstraction in CoMingle (e.g., GPS, speedometer)
  - More platforms (iOS, Raspberry Pi, Arduino, backend servers)
- Going beyond toy applications
  - Augmenting event/conference applications
  - Social interactive mobile applications

Questions?

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